### Cooperative Execution on Heterogeneous Multi-core Systems

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### COMMODITY COMPUTERS = HETEROGENEOUS SYSTEMS

- Multi-core General-Purpose Processors (CPUs)
- Many-core Graphic Processing Units (GPUs)
- ...
- Special accelerators, co-processors, FPGAs
- => HUGE COMPUTING POWER
- Not yet completely explored for **COLLABORATIVE COMPUTING**

### HETEROGENEITY MAKES PROBLEMS MUCH MORE COMPLEX!

- Performance modeling and load balancing
- Different programming models and languages





# Collaborative Environment for Heterogeneous Computers

### PERFORMANCE MODELING AND LOAD BALANCING

- State of the art for heterogeneous systems
- for CPU+GPU

CASE STUDY: 2D BATCH FAST FOURIER TRANSFORM

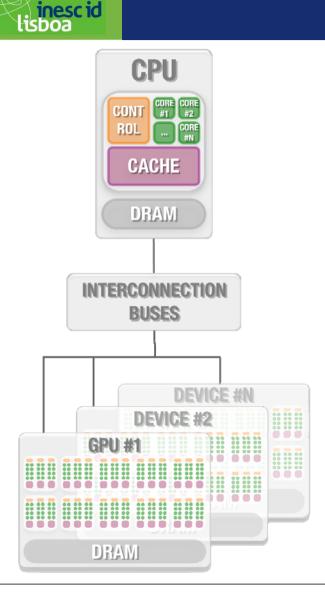
# CONCLUSIONS AND FUTURE WORK



### **Desktop Heterogeneous Systems**

### **MASTER-SLAVE** paradigm

- CPU (Master)
  - Global execution controller
  - Access the whole global memory
- INTERCONNECTION BUSSES
  - Limited and asymmetric communication bandwidth
  - Potential execution bottleneck
- UNDERLYING DEVICES (Slaves)
  - Different architectures and programming models
  - Computation performed using local memories



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### **Redefining Tasks and Primitive Jobs**

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### TASK – basic programming unit (coarser-grained)

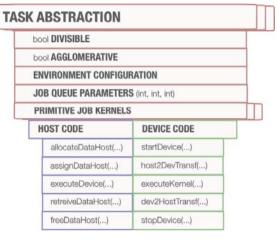
- CONFIGURATION PARAMETERS
  - Task: application and task dependency information
  - Environment: device type, number of devices...
- PRIMITIVE JOB WRAPPER
  - DIVISIBLE TASK comprise several finer-grained Primitive Jobs
  - AGGLOMERATIVE TASK allows grouping of Primitive Jobs

# PRIMITIVE JOB – minimal program portion for parallel execution

- CONFIGURATION PARAMETERS
  - I/O and performance specifics, ...
- CARRIES PER-DEVICE-TYPE IMPLEMENTATIONS
  - Vendor-specific programming models and tools
  - Specific optimization techniques

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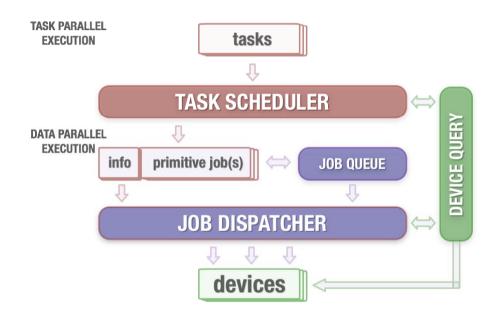
Primitive Job	Task Type	
Granularity	Divisible	Agglomerative
Coarser- grained	NO	
Balanced	YES	NO
Finer/Balanced	YES	YES

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Task Type		
Divisible	Agglomerative	
NO		
YES	NO	
YES	YES	

### Task Level Parallelism

TASK SCHEDULER submits independent tasks to JOB
 DISPATCHER in respect to task and environment
 configuration parameters and current platform state
 from DEVICE QUERY structure

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### Data Level Parallelism

- PRIMITIVE JOBS may be arranged into JOB QUEUES (currently, 1D-3D grid organization) for DIVISIBLE (AGGLOMERATIVE) TASKS
- JOB DISPATCHER USES DEVICE QUERY and JOB QUEUE information to map (agglomerated) PRIMITIVE JOBS to the requested devices; then initiates and controls further execution;

### **Nested Parallelism**

 If provided, JOB DISPATCHER can be configured to perceive certain number of cores of a multi-core device as a single device

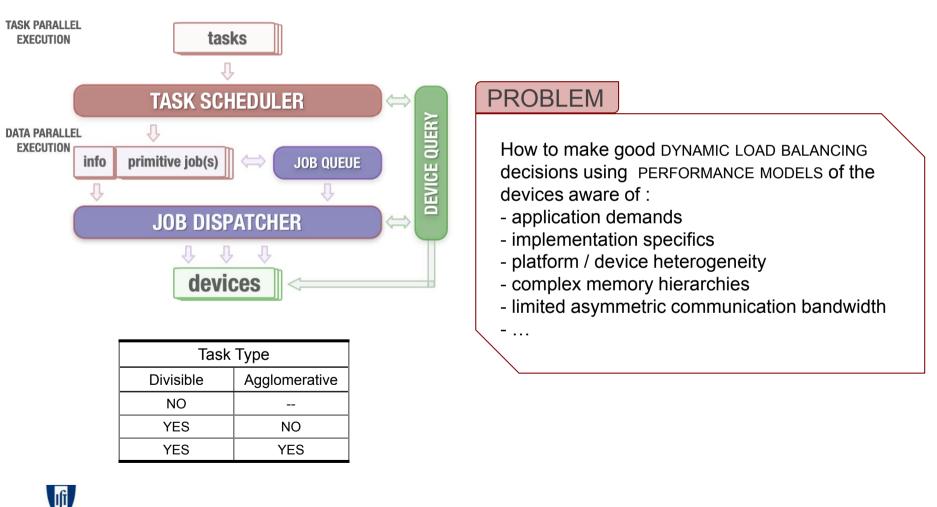


\*Aleksandar Ilic and Leonel Sousa. "Collaborative Execution Environment for Heterogeneous Parallel Systems", In 12th Workshop on Advances in Parallel and Distributed Computational Models (APDCM/IPDPS 2010), April 2010.

### Collaborative Execution Environment for Heterogeneous Systems\*



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# CONSTANT PERFORMANCE MODELS (CPM)

- DEVICE PERFORMANCE (SPEED) : constant positive number
  - Typically represents relative speed when executing a serial benchmark of a given size
- COMPUTATION DISTRIBUTION : proportional to the speed of device

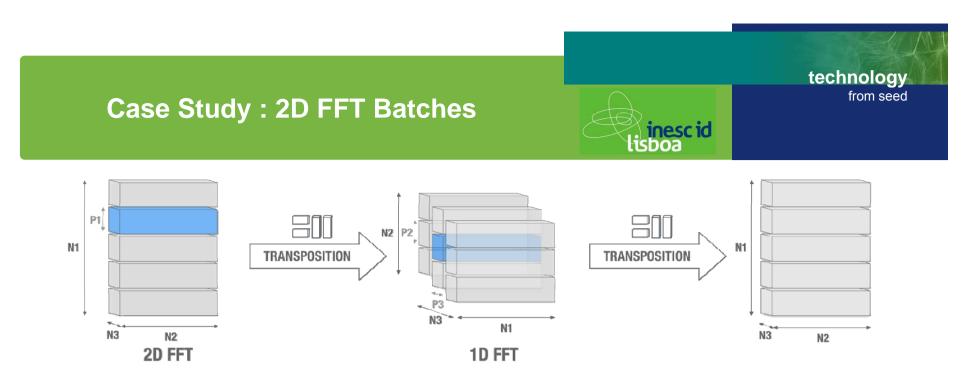
## FUNCTIONAL PERFORMANCE MODELS (FPM)

- DEVICE PERFORMANCE (SPEED) : continuous function of the problem size
  - Typically require several benchmark runs and significant amount of time for building it
- COMPUTATION DISTRIBUTION : relies on the "functional speed" of the processor

# FPM vs. CPM

- MORE REALISTIC : integrates features of heterogeneous processor
  - Processor heterogeneity, the heterogeneity of memory structure, and the other effects (such as paging)
- MORE ACCURATE DISTRIBUTION of computation across heterogeneous devices
- APPLICATION-CENTRIC approach characterize speed for different applications with different functions





