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Enhanced Coalgebraic Bisimulation

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We present a systematic study of bisimulation-up-to techniques for coalgebras. This enhances the bisimulation proof method for a large class of state based systems, including labelled transition systems but also stream systems and weighted automata. Our approach allows for compositional reasoning about the soundness of enhancements. Applications include the soundness of bisimulation up to bisimilarity, up to equivalence and up to congruence. All in all, this gives a powerful and modular framework for simplified coinductive proofs of equivalence.

1. Introduction

In the quest for good models of computation, the challenge of finding canonical notions of equivalence and corresponding proof methods has occupied the mind of many researchers. The pioneering work of Milner and Park (Milner, 1980; Park, 1981) on bisimulation has resulted in a vast amount of follow-up notions and improvements. Milner himself has proposed a powerful technique for modular reasoning about bisimilarity – bisimulation-up-to – which allows the re-use of existing bisimulation proofs and the construction of smaller relations to prove equivalence (Milner, 1983). Sangiorgi (Sangiorgi, 1998; Pous and Sangiorgi, 2012) has followed up on Milner's idea and proposed many enhancements to the theory of bisimulation-up-to for labelled transition systems. The gain of using bisimulations-up-to lies in the fact that they are smaller relations than usual bisimulations, thereby in many cases substantially reducing the amount of work and thus making the method more efficient. Bisimulation up to context is an example of an enhanced technique in which one can use the algebraic structure (syntax) of processes. Other examples are the notions of bisimulation up to union and bisimulation up to equivalence as well as combinations of any of these, which enable compositional, succinct reasoning on equivalence, combining both inductive and coinductive techniques.

In fact, some of the most useful up-to techniques are based on combinations of other

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enhancements. One of the difficulties in proving such up-to techniques to be sound is that the combination of sound enhancements is not necessarily sound. The first systematic study which addressed the issue of when such techniques can be safely combined is due to Sangiorgi (Sangiorgi, 1998). While this work focused on labelled transition systems, a more general, abstract *algebra of enhancements* in terms of lattices and monotone functions has been introduced by Pous and Sangiorgi (Pous, 2007; Pous and Sangiorgi, 2012). An important feature there is the notion of *compatible* functions, defining a class of sound enhancements that is closed under composition.

Enhancements of the bisimulation proof method are interesting not only for labelled transition systems but also for other types of state-based systems; for example, recently an efficient algorithm for checking equivalence of non-deterministic automata was introduced, based on bisimulation up to congruence (Bonchi and Pous, 2013). Another recent example is the application of a different kind of up-to techniques for deterministic automata to proving language equivalence (Rot et al., 2013b). Orthogonally to enhancements of the bisimulation proof method there is the theory of *coalgebra* in which the notion of bisimulation is extended to other models of computation, including all kinds of infinite data types, automata, and dynamical systems from a unifying perspective. By generalizing the theory of bisimulation-up-to to coalgebras one can study these techniques at a general level, with applications to many different types of state-based systems.

In the present paper we establish the connection between coalgebraic bisimulation-up-to and the algebra of enhancements by using the characterization of bisimulation in terms of monotone functions. This allows us to reason compositionally about the soundness of enhancements at the level of coalgebras. By showing that an up-to-technique is *compatible* one can now safely compose it with other compatible enhancements of coalgebraic bisimulation. We show that the most important enhancements are compatible.

In general many important instances of bisimulation-up-to, such as bisimulation up to equivalence and bisimulation up to bisimilarity, are not sound at the general level of coalgebras. We address this problem by a restriction to coalgebras for functors which preserve weak pullbacks; we prove the compatibility of such composition-based enhancements by using the theory of *relators* (Trnková, 1980; Rutten, 1998).

We show that bisimulation up to context is compatible whenever the system under consideration is a so-called λ-*bialgebra* for a distributive law λ (see, e.g., (Turi and Plotkin, 1997; Bartels, 2004; Klin, 2011)). Examples of such λ-bialgebras include non-deterministic and weighted automata but also operational models of specifications adhering to the *abstract GSOS* format (Turi and Plotkin, 1997), which generalizes the well-known GSOS format (Bloom et al., 1995) for labelled transition systems. So even in the more classical case of labelled transition systems this generalizes the result of Sangiorgi (Sangiorgi, 1998), who proved compatibility for the strictly less expressive De Simone format. Examples of operations which are expressible in GSOS but not in De Simone are the Kleene star and the priority operator (Aceto et al., 2001).

Most coalgebras considered in practice, such as labelled transition systems, stream systems and (non)-deterministic automata, are modeled by type functors which preserve weak pullbacks. However there are important instances where this is not the case, including certain weighted transition systems (Klin, 2009). In such cases one can consider *behavioural equivalence*, which is a weaker notion of equivalence. To accommodate proofs of behavioural equivalence, in this paper we additionally introduce a compositional theory of up-to techniques for behavioural equivalence, most of which are sound independently of the type functor.

Related work. The first account of bisimulation-up-to at the level of coalgebras was given by Lenisa (Lenisa, 1999; Lenisa et al., 2000). In (Lenisa, 1999), Lenisa considers a set-theoretic notion of coinduction and coinduction-up-to for abstract monotone operators, working in the direction of (Pous and Sangiorgi, 2012), and defines coalgebraic bisimulation-up-to-T for a monad T. However, in (Lenisa, 1999, page 22) the treatment of instances such as bisimulation up to bisimilarity are explicitly mentioned as an open problem. Interestingly, she conjectured that "the theory of functors and relators could shed some light on this problem" which is indeed precisely the successful approach taken in the present work.

The up-to-context technique for coalgebraic bisimulation was later derived as a special case of so-called λ-coinduction (Bartels, 2004). However, (Bartels, 2004, pages 126, 129) mentions already that it would be ideal to combine the up-to-context technique with other enhancements. Indeed, combining up-to-context with up-to-bisimilarity or up-to-equivalence yields powerful proof techniques (see, e.g., (Pous and Sangiorgi, 2012) and this paper for examples). In this paper we strengthen the soundness result of (Bartels, 2004) to *compatibility* of up-to-context, allowing for such combinations.

The recent paper (Zhou et al., 2010) introduces bisimulation-up-to where the notion of bisimulation is based on a specification language for polynomial functors (which does not include, for example, labelled transition systems). In contrast, we base our work on the standard notion of bisimulation, and only need to restrict to weak pullback preserving functors to obtain our soundness results. In the paper (Luo, 2006) coalgebraic bisimulation-up-to techniques are studied based on relation lifting. There, a concrete coalgebraic notion of compatibility, based on the notion of *consistency* proposed by Sangiorgi (Sangiorgi, 1998) is introduced, and it is used to prove soundness of bisimulation up to context and of bisimulation up to bisimilarity (the latter is actually false in general, as we show in this paper). However, in (Luo, 2006) combinations of enhancements are not considered.

Recently, a new generalization of bisimulation-up-to to coalgebras was introduced by a subset of the authors in (Rot et al., 2013a). In the present paper we take this generalization as our starting point. The solution of (Rot et al., 2013a) to the problem of unsoundness of bisimulation up to bisimilarity was, similarly to the present paper, to restrict to functors which preserve weak pullbacks. For such systems, bisimulation coincides with behavioural equivalence, and for the latter, the problematic up-to techniques were shown to be sound. In (Rot et al., 2013a), the soundness of each of the enhancements and of their combinations had to be shown separately. Indeed, the problem of compositionality of enhancements, which we solve in the present paper, was left as the main open problem.

Outline. In Section 2 we recall coalgebras and bisimulations. Then in Section 3 we introduce bisimulation-up-to, together with the main instances and a number of examples. In Section 4 we recall the algebra of enhancements; Section 5 then rephrases bisimulation-up-to in terms of this theory. In Section 6 we prove compatibility results for the instances of bisimulation-up-to introduced in Section 3. Section 7 contains a similar development of up-to techniques for behavioural equivalence, and we conclude in Section 8.

Notation. Let Set be the category of sets and functions. Sets are denoted by capital letters $X, Y, \ldots$ and functions by lower case $f, g, \ldots$. We write id for the identity function and $g \circ f$ for function composition, defined by $(g \circ f)(x) = g(f(x))$. We write $f[S]$, for a function $f\colon X \to Y$ and a set $S \subseteq X$, to denote the image of S under f. Given sets X and Y, $X \times Y$ is the cartesian product of X and Y (with the usual projection maps π_1 and π_2), X^Y is the set of functions $f\colon Y \to X$ and $\mathcal{P}(X)$ is the set of subsets of X. These operations, defined on sets, can analogously be defined on functions, yielding (bi-)functors. We write 2 for the two elements set $2 = \{0, 1\}$, ω for the set of natural numbers and $\mathbb{R}$ for the set of real numbers. By $\mathbb{R}^X_\omega$ we denote the set of functions $f\colon X \to \mathbb{R}$ with *finite support*, i.e., such that $f(x) \neq 0$ for finitely many elements x. We will write the elements v of $\mathbb{R}^X_\omega$ as a formal sum $v = f(x_1)x_1 + \cdots + f(x_n)x_n$. $\mathbb{R}^X_\omega$ carries a vector space structure where sum and scalar product (denoted by $+$ and $\cdot$) are defined pointwise: we call it the *free vector space* generated by X.

