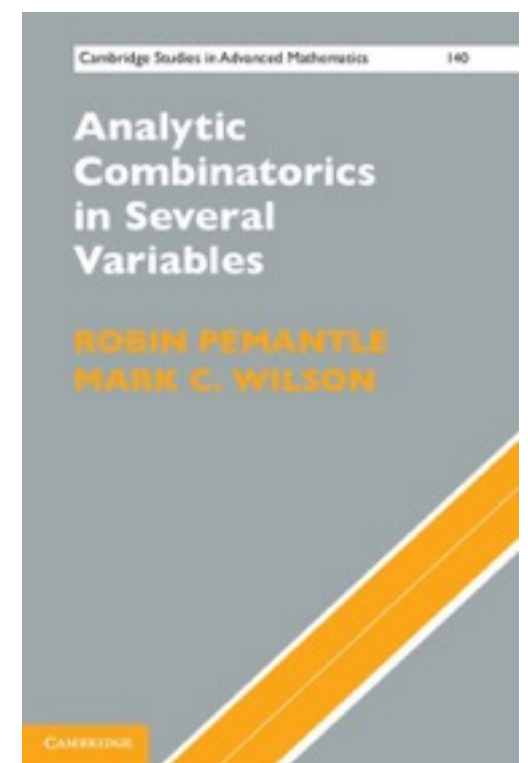
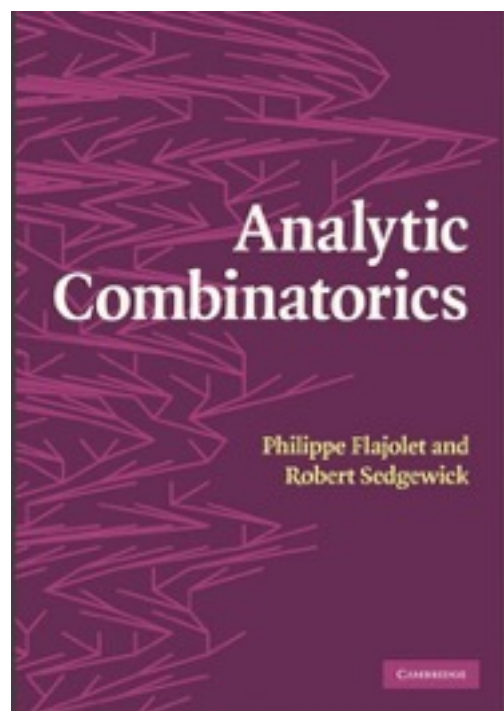


Symbolic-Numeric Tools for Multivariate Asymptotics

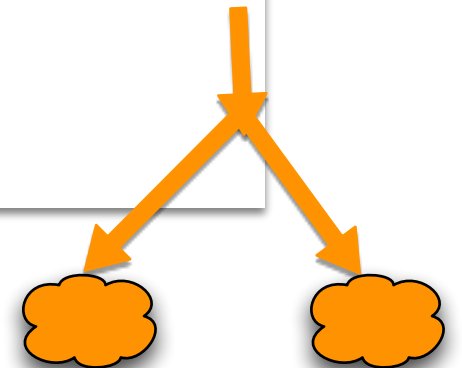
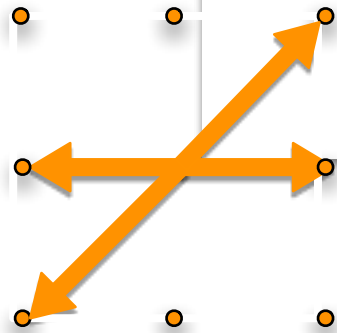
Bruno Salvy

