

Sparse matrix partitioning, ordering, and visualisation by Mondriaan 3.0

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Workshop Scheduling in Aussois, June 4, 2010

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Partitioning problems

Parallel sparse matrix–vector multiplication

Visualisation by MondriaanMovie

Hypergraphs

2D matrix partitioning

Ordering problems

Separated Block Diagonal structure

Parallel computing revolution

Conclusions and future work

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

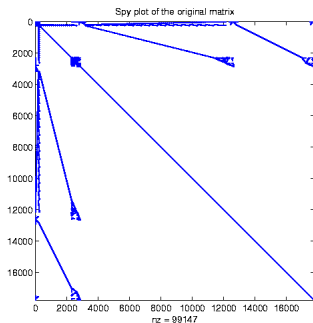
SBD

Revolution

Conclusions



Motivation: sparse matrix memplus



17758×17758 matrix with 126150 nonzeros.

Contributed to MatrixMarket in 1995 by Steve Hamm (Motorola). Represents the design of a [memory circuit](#).

Iterative solver multiplies matrix repeatedly with a vector



Motivation: high-performance computer



- ▶ National supercomputer Huygens named after Christiaan Huygens. Wikipédia: “En 1655, Huygens découvre Titan, la lune de Saturne. Il examina également les anneaux de Saturne et établit qu’il s’agissait bien d’un anneau entourant la planète”
- ▶ Huygens, the machine, has 104 nodes
- ▶ Each node has 16 processors
- ▶ Each processor has 2 cores and a shared L3 cache
- ▶ Each core has a local L1 and L2 cache

Outline

Partitioning

Matrix-vector
Movies
Hypergraphs
2D

Ordering

SBD

Revolution

Conclusions

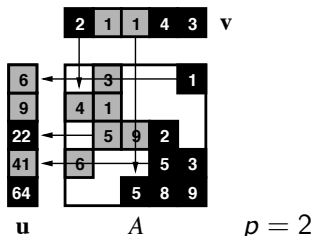


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Parallel sparse matrix–vector multiplication $\mathbf{u} := \mathbf{A}\mathbf{v}$

A sparse $m \times n$ matrix, \mathbf{u} dense m -vector, \mathbf{v} dense n -vector

$$u_i := \sum_{j=0}^{n-1} a_{ij}v_j$$



4 supersteps: **communicate**, compute, **communicate**, compute

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Divide evenly over 4 processors

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

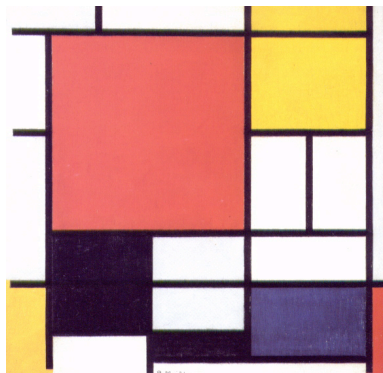
Revolution

Conclusions



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Composition with Red, Yellow, Blue and Black



Piet Mondriaan 1921

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

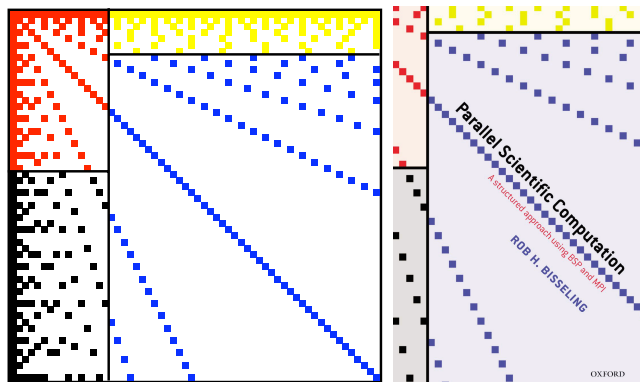
Revolution

Conclusions



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Matrix prime60



- ▶ Mondriaan block partitioning of 60×60 matrix prime60 with 462 nonzeros, for $p = 4$
- ▶ $a_{ij} \neq 0 \iff i|j$ or $j|i$ ($1 \leq i, j \leq 60$)

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Avoid communication completely, if you can

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

All nonzeros in a row or column have the same colour.



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Permute the matrix rows/columns

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

First the **green** rows/columns, then the **blue** ones.