We denote the category of sets and relations by Rel. Relations are denoted by capital letters $R, S, \ldots$ We write Δ for the identity relation and $R \circ S$ for relation composition, defined as usual: $R \circ S = \{(x, z) \in X \times Z \mid \exists y \text{ s.t. } xRy \text{ and } ySz\}$.

2. Coalgebra and bisimulation

We recall coalgebras and bisimulations, and make explicit the underlying notion of progression, which we need in the sequel. A *coalgebra* for a functor $F\colon \mathsf{Set} \to \mathsf{Set}$ is a pair (X, α) consisting of a set X and a function $\alpha\colon X \to FX$. A function $f\colon X \to Y$ is an *(F-coalgebra) homomorphism* between (X, α) and (Y, β) if $Ff \circ \alpha = \beta \circ f$.

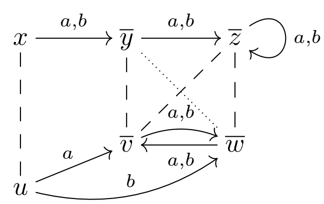
Definition 1. For a coalgebra $\alpha\colon X \to FX$ and relations $R, S \subseteq X \times X$, we say R *progresses to* S, denoted $R \rightarrowtail S$, if there exists a $\gamma\colon R \to FS$ making the following diagram commute:

$$\begin{array}{ccccc} X & \xleftarrow{\pi_1} & R & \xrightarrow{\pi_2} & X \\ \downarrow{\scriptstyle \alpha} & & \downarrow{\scriptstyle \gamma} & & \downarrow{\scriptstyle \alpha} \\ FX & \xleftarrow{F\pi_1} & FS & \xrightarrow{F\pi_2} & FX \end{array}$$

A *bisimulation* is a relation R such that $R \rightarrowtail R$.

Bisimulations are usually defined between two different systems, which however can be reduced to bisimulations on a single system by using coproducts (c.f., Appendix A). We use bisimulations on single systems for technical convenience and notational clarity. *Bisimilarity*, denoted by $\sim$, is defined as the largest bisimulation. Bisimulations can be seen as a proof technique for bisimilarity: in order to prove that $x \sim y$ (for any two states $x, y \in X$) it suffices to exhibit a bisimulation R such that $x \; R \; y$.

Example 1. Deterministic automata on the alphabet A are coalgebras for the functor $FX = 2 \times X^A$. Indeed, a deterministic automaton is a pair $(X, \langle o, t \rangle)$, where X is a set of states and $\langle o, t \rangle \colon X \to 2 \times X^A$ is a function with two components: o, the output function, determines if a state x is final ($o(x) = 1$) or not ($o(x) = 0$); and t, the transition function, returns for each input letter $a \in A$ the next state. Bisimilarity coincides with the standard notion of language equivalence, which can thus be proved by providing a suitable bisimulation. Unfolding the definition, a relation $R \subseteq X \times X$ is a bisimulation provided that for all $(x, y) \in R$: $o(x) = o(y)$ and, for all $a \in A$, $(t(x)(a), t(y)(a)) \in R$. As an example consider the automaton below, with final states y, z, v, w and transitions given by the solid arrows. The relation given by the four dashed lines together with the dotted line is a bisimulation.



Example 2. Labelled transition systems over a set of labels A are coalgebras for the functor $FX = \mathcal{P}(A \times X)$. An F-coalgebra (X, α) consists of a set of states X and a function $\alpha \colon X \to \mathcal{P}(A \times X)$ that maps each state $x \in X$ into a set of possible transitions (a, x'), where a is the label and x' is the arriving state. We write $x \xrightarrow{a} x'$ iff $(a, x') \in \alpha(x)$. Bisimilarity and bisimulation instantiate to the classical notions by Milner and Park (Milner, 1980; Park, 1981). A relation $R \subseteq X \times X$ is called a *bisimulation* provided that for all $(x, y) \in R$: if $x \xrightarrow{a} x'$ then there exists a state y' such that $y \xrightarrow{a} y'$ and $(x', y') \in R$, and vice versa.

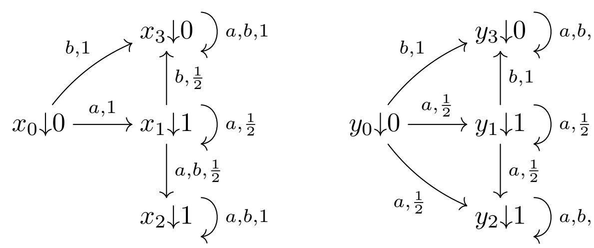
Example 3. A weighted automaton with input alphabet A is a pair $(X, \langle o, t \rangle)$, where X is a set of states, $o \colon X \to \mathbb{R}$ is an output function associating to each state its output weight and $t \colon X \to (\mathbb{R}^X_\omega)^A$ is the transition relation that associates a weight to each transition. We shall use the following notation: $x \xrightarrow{a,r} y$ means that $t(x)(a)(y) = r$. Weight 0 means no transition.

Every weighted automaton induces a coalgebra for the functor $FX = \mathbb{R} \times X^A$ that is defined as $(\mathbb{R}^X_\omega, \langle o^\sharp, t^\sharp \rangle)$ where $\mathbb{R}^X_\omega$ is the free vector space generated by X and $o^\sharp \colon \mathbb{R}^X_\omega \to \mathbb{R}$ and $t^\sharp \colon \mathbb{R}^X_\omega \to (\mathbb{R}^X_\omega)^A$ are the linear extensions of o and t. For a detailed explanation see (Bonchi et al., 2012, Section 3).

For an example consider the weighted automaton $(X, \langle o, t \rangle)$ depicted below (1), where we use $x{\downarrow}r$ to denote $o(x) = r$ and, as usual, arrows represent transitions. Part of the infinite corresponding F-coalgebra is depicted in (2). Note that now states are vectors in $\mathbb{R}^X_\omega$ and that transitions are only labeled by symbols in A: the vector $v = \frac{1}{2}y_1 + \frac{1}{2}y_2 \in \mathbb{R}^X_\omega$ goes with a into $t^\sharp(\frac{1}{2}y_1 + \frac{1}{2}y_2)(a) = \frac{1}{2}t(y_1)(a) + \frac{1}{2}t(y_2)(a) = \frac{1}{4}y_1 + \frac{1}{4}y_2 + \frac{1}{2}y_2$.

In (Bonchi et al., 2012) it is shown that bisimilarity on $(\mathbb{R}^X_\omega, \langle o^\sharp, t^\sharp \rangle)$ coincides with standard weighted language equivalence (Salomaa and Soittola, 1978; Berstel and Reutenauer, 1988) which can therefore be proved by means of bisimulations. A relation $R \subseteq \mathbb{R}^X_\omega \times \mathbb{R}^X_\omega$ is a bisimulation provided that for all $(v, w) \in R$: $o^\sharp(v) = o^\sharp(w)$ and, for all $a \in A$, $(t^\sharp(v)(a), t^\sharp(w)(a)) \in R$.

For an example, consider the weighted automaton below.



(1)

The states x_0 and y_0 are weighted language equivalent. To formally prove it we exhibit a bisimulation $R \subseteq \mathbb{R}^X_\omega \times \mathbb{R}^X_\omega$ such that $(x_0, y_0) \in R$. Note that this relation is infinite since it must contain at least all the pairs given by the dotted lines below.

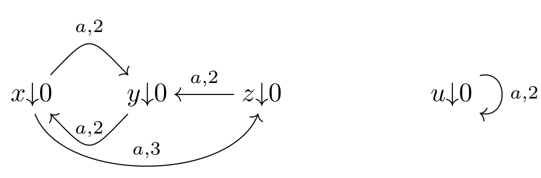
$$\begin{array}{ccccccccc} x_0{\downarrow}0 & \xrightarrow{a} & x_1{\downarrow}1 & \xrightarrow{a} & \frac{1}{2}x_1 + \frac{1}{2}x_2{\downarrow}1 & \xrightarrow{a} & \frac{1}{4}x_1 + \frac{3}{4}x_2{\downarrow}1 & \xrightarrow{a} & \dots \\ \vdots & & \vdots & & \vdots & & \vdots & & \vdots \\ y_0{\downarrow}0 & \xrightarrow{a} & \frac{1}{2}y_1 + \frac{1}{2}y_2{\downarrow}1 & \xrightarrow{a} & \frac{1}{4}y_1 + \frac{3}{4}y_2{\downarrow}1 & \xrightarrow{a} & \frac{1}{8}y_1 + \frac{7}{8}y_2{\downarrow}1 & \xrightarrow{a} & \dots \end{array} \tag{2}$$

In Section 3 we will show that there exists a finite bisimulation up to context proving that x_0 and y_0 are bisimilar and therefore language equivalent.