Joint work with Stephen Melczer
to appear in Proc. ISSAC 2016

FastRelax meeting
May 25, 2016



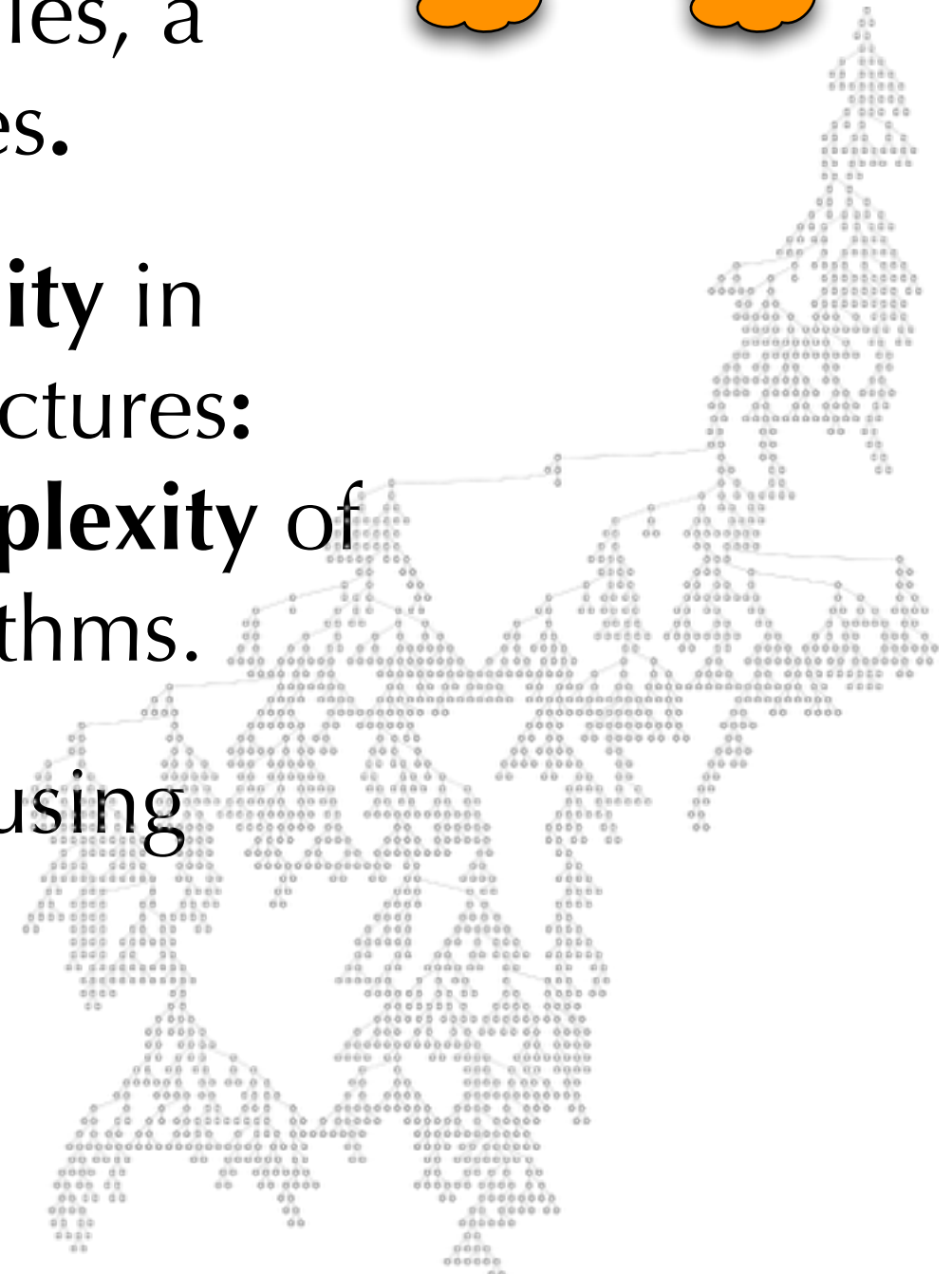
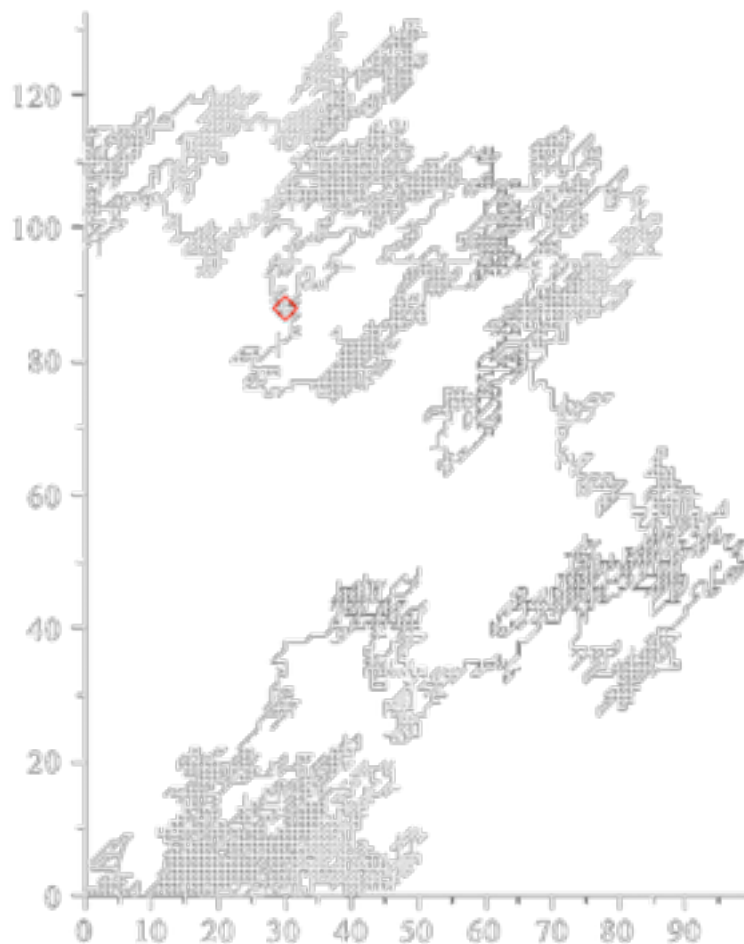
Combinatorics, Randomness and Analysis



From **simple local** rules, a **global** structure arises.

A quest for **universality** in random discrete structures:
➔ **probabilistic complexity** of structures and algorithms.

Quantitative results using **complex analysis**.



Overview

Ex.: binary trees

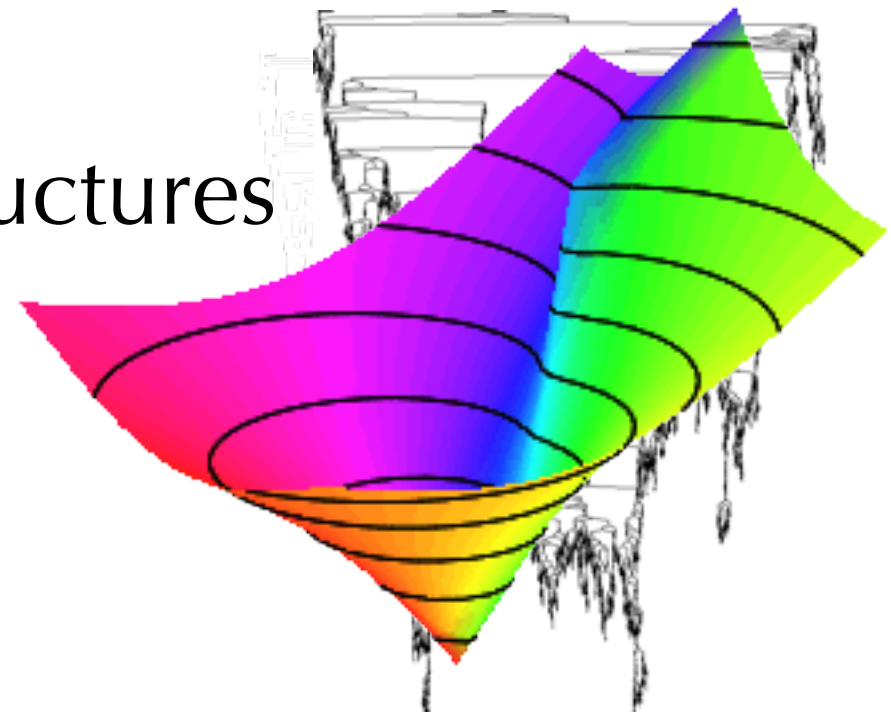
- Equations over combinatorial structures
- Generating functions

$$F(z) = \sum_{n=0}^{\infty} f_n z^n$$

- Complex analysis

$$f_n \sim \dots, \quad n \rightarrow \infty$$

Aim: several vars, automatically.



