

Proofs and Programs

TD 3 - The heart of Curry-Howard

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HW is due on 27 February 8am

Highlights

- Working on basics of type theory (typability - ex 1, inhabitation - ex 2).
- Understanding the C.-H. Correspondence between NJ(\Rightarrow) derivation trees, and λ_{\rightarrow} terms (ex 3,5);
- Taking advantage of that understanding, by dealing with other logical connectors (ex-4).

Let X, Y, \dots be *type variables* and x, y, \dots be *term variables*. Types (capital letter) and pure terms (small letters) are inductively defined by:

$$\begin{aligned} S, T, \dots &::= X \mid S \rightarrow T \\ a, b, \dots &::= x \mid \lambda x.a \mid ab \end{aligned}$$

The *simply typed λ -calculus* λ_{\rightarrow} is defined as the set of *typable* λ -terms in the following *type system*:

$$\frac{}{\Delta, x : A \vdash x : A} \text{ (Hyp)} \quad \frac{\Delta, x : A \vdash m : B}{\Delta \vdash \lambda x.m : A \rightarrow B} (\rightarrow_I) \quad \frac{\Delta \vdash f : A \rightarrow B \quad \Delta \vdash n : A}{\Delta \vdash fn : B} (\rightarrow_E)$$

where every variable appears at most once in the *context* (hence viewed as a set of typed variables). Simply typed λ -terms correspond to annotations of proof derivations in the implicative fragment of NJ, NJ(\Rightarrow). Then, for every proposition P :

$$P \text{ is provable in NJ}(\Rightarrow) \quad \text{iff} \quad P \text{ (viewed as a type) is inhabited in } \lambda_{\rightarrow}$$

Exercise 1 ($\vdash t : ?$). Provide the “most general types” that solve the following *type inference problems* or explain why it is not possible:

- | | |
|---|--|
| 1. $t \equiv \lambda f.\lambda g.g(f \mathbf{F})(f \mathbf{T})$; | $t \equiv \lambda f.\lambda g.g(f \mathbf{I})(f \mathbf{T})$; |
| 2. $t \equiv \lambda f.\lambda x.f(f x)$; | $t \equiv \lambda f.\lambda x.f(x x)$; |
| 3. HW $t \equiv \lambda x.x(\mathbf{II})$; | $t \equiv \mathbf{Y}$. |

Exercise 2 ($\vdash ? : A$). Recall that λ -terms in β -normal form are described by the following grammar:

$$n := a \mid \lambda x.n \quad a := x \mid a n$$

Solve the following *inhabitation problems* using terms in normal form or explain why it is not possible:

1. $A \equiv (S \rightarrow T) \rightarrow (T \rightarrow U) \rightarrow S \rightarrow U$;
2. $A \equiv (S \rightarrow T) \rightarrow ((S \rightarrow U) \rightarrow U) \rightarrow (T \rightarrow U) \rightarrow U$.
3. $A \equiv ((S \rightarrow T) \rightarrow S) \rightarrow S$ (Pierce’s law)

Exercise 3 (Avoiding detours). 1. Give a term associated to the following NJ(\Rightarrow) proof

$$\frac{\frac{\frac{}{\vdash P} \text{ (Hyp)}}{\vdash P} \text{ (Hyp)}}{\vdash P} \text{ (Hyp)} \quad \frac{\frac{\frac{\frac{}{P \Rightarrow P, Q \vdash P \Rightarrow P} \text{ (Hyp)}}{P \Rightarrow P \vdash Q \Rightarrow P \Rightarrow P} \text{ (Hyp)}}{\vdash (P \Rightarrow P) \Rightarrow Q \Rightarrow P \Rightarrow P} \text{ (Hyp)}}{Q \Rightarrow P \Rightarrow P} \text{ (Hyp)}}{\vdash P} \text{ (Hyp)} \text{ (Hyp)}$$