Example 4. The notion of weighted automata from Example 3 can be generalized by replacing the field of reals $\mathbb{R}$ with any commutative semiring $\mathbb{S}$. As discussed in (Bonchi et al., 2012), the coalgebraic characterization can be easily extended by taking the free semi-module $\mathbb{S}^X_\omega$ rather than the free vector space $\mathbb{R}^X_\omega$.

We now exhibit an example of a weighted automaton for the tropical semiring $\mathbb{T} = \langle \mathbb{R} \cup \{\infty\}, min, \infty, +, 0\rangle$. In this semiring, the addition operation is given by the function *min* (defined in the obvious way) having ∞ as neutral element. The multiplication is given by the function + (defined in the obvious way) having 0 as neutral element.

The weighted automaton $(X, \langle o, t\rangle)$ below



(3)

induces the coalgebra $(\mathbb{T}^X_\omega, \langle o^\sharp, t^\sharp\rangle)$ which is partially depicted below.

$$\begin{array}{ccccccc} x{\downarrow}0 & \xrightarrow{a} & min(2+y, 3+z){\downarrow}2 & \xrightarrow{a} & min(4+x, 5+y){\downarrow}4 & \xrightarrow{a} & \dots \\ \vdots & & \vdots & & \vdots & & \vdots \\ u{\downarrow}0 & \xrightarrow{a} & 2+u{\downarrow}2 & \xrightarrow{a} & 4+u{\downarrow}4 & \xrightarrow{a} & \dots \end{array} \tag{4}$$

The states x and u are weighted language equivalent. To prove it we would need an infinite bisimulation, since it should relate all the pairs of states linked by the dotted lines in the above figure. In Section 3, we will exhibit a finite bisimulation up to congruence proving that x and u are language equivalent.

Example 5. We now consider *stream systems* (over the reals), which are coalgebras for the functor $FX = \mathbb{R} \times X$. At first, we take the set $\mathbb{R}^\omega = \{\sigma \mid \sigma\colon \omega \to \mathbb{R}\}$ of all streams (infinite sequences) of elements of $\mathbb{R}$ and we define $(-)_0\colon \mathbb{R}^\omega \to \mathbb{R}$ and $(-)'\colon \mathbb{R}^\omega \to \mathbb{R}^\omega$ as $(\sigma)_0 = \sigma(0)$ and $(\sigma)'(n) = \sigma(n+1)$. The F-coalgebra $(\mathbb{R}^\omega, \langle(-)_0, (-)'\rangle)$ is called *final*, which means that from any F-coalgebra there exists a unique homomorphism into it (Rutten, 2000).

Then, we define operations on $\mathbb{R}^\omega$ by means of *behavioural differential equations* (Rutten, 2003), in which an operation is defined by specifying its initial value $(-)_0$ and its derivative $(-)'$. These operations will become relevant in the examples in Section 3.

Differential equation	Initial value	Name
$(\sigma+\tau)' = \sigma' + \tau'$	$(\sigma+\tau)_0 = \sigma_0 + \tau_0$	sum
$(\sigma \otimes \tau)' = \sigma' \otimes \tau + \sigma \otimes \tau'$	$(\sigma \otimes \tau)_0 = \sigma_0 \times \tau_0$	shuffle product
$(\sigma^{-1})' = -\sigma' \otimes (\sigma^{-1} \otimes \sigma^{-1})$	$(\sigma^{-1})_0 = (\sigma_0)^{-1}$	shuffle inverse

In the second column, the operations $+$, $\times$ and $(-)^{-1}$ on the right of the equations are the standard operations on $\mathbb{R}$. The inverse is only defined on streams σ for which $\sigma_0 \neq 0$. With every real number r we associate a stream $\mathsf{r} = (r, 0, 0, 0, \ldots)$, and we abbreviate $(-\mathsf{1}) \otimes \sigma$ by $-\sigma$. The set of *terms* $T(\mathbb{R}^\omega)$ is defined by the grammar $t ::= \ \sigma \mid t_1 + t_2 \mid t_1 \otimes t_2 \mid t_1^{-1}$ where σ ranges over $\mathbb{R}^\omega$. We call a term *well-formed* if the inverse is never applied to a subterm with initial value 0; this notion can be straightforwardly defined by induction. We can turn the set $T_{wf}(\mathbb{R}^\omega)$ of well-formed terms into an F-coalgebra $\mathcal{S} = (T_{wf}(\mathbb{R}^\omega), \langle(-)_0, (-)'\rangle)$ by defining $(-)_0\colon T_{wf}(\mathbb{R}^\omega) \to \mathbb{R}$ and $(-)'\colon T_{wf}(\mathbb{R}^\omega) \to T_{wf}(\mathbb{R}^\omega)$ by induction according to the final coalgebra (for the base case σ) and the above specification (for the other terms).

In (Rutten, 2003) it is shown that every term $t \in T_{wf}(\mathbb{R}^\omega)$ denotes a stream in $\mathbb{R}^\omega$ and that two terms t_1 and t_2 denote the same stream iff $t_1 \sim t_2$. As a result, in order to prove that two terms denote the same stream it is enough to exhibit a bisimulation relating them. A relation $R \subseteq T_{wf}(\mathbb{R}^\omega) \times T_{wf}(\mathbb{R}^\omega)$ is called a bisimulation provided that for all $(t_1, t_2) \in R$ it holds that $(t_1)_0 = (t_2)_0$ and $((t_1)', (t_2)') \in R$.

3. Bisimulation-up-to

The following definition generalizes the notions of bisimulation-up-to (Sangiorgi, 1998; Pous, 2007; Pous and Sangiorgi, 2012) from labelled transition systems to coalgebras.

Definition 2. Let (X, α) be a coalgebra and $f\colon \mathcal{P}(X \times X) \to \mathcal{P}(X \times X)$ be a function on relations. A *bisimulation up to* f is a relation R such that $R \rightarrowtail f(R)$. We say that f is *sound* if $R \subseteq \sim$ for all R such that $R \rightarrowtail f(R)$.

If a function f is sound then giving a bisimulation up to f relating two states x and y is enough to prove that $x \sim y$. We now exhibit some functions that we will prove to be sound, under certain conditions, in Section 6. These conditions are satisfied in all the examples presented in this section.

Up-to-equivalence. Consider the function e mapping a relation R to its equivalence closure $e(R)$. A bisimulation up to e is called a bisimulation *up to equivalence*. Similarly one can define *up-to-transitivity* and *up-to-symmetry*.

Example 6. The relation R denoted by the four dashed lines in the automaton of Example 1 is *not* a bisimulation, since $((t(x)(b), t(u)(b)) = (y, w) \notin R$. However R *is* a bisimulation up to equivalence, since the pair (y, w) is in $e(R)$. Hopcroft and Karp's algorithm (Hopcroft and Karp, 1971) exploits this technique for checking equivalence of deterministic automata: rather than exploring n^2 pairs of states (where n is the number of states), the algorithm visits at most n pairs (that is the number of equivalence classes).

Up-to-union. For a fixed relation $S \subseteq X \times X$ consider the function $u_S(R) = R \cup S$. We call a bisimulation up to u_S a bisimulation *up to union with* S. Intuitively the successor states may be related either by R again or by S.

Up-to-union-and-equivalence. By composing the above functions e and u_S we obtain a new interesting up-to technique. If R is a bisimulation up to $e \circ u_S$ then we say R is a bisimulation *up to S-union and equivalence.*

Example 7. Recall Example 5 and suppose that we want to prove that the stream $\mathsf{1} = (1, 0, 0, \ldots)$ is the unit for the shuffle product $\otimes$, that is, $\sigma \otimes \mathsf{1} \sim \sigma$. We make use of the relation $R = \{(\sigma \otimes \mathsf{1}, \sigma) \mid \sigma \in T_{wf}(\mathbb{R}^\omega)\}$. For any $\sigma \in T_{wf}(\mathbb{R}^\omega)$, we have $(\sigma \otimes \mathsf{1})_0 = \sigma_0 \times \mathsf{1}_0 = \sigma_0$. Further $(\sigma \otimes \mathsf{1})' = \sigma' \otimes \mathsf{1} + \sigma \otimes \mathsf{1}' = \sigma' \otimes \mathsf{1} + \sigma \otimes \mathsf{0}$; this element is not in relation with σ', so R is not a bisimulation. However given some basic laws of stream calculus, in particular $\sigma \otimes \mathsf{0} \sim \mathsf{0}$, $\sigma + \mathsf{0} \sim \sigma$ and the fact that $\sim$ is a congruence, we obtain

$$\sigma' \otimes \mathsf{1} + \sigma \otimes \mathsf{0} \sim \sigma' \otimes \mathsf{1} + \mathsf{0} \sim \sigma' \otimes \mathsf{1} \; R \; \sigma'$$

so R is a bisimulation up to $\sim$-union and equivalence and it proves that $\sigma \otimes \mathsf{1} \sim \sigma$.