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Combinatorial problem: sparse matrix partitioning

Problem: Split the set of nonzeros A of the matrix into p subsets, A_0, A_1, \dots, A_{p-1} , minimising the communication volume $V(A_0, A_1, \dots, A_{p-1})$ under the load imbalance constraint

$$\text{nz}(A_i) \leq \frac{\text{nz}(A)}{p}(1 + \epsilon), \quad 0 \leq i < p.$$

The **maximum** amount of work should not exceed the **average** amount by more than a fraction ϵ .

- ▶ $p = 2$ problem is already NP-complete (Lengauer 1990, circuit layout)

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

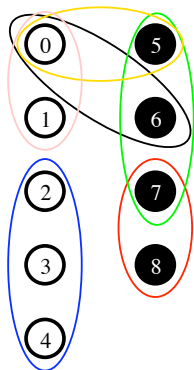
Revolution

Conclusions



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The hypergraph connection



Hypergraph with 9 vertices and 6 hyperedges (nets),
partitioned over 2 processors, black and white

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

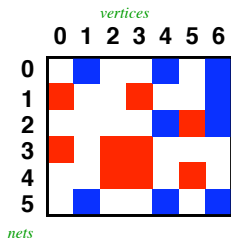
Revolution

Conclusions



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1D matrix partitioning using hypergraphs



- ▶ Hypergraph $\mathcal{H} = (\mathcal{V}, \mathcal{N}) \Rightarrow$ exact communication volume in sparse matrix–vector multiplication.
- ▶ Columns \equiv Vertices: 0, 1, 2, 3, 4, 5, 6.
Rows \equiv Hyperedges (nets, subsets of \mathcal{V}):

$$n_0 = \{1, 4, 6\}, \quad n_1 = \{0, 3, 6\}, \quad n_2 = \{4, 5, 6\},$$

$$n_3 = \{0, 2, 3\}, \quad n_4 = \{2, 3, 5\}, \quad n_5 = \{1, 4, 6\}.$$

- ▶ **Cut** nets n_1, n_2 cause 1 horizontal **communication**.



$(\lambda - 1)$ -metric for hypergraph partitioning

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

- ▶ 138×138 symmetric matrix bcsstk22, $nz = 696$, $p = 8$
- ▶ Reordered to **Bordered Block Diagonal** (BBD) form
- ▶ Split of row i over λ_i processors causes a communication volume of $\lambda_i - 1$ data words



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Cut-net metric for hypergraph partitioning

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

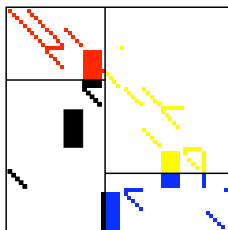
Conclusions

- ▶ Row split has **unit cost**, irrespective of λ_i



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Mondriaan 2D matrix partitioning



- ▶ Block partitioning (without row/column permutations) of 59×59 matrix `impcol_b` with 312 nonzeros, for $p = 4$
- ▶ Mondriaan package v1.0 (May 2002). Originally developed by Vastenhouw and Bisseling for partitioning term-by-document matrices for a parallel web search machine.

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Mondriaan 2D matrix partitioning

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

- ▶ $p = 4$, $\epsilon = 0.2$, global non-permuted view



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Fine-grain 2D matrix partitioning

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

- ▶ Each individual nonzero is a vertex in the hypergraph
Çatalyürek and Aykanat, 2001.



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Mondriaan 2.0, Released July 14, 2008



- ▶ New algorithms for **vector partitioning**.
- ▶ Much **faster**, by a factor of 10 compared to version 1.0.
- ▶ 10% better **quality** of the matrix partitioning.
- ▶ Inclusion of **fine-grain** partitioning method
- ▶ Inclusion of **hybrid** between original Mondriaan and fine-grain methods.
- ▶ Can also handle $p \neq 2^q$.

Outline

Partitioning

Matrix-vector
Movies
Hypergraphs
2D

Ordering

SBD

Revolution

Conclusions



Matrix 1ns3937 (Navier–Stokes, fluid flow)

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

Splitting the sparse matrix 1ns3937 into 5 parts.