Part of a PARALLEL 3D FFT PROCEDURE :  $H = FFT_{1D} (FFT_{2D} (h))$ 

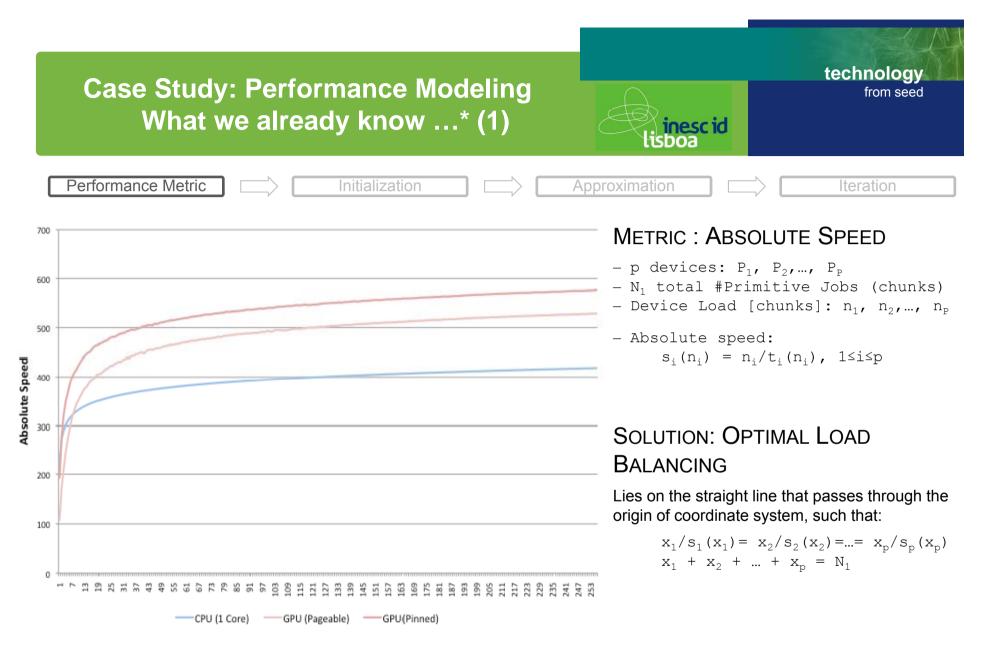
– Very HIGH COMMUNICATION-TO-COMPUTATION ratio

PROBLEM DEFINITION

 $- N_{in} = N_{out} = N_1 N_2 N_3 * sizeof(data)$ 

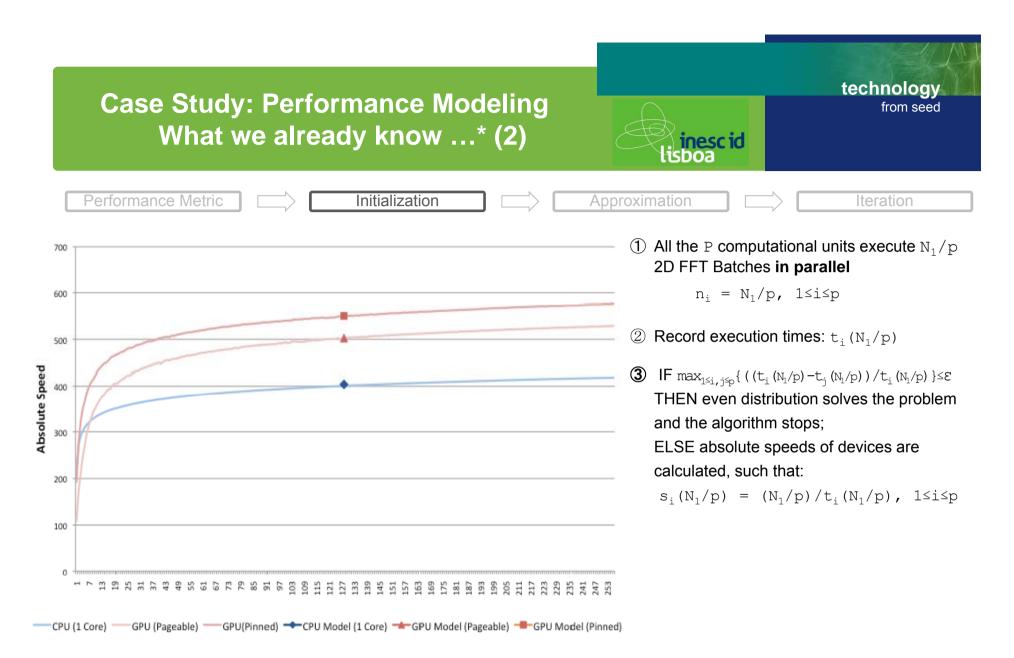
- Performance [FLOPS] :  $\mathrm{N_1N_2N_3}$  \* log ( $\mathrm{N_1N_2N_3}$ )
- $-N_2 = const; N_3 = const;$
- N<sub>1</sub> total number of computational chunks (PRIMITIVE JOBS)





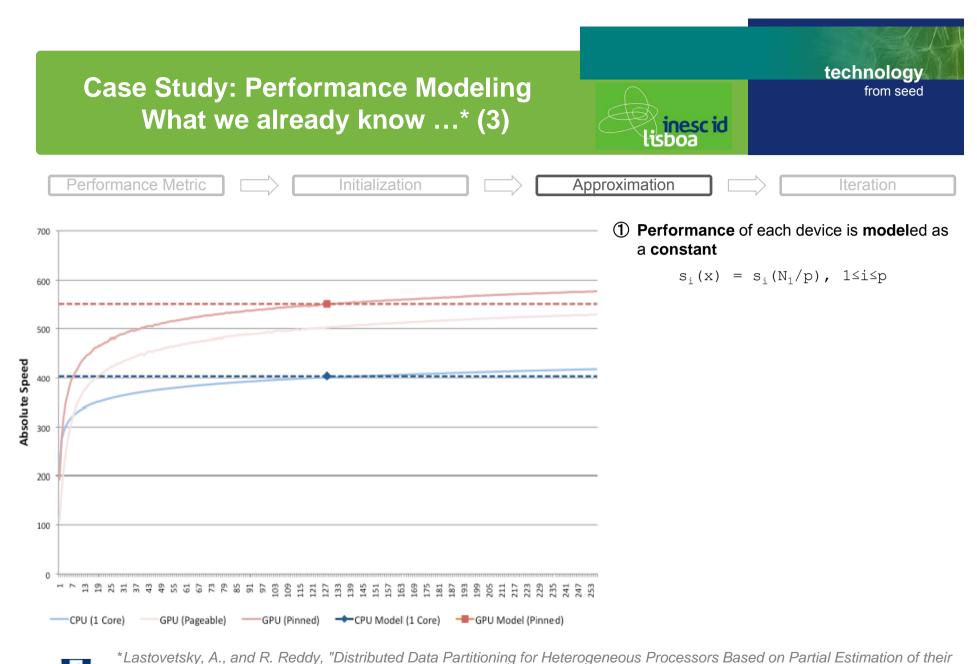


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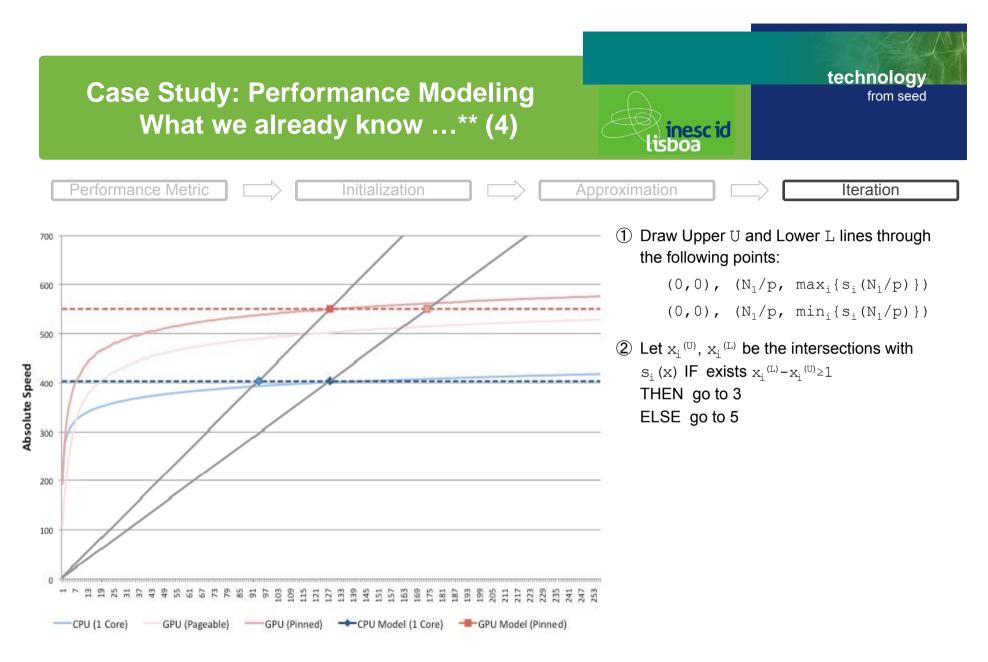
\*Lastovetsky, A., and R. Reddy, "Distributed Data Partitioning for Heterogeneous Processors Based on Partial Estimation of their Functional Performance Models", HeteroPar 2009, Netherlands, Lecture Notes in Computer Science, vol. 6043, Springer, pp. 91-101, 25/9/2009, 2010.