Up-to-bisimilarity. Consider the function $b(R) = \sim \circ R \circ \sim$ which composes a relation on both sides with bisimilarity. A bisimulation up to b corresponds to the well-known concept of bisimulation *up to bisimilarity*, in which derivatives (i.e., the arriving states) do not need to be related directly but may be bisimilar to elements that are. Notice that every bisimulation up to bisimilarity is also a bisimulation up to $\sim$-union and equivalence. Since $\sim$ is transitive on stream systems, the relation R in Example 7 is also a bisimulation up to bisimilarity.

Up-to-context. When the state space of a coalgebra carries some kind of algebraic structure (as it is the case, for instance, with process algebras and regular expressions) it can be interesting to consider bisimulation up to *context*. In order to achieve the desired level of generality, we define the notion of contextual closure of a relation with respect to an algebra for a monad T; the uninitiated reader can safely skip to the examples below.

Recall that a *monad* is a triple (T, μ, η) where T is an endofunctor, $\mu\colon TT \Rightarrow T$ and $\eta\colon Id \Rightarrow T$ are natural transformations such that $\mu \circ T\eta = id = \mu \circ \eta T$ and $\mu \circ \mu T = \mu \circ T\mu$. A T-algebra is a pair (X, β) where X is a set and $\beta\colon TX \to X$ is a function such that $\beta \circ \eta_X = id$ and $\beta \circ \mu_X = \beta \circ T\beta$. A function $f\colon X \to Y$ is a *(T-algebra) homomorphism* between (X, β) and (Y, γ) if $f \circ \beta = \gamma \circ Tf$.

For a T-algebra (X, β), the *contextual closure* of a relation $R \subseteq X \times X$ is defined as

$$c_\beta(R) = \langle \pi_1^\sharp, \pi_2^\sharp \rangle [TR] = \{(\pi_1^\sharp(t), \pi_2^\sharp(t)) \mid t \in TR\}$$

where $\pi_i^\sharp = \beta \circ T\pi_i$. Whenever β is clear from the context we will simply write $c(R)$. If R is a bisimulation up to c then we call R a *bisimulation up to context.* It can be verified that since β is a T-algebra (for a monad), it holds that $R \subseteq c_\beta(R)$ for any relation R.

Example 8. Given a signature Σ, i.e., a set of operations with associated arities, we consider the free T_Σ-algebra $\mu\colon T_\Sigma T_\Sigma X \to T_\Sigma X$. Intuitively, $T_\Sigma X$ consists of all Σ-terms with variables in X. Now, given a relation $R \subseteq T_\Sigma X \times T_\Sigma X$ on these terms, the contextual closure $c(R) \subseteq T_\Sigma X \times T_\Sigma X$ can be inductively characterized by the following rules, where g is any operator of Σ with arity n.

$$\frac{s\ R\ t}{s\ c(R)\ t} \qquad \frac{s_i\ c(R)\ t_i \qquad i = 1 \ldots n}{g(s_1, \ldots, s_n)\ c(R)\ g(t_1, \ldots, t_n)}$$

This slightly differs from the definition in (Pous and Sangiorgi, 2012) where the contextual closure is defined as $c'(R) = \{(C[s_1, \ldots s_n], C[t_1, \ldots t_n]) \mid C$ a context and for all i, $(s_i, t_i) \in R\}$ (a context C is a term with $n \geq 0$ holes $[\cdot]_i$ in it). In our case c' can be obtained as $c \circ r$, i.e., by precomposing c with the *reflexive closure* function r.

Example 9. Recall from Example 3 that every weighted automaton $(X, \langle o, t\rangle)$ induces a coalgebra whose state space is the free vector space $\mathbb{R}_\omega^X$, that is, an algebra for the monad $\mathbb{R}_\omega^-$. Given a relation $R \subseteq \mathbb{R}_\omega^X \times \mathbb{R}_\omega^X$ its contextual closure $c(R) \subseteq \mathbb{R}_\omega^X \times \mathbb{R}_\omega^X$ can be inductively characterized by the following rules.

$$\frac{v\ R\ w}{v\ c(R)\ w} \qquad \frac{-}{0\ c(R)\ 0} \qquad \frac{v_1\ c(R)\ w_1 \qquad v_2\ c(R)\ w_2}{v_1 + v_2\ c(R)\ w_1 + w_2} \qquad \frac{v\ c(R)\ w \qquad r \in \mathbb{R}}{r \cdot v\ c(R)\ r \cdot w}$$

With the above characterization, it is easy to introduce bisimulation up to context for weighted automata: a relation $R \subseteq \mathbb{R}_\omega^X \times \mathbb{R}_\omega^X$ is a bisimulation up to context provided that for all $(v, w) \in R$ it holds that $o_1^\sharp(v) = o_2^\sharp(w)$ and for all $a \in A$, $(t_1^\sharp(v)(a), t_2^\sharp(w)(a)) \in c(R)$.

As an example consider the weighted automaton in (1). It is easy to see that the relation $R = \{(x_2, y_2), (x_3, y_3), (x_1, \frac{1}{2}y_1 + \frac{1}{2}y_2), (x_0, y_0)\}$ is a bisimulation up to context: consider $(x_1, \frac{1}{2}y_1 + \frac{1}{2}y_2)$ (the other pairs are trivial) and observe that

$$\begin{array}{ccc} x_1 & \xrightarrow{a} & \frac{1}{2}x_1 + \frac{1}{2}x_2 \\ \vdots R & & \vdots c(R) \\ \frac{1}{2}y_1 + \frac{1}{2}y_2 & \xrightarrow[a]{} & \frac{1}{4}y_1 + \frac{3}{4}y_2 \end{array} \qquad \begin{array}{ccc} x_1 & \xrightarrow{b} & \frac{1}{2}x_3 + \frac{1}{2}x_2 \\ \vdots R & & \vdots c(R) \\ \frac{1}{2}y_1 + \frac{1}{2}y_2 & \xrightarrow[b]{} & \frac{1}{2}y_3 + \frac{1}{2}y_2 \end{array}$$

It is worth noting that the above bisimulation up to context is finite, while one would need an infinite bisimulation to prove the equivalence of x_0 and y_0.

Example 10 ((Rot et al., 2013b)). The set $\mathcal{P}(A^*)$ of all languages forms a deterministic automaton as follows: the set of states is precisely the set of languages $\mathcal{P}(A^*)$ itself; a state $L \in \mathcal{P}(A^*)$ is accepting, i.e., $o(L) = 1$, if and only if the empty word ε is in L, and for every $a \in A$, the next state after an a-transition is given by the language derivative $t(L)(a) = \{w \mid aw \in L\}$. One can readily show that the language accepted by a state L is precisely L itself, and so whenever two languages L and K are bisimilar, they are in fact equal. The operations of language union $+$, composition $\cdot$ and Kleene star *, defined as usual, define an algebra on $\mathcal{P}(A^*)$. We have the following properties of derivatives of these operations due to Brzozowski; we formulate this in terms of languages (e.g., (Conway, 1971, page 41)):

$$
\begin{array}{ll}
t(\mathsf{0})(a) = \mathsf{0} & o(\mathsf{0}) = \mathsf{0} \\
t(\mathsf{1})(a) = \mathsf{0} & o(\mathsf{1}) = \mathsf{1} \\
t(b)(a) = \begin{cases} \mathsf{1} & \text{if } b = a \\ \mathsf{0} & \text{otherwise} \end{cases} & o(b) = \mathsf{0} \\
t(L + K)(a) = t(L)(a) + t(K)(a) & o(L + K) = o(L) \vee o(K) \\
t(L \cdot K)(a) = t(L)(a) \cdot K + o(L) \cdot t(K)(a) & o(L \cdot K) = o(L) \wedge o(K) \\
t(L^*)(a) = t(L)(a) \cdot L^* & o(L^*) = \mathsf{1}
\end{array}
$$

for any languages L, K.

Arden's rule states that if $L = KL + M$ for some languages L, K and M, and K does not contain the empty word, then $L = K^*M$. In order to prove its validity coinductively, let L, K, M be languages such that $\varepsilon \notin K$ and $L = KL + M$, and let $R = \{(L, K^*M)\}$. Then $o(K) = 0$ since by assumption $\varepsilon \notin K$, so

$$
\begin{aligned}
o(L) &= o(KL + M) = (o(K) \wedge o(L)) \vee o(M) = (0 \wedge o(L)) \vee o(M) \\
&= o(M) = 1 \wedge o(M) = o(K^*) \wedge o(M) = o(K^*M)
\end{aligned}
$$

and for any $a \in A$:

$$
\begin{aligned}
t(L)(a) &= t(KL + M)(a) = t(K)(a) \cdot L + o(K) \cdot t(L)(a) + t(M)(a) \\
&= t(K)(a) \cdot L + t(M)(a) \;\; c(R) \;\; t(K)(a) \cdot K^*M + t(M)(a) = t(K^*M)(a)
\end{aligned}
$$

So R is a bisimulation up to context, where the contextual closure is taken with respect to the operators of union and composition.