2. Reduce the term from question 1. and draw its type derivation tree. What is changed compare to the tree above?
3. Assuming the *subject reduction* and *strong normalization* properties of λ_{\rightarrow} (cf. Lecture notes: Lemma 4 and Theorem 6). Prove the following equivalences:

$$\begin{aligned}
 P \text{ is provable in NJ}(\Rightarrow) & \quad \text{iff } P \text{ (viewed as a type) is inhabited by a } \beta\text{-normal form of } \lambda_{\rightarrow} \\
 & \quad \text{iff } P \text{ is provable without (I/E)-detours}
 \end{aligned}$$

4. Deduce that Pierce's law is not derivable in NJ.

Exercise 4 (Product Type). In this exercise we are interested in extending the CH correspondence to $\text{NJ}(\Rightarrow, \wedge)$, the NJ fragment with implication and conjunction. Following the BHK interpretation, a witness for the *conjunction* $A \wedge B$ will correspond to a *pair* of witnesses for A and B .

1. Recall the encoding of **pair**, π_1 and π_2 *combinators* seen in TD1. Can we use them to encode *general* pairs and projection in λ_{\rightarrow} ?
2. Instead, we extend types with *products* (\times) and pure λ -terms with three new *constants* **pair**, π_1 , π_2 (hence making the calculus unpure...):

$$\begin{aligned}
 S, T, \dots & ::= X \mid S \rightarrow T \mid S \times T \\
 a, b, \dots & ::= x \mid \lambda x.a \mid ab \mid \mathbf{pair}(a, b) \mid \pi_i(a)
 \end{aligned}$$

Give typing rules for the new constants so that they annotate the intro-elim rules of \wedge in NJ.

3. Similarly to the *I/E*-detour of \Rightarrow in $\text{NJ}(\Rightarrow)$, explain what detours can be created by \wedge in proof derivations and how to eliminate them.
4. Deduce new reductions (β -rules) for the extended λ -calculus. We will denote $\lambda_{\rightarrow, \times}$ this new calculus.
5. Inhabit the following types of $\lambda_{\rightarrow, \times}$:
 - a) $(A \rightarrow B \rightarrow C) \rightarrow A \times B \rightarrow C$;
 - b) $(A \times B \rightarrow C) \rightarrow A \rightarrow B \rightarrow C$.
6. What is the grammar for β -normal terms in $\lambda_{\rightarrow, \times}$?
7. Let $\Delta \vdash p : A \times B$, what proof simplification in $\text{NJ}(\Rightarrow, \wedge)$ corresponds to the following reduction?

$$\mathbf{pair}(\pi_1 p)(\pi_2 p) \rightarrow_{\eta} p$$

Exercise 5 (Making a choice). **HW** Let A be a propositional variable and $P \equiv (A \Rightarrow A) \Rightarrow (A \Rightarrow A)$

1. Give a term corresponding to the following proof π_1 :

$$\overline{\vdash P}(\pi_1) = \frac{\frac{\overline{A \Rightarrow A, A \vdash A \Rightarrow A} \text{ (Hyp)} \quad \overline{A \Rightarrow A, A \vdash A} \text{ (Hyp)}}{A \Rightarrow A, A \vdash A} (\Rightarrow_E)}{A \Rightarrow A \vdash A \Rightarrow A} (\Rightarrow_I)}{\vdash P} (\Rightarrow_I)$$

2. Give a term corresponding to the following proof π_2 :

$$\overline{\vdash P}(\pi_2) = \frac{\frac{\overline{P, P \vdash P} \text{ (Hyp)}}{P \vdash P \Rightarrow P} (\Rightarrow_I)}{\vdash P \Rightarrow P \Rightarrow P} (\Rightarrow_I)}{\vdash P \Rightarrow P} (\Rightarrow_E)}{\vdash P} (\Rightarrow_I)}{\overline{\vdash P}(\pi_1)} (\Rightarrow_E)$$

3. Reduce the previous term to its β -normal form and give at each step of β reduction the corresponding proof simplification.
4. What would have happened if you had choose a different annotation/witness for the first hypothesis $\overline{P, P \vdash P} \text{ (Hyp)}$?