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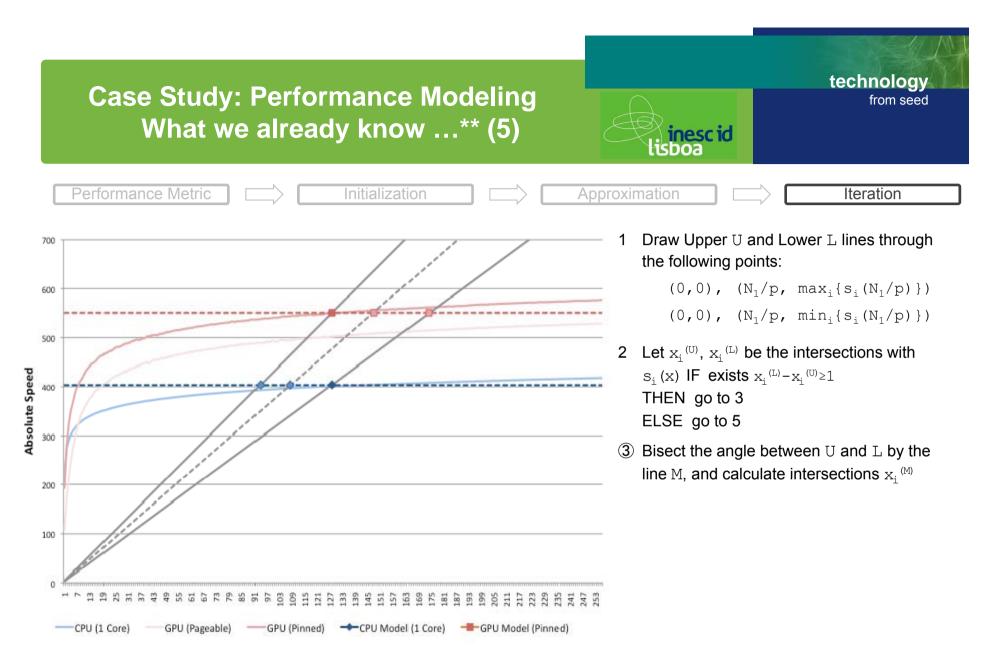


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Functional Performance Models", HeteroPar 2009, Netherlands, Lecture Notes in Computer Science, vol. 6043, Springer, pp. 91-101, 25/9/2009, 2010.

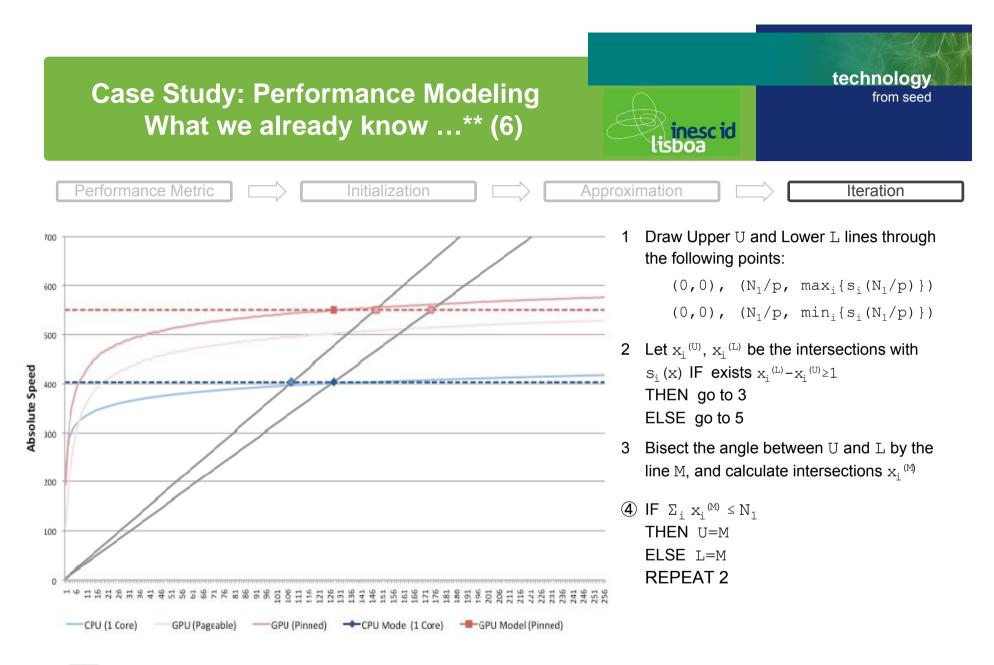


\*\*Lastovetsky, A., and R. Reddy, "Data Partitioning with a Functional Performance Model of Heterogeneous Processors", ıſi International Journal of High Performance Computing Applications, vol. 21, issue 1: Sage, pp. 76-90, 2007 INSTITUTO SUPERIOR TÉCNICO



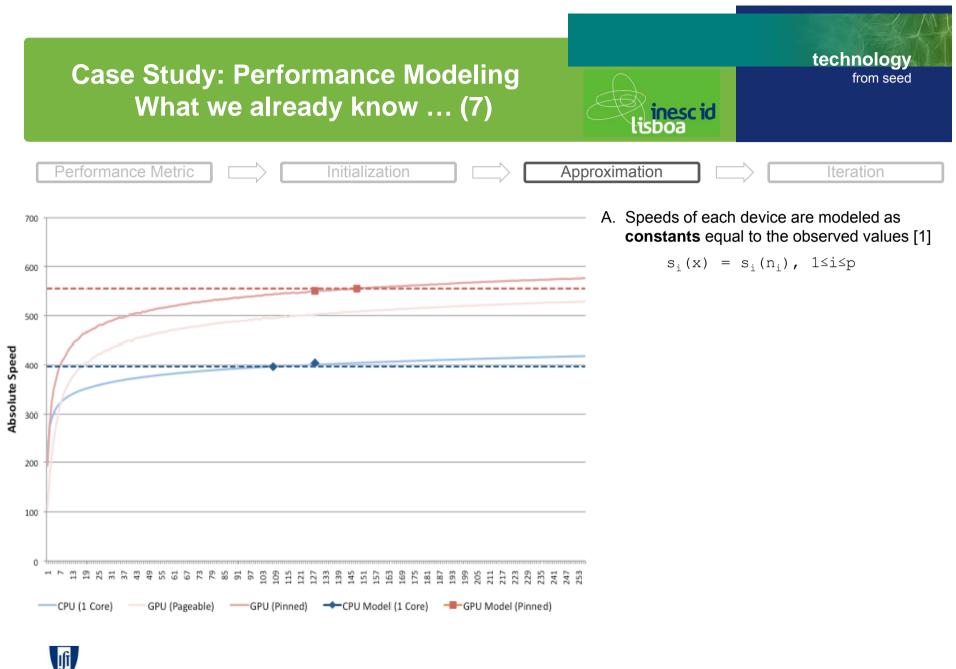
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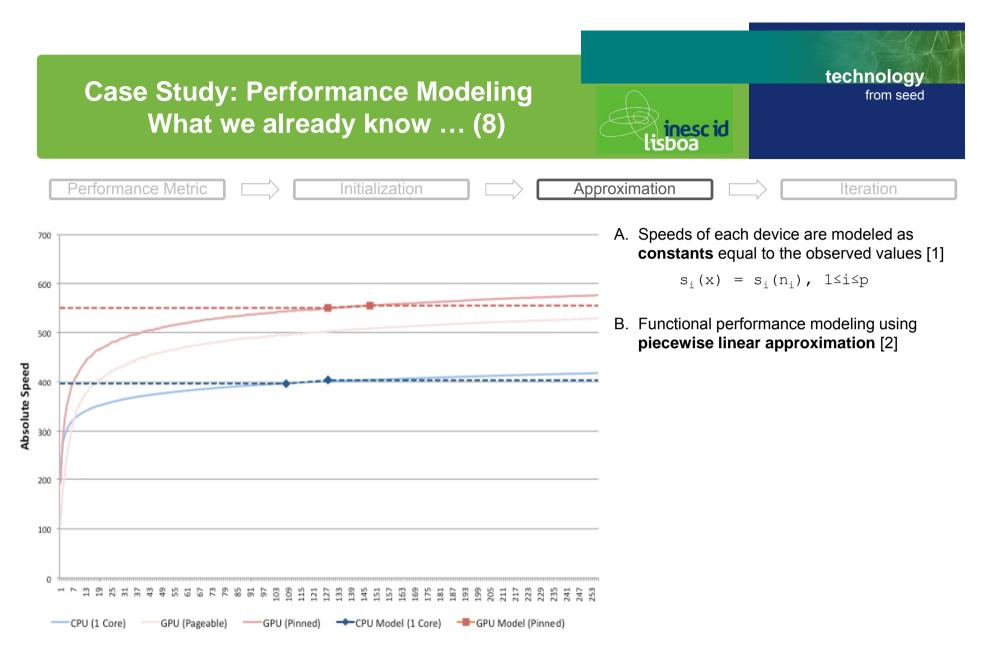


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INSTITUTO SUPERIOR TÉCNICO  Galindo I., Almeida F., and Badía-Contelles J.M., "Dynamic Load Balancing on Dedicated Heterogeneous Systems", In EuroPVM/MPI 2008, Springer, pp. 64-74, 2008.
 Lastovetsky, A., and R. Reddy, "Distributed Data Partitioning for Heterogeneous Processors Based on Partial Estimation of

their Functional Performance Models", HeteroPar 2009, vol. 6043, Springer, pp. 91-101, 25/9/2009, 2010



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### **ONLINE PERFORMANCE MODELING**

- PERFORMANCE ESTIMATION of all heterogeneous devices DURING THE EXECUTION
  - No prior knowledge on the performance of an application is available on any of the devices
  - Modeling of the overall CPU and GPU performance for different problem sizes (+ kernel-only GPU performance)