Up-to-congruence. By composing the functions e, c and r described above, we obtain another interesting up to technique. A bisimulation up to $e \circ c \circ r$ is called a bisimulation *up to congruence.* A recently introduced algorithm (Bonchi and Pous, 2013), for checking the equivalence of non-deterministic automata, exploits this technique. The bisimulations up to congruence built by this algorithm can be exponentially smaller than bisimulation up to context. This is due to the use of *transitivity.*

Example 11. Recall from Example 4 the tropical semiring $\mathbb{T}$. Given a relation $R \subseteq \mathbb{T}_\omega^X \times \mathbb{T}_\omega^X$, its congruence closure can be inductively characterized by the following rules.

$$\frac{v\ R\ w}{v\ ecr(R)\ w} \qquad \frac{-}{v\ ecr(R)\ v} \qquad \frac{v\ ecr(R)\ w}{w\ ecr(R)\ v} \qquad \frac{u\ ecr(R)\ v\ ecr(R)\ w}{u\ ecr(R)\ w}$$

$$\frac{v_1\ ecr(R)\ w_1 \qquad v_2\ c(R)\ w_2}{min(v_1, v_2)\ ecr(R)\ min(w_1, w_2)} \qquad \frac{v\ ecr(R)\ w \qquad r \in \mathbb{R} \cup \{\infty\}}{r + v\ ecr(R)\ r + w}$$

For an example of bisimulation up to congruence consider the relation $R = \{(x, u),\ (min(2+y, 3+z), 2+u)\}$ and the weighted automaton depicted in (3). To prove that R is a bisimulation up to congruence we only have to show that $(min(4 + x, 5 + y), 4 + u) \in ecr(R)$:

$$\begin{aligned}
min(4 + x, 5 + y)\ ecr(R)\ & min(4 + u, 5 + y) && ((x, u) \in R)\\
ecr(R)\ & min(4 + y, 5 + z), 5 + y) && ((min(2 + y, 3 + z), 2 + u) \in R)\\
=\ & 2 + min(2 + y, 3 + z) &&\\
ecr(R)\ & 4 + u && ((min(2 + y, 3 + z), 2 + u) \in R)
\end{aligned}$$

Note that R is not a bisimulation up to context, since $(min(4 + x, 5 + y), 4 + u) \notin c(R)$. Here transitivity is really needed.

Up-to-union-context-and-equivalence. A bisimulation up to $e \circ c \circ u_S$ is called a *bisimulation up to S-union, context and equivalence.* This is an important extension of bisimulation up to context because the equivalence closure allows us to relate derivatives of R using $c(R \cup S)$ in "multiple steps".

Example 12. Recall the operations of shuffle product and inverse from Example 5 and suppose that we want to prove that the inverse operation is really the inverse of shuffle product, that is, $\sigma \otimes \sigma^{-1} \sim 1$ for all $\sigma \in T_{wf}(\mathbb{R}^\omega)$ such that $\sigma_0 \neq 0$. Suppose we are given that $\otimes$ is associative and commutative (so $\sigma \otimes \tau \sim \tau \otimes \sigma$, etc.) and that $\sigma + (-\sigma) \sim \mathsf{0}$ (note that these assumptions actually hold in general (Rutten, 2003)). Let

$$R = \{(\sigma \otimes \sigma^{-1}, 1) \mid \sigma \in T_{wf}(\mathbb{R}^\omega), \sigma_0 \neq 0\}\,.$$

We can now establish that R is a bisimulation up to $\sim$-union, context and equivalence. First consider the initial values:

$$(\sigma \otimes \sigma^{-1})_0 = \sigma_0 \times (\sigma^{-1})_0 = \sigma_0 \times (\sigma_0)^{-1} = 1 = \mathsf{1}_0$$

Next, we relate the derivatives by $e(c(R \cup \sim))$:

$$\begin{aligned}
(\sigma \otimes \sigma^{-1})' &= \sigma' \otimes \sigma^{-1} + \sigma \otimes (\sigma^{-1})'\\
&= \sigma' \times \sigma^{-1} + \sigma \otimes (-\sigma' \otimes (\sigma^{-1} \otimes \sigma^{-1}))\\
&t(c(\sim))\ (\sigma' \otimes \sigma^{-1}) + (-(\sigma' \otimes \sigma^{-1}) \otimes (\sigma \otimes \sigma^{-1}))\\
&c(R \cup \sim)\ (\sigma' \otimes \sigma^{-1}) + (-(\sigma' \otimes \sigma^{-1}) \otimes 1)\\
&t(c(\sim))\ \mathsf{0} = \mathsf{1}'
\end{aligned}$$

where $t(c(\sim))$ denotes the transitive closure of $c(\sim)$; in the above we apply multiple substitutions of terms for bisimilar terms. Since $t(c(\sim)) \subseteq e(c(R \cup \sim))$ and $c(R \cup \sim) \subseteq e(c(R \cup \sim))$ we may conclude that R is a bisimulation up to $\sim$-union, context, and equivalence. Notice that R is not a bisimulation; establishing that it is a bisimulation-up-to is much easier than finding a bisimulation which contains R.

In the above, rather than $c(R \cup \sim)$ we could have used $c(r(R))$. Moreover, since in this example $\sim\, = t(c(\sim))$, the above is also an example of *bisimulation up to context, reflexivity and bisimilarity*, that is, a bisimulation up to $b \circ c \circ r$. (Any bisimulation up to context, reflexivity and bisimilarity is also a bisimulation up to $\sim$-union, context and equivalence.)

4. An algebra of enhancements

The above examples illustrate the large range of enhancements available for bisimilarity, and the need to combine such enhancements. For instance, up-to-union is rarely used on its own: it needs to be combined with up-to-equivalence or up-to-context. However, the soundness of such a combination does not necessarily follow from the soundness of its basic constituents, and it can be hard to prove it from scratch. This calls for a theory of enhancements which would allow one to freely combine them. Such a theory was developed at the rather abstract level of complete lattices (Pous, 2007; Pous and Sangiorgi, 2012). We rephrase it here at the level of binary relations, for the sake of clarity. We instantiate it in the following sections to obtain our general theory of coalgebraic bisimulations and behavioural equivalences up-to.

Let b be a monotone function on binary relations. By the Knaster-Tarski theorem b has a greatest fixpoint, denoted by $\mathsf{gfp}(b)$, which is also the greatest post-fixpoint: $\mathsf{gfp}(b) = \bigcup\{R \mid R \subseteq b(R)\}$. The intuition is that by choosing b in an appropriate way, $\mathsf{gfp}(b)$ will be the desired notion of bisimilarity. This motivates the following terminology:

— A *b-simulation* is a relation R such that $R \subseteq b(R)$.
— *b-similarity* is the greatest b-simulation, i.e., $\mathsf{gfp}(b)$.

The bisimulation proof method can now be rephrased as follows: to prove that some states x, y are b-similar it suffices to exhibit a b-simulation R such that $x\ R\ y$. Enhancements of the bisimulation proof method allow one to weaken the requirement that R is a b-simulation: rather than checking $R \subseteq b(R)$, we would like to check $R \subseteq b(S)$ for a relation S which is possibly larger than R. The key idea consists in using a function f to obtain this larger relation out of R: $S = f(R)$.

Definition 3. Let f be a function on binary relations.

— A *b-simulation up to f* is a relation R such that $R \subseteq b(f(R))$.
— f is *b-sound* if all b-simulations up to f are contained in b-similarity.
— f is *b-compatible* if it is monotone and $f \circ b \subseteq b \circ f$.

The notion of b-compatible function is introduced to get around the fact that b-sound functions cannot easily be composed: b-compatible functions are b-sound and they enjoy nice compositionality properties:

Theorem 1. All b-compatible functions are b-sound.

Proof. Suppose that R is a bisimulation up to f, i.e., that $R \subseteq b(f(R))$. Using compatibility of f and by a simple induction on n, we get $\forall n, f^n(R) \subseteq b(f^{n+1}(R))$. Therefore, we have

$$\bigcup_n f^n(R) \subseteq \bigcup_n b(f^n(R)) \subseteq b\left(\bigcup_n f^n(R)\right) \ .$$

(The second inclusion holds by monotonicity of b.) In other words, $f^\omega(R) = \bigcup_n f^n(R)$ is a b-simulation. This latter relation trivially contains R, by taking $n = 0$, so that we can conclude that R is contained in b-similarity. □

Proposition 1. The following functions are b-compatible:

1 id — the identity function;
2 con_S — the constant-to-S function, for any b-simulation S;
3 $f \circ g$ for any b-compatible functions f and g;
4 $\bigcup F$ for any set F of b-compatible functions.

The last two items allow one to freely combine b-compatible functions using functional composition and pointwise union. There is a third way of combining two functions f, g on relations, using relational composition: $f \bullet g(R) = f(R) \circ g(R)$. This composition operator does *not* always preserve b-compatible functions; the following lemma gives a sufficient condition:

Proposition 2. If b satisfies the following condition:

$$\text{for all relations } R, S, \quad b(R) \circ b(S) \subseteq b(R \circ S) \ , \tag{†}$$

then $f \bullet g$ is b-compatible for all b-compatible functions f and g.