$$\mathcal{B} = \mathcal{Z} \cup \mathcal{B} \times \mathcal{B}$$

$$B(z) = z + B(z)^2$$

$$B_n \sim \frac{4^{n-1} n^{-3/2}}{\sqrt{\pi}}$$

I. Families of Generating Functions

Rational Generating Functions

(regular languages, linear recurrences with const. coeffs.)

Example

Generating function of $1, 2, 4, 8, \dots$ is $\frac{1}{1 - 2z}$.

Example

Fibonacci numbers $1, 1, 2, 3, 5, 8, \dots$ have gf $\frac{1}{1 - z - z^2}$.

Example (Simple lattice walks)

$$\begin{aligned} F(x, y) &= \frac{1}{1 - x - y} = \sum_{i,j} c_{i,j} x^i y^j = \sum_{i,j} \binom{i+j}{i} x^i y^j \\ &= 1 + x + 2xy + x^2 + y^2 + x^3 + y^3 + 3x^2y + 3xy^2 + 6x^2y^2 + \dots \end{aligned}$$

Here $c_{i,j}$ counts the number of ways to walk from $(0,0)$ to (i,j) using i steps $(0,1)$ and j steps $(1,0)$.

Rational Generating Functions

(regular languages, linear recurrences with const. coeffs.)

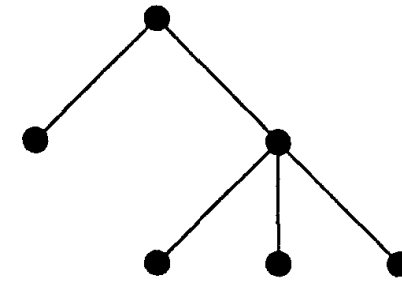
Example (Restricted factors in words)

$$F(x, y) = \frac{1 - x^3y^6 + x^3y^4 + x^2y^4 + x^2y^3}{1 - x - y + x^2y^3 - x^3y^3 - x^4y^4 - x^3y^6 + x^4y^6}$$

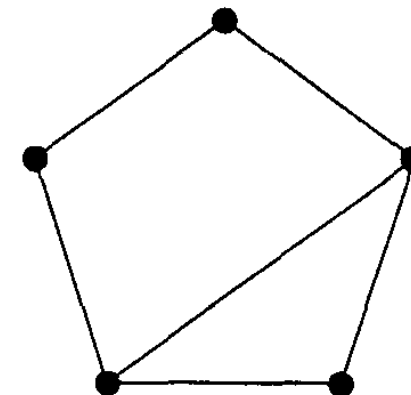
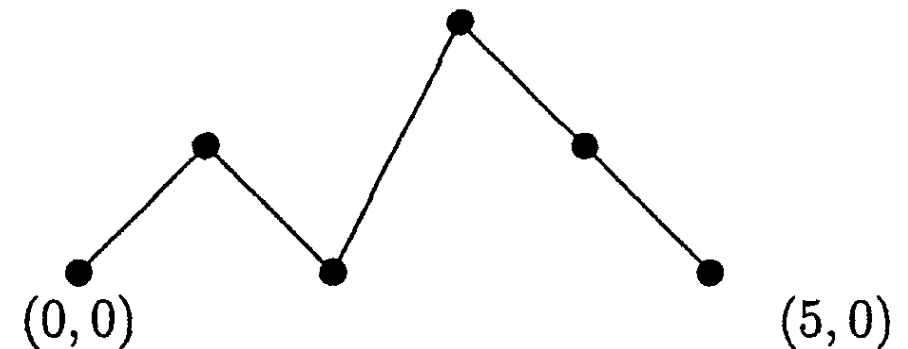
Here $c_{i,j}$ counts the number of binary words with i zeroes and j ones that do not contain 10101101 or 1110101.

Algebraic Generating Functions

- Plane trees with n nodes and given arity set;
- Bracketings of a word of length n ;
- Paths from $(0,0)$ to $(n,0)$ using steps $(1,k)$ and $(1,-1)$;
- Dissections of a convex n -gon...



$(x(xxx))$



Diagonals of Rational Functions

Input: Rational function $F(\mathbf{z}) = G(\mathbf{z})/H(\mathbf{z})$ with power series

$$F(\mathbf{z}) = \sum_{\mathbf{i} \in \mathbb{N}^n} c_{\mathbf{i}} \mathbf{z}^{\mathbf{i}}$$