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Recursive, adaptive bipartitioning algorithm

MatrixPartition(A, p, ϵ)

input: p = number of processors, $p = 2^q$

ϵ = allowed load imbalance, $\epsilon > 0$.

output: p -way partitioning of A with imbalance $\leq \epsilon$.

if $p > 1$ **then**

$q := \log_2 p$;

$(A_0^r, A_1^r) := h(A, \text{row}, \epsilon/q)$; **hypergraph splitting**

$(A_0^c, A_1^c) := h(A, \text{col}, \epsilon/q)$;

$(A_0^f, A_1^f) := h(A, \text{fine}, \epsilon/q)$;

$(A_0, A_1) := \text{best of } (A_0^r, A_1^r), (A_0^c, A_1^c), (A_0^f, A_1^f)$;

$\text{maxnz} := \frac{\text{nz}(A)}{p}(1 + \epsilon)$;

$\epsilon_0 := \frac{\text{maxnz}}{\text{nz}(A_0)} \cdot \frac{p}{2} - 1$; **MatrixPartition**($A_0, p/2, \epsilon_0$);

$\epsilon_1 := \frac{\text{maxnz}}{\text{nz}(A_1)} \cdot \frac{p}{2} - 1$; **MatrixPartition**($A_1, p/2, \epsilon_1$);

else output A ;

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Mondriaan version 1 vs. 3

Name	p	v1.0	v3.0
df1001	4	1484	1406
	16	3713	3640
	64	6224	6022
cre_b	4	1872	1491
	16	4698	4158
	64	9214	9095
tbdmatlab	4	10857	10060
	16	28041	24910
	64	52467	50020
nug30	4	55924	58770
	16	126255	137200
	64	212303	267200
tbdlinux	4	30667	30240
	16	73240	68890
	64	146771	140500

Outline

Partitioning

Matrix-vector
Movies
Hypergraphs
2D

Ordering

SBD

Revolution

Conclusions



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Mondriaan, default values (v1 localbest, v3 hybrid), $\epsilon = 0.03$

Mondriaan 3.0 coming this month



- ▶ Ordering to **SBD** and **BBD** structure: cut rows are placed in the middle, and at the end, respectively.
- ▶ Visualisation through **Matlab** interface, **MondriaanPlot**, and **MondriaanMovie**
- ▶ Metrics: $\lambda - 1$ for parallelism, and **cut-net** for other applications
- ▶ **Library-callable**, so you can link it to your own program
- ▶ Interface to **PaToH** hypergraph partitioner

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

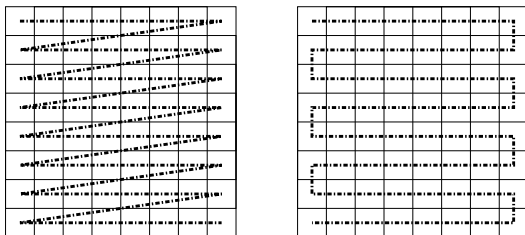
Revolution

Conclusions



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Ordering a sparse matrix to improve cache use



- ▶ Compressed Row Storage (CRS, left) and **zig-zag** CRS (right) orderings.
- ▶ Zig-zag CRS avoids unnecessary end-of-row jumps in cache, thus improving access to the input vector in a matrix–vector multiplication.
- ▶ Yzelman and Bisseling, *SIAM Journal on Scientific Computing* 2009.

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

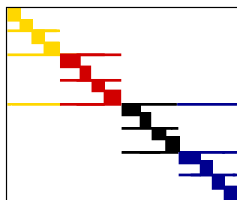
Revolution

Conclusions



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Separated block-diagonal (SBD) structure



- ▶ SBD structure is obtained by recursively partitioning the columns of a sparse matrix, each time moving the cut (mixed) rows to the middle. Columns are permuted accordingly.
- ▶ Mondriaan is used in one-dimensional mode, splitting only in the column direction.
- ▶ The cut rows are sparse and serve as a **gentle transition** between accesses to two different vector parts.

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Partition the columns till the end, $p = n = 59$

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions

- ▶ The recursive, fractal-like nature makes the ordering method work, irrespective of the actual cache characteristics (e.g. sizes of L1, L2, L3 cache).
- ▶ The ordering is **cache-oblivious**.



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Try to forget it all

- ▶ Ordering the matrix in SBD format makes the matrix-vector multiplication **cache-oblivious**. Forget about the exact cache hierarchy. It will always work.
- ▶ We also like to forget about the cores: **core-oblivious**. And then processor-oblivious, node-oblivious.
- ▶ All that is needed is a good ordering of the **rows** and **columns** of the matrix, and subsequently of its **nonzeros**.

