

TD de Sémantique et Vérification

VIII– LTL and Büchi Automata Monday 15th March 2021

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- We denote the words characterised by an LTL formula φ by $[\![\varphi]\!] = \{\sigma \mid \sigma \vDash \varphi\}$.
- For a set of finite words W, we denote by \overrightarrow{W} the set of words having an infinite number of prefixes in $W \colon \overrightarrow{W} = \{ \sigma \in \Sigma^{\omega} \mid \exists^{\infty} \hat{\sigma} \in W, \hat{\sigma} \subseteq \sigma \}$
- A regular safety property P is a safety property such that P_{bad} is regular (i.e., recognizable by a DFA)

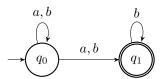
Let A be a set. Recall that

- For $u \in A^*$, $ext(u) = \{ \sigma \in A^\omega \mid u \subseteq \sigma \}$
- We can equip A^{ω} with the topology $\mathcal{U} = \{ \text{ext}(U) \mid U \subseteq A^* \}$

Decomposition of ω -regular Linear Time Properties

Exercise 1.

Let $L = (a+b)b^{\omega}$ and the following NBA \mathcal{A} recognizing L:



Show that there is a word $\sigma \notin \mathcal{L}_{\omega}(\mathcal{A})$ such that σ has an infinite number of prefixes recognized by \mathcal{A} as a DFA (i.e., $\exists^{\infty} \hat{\sigma} \subseteq_{\text{finite}} \sigma, \hat{\sigma} \in \mathcal{L}(\mathcal{A})$).

Exercise 2.

We want to show the decomposition theorem for ω -regular properties. Let $U \subseteq \Sigma^*$.

- 1. Show that if $A \subseteq \Sigma^{\omega}$ is ω -regular, then $\operatorname{pref}(A) \subseteq \Sigma^*$ is regular. (You may use that $\operatorname{pref}(U)$ is regular if $U \subseteq \Sigma^*$ is regular.)
- 2. Show that cl(U) is a safety property induced by $P_{bad} = U^c$.
- 3. Deduce that if $P \subseteq \mathcal{P}(AP)^{\omega}$ is an ω -regular safety property, then P is a regular safety property.
- 4. Let P_{bad} be a set of bad prefix of cl(U) and \mathcal{A} be a complete deterministic automaton recognizing P_{bad} such that for all final state q_F of \mathcal{A} and all $a \in \Sigma$, $\delta(q_F, a) = q_F$. Show that $\mathcal{L}_{\omega}(\mathcal{A}) = (2^{AP})^{\omega} \backslash cl(U)$.
- 5. Show that if U is regular then cl(U) is ω -regular (we recall that if cl(U) is a regular safety property, we can always find P_{bad} and A satisfying the properties of the previous question. Try to adapt A to recognize cl(U)).
- 6. Show that for every ω -regular linear time property P, there is a ω -regular safety property P_{safe} and a ω -regular liveness property P_{live} such that $P = P_{safe} \cap P_{live}$.

From LTL to NBAs

Exercise 3.

Let φ be a the LTL formula over AP = $\{a,b,c\}$ given by $\varphi = \Box a \wedge (b \ \mathcal{U} \neg c)$. Construct an (G)NBA \mathcal{A} , such that, $\mathcal{L}_{\omega}(\mathcal{A}) = \llbracket \varphi \rrbracket$.

Characterization of DBAs

Exercise 4.

- 1. Let AP = $\{a,b\}$. Show that $\{\sigma \mid \exists^{\infty}t, a \in \sigma(t)\} \subseteq (2^{AP})^{\omega}$ is not closed.
- 2. Show that there is a DBA \mathcal{A} such that $\mathcal{L}_{\omega}(\mathcal{A})$ is not a safety property.
- 3. Show that for every regular safety property P, there is a DBA \mathcal{A} such that $P = \mathcal{L}_{\omega}(\mathcal{A})$.