We show in the following section that for all functors F there exists a function φ such that the F-bisimulations are the φ-simulations. Any such function is monotone and the property (†) holds iff the functor F preserves weak pullbacks.

We conclude this section with two lemmas which will be useful in the sequel: the first one gives an alternative characterisation of b-compatible functions; the second one shows that b-similarity is closed under any b-compatible function.

Lemma 1. A monotone function f is b-compatible iff for all relations R, S, $R \subseteq b(S)$ implies $f(R) \subseteq b(f(S))$.

Lemma 2. For all b-compatible functions f, $f(\mathsf{gfp}(b)) \subseteq \mathsf{gfp}(b)$.

5. Bisimulation and φ-simulation

In this section we show how to characterize bisimulation and bisimulation-up-to in terms of monotone functions. This allows us to study bisimulation-up-to, as introduced in Section 3, in terms of the abstract framework of Section 4.

Let (X, α) be an F-coalgebra. We define an endofunction φ_α on the complete lattice of relations on X ordered by inclusion $(\mathcal{P}(X \times X), \subseteq)$ as follows (Rutten, 1998; Hermida and Jacobs, 1998):

$$\begin{aligned}\varphi_\alpha(R) &= \{(x,y) \mid (\alpha(x), \alpha(y)) \in F(\pi_1^R)^{-1} \circ F(\pi_2^R)\} \\ &= \{(x,y) \mid \exists z \in FR \text{ s.t. } F(\pi_1^R)(z) = \alpha(x) \text{ and } F(\pi_2^R)(z) = \alpha(y)\}\end{aligned}$$

We write φ instead of φ_α if α is clear from the context.

Example 13. We describe φ for several concrete types of systems.

1 For deterministic automata, φ corresponds to the classical functional exploited by the Hopcroft minimization algorithm:

$$\varphi(R) = \{(x,y) \mid o(x) = o(y) \text{ and, for all } a \in A,\ (t(x)(a), t(y)(a)) \in R\}$$

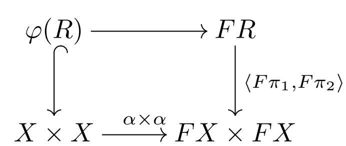
2 In the case of labelled transition systems, φ corresponds to the well-known functional of bisimilarity (e.g., (Sangiorgi, 2012)):

$$\begin{aligned}\varphi(R) = \{(x,y) \mid &\text{ if } x \xrightarrow{a} x' \text{ then there exists } y' \text{ s.t. } y \xrightarrow{a} y' \text{ and } x'Ry', \text{ and} \\ &\text{if } y \xrightarrow{a} y' \text{ then there exists } x' \text{ s.t. } x \xrightarrow{a} x' \text{ and } x'Ry'\}\end{aligned}$$

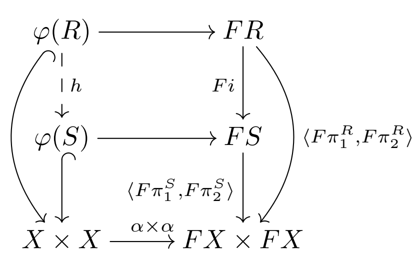
3 For stream systems, i.e., coalgebras for the functor $FX = \mathbb{R} \times X$, φ instantiates to $\varphi(R) = \{(x,y) \mid x_0 = y_0 \text{ and } x'Ry'\}$.

Lemma 3. For any coalgebra (X, α): φ_α is monotone.

Proof. Notice that for any R, $\varphi(R)$ can be characterized as a pullback (cf. (Staton, 2011)):



Suppose $R \subseteq S$; denote the corresponding inclusion map by $i\colon R \hookrightarrow S$. Then one can show that the following holds: $\langle F\pi_1^R, F\pi_2^R\rangle = \langle F\pi_1^S, F\pi_2^S\rangle \circ Fi$. Because of that and the fact that $\varphi(R)$ and $\varphi(S)$ are pullbacks the following diagram commutes (not including h, which will be introduced below):



By the fact that $\varphi(S)$ is a pullback, the map h, making the above diagram commute, exists; commutativity of the triangle on the left shows that h is an inclusion map, i.e., $\varphi(R) \subseteq \varphi(S)$. $\square$

The following lemma establishes the connection of the above monotone functions to bisimulation and bisimulation-up-to.

Lemma 4. For any coalgebra (X, α) and for any relations $R, S \subseteq X \times X$:

$$R \subseteq \varphi_\alpha(S) \text{ iff } R \rightarrowtail S.$$

Proof. Follows easily from the second characterization of φ as given above. □

From the above lemma we directly obtain the following known result (Rutten, 1998):

Corollary 1. For any coalgebra (X, α): R is a bisimulation iff $R \subseteq \varphi_\alpha(R)$.

In other words, a φ-simulation (Section 4) is the same as a bisimulation. Thus, the greatest fixpoint of φ is precisely $\sim$. Lemma 4 also establishes a tight connection between bisimulation-up-to and φ-simulation-up-to.

Corollary 2. Let $f\colon \mathcal{P}(X \times X) \to \mathcal{P}(X \times X)$ be monotone. For any coalgebra (X, α):

1 $R \subseteq X \times X$ is a bisimulation up to f iff it is a φ_α-simulation up to f;
2 If f is φ_α-compatible (Def. 3), then f is sound (Def. 2).

Proof.

1 Follows directly from Lemma 4: $R \subseteq \varphi(f(R))$ iff $R \rightarrowtail f(R)$.
2 Suppose $R \rightarrowtail f(R)$; then $R \subseteq \varphi(f(R))$ by (1). If f is φ-compatible, then by Theorem 1 it is φ-sound. So $R \subseteq \mathsf{gfp}(\varphi) = \sim$.

□

Via the above results we can apply the general theory of Section 4 to reason about coalgebraic bisimulation-up-to.

6. Compatibility

In this section we study the φ-compatibility of the instances of bisimulation-up-to introduced in Section 3. By proving the compatibility of a function f we obtain the soundness of bisimulation up to f and we can compose it to other compatible functions, knowing that the result is again compatible.

Theorem 2. Let (X, α) be a coalgebra for a functor F. The following functions are φ_α-compatible:

1 r — the reflexive closure;
2 s — the symmetric closure;
3 u_S — union with S (for a bisimulation S);

If F preserves weak pullbacks, then the following are φ_α-compatible:

4. t — the transitive closure;
5. e — the equivalence closure;
6. b — bisimilarity;
7. $e \circ u_S$ — S-union and equivalence (for a bisimulation S).

The functions exploiting the contextual closure c will be considered later (Section 6.1). We will prove the above theorem below. But first, notice that for the compatibility of several functions we require the functor to preserve weak pullbacks. Indeed, bisimulation up to bisimilarity and bisimulation up to equivalence are not sound in general, and consequently not compatible either. This is illustrated by the following example, which is strongly inspired by an example from (Aczel and Mendler, 1989).

Example 14. Define the functor $F\colon \mathsf{Set} \to \mathsf{Set}$ as

$$FX = \{(x_1, x_2, x_3) \in X^3 \mid |\{x_1, x_2, x_3\}| \leq 2\}$$
$$F(f)(x_1, x_2, x_3) = (f(x_1), f(x_2), f(x_3))$$

Consider the F-coalgebra with states $X = \{0, 1, 2, \tilde{0}, \tilde{1}\}$ and transition structure

$$\begin{array}{lll} 0 \mapsto (0,1,0) & \tilde{0} \mapsto (0,0,0) & 2 \mapsto (2,2,2) \\ 1 \mapsto (0,0,1) & \tilde{1} \mapsto (1,1,1) & \end{array}$$

Then $0 \not\sim 1$. To see this, note that in order for the pair $(0,1)$ to be contained in a bisimulation R, there must be a transition structure on this relation which maps $(0,1)$ to $((0,0),(1,0),(0,1))$. But this triple can not be in FR, because it contains three different elements. However, it is easy to show that $0 \sim 2$ and $1 \sim 2$: the relation $\{(0,2),(1,2)\}$ is a bisimulation.

Now consider the relation $S = \{(\tilde{0}, \tilde{1}), (2,2)\}$. S is not a bisimulation, since for that there should be a function from S to FS mapping the elements as follows:

$$(\tilde{0}, \tilde{1}) \mapsto ((0,1),(0,1),(0,1)) \qquad (2,2) \mapsto ((2,2),(2,2),(2,2))$$

and neither $((0,1),(0,1),(0,1))$ nor $((2,2),(2,2),(2,2))$ are contained in FS. However, since $0 \sim 2 \; S \; 2 \sim 1$ (and $2 \sim S \sim 2$), they *are* contained in $F(\sim S \sim)$; so S is a bisimulation up to bisimilarity. Thus if up-to-bisimilarity is sound, then $S \subseteq \sim$ so $0 \sim 1$, which is a contradiction.