Outline

Partitioning

Matrix-vector
Movies
Hypergraphs
2D

Ordering

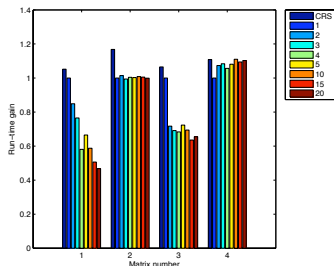
SBD

Revolution

Conclusions



Wall clock timings of SpMV on Huygens



Splitting into 1–20 parts

- ▶ Experiments on 1 core of the dual-core 4.7 GHz Power6+ processor of the Dutch national supercomputer Huygens.
- ▶ 64 kB L1 cache, 4 MB L2, 32 MB L3.
- ▶ Test matrices: 1. stanford; 2. stanford_berkeley; 3. wikipedia-20051105; 4. cage14

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

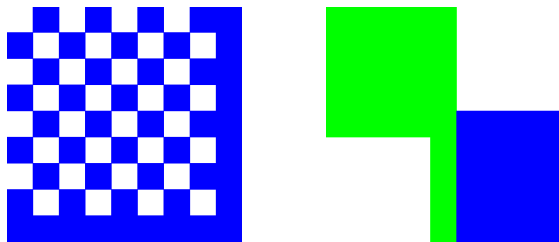
Revolution

Conclusions



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Doubly Separated Block-Diagonal structure



- ▶ 9×9 chess-arrowhead matrix, $nz = 49$, $p = 2$, $\epsilon = 0.2$.
- ▶ DSBD structure is obtained by recursively partitioning the sparse matrix, each time moving the cut rows and columns to the middle.
- ▶ The nonzeros must also be reordered by a [Z-like ordering](#).
- ▶ Mondriaan is used in two-dimensional mode.

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

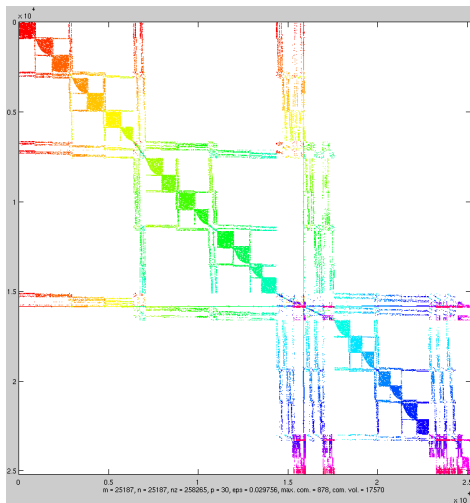
Revolution

Conclusions



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Screenshot of Matlab interface



► Matrix rhpentium, split over 30 processors

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Pictures of a revolution: the guillotine



King Louis XVI of France executed at the Place de la Concorde in Paris, January 23, 1793. Source:

http://www.solarnavigator.net/history/french_revolution.htm

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

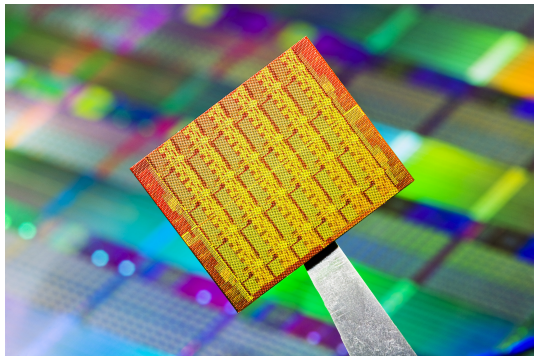
Revolution

Conclusions



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The parallel computing revolution



Intel Single-Chip Cloud computer with 48 cores, available Q3 2010. Energy consumption from 25 to 125 Watt, depending on use. Each pair of cores has a variable clock frequency. Source: <http://techresearch.intel.com>

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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Conclusions

- ▶ **Flop counts** become less and less important.
- ▶ It's all about **restricting movement**: moving less data, moving fewer electrons.
- ▶ We have presented two combinatorial problems: **partitioning** and **ordering**. Solution of these is an enabling technology for high-performance computing.
- ▶ **Reordering** is a promising method for oblivious computing. We have shown its utility in enhancing cache performance.
- ▶ Mondriaan 3.0, to be released soon, provides new reordering methods, based on hypergraph partitioning.
- ▶ **Visualisation can help in designing new algorithms!**

Outline

Partitioning

Matrix-vector

Movies

Hypergraphs

2D

Ordering

SBD

Revolution

Conclusions



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