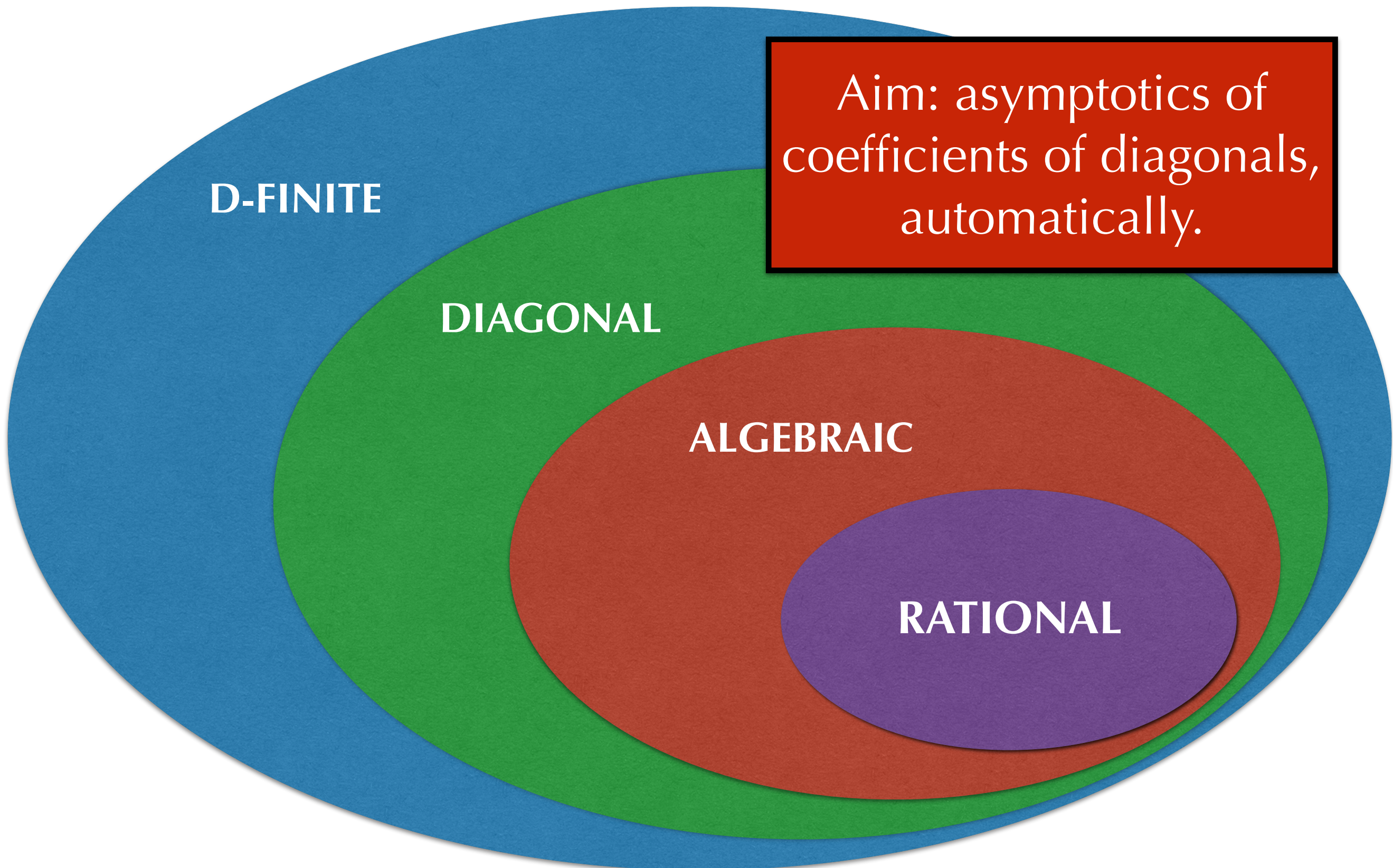
Goal: Asymptotics of the *diagonal sequence* $c_{k,k,\dots,k}$ as $k \rightarrow \infty$

Example (Apéry)

$$F(a, b, c, z) = \frac{1}{1 - z(1 + a)(1 + b)(a + c)(1 + b + c + bc + abc)}$$

Here $(c_{k,k,k,k,k})_{k \geq 0}$ determines Apéry's sequence, related to his celebrated proof of the irrationality of $\zeta(3)$.

Univariate Generating Functions



II. Analytic Combinatorics

Cauchys' formula

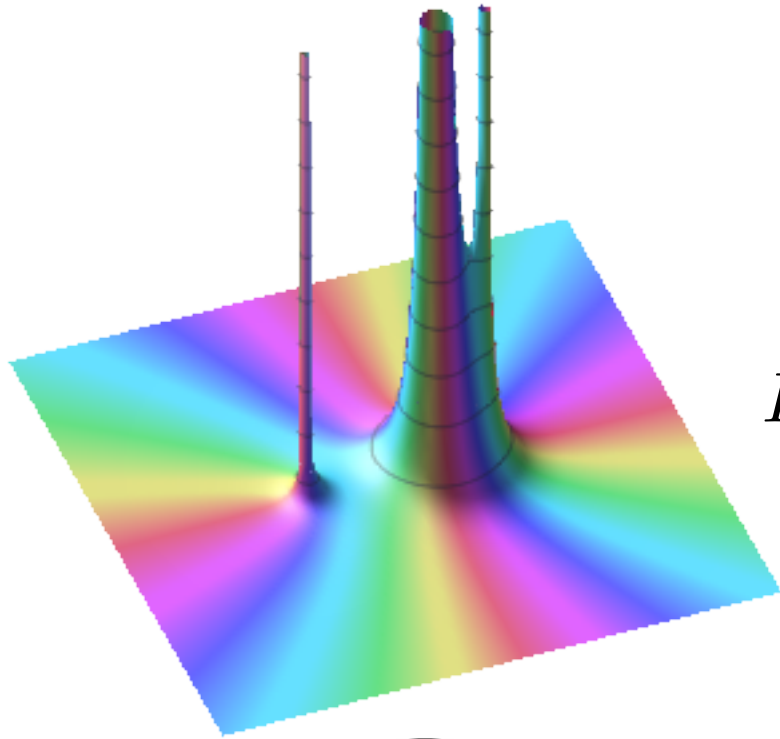
Thm. If $f = \sum_{i_1, \dots, i_n \geq 0} c_{i_1, \dots, i_n} z_1^{i_1} \cdots z_n^{i_n}$ is convergent in the neighborhood of 0, then

$$c_{i_1, \dots, i_n} = \left(\frac{1}{2\pi i} \right)^n \int_T f(z_1, \dots, z_n) \frac{dz_1 \cdots dz_n}{z_1^{i_1+1} \cdots z_n^{i_n+1}}$$

for any small torus T ($|z_j| = r e^{i\theta_j}$) around 0.