## DYNAMIC LOAD BALANCING

- OPTIMAL DISTRIBUTION OF COMPUTATIONS (PRIMITIVE JOBS)
  - Partial estimations of the performance should be built and used to decide on optimal mapping
  - Returned solution should provide load balancing within a given accuracy

### **COMMUNICATION AWARENESS**

- MODELING THE BANDWIDTH for interconnection busses DURING THE EXECUTION
  - To select problem sizes that maximize the interconnection bandwidth
  - The algorithm should be aware of asymmetric bandwidth for Host-To-Device and Device-To-Host transfers

# **CPU+GPU ARCHITECTURAL SPECIFICS**

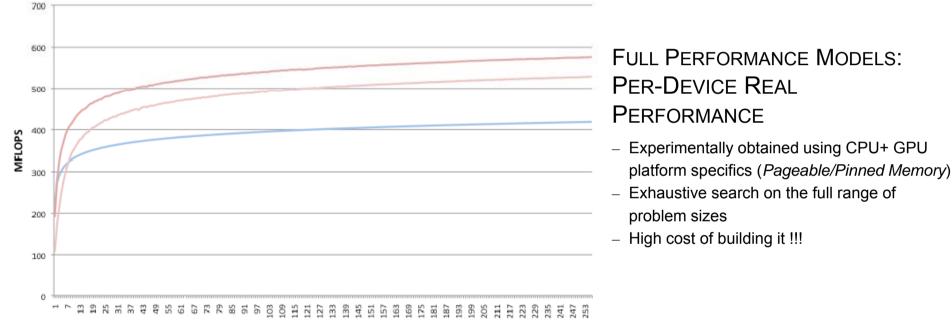
- Make use of ENVIRONMENT-SPECIFIC FUNCTIONS to ease performance modeling
  - Asynchronous transfers and CUDA streams to overlap communication with computation
  - Be aware of diverse capabilities of different devices, but also for devices of the same type (e.g. GT200 vs. Fermi)



### Case Study: Building Full Performance Models



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### Load (Nin = Nout)

— CPU (1 Core) — GPU (Pageable) — GPU (Pinned)

	CPU	GPU	
Experimental Setup	Intel Core 2 Quad	nVIDIA GeForce 285GTX	
Speed/Core (GHz)	2.83	1.476	
Global Memory (MB)	4096	1024	

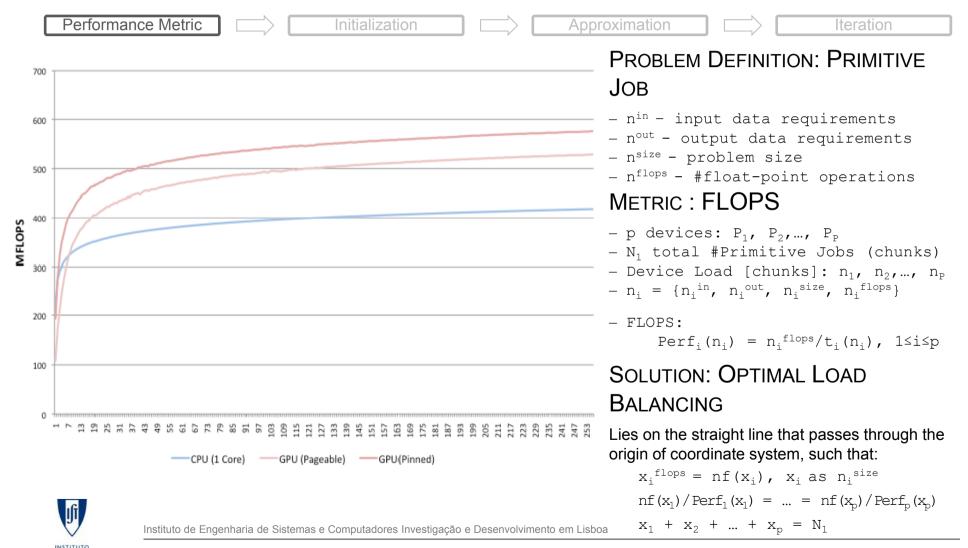
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2D FFT Batch	CPU GPU		
FFT Type	Complex; Double Precision		
Total Size (N <sub>1</sub> xN <sub>2</sub> xN <sub>3</sub> )	256x256x256		
High Performance Soft	Software		
FFT	Intel MKL 10.2 CUFFT 3.1		

### Case Study: Performance Modeling in CPU + GPU Environment (1)

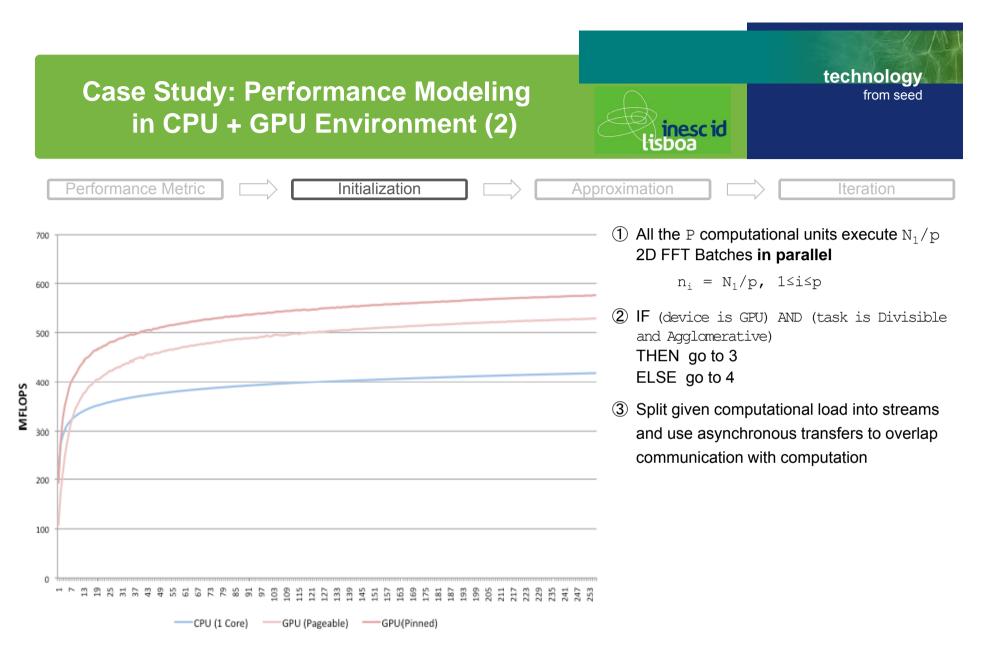


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# Case Study: Performance Modeling in CPU + GPU Environment (3) inescid initialization Approximation Iteration Streaming Approximation Iteration Streaming Approximation Iteration Iteration Iteration Approximation Iteration Iteration Iteration Iteration Iteration Iteration

- SUBDIVIDE  $n_i$  computational chunks using DIV2 STRATEGY
  - No prior knowledge on the performance of an application!
  - The next stream has half the load of a previous stream
  - Algorithm may continue splitting the workload until the last stream is assigned with load equal to 1
- BANDWIDTH-AWARE DIV2 STRATEGY
  - Interconnection bandwidth is a subject to the amount of data that should be transferred and not to the application-specific demands
  - Run small pre-calibration tests for HOST-TO-DEVICE AND DEVICE-TO-HOST transfers
  - Tests can be stopped when saturation points are detected, or when transfers reach desired value (e.g. 60% of its theoretical)
  - CASE STUDY: n<sup>min\_size</sup> = 4



1 All the P computational units execute  $N_1/P$ 2D FFT Batches in parallel

 $n_i = N_1/p$ ,  $1 \le i \le p$ 

- 2 IF (device is GPU) AND (task is Divisible and Agglomerative) THEN go to 3 ELSE go to 4
- ③ Split given computational load into streams and use asynchronous transfers to overlap communication with computation



