

# ENS Lyon. Day 3. Basic group: Problem Analysis

October 28, 2015

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Questions

# Problem A. K-th Maximum

## Statement

- ▶ Implement the data structure that allows the following queries:
  - ▶ Add and remove elements;
  - ▶ Find the  $k$ -th maximum among elements;

## Solution

- ▶ Store size field;
- ▶ If size of the left subtree is more than  $k$ , go to the left;
- ▶ Otherwise, go to the right with  $k = k - L.size + 1$ .

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# Problem B. Swapper

## Statement

- ▶ Given an array of  $n$  elements;
- ▶ Queries of two kinds:
  - ▶ Swap adjacent elements in the given range:  $x$  with  $x + 1$ ,  $x + 2$  with  $x + 3$  and so on;
  - ▶ Calculate the sum in the given range.

## Solution

- ▶ Keep two treaps with implicit key: elements with odd and even indices;
- ▶ Sum: find the sum in both trees;
- ▶ Swap: cut out corresponding ranges in both tree, swap it.

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# Problem C. Reverse

## Statement

- ▶ Given an array of  $n$  elements;
- ▶ Queries: reverse arbitrary range, calculate sum in the given range.

## Solution

- ▶ Treap with implicit key;
- ▶ Push: swap left and right subtrees; flip inconsistency in subtrees;
- ▶ See lecture notes.

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# Problem D. Cartesian Tree (Treap)

## Statement

- ▶ Given  $n$  pairs  $(a_i, b_i)$ ;
- ▶ Construct cartesian tree.

## Solution

- ▶ Sort pairs in increasing order of  $a_i$ ;
- ▶ Iteratively add  $(a_{i+1}, a_{i+1})$ 
  - ▶  $b_i > b_{i+1}$ , add  $(a_{i+1}, b_{i+1})$  as the right child  $(a_i, a_i)$
  - ▶ Otherwise, find parent  $v = (a_k, b_k)$  such that  $b_k > b_{i+1}$ ; new node is the right child of  $v$ , right child of  $v$  is the left child of new node;
  - ▶ See lecture notes.

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# Problem E. Sum

## Statement

- ▶ Consider set of integers;
- ▶ Queries: add element, find sum of all element  $x$  s.t.  $l \leq x \leq r$ .

## Solution

- ▶ Additional field “sum”: sum of elements in subtree with root in this vertex;
- ▶ Split the given range, return sum stored in the root;
- ▶ Update sum in merge and split operations (in a similar way as size of subtree).

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# Problem F. Move to front

## Statement

- ▶ Consider an array of  $n$  elements;
- ▶ Queries: Move given range in the beginning of the array.

## Solution

- ▶ Treap with implicit key;
- ▶ Move: cut required range;
- ▶ Merge trees in the order: range, left part, right part.

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# Problem G. Key Values

## Statement

- ▶ Consider array  $A$  of infinity length;
- ▶ Query  $Insert(L, K)$ :
  - ▶ If  $A[i]$ , then  $A[L] := K$ ;
  - ▶ Otherwise:  $Insert(L + 1, A[L])$  and after  $A[L] := K$ .

## Solution

- ▶ Treap with implicit key;
- ▶ To process “insert”: shift suffix of the array to the right.

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# Problem H. Implicit Key

## Statement

- ▶ Consider an array of  $n$  elements;
- ▶ Query: insert element in arbitrary position, remove element from arbitrary position.

## Solution

- ▶ Treap with implicit key;
- ▶ See lecture notes.

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# Problem I. Binary Search Tree

## Statement

- ▶ Implement Binary Search Tree;
- ▶ Effective implementation: all queries process at  $O(\log n)$  time.

## Statement

- ▶ Treap;
- ▶ Priorities are choosed randomly.

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# Problem J. Range Minimum Query

## Statement

- ▶ Consider an array of  $n$  elements;
- ▶ Queries: insert element in arbitrary position, find minimum in the given range.

## Solution

- ▶ Treap with implicit key;
- ▶ Store field for minimum in the subtree; update it in the similar way as size of subtree.

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