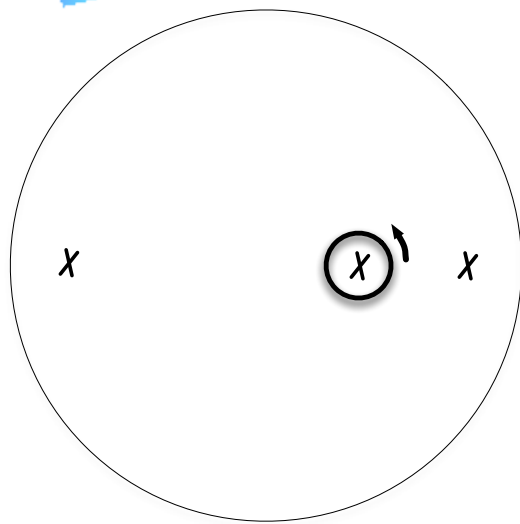
Asymptotics: deform the torus to pass where the integral concentrates asymptotically.

Coefficients of Univariate Rational Functions

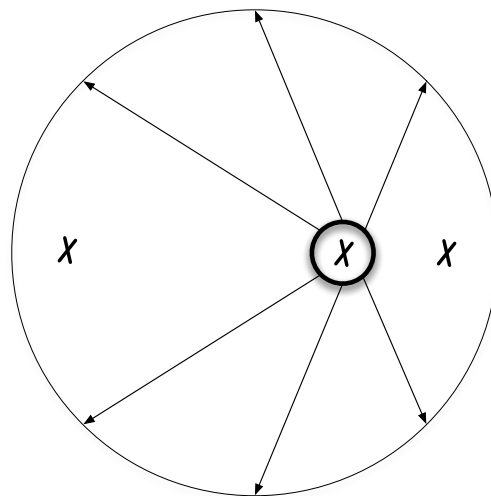


$$F_1 = 1 = \frac{1}{2\pi i} \oint \frac{1}{1-z-z^2} \frac{dz}{z^2}$$

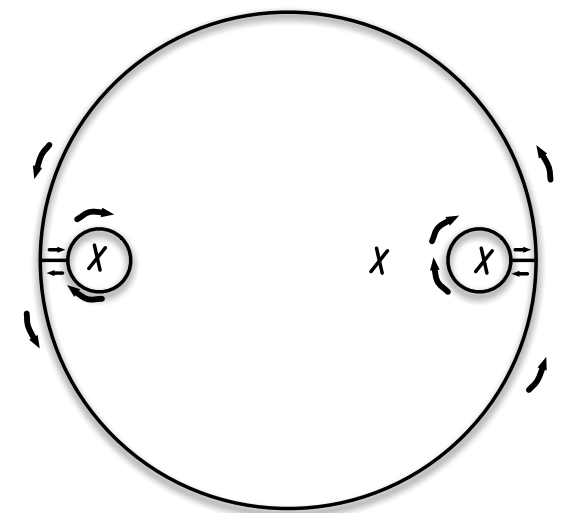
$$a_n = \frac{1}{2\pi i} \int_{\gamma} \frac{f(z)}{z^{n+1}} dz$$



=



=



As n increases, the smallest singularities dominate.

$$F_n = \frac{\phi^{-n-1}}{1+2\phi} + \frac{\overline{\phi}^{-n-1}}{1+2\overline{\phi}}$$

Conway's sequence

1,11,21,1211,111221,...

Generating function for lengths:

$$f(z) = P(z)/Q(z)$$

with $\deg Q = 72$.

Smallest singularity:

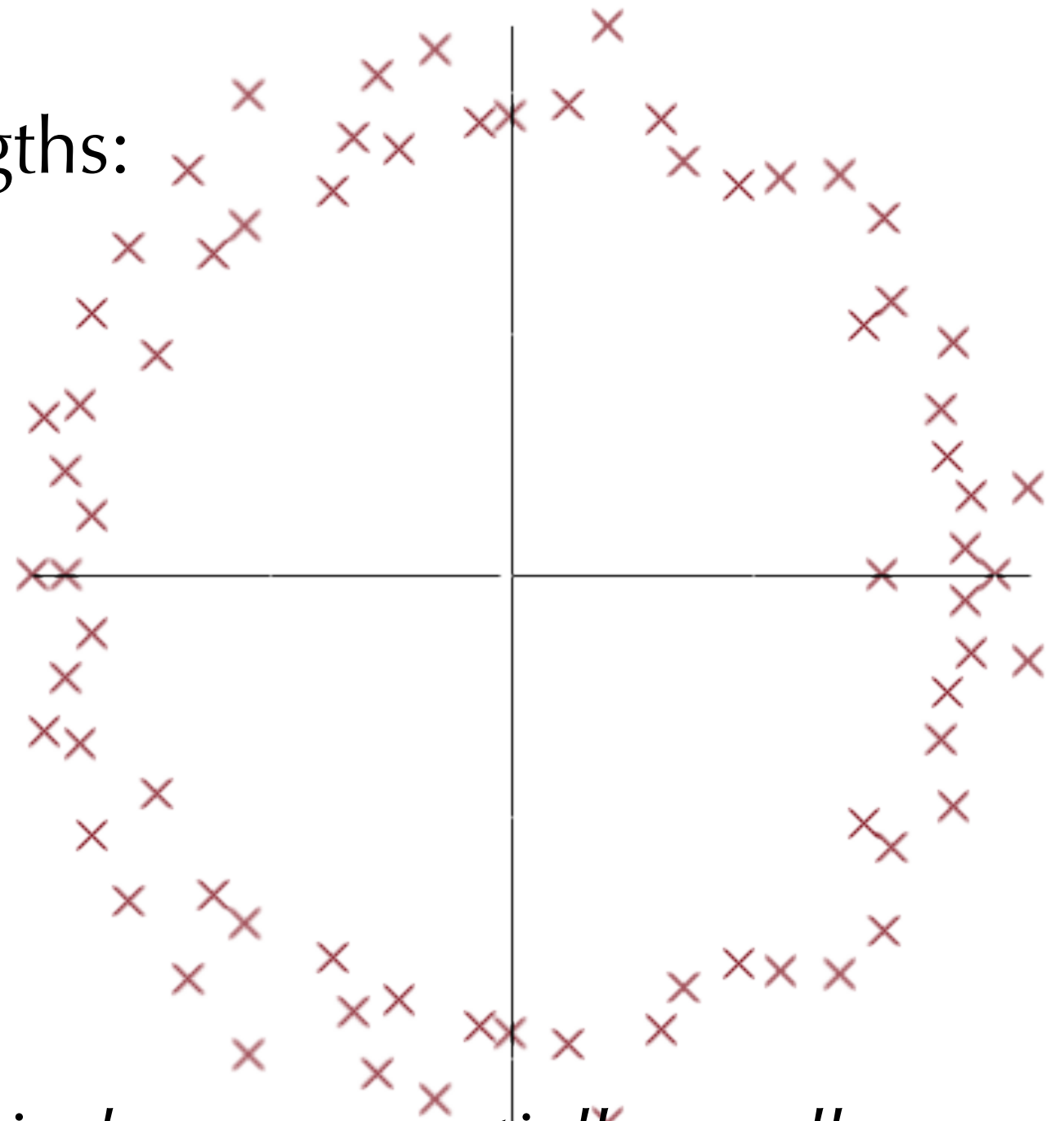
$$\delta(f) \approx 0.7671198507$$

$$\rho = 1/\delta(f) \approx 1.30357727$$

$$\ell_n \in 2.04216 \rho^n$$

$\rho \operatorname{Res}(f, \delta(f))$

remainder exponentially small



Coefficients of Multivariate Rational Functions

Def. $F(z_1, \dots, z_n)$ is **combinatorial** if every coefficient is ≥ 0 .

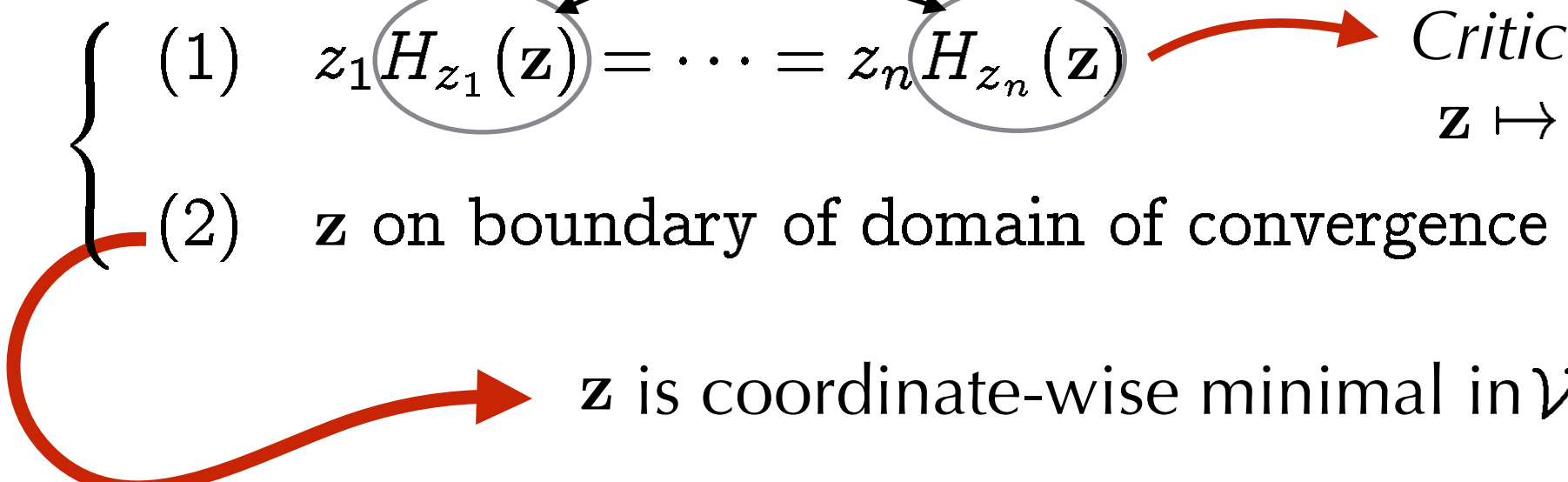
Assuming $F=G/H$ is combinatorial + generic conditions,
asymptotics are determined by the points in $\mathcal{V} = \{z : H(z) = 0\}$
such that

partial derivatives

$$\left\{ \begin{array}{l} (1) \quad z_1 H_{z_1}(z) = \dots = z_n H_{z_n}(z) \\ (2) \quad z \text{ on boundary of domain of convergence} \end{array} \right.$$

Critical Points
 $z \mapsto (z_1 \cdots z_n)$

z is coordinate-wise minimal in \mathcal{V}



(1) is an **algebraic** condition.

(2) is a **semi-algebraic** condition (can be expensive).

Aim: find these points,
in good complexity.



III. Symbolic-Numeric Computation

Kronecker Representation

Algebraic part: “compute” the solutions of the system

Suppose $H(\mathbf{z}) = 0, \quad z_1 H_{z_1}(\mathbf{z}) = \cdots = z_n H_{z_n}(\mathbf{z})$

$$\deg(H) = d, \quad \max \text{coeff}(H) \leq 2^h \quad D := d^n$$

Under genericity assumptions, in $\tilde{O}(hD^3)$ bit ops there is a prob. algorithm to find:

$$\left. \begin{array}{l} P(u) = 0 \\ P'(u)z_1 - Q_1(u) = 0 \\ \vdots \\ P'(u)z_n - Q_n(u) = 0 \end{array} \right\} \begin{array}{l} \text{Degree} \leq D \\ \text{Height} \leq \tilde{O}(hD) \end{array}$$

History and Background:
see Castro, Pardo, Hägele,
and Morais (2001)

Giusti, Lecerf, Salvy (2001)
Schost (2001)

System reduced to
a univariate polynomial.