### Case Study: Performance Modeling in CPU + GPU Environment (4)

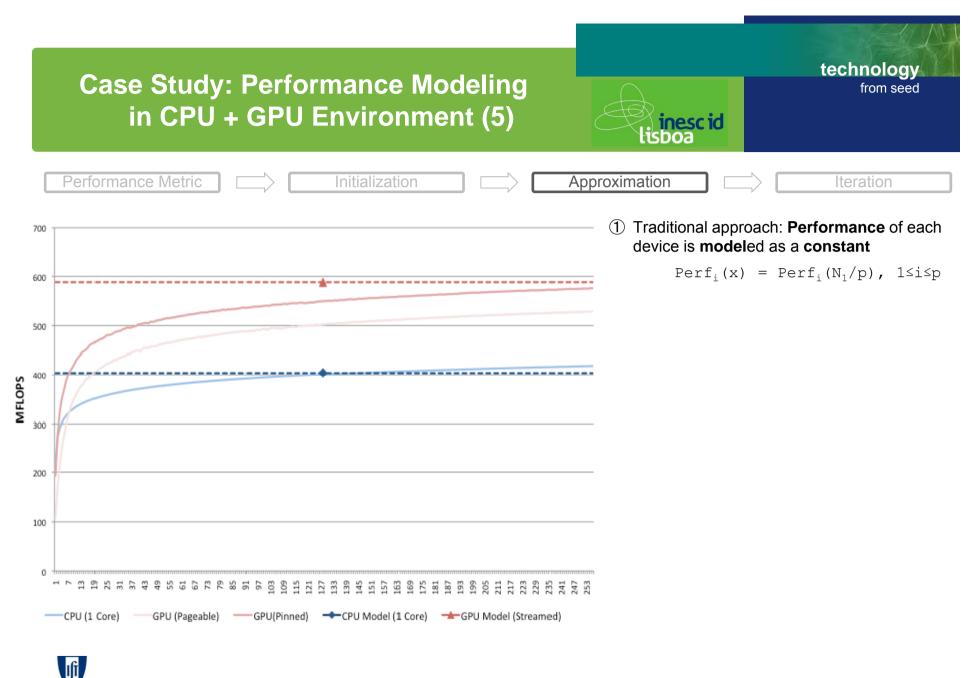
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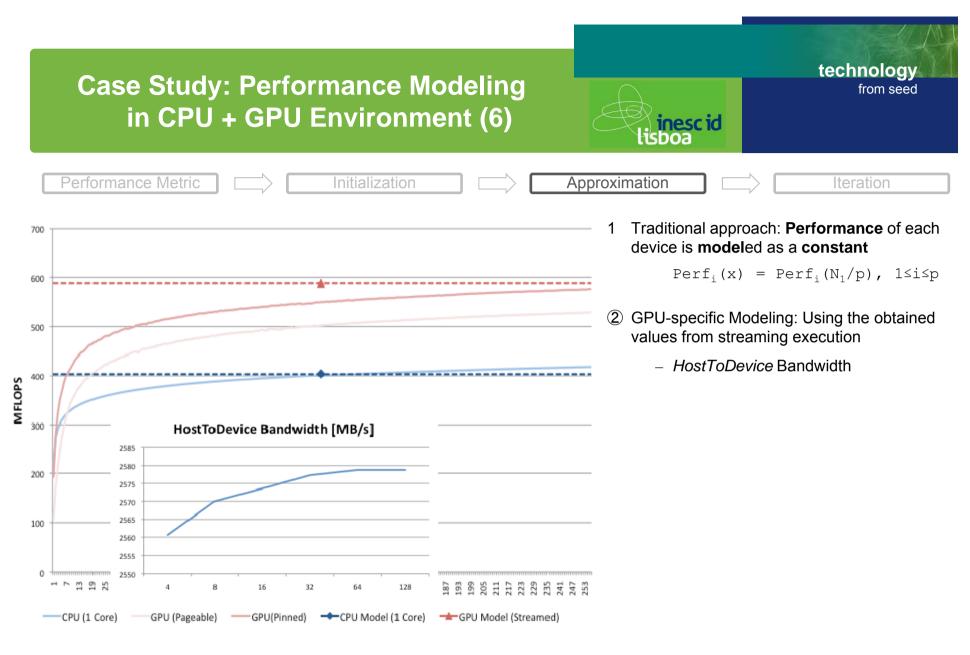
	Performance Metric App	roximation
70		1 All the P computational units execute $N_1/p$ 2D FFT Batches in parallel
60		n <sub>i</sub> = N <sub>1</sub> /p, 1≤i≤p
50 <b>Sd</b>		2 IF (device is GPU) AND (task is Divisible and Agglomerative) THEN go to 3 ELSE go to 4
SAOTJJW 30		3 Split given computational load into streams and use asynchronous transfers to overlap communication with computation
		(4) Execute & record execution times: $t_i (N_1/p)$
10		(5) IF $\max_{1 \le i, j \le p} \{ ((t_i(N_1/p) - t_j(N_1/p))/t_i(N_1/p)) \le t $ THEN even distribution solves the problem
		and the algorithm stops; ELSE performance of devices is calculated, such that: $Perf_i (N_1/p) = nf (N_1/p) / t_i (N_1/p), 1 \le i \le p$

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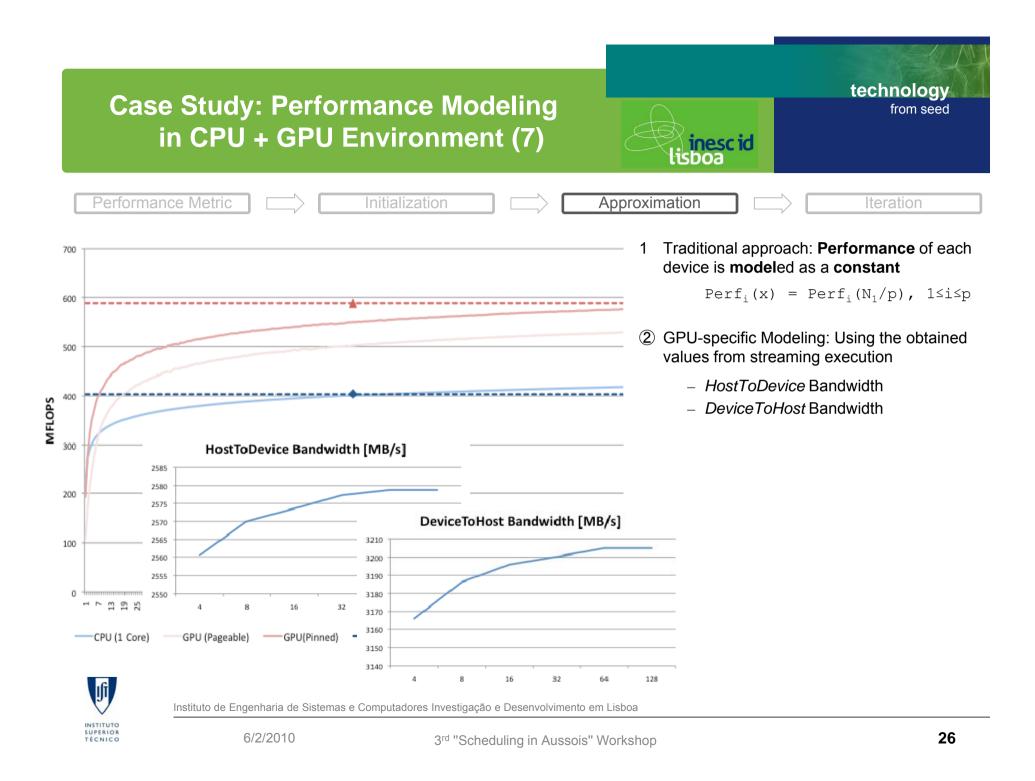


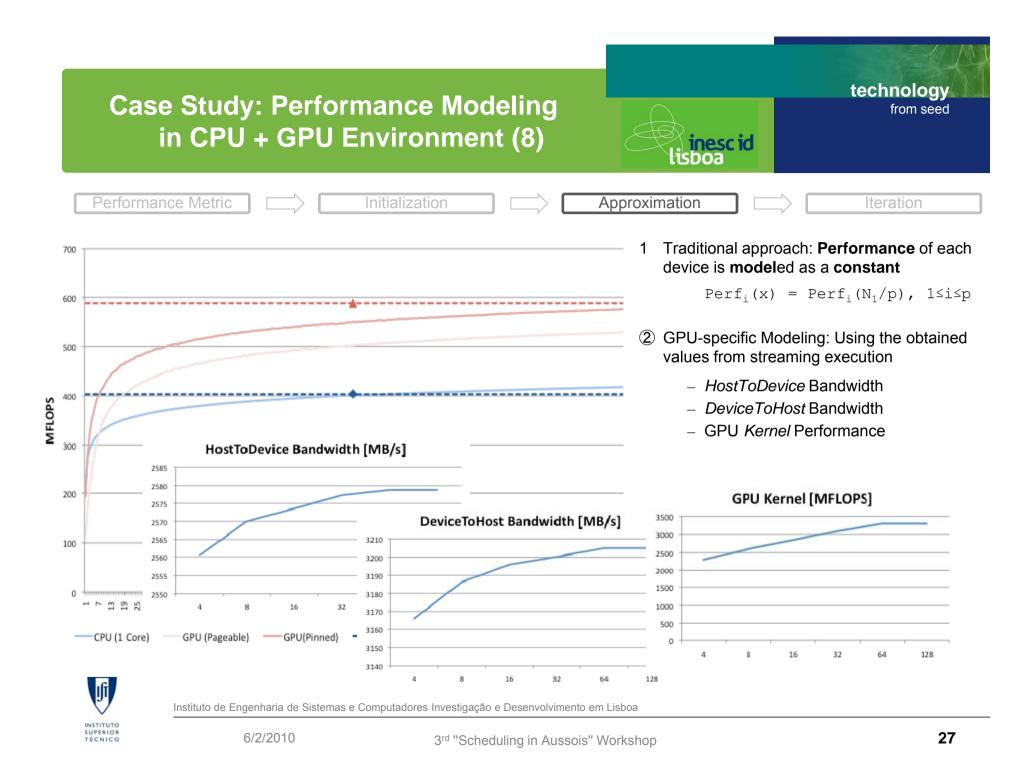
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### technology **Case Study: Performance Modeling** from seed in CPU + GPU Environment (9) linescid **Performance Metric** Approximation Initialization Iteration Traditional approach: Performance of each 700 device is modeled as a constant $Perf_{i}(x) = Perf_{i}(N_{1}/p), 1 \le i \le p$ 600 GPU-specific Modeling: Using the obtained 2 500 values from streaming execution $\Delta$ - HostToDevice Bandwidth 400 MFLOPS DeviceToHost Bandwidth \_ 300 GPU Kernel Performance \_ HostToDevice Bandwidth [MB/s] ③ Incorporate streaming results 2585 200 2580 **GPU Kernel [MFLOPS]** 2575 100 3500 DeviceToHost Bandwidth [MB/s] 2570 3000 2565 3210 2500 2560 3200 0 2000 6 52 2555 3190 1500 2550 3180 32 1000 16 3170 CPU (1 Core) G 500 3160 0 3150 16 32 64 128 3140 л 8 16 32 64 128 ıſī Instituto de Engenharia de Sistemas e Computadores Investigação e Desenvolvimento em Lisboa INSTITUTO SUPERIOR