Below we will show that if the functor preserves weak pullbacks then φ-compatible functions are closed under the operation $\bullet$ (defined in Section 4), which allows to prove items 4,5,6 and 7. In order to proceed we recall some fundamental results relating preservation of weak pullbacks to composition of relations.

Theorem 3. Let $F\colon \mathsf{Set} \to \mathsf{Set}$ be a functor. The following are equivalent:

1 F preserves weak pullbacks.

2 $\tilde{F}\colon \mathsf{Rel} \to \mathsf{Rel}$, defined as

$$\tilde{F}X = FX$$
$$\tilde{F}R = F(\pi_1^R)^{-1} \circ F(\pi_2^R)$$

is a functor (i.e., $\tilde{F}$ preserves composition).

3 The composition of two F-bisimulations is again a bisimulation.

The equivalence of (1) and (2) is due to Trnková (Trnková, 1980). Notice that φ is in fact defined in terms of the action of $\tilde{F}$ on relations: $\varphi_\alpha(R) = \{(x,y) \mid (\alpha(x), \alpha(y)) \in$

$\tilde{F}R\}$ (Rutten, 1998). Rutten (Rutten, 2000) established the implication from (1) to (3). The reverse implication is due to Gumm and Schröder (Gumm and Schröder, 2000). Their result is based on bisimulations on two coalgebras (X, α) and (Y, β) but for our notion of bisimulation (restricted to one coalgebra) the implication still holds, as we show in Appendix A.

Using Theorem 3 we show that preservation of weak pullbacks coincides precisely with the property (†) of Section 4. Then by Proposition 2 we obtain that φ-compatible functions are closed under $\bullet$ in the case of a functor which preserves weak pullbacks.

Proposition 3. *F preserves weak pullbacks iff for any F-coalgebra (X, α), φ_α satisfies (†), i.e., for all relations $R, S\colon \varphi_\alpha(R) \circ \varphi_\alpha(S) \subseteq \varphi_\alpha(R \circ S)$.*

Proof. Suppose F preserves weak pullbacks. Let (X, α) be an F-coalgebra, $R, S \subseteq X \times X$ relations, and $(x, z) \in \varphi_\alpha(R) \circ \varphi_\alpha(S)$, so there is some y such that $(x, y) \in \varphi_\alpha(R)$ and $(y, z) \in \varphi_\alpha(S)$. Then $(\alpha(x), \alpha(y)) \in \tilde{F}(R)$ and $(\alpha(y), \alpha(z)) \in \tilde{F}(S)$, so $(\alpha(x), \alpha(z)) \in \tilde{F}(R) \circ \tilde{F}(S)$. But by assumption and Theorem 3 $\tilde{F}$ is functorial, so $\tilde{F}(R) \circ \tilde{F}(S) = \tilde{F}(R \circ S)$. Consequently $(x, z) \in \varphi_\alpha(R \circ S)$ as desired.

Conversely, suppose that (†) holds; then by Proposition 2, compatible functions are closed under $\bullet$. Let R, S be bisimulations, so con_R and con_S are compatible by Proposition 1. By assumption $con_R \bullet con_S$ is compatible, so by Lemma 1 we have $R \circ S \subseteq \varphi(R \circ S)$. Now by Corollary 1, $R \circ S$ is a bisimulation. From Theorem 3 we conclude that F preserves weak pullbacks. □

The inverse function is compatible as well, which will be useful to prove compatibility of the equivalence closure:

Proposition 4. *Let (X, α) be a coalgebra. The inverse map $i(R) = R^{-1}$ is φ_α-compatible.*

Proof. Suppose $R \subseteq \varphi(S)$, and let $(x, y) \in R^{-1}$, so $(y, x) \in R$. Then $(\alpha(y), \alpha(x)) \in (F(\pi_1^S))^{-1} \circ F(\pi_2^S)$. But $\pi_1^S = \pi_2^{S^{-1}}$ and $\pi_2^S = \pi_1^{S^{-1}}$, so $(\alpha(y), \alpha(x)) \in (F(\pi_2^{S^{-1}}))^{-1} \circ F(\pi_1^{S^{-1}})$. Consequently

$$(\alpha(x), \alpha(y)) \in ((F(\pi_2^{S^{-1}}))^{-1} \circ F(\pi_1^{S^{-1}}))^{-1} = (F(\pi_1^{S^{-1}}))^{-1} \circ F(\pi_2^{S^{-1}})$$

so $x, y \in \varphi(S^{-1})$. By Lemma 1, i is φ-compatible. □

We proceed with the proof of Theorem 2. Below we use the general compatibility results of Proposition 1 without further reference.

1 The identity relation Δ is a bisimulation (Rutten, 2000) and thus, by Proposition 1, con_Δ is compatible and thus $r = id \cup con_\Delta$ is compatible.

2 The inverse function is compatible by Proposition 4. Compatibility of $s = id \cup i$ then follows directly.

3 $u_S = id \cup con_S$ is compatible for a bisimulation S, since con_S is compatible.

4 First, we define $(-)^n$ as $(-)^1 = id$ and $(-)^{n+1} = id \bullet (-)^n$. We prove that for all $n \geq 1$, $(-)^n$ is compatible, by induction on n. For the base case, notice that id is compatible. Now suppose $(-)^n$ is compatible. Then, by Proposition 3 and Proposition 2, $(-)^{n+1} = id \bullet (-)^n$ is also compatible. Now notice that $t = \bigcup_{n \geq 1} (-)^n$; so the function t is

the (infinite) union of compatible functions, and consequently by Proposition 1 it is compatible.

5 $e = t \circ s \circ r$ is compatible, since r, s, and t are compatible.

6 $con_\sim$ is compatible since $\sim$ is a bisimulation. By Proposition 3 and Proposition 2, the function $b = con_\sim \bullet id \bullet con_\sim$ is compatible.

7 $e \circ u_S$ is compatible since e and u_S are compatible.

6.1. *Bisimulation up to context*

In order to define the contextual closure c, we need a T-algebra $\beta\colon TX \to X$ on the states of an F-coalgebra (X, α). For compatibility of c one might expect that it is enough to know that bisimilarity is a congruence with respect to this algebra; however, it is known that this is not even enough for soundness of bisimulation up to context (Pous and Sangiorgi, 2012). As we will show below, in order to prove that c is compatible, it is sufficient to assume that (X, β, α) is a λ-*bialgebra* for some natural transformation $\lambda\colon TF \Rightarrow FT$. Such a natural transformations is typically called a *distributive law* of the functor T over the functor F, but we will not use that term to avoid confusion with the stronger notion of distributive laws of *monads* over functors. We refer to (Klin, 2011) for a nice overview on λ-bialgebras and report their formal definition below.

Definition 4. Let (T, μ, η) be a monad and F an endofunctor, both on Set.

— An (F, T)-*bialgebra* is a triple (X, β, α) where X is a set, (X, β) is a T-algebra and (X, α) is an F-coalgebra.

— Given a natural transformation $\lambda\colon TF \Rightarrow FT$ we say (X, β, α) is a λ-*bialgebra* if $\alpha \circ \beta = F\beta \circ \lambda_X \circ T\alpha$.

Below, we give an intuition on λ-bialgebras by exhibiting some examples.

Example 15. The coalgebra $(\mathbb{R}_\omega^X, \langle o^\sharp, t^\sharp \rangle)$ induced by a weighted automaton (Example 3) is a λ-bialgebra, where λ is a certain natural transformation law of the (underlying functor of the) free vector space monad $\mathbb{R}_\omega^-$ over the functor $FX = \mathbb{R} \times X^A$.

Recall the coalgebra $(\mathcal{P}(A^*), \langle o, t \rangle)$ introduced in Example 10, where $\mathcal{P}(A^*)$ is the set of languages over an alphabet A. The operations of union, concatenation and Kleene star induce an algebra on $\mathcal{P}(A^*)$. Together, this algebra and coalgebra form a λ-bialgebra (Jacobs, 2006).

Other important examples include certain types of process algebras and stream coalgebras induced by behavioural differential equations, but these involve a technicality treated in Section 6.2.

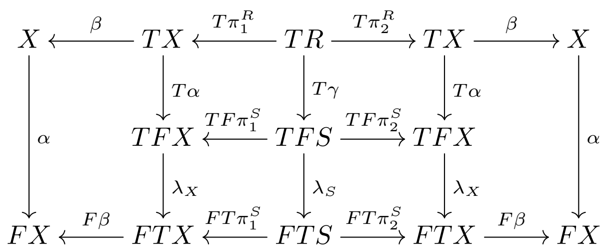
Theorem 4. Let (X, β, α) be a λ-bialgebra for $\lambda\colon TF \Rightarrow FT$. The contextual closure function c_β is φ_α-compatible. If F preserves weak pullbacks then the following are φ_α-compatible as well:

1 $e \circ c_\beta \circ r$ — congruence;

2 $e \circ c_\beta \circ u_S$ — context, S-union and equivalence;

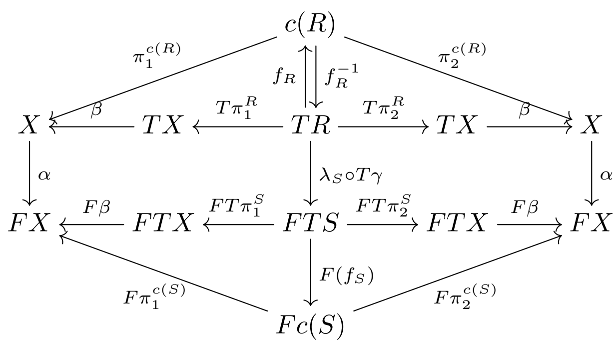
3 $b \circ c_\beta \circ r$ — context, reflexivity and bisimilarity.