Example (Lattice Path Model)

The number of walks from the origin taking steps $\{NW, NE, SE, SW\}$ and staying in the first quadrant has

$$F(x, y, t) = \frac{(1+x)(1+y)}{1 - t(x^2 + y^2 + x^2y^2 + 1)}$$

One can calculate the Kronecker representation

$$P(u) = 4u^9 + 39u^8 - (4339/2)u^7 + 4669/2u^6 + 4669/2u^5 - (4339/2)u^4 + 39u^3 + 4u^2 + 4$$

Next, find the *minimal* critical points, that lie on the boundary of the domain of convergence.

$$Q_t(u) = 4u^9 + 39u^8 - (4339/2)u^7 + 4669/2u^6 + 4669/2u^5 - (4339/2)u^4 + 39u^3 + 4u^2 + 4$$

so that the critical points are given by:

$$P(u) = 0, \quad x = \frac{Q_x(u)}{P'(u)}, \quad y = \frac{Q_y(u)}{P'(u)}, \quad t = \frac{Q_t(u)}{P'(u)}$$

Testing Minimality

Semi-algebraic part: In the combinatorial case, it is enough to examine only the positive real points on the variety to determine *minimality*.

Thus, we add the equation $H(tz_1, \dots, tz_n) = 0$ for a new variable t and select the positive real point(s) \mathbf{z} with no $t \in (0, 1)$ from a new Kronecker representation:

$$\tilde{P}(v) = 0$$

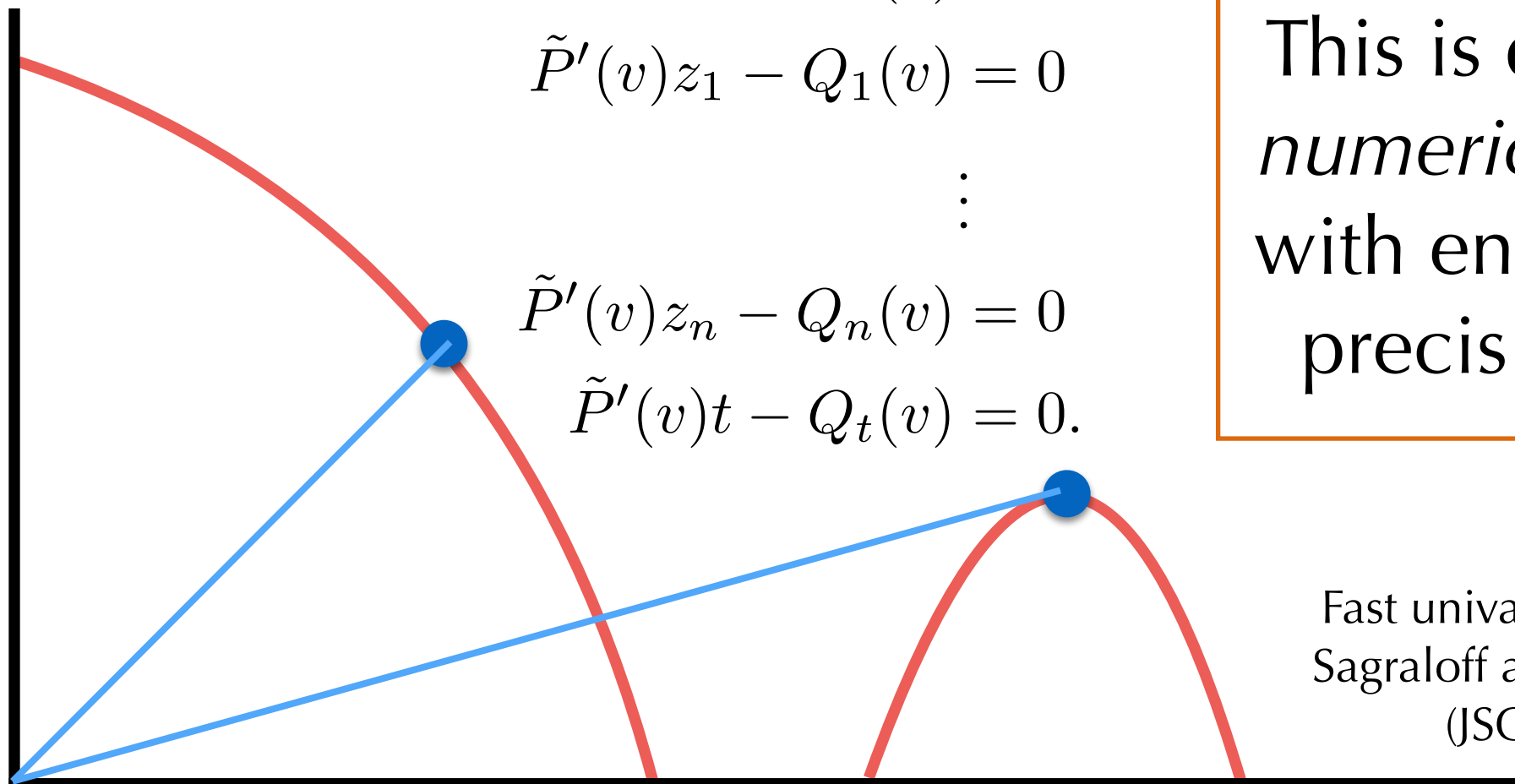
$$\tilde{P}'(v)z_1 - Q_1(v) = 0$$

$$\vdots$$

$$\tilde{P}'(v)z_n - Q_n(v) = 0$$

$$\tilde{P}'(v)t - Q_t(v) = 0.$$

This is done
numerically,
with enough
precision.