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### technology **Case Study: Performance Modeling** from seed in CPU + GPU Environment (10) linescid **Performance Metric** Approximation Initialization Iteration Traditional approach: Performance of each 700 device is modeled as a constant ---- $Perf_{i}(x) = Perf_{i}(N_{1}/p), 1 \le i \le p$ 600 GPU-specific Modeling: Using the obtained 2 500 values from streaming execution - HostToDevice Bandwidth 400 MFLOPS - DeviceToHost Bandwidth 300 - GPU Kernel Performance ③ Incorporate streaming results 200 100 0 £ 6 5 5 5 2 2 88 5 21 27 23 8 8 22 5 5 CPU (1 Core) GPU (Pageable) — GPU(Pinned) + CPU Model (1 Core) - GPU Model GPU Model (Streamed) [Streamed Values]



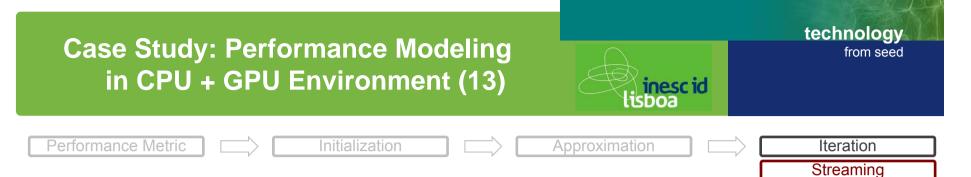
### technology **Case Study: Performance Modeling** from seed in CPU + GPU Environment (11) linescid **Performance Metric** Approximation Iteration ① Draw Upper U and Lower L lines through 700 the following points: 600 (0,0), $(N_1/p, max_i \{ Perf_i (N_1/p) \} )$ (0,0), $(N_1/p, min_i \{Perf_i (N_1/p)\})$ 500 (2) Let $x_i^{(U)}$ and $x_i^{(L)}$ be the intersections with $Perf_{i}(x)$ MFLOPS 400 IF exists $x_i^{(L)} - x_i^{(U)} \ge 1$ 300 THEN go to 3 ELSE go to 5 200 ③ Bisect the angle between U and L by the line M, and calculate intersections $x_i^{(M)}$ 100 (4) IF $\Sigma_i X_i^{(M)} \leq N_1$ 0 THEN U=M 211 217 223 223 223 223 223 223 2241 2241 2247 253 6 5 ELSE L=M **REPEAT 2** CPU (1 Core) GPU (Pageable) — GPU(Pinned) - CPU Model (1 Core) - GPU Model [Streamed Values]



### technology **Case Study: Performance Modeling** from seed in CPU + GPU Environment (12) linescid **Performance Metric** Approximation Iteration Draw Upper U and Lower L lines through 700 the following points: 600 (0,0), $(N_1/p, max_i \{ Perf_i (N_1/p) \} )$ (0,0), $(N_1/p, min_i \{Perf_i (N_1/p)\})$ 500 2 Let $x_i^{(U)}$ and $x_i^{(L)}$ be the intersections with $Perf_{i}(x)$ 400 MFLOPS IF exists $x_{i}^{(L)} - x_{i}^{(U)} \ge 1$ 300 THEN go to 3 ELSE go to 5 200 Bisect the angle between U and L by the 3 line M, and calculate intersections $x_i^{(M)}$ 100 4 IF $\Sigma_i X_i^{(M)} \leq N_1$ 0 THEN U=M 211 217 223 229 229 235 235 247 247 253 ELSE L=M **REPEAT 2** CPU (1 Core) GPU (Pageable) — GPU(Pinned) - CPU Model (1 Core) - GPU Model [Streamed Values] (5) Employ streaming strategy on the calculated workload value ıſi

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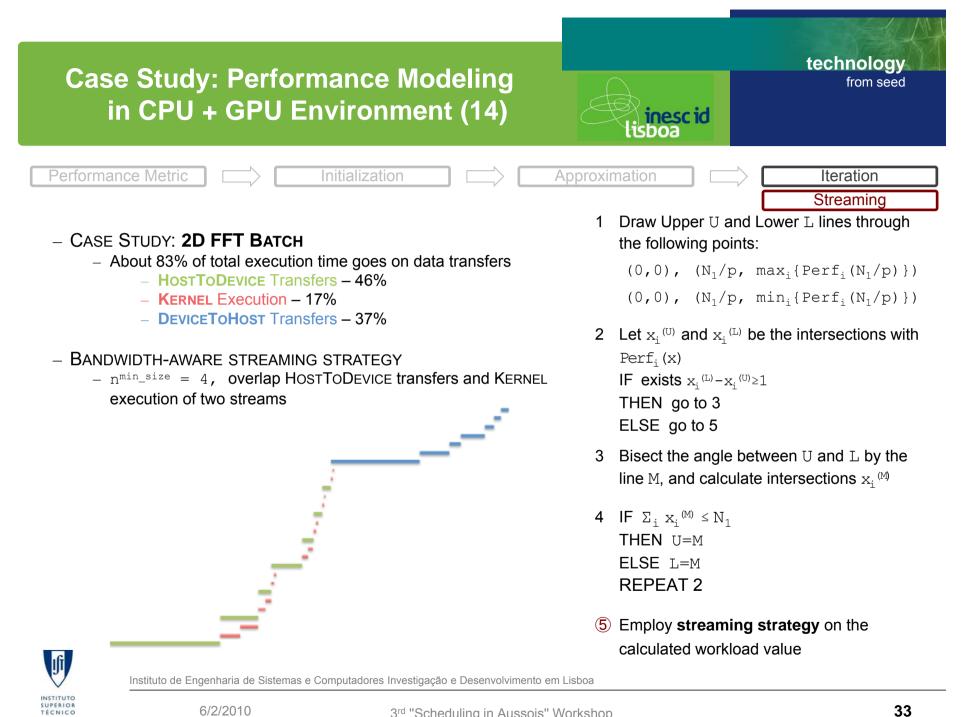
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- STREAMING STRATEGY
  - Results obtained using DIV2 STRATEGY consider application characterization (e.g. communication-to-computation ratio)
  - Workload size for the next stream should be chosen in order to OVERLAP TRANSFERS WITH COMPUTATION in the previous stream
- BANDWIDTH-AWARE STREAMING STRATEGY
  - Reuses the MINIMAL WORKLOAD SIZE FROM DIV2 STRATEGY (obtained via HOST-TO-DEVICE and DEVICE-TO-HOST tests)
    - IF (n<sup>curr</sup> ≥ n<sup>min\_size</sup>)
      THEN use strategy (cont. overlapping)
      ELSE restart strategy on n<sup>curr</sup> load
  - In case that load drops under  $n^{\texttt{min}\_\texttt{size}},$  strategy is restarted on remaining load

- 1 Draw Upper U and Lower L lines through the following points:
  - $(0,0), (N_1/p, max_i \{Perf_i(N_1/p)\})$
  - (0,0), (N<sub>1</sub>/p, min<sub>i</sub>{Perf<sub>i</sub>(N<sub>1</sub>/p)})
- 2 Let  $x_i^{(U)}$  and  $x_i^{(L)}$  be the intersections with Perf<sub>i</sub> (x) IF exists  $x_i^{(L)} - x_i^{(U)} \ge 1$ THEN go to 3 ELSE go to 5
- 3 Bisect the angle between U and L by the line M, and calculate intersections  $x_i^{(M)}$
- 4 IF  $\Sigma_{i} X_{i}^{(M)} \leq N_{1}$ THEN U=M ELSE L=M REPEAT 2
- (5) Employ streaming strategy on the calculated workload value





### **Case Study: Performance Modeling** in CPU + GPU Environment (15)

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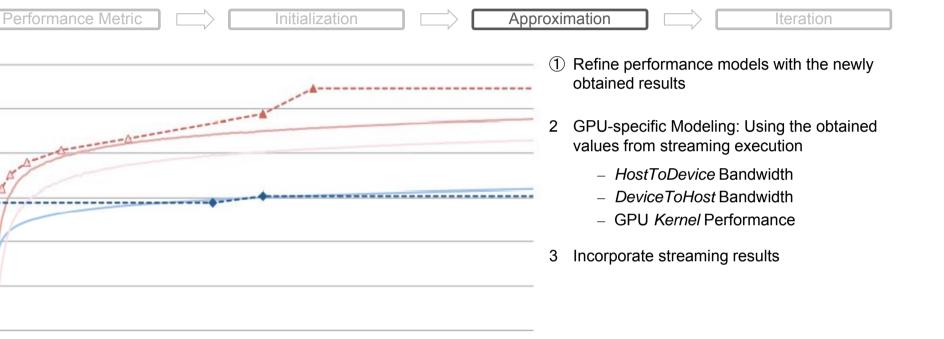
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	Performance Metric Appr	oxi	mation Iteration
700	▲	1	Draw Upper $\mathbb U$ and Lower $\mathbb L$ lines through the following points:
600 500	AA		$(0,0), (N_1/p, max_i {Perf_i(N_1/p)})$ $(0,0), (N_1/p, min_i {Perf_i(N_1/p)})$
400 300	A A A A A A A A A A A A A A A A A A A	2	Let $x_i^{(U)}$ and $x_i^{(L)}$ be the intersections with Perf <sub>i</sub> (x) IF exists $x_i^{(L)} - x_i^{(U)} \ge 1$
200		3	THEN go to 3 ELSE go to 5 Bisect the angle between $U$ and $L$ by the
100			line M, and calculate intersections $x_{\underline{i}}^{\mbox{ (M)}}$
о СРU (	<sup></sup> 다 다 다 다 다 다 다 다 다 다 다 다 다 다 다 다 다 다	4	IF $\Sigma_{i} \times_{i}^{(M)} \leq N_{1}$ THEN U=M ELSE L=M REPEAT 2
	[Streamed Values]	5	Employ streaming strategy on the calculated workload value



### Case Study: Performance Modeling in CPU + GPU Environment (16)

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700

600

500

400

300

200

100

0

CPU (1 Core)

m

MFLOPS

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[Streamed Values]

GPU (Pageable) — GPU(Pinned) - CPU Model (1 Core) 🛆 GPU Model

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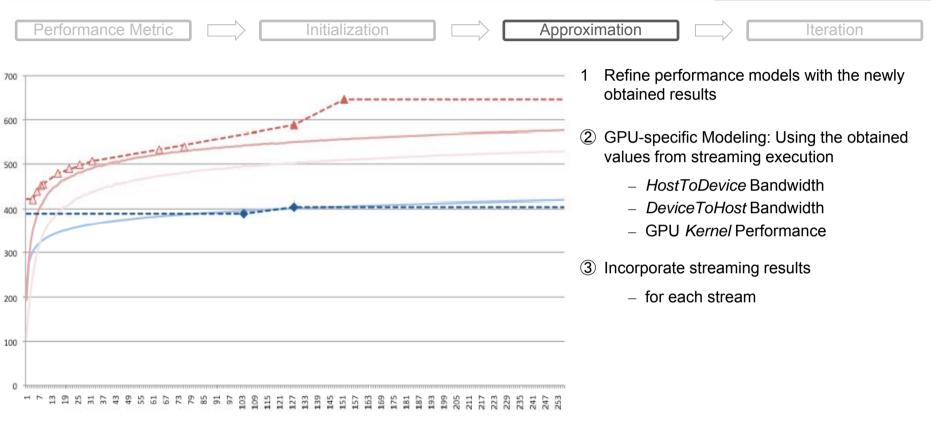
217 223 229 229 235 247 247 253

-----GPU Model (Streamed)

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-CPU (1 Core) — GPU (Pageable) — GPU(Pinned) ← CPU Model (1 Core) △ GPU Model [Streamed Values]

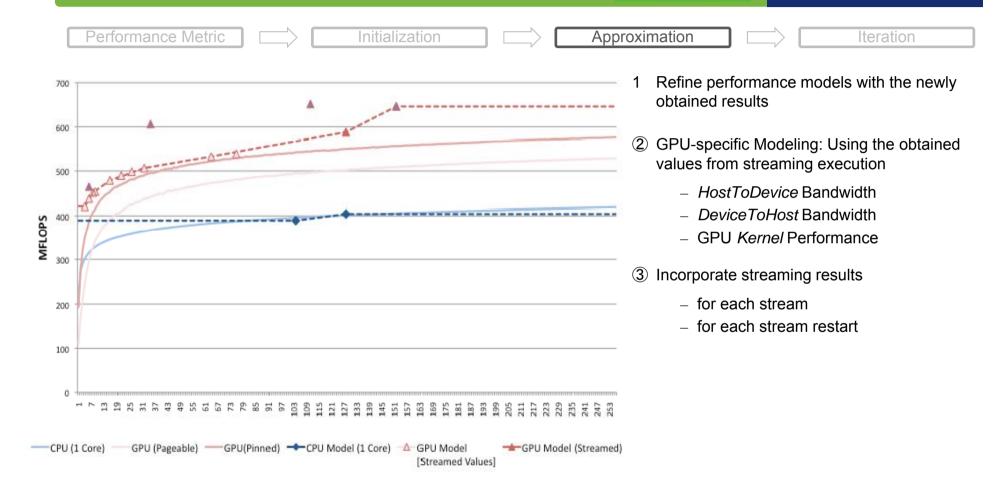


MFLOPS

### Case Study: Performance Modeling in CPU + GPU Environment (18)

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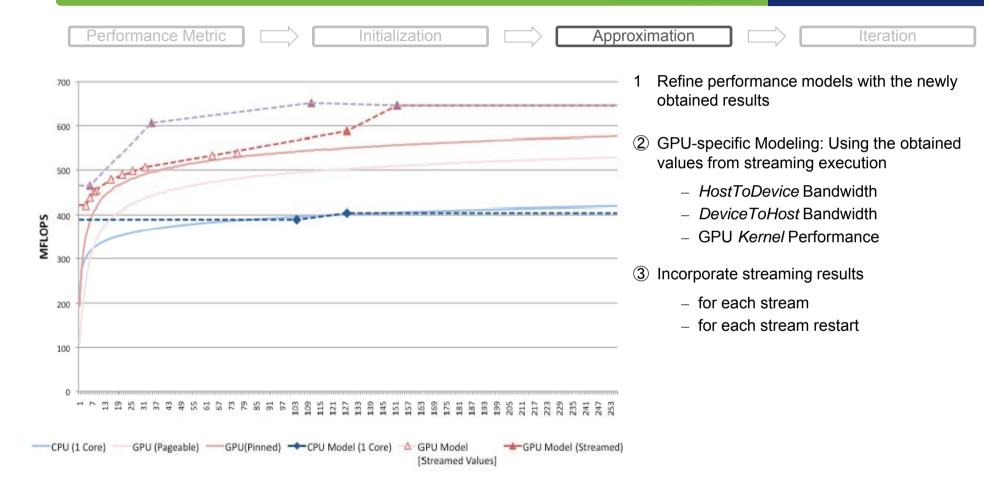
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### Case Study: Performance Modeling in CPU + GPU Environment (19)

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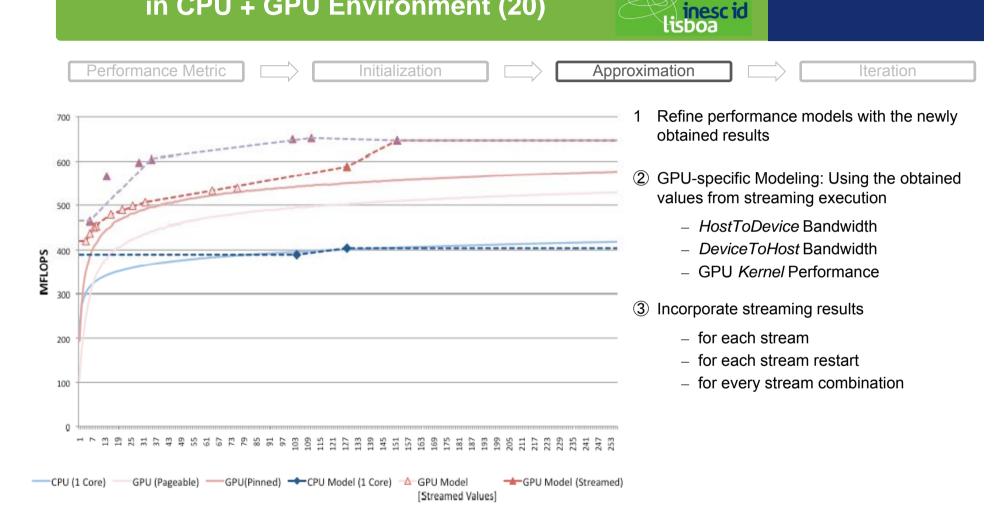
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### Case Study: Performance Modeling in CPU + GPU Environment (20)