Fast univariate solving:
Sagraloff and Mehlhorn
(JSC 2016)

How many digits do we need?

Suppose $\alpha \neq \beta$ are roots of square-free poly $A(z)$ of degree δ and height σ .

For the Kronecker Representation, we have polynomials of

degree $D = d^n$

height $\tilde{O}(hD)$

so we need a precision of $2^{-\tilde{O}(hD^2)}$ to rigorously decide signs, find exact zeroes, and identify equal coordinates.

First Complexity Results for ACSV

Theorem (M. and Salvy, 2016)

Under generic conditions, and assuming $F(\mathbf{z})$ is combinatorial, the points contributing to dominant diagonal asymptotics can be determined in $\tilde{O}(hD^4)$ bit operations. Each contribution has the form

$$A_k = \left(T^{-k} k^{(1-n)/2} \cdot (2\pi)^{(1-n)/2} \right) (C + O(1/k))$$

and can be found to $2^{-\kappa}$ precision in $\tilde{O}(h(dD)^3 + D\kappa)$ bit ops.

The genericity assumptions heavily restrict the form of the asymptotic growth. Removing more of these assumptions is ongoing work.

Example 1

Example (Apéry)

$$F(a, b, c, z) = \frac{1}{1 - z(1 + a)(1 + b)(a + c)(1 + b + c + bc + abc)}$$

Here a Kronecker representation is

$$P(u) = u^2 - 366u - 17711$$

$$a = \frac{2u - 1006}{P'(u)}, \quad b = c = \frac{-320}{P'(u)}, \quad z = \frac{-164u - 7108}{P'(u)}$$

There are two real critical points, and one is positive. After testing minimality, one has proven asymptotics

```
> A, U, PRINT := DiagonalAsymptotics(numer(F),denom(F),[a,b,c,z],u,k, useFGb):  
A, U;
```

$$\frac{1}{4} \frac{\left(\frac{2u - 366}{34u + 1458} \right)^k \sqrt{2} \sqrt{\frac{2u - 366}{-96u - 4192}}}{k^{3/2} \pi^{3/2}}, [RootOf(_Z^2 - 366_Z - 17711, -43.27416997969$$

Example 2

Example (Restricted Words in Factors)

$$F(x, y) = \frac{1 - x^3 y^6 + x^3 y^4 + x^2 y^4 + x^2 y^3}{1 - x - y + x^2 y^3 - x^3 y^3 - x^4 y^4 - x^3 y^6 + x^4 y^6}$$

```
> ASM, U := DiagonalAsymptotics(numer(F),denom(F),indets(F),u,k,true,u-T,T):
ASM;
```

$$\frac{1}{2} \left(\left(\frac{84 u^{20} + 240 u^{19} - 285 u^{18} - 1548 u^{17} - 2125 u^{16} - 1408 u^{15} + 255 u^{14} + 756 u^{13} + 2509 u^{12} + 2856 u^{11} + 605 u^{10} + 2020 u^9 + 1233 u^8 - 1760 u^7 + 1233 u^6 - 162 u^5 - 612 u^4 - 902 u^3 - 616 u^2 + 254 u + 548}{-12 u^{20} + 30 u^{19} + 258 u^{18} + 500 u^{17} + 440 u^{16} - 102 u^{15} - 378 u^{14} - 1544 u^{13} - 2142 u^{12} - 550 u^{11} - 2222 u^{10} - 1644 u^9 + 2860 u^8 - 1848 u^7 + 1233 u^6 - 162 u^5 - 612 u^4 - 902 u^3 - 616 u^2 + 254 u + 548} \right. \right. \\ \left. \sqrt{\frac{84 u^{20} + 240 u^{19} - 285 u^{18} - 1548 u^{17} - 2125 u^{16} - 1408 u^{15} + 255 u^{14} + 756 u^{13} + 2509 u^{12} + 2856 u^{11} + 605 u^{10} + 2020 u^9 + 1233 u^8 - 1760 u^7 + 1233 u^6 - 162 u^5 - 612 u^4 - 902 u^3 - 616 u^2 + 254 u + 548}{-162 u^{18} - 612 u^{17} - 902 u^{16} - 616 u^{15} + 254 u^{14} + 548 u^{13} + 2054 u^{12} + 2156 u^{11} + 898 u^{10} + 2268 u^9 + 2462 u^8 - 2088 u^7 + 1312 u^6 - 255 u^5 - 190 u^4 - 19 u^3 + 46 u^2 + 461 u + 628}} \right. \\ \left. - 255 u^{16} - 190 u^{15} - 19 u^{14} + 46 u^{13} + 461 u^{12} + 628 u^{11} + 133 u^{10} + 374 u^9 + 161 u^8 - 384 u^7 + 146 u^6 - 138 u^5 - 285 u^4 - 40 u^3 + 91 u^2 - 30 u + 16 \right) \\ + 756 u^{13} + 2509 u^{12} + 2856 u^{11} + 605 u^{10} + 2020 u^9 + 1233 u^8 - 1760 u^7 + 924 u^6 - 492 u^5 - 675 u^4 + 632 u^3 - 249 u^2 + 24 u + 16) \Big)$$

```
> U;
[RootOf(4 _Z^21 + 12 _Z^20 - 15 _Z^19 - 86 _Z^18 - 125 _Z^17 - 88 _Z^16 + 17 _Z^15 + 54 _Z^14 + 193 _Z^13 + 238 _Z^12 + 559 _Z^11 + 6507895675196348205194331680572580787580891704416439705170142841771704333646705442481154976482193290559798252115981422944236560602741008070812421670189339690141212064002448936239379373317765301730899734374355850189724258427758618089253738395794473630503792093715946038342747035565025302472725126
```

Conclusion

We

- Give the first complexity bounds for methods in analytic combinatorics in several variables
- Combine strong symbolic results on the Kronecker representation with fast algorithms on univariate polynomials

Work in progress: extend beyond some of the assumptions.