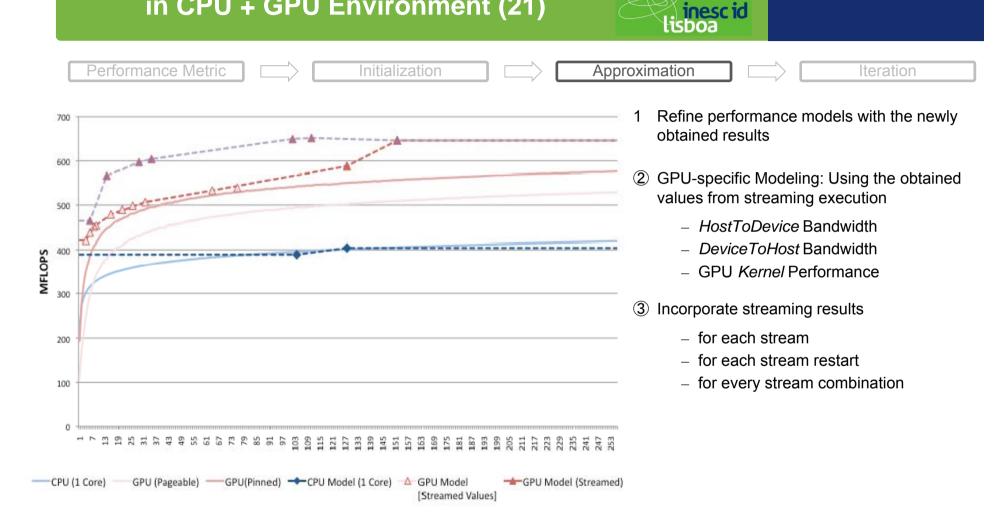
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### Case Study: Performance Modeling in CPU + GPU Environment (21)

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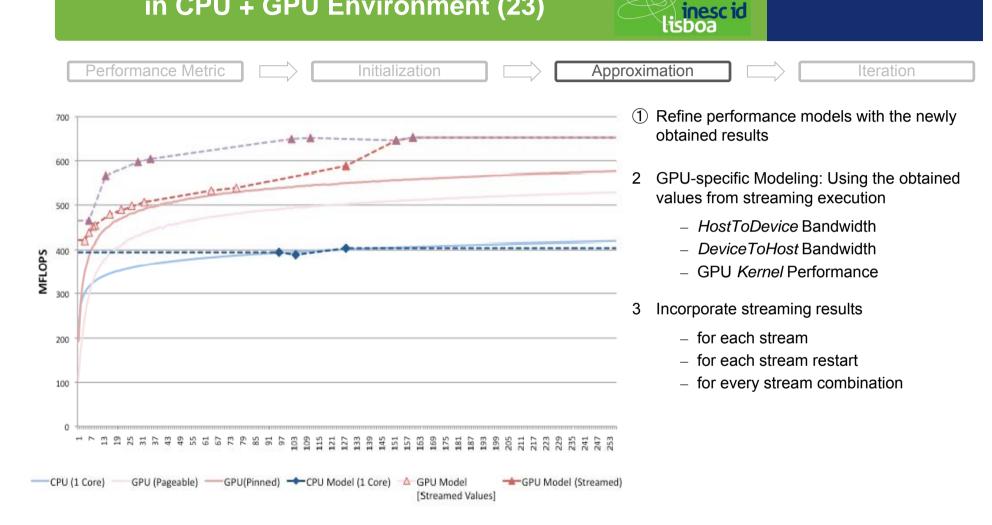


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### Case Study: Performance Modeling in CPU + GPU Environment (23)

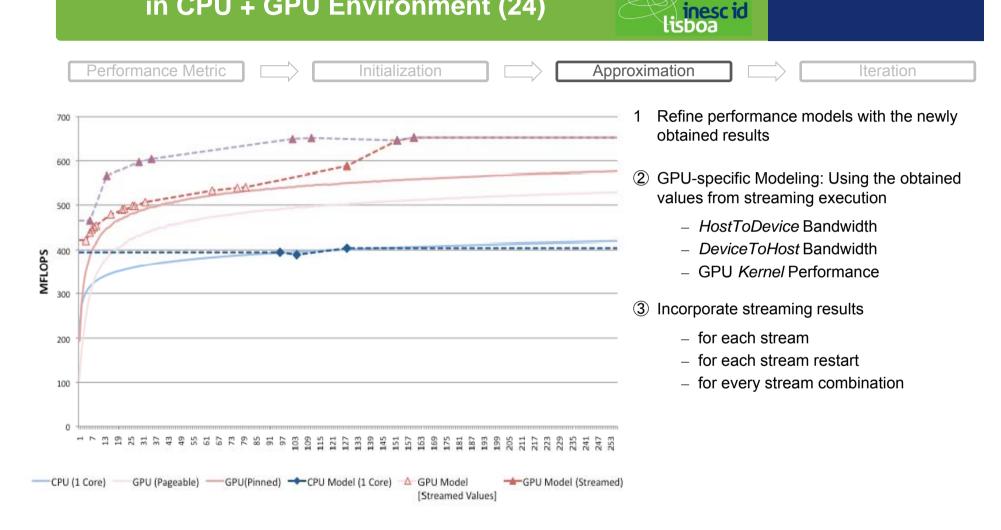
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### Case Study: Performance Modeling in CPU + GPU Environment (24)

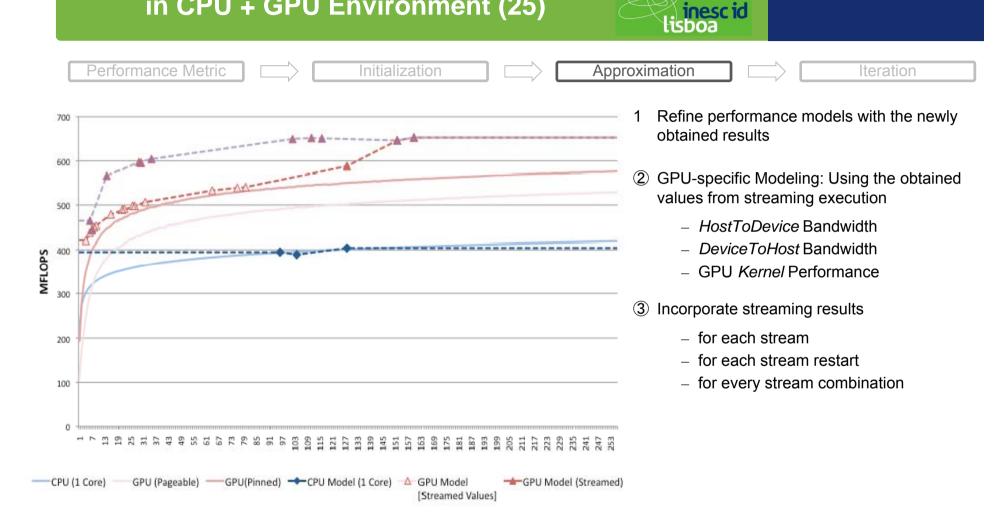
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### Case Study: Performance Modeling in CPU + GPU Environment (25)

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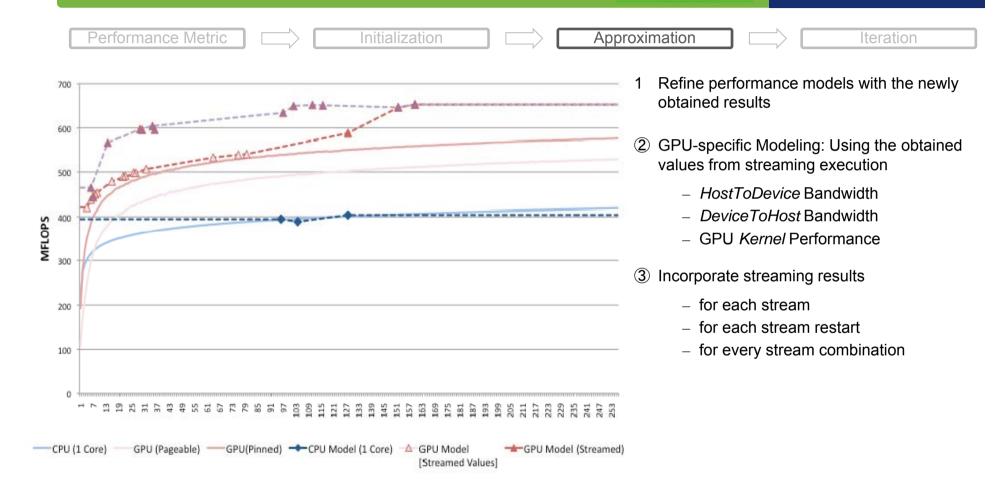




### Case Study: Performance Modeling in CPU + GPU Environment (26)

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### COLLABORATIVE ENVIRONMENT FOR HETEROGENEOUS COMPUTERS

### TRADITIONAL APPROACHES FOR PERFORMANCE MODELING

- Approximate the performance using number of points equal to the number of iterations
- In this case, **3 POINTS** per each device

### PRESENTED APPROACH FOR PERFORMANCE MODELING

- Models the performance using **MORE THAN 30 POINTS**, in this case
- COMMUNICATION-AWARE schedules in respect to limited and asymmetric interconnection bandwidth
- Employs STREAMING STRATEGIES to overlap communication with computation across devices
- BUILDS SEVERAL PER-DEVICE MODELS AT THE SAME TIME
  - OVERALL PERFORMANCE for each device + STREAMING GPU PERFORMANCE
  - HOSTTODEVICE BANDWIDTH Modeling
  - DEVICETOHOST BANDWIDTH Modeling
  - GPU KERNEL PERFORMANCE Modeling