Proof. We prove compatibility of c; then items 1,2 and 3 follow directly from Theorem 2 and Proposition 1. Suppose $R \subseteq \varphi(S)$ for some R and S. Consider the following diagram:



The existence of γ and commutativity of the upper squares follow from Lemma 4 and an application of T. The lower squares commute by naturality. Finally the outer rectangles commute since (X, β, α) is a λ-bialgebra.

Let $f_R : TR \to c(R)$ be the corestriction of $\langle \beta \circ T\pi_1^R, \beta \circ T\pi_2^R \rangle : TR \to X \times X$ to its range, so that $f_R[TR] = c(R)$. Let $f_S : TS \to c(S)$ be defined analogously, and take f_R^{-1} to be any right inverse of f_R. Then the following diagram commutes:



So $c(R)$ progresses to $c(S)$, and consequently $c(R) \subseteq \varphi(c(S))$ by Lemma 4. By Lemma 1 we conclude that c is φ-compatible. $\square$

Remark 1. The greatest bisimulation on a λ-bialgebra is closed under the algebraic operations. This was first shown by Turi and Plotkin (Turi and Plotkin, 1997) under the assumption that F preserves weak pullbacks; Bartels (Bartels, 2004) showed that this assumption is not necessary. We obtain the same result as a direct consequence of the above theorem and Lemma 2.

6.2. *Coalgebras for copointed functors*

As was first shown by Turi and Plotkin (Turi and Plotkin, 1997) one can obtain process algebras whose operational rules conform to the GSOS format (Bloom et al., 1995) as λ-bialgebras. Every GSOS specification over some signature Σ induces an operational model

$$T_\Sigma T_\Sigma \emptyset \xrightarrow{\beta} T_\Sigma \emptyset \xrightarrow{\alpha} \mathcal{P}_f(A \times T_\Sigma \emptyset)$$

on closed terms, where β is the initial algebra and α is a transition structure induced by the specification. However, in several concrete cases there is no natural transformation λ such that $(T_\Sigma\emptyset, \beta, \alpha)$ is a λ-bialgebra. Instead, one needs to consider the bialgebra $(T_\Sigma\emptyset, \beta, \langle \alpha, id\rangle)$; $\langle \alpha, id\rangle$ now is a coalgebra for the so-called *cofree copointed endofunctor* $\mathcal{P}_f(A \times Id) \times Id$ (see, e.g., (Klin, 2011)). Analogously, any (non-partial)[†] specification of operations on streams in terms of behavioural differential equations (Rutten, 2003) corresponds to a natural transformation involving not the functor $FX = \mathbb{R} \times X$ but the functor $F \times Id$. Yet another example is given by the coalgebraic characterization of context-free grammars as in (Winter et al., 2011); as discussed in (Bonsangue et al., 2013), this construction involves a bialgebra which is a λ-bialgebra when the coalgebra is paired with the identity function.

All of the above are examples of bialgebras (X, β, α) such that $(X, \beta, \langle \alpha, id\rangle)$ is a λ-bialgebra. In such cases one wants to consider bisimulation(-up-to) on the coalgebra α and not on $\langle \alpha, id\rangle$. However, while Theorem 4 gives us $\varphi_{\langle \alpha, id\rangle}$-compatibility of the contextual closure c_β, it does *not* provide φ_α-compatibility. For the convenience of the reader, we recall a counterexample from (Pous and Sangiorgi, 2012).

Example 16 ((Pous and Sangiorgi, 2012)). Consider the following specification of the prefix and the replication operation on labelled transition systems:

$$\frac{}{a.x \xrightarrow{a} x} \qquad\qquad \frac{x \xrightarrow{a} x'}{!x \xrightarrow{a} !x \mid x'}$$

together with the standard definition of the parallel operator $x \mid y$. While this is arguably not the best way to specify replication in the context of CCS (Pous and Sangiorgi, 2012) it suffices for our purposes. This specification induces a coalgebra $\alpha \colon T\emptyset \to \mathcal{P}_f(A \times T\emptyset)$ on terms. Now consider the relations $R = \{(!a.b, !a.c)\}$ and $S = \{(!a.b \mid b, !a.c \mid c)\}$. Then R progresses to S, but $c(R)$ does not progress to $c(S)$. For example, $d.!a.b\ c(R)\ d.!a.c$ but $!a.b$ and $!a.c$ are not related by $c(S)$. Thus, by Lemma 1 the contextual closure c is not φ_α-compatible.

The solution is to define a different function φ' on the complete lattice of relations of a coalgebra α as

$$\varphi'(R) = \varphi(R) \cap R\,.$$

Clearly φ' is monotone, and the greatest fixpoint of φ' is bisimilarity. In terms of progression, we have $R \subseteq \varphi'(S)$ if and only if R progresses to S and $R \subseteq S$. Thus if R progresses to $f(R)$ for a function satisfying $R \subseteq f(R)$, then $R \subseteq \varphi'(f(R))$. But notice that for every function f considered in Theorem 2 and Theorem 4 we have $R \subseteq f(R)$ for any R. Thus to obtain soundness, we only need to show that all these functions are φ'-compatible.

Theorem 5. Theorem 2 holds when every occurence of φ is replaced by φ'. Moreover, if $(X, \beta, \langle \alpha, id\rangle)$ is a λ-bialgebra then Theorem 4 holds for φ'_α as well.

[†] The partial specification of Example 5 can be completed by fixing the initial value of 0^{-1} to some arbitrary constant.

Proof. If $R \subseteq f(R)$ then φ-compatibility of f implies φ' compatibility of f, and thus the first part of the claim follows from Theorem 2. By Theorem 4 we obtain $\varphi'_{\langle\alpha,id\rangle}$-compatibility of the contextual closure function; and one can readily show that $\varphi'_{\langle\alpha,id\rangle} = \varphi'_\alpha$, obtaining the desired result. ☐

7. Behavioural equivalence-up-to

Whenever the functor F does not preserve weak pullbacks (as it is the case, for instance, with certain types of weighted transition systems (Klin, 2009)) one can consider *behavioural equivalence*, rather than bisimilarity.

Definition 5. For a coalgebra $\alpha \colon X \to FX$ and relations $R, S \subseteq X \times X$, we say R *progresses to* S (with respect to behavioural equivalence), denoted $R \rightsquigarrow S$, if the following diagram commutes:

$$R \underset{\pi_2}{\overset{\pi_1}{\rightrightarrows}} X \xrightarrow{\ \alpha\ } FX \xrightarrow{\ Fq\ } F(X/e(S))$$

where q is the quotient map of $e(S)$. If $R \rightsquigarrow R$ then R is called a *behavioural equivalence*.

Equivalently, R progresses to S if $Fq \circ \alpha$ factors through the quotient map of $e(R)$. In particular, R is a behavioural equivalence iff the quotient map of R is a coalgebra homomorphism.

The largest behavioural equivalence is denoted by $\approx$. An equivalent definition of $\approx$ is: $x \approx y$ iff there exists some homomorphism f from (X, α) to some coalgebra (Y, β) such that $f(s) = f(t)$ (Gumm, 1999; Rot et al., 2013a).

The relation R of Example 6 is a behavioural equivalence: note that, intuitively, behavioural equivalences are implicitly "up-to-equivalence", since the arriving states can be related by $e(R)$. Note also that in (Aczel and Mendler, 1989) behavioural equivalences are called *pre-congruences*.

Definition 6. If $R \rightsquigarrow f(R)$ for a function $f \colon \mathcal{P}(X \times X) \to \mathcal{P}(X \times X)$ then we say R is a *behavioural equivalence up to f*. We say that f is *sound* (w.r.t. behavioural equivalence) if $R \subseteq\, \approx$ for all R such that $R \rightsquigarrow f(R)$.

We proceed with a similar development as for bisimulation-up-to: first, we characterize behavioural equivalence as a fixed point of a monotone function, as done already in (Aczel and Mendler, 1989). Define the function $\psi_\alpha \colon \mathcal{P}(X \times X) \to \mathcal{P}(X \times X)$ as

$$\psi_\alpha(R) = \{(x, y) \mid Fq \circ \alpha(x) = Fq \circ \alpha(y)\}$$

i.e., as the kernel of $Fq \circ \alpha$, where $q \colon X \to X/e(R)$ is the quotient map of $e(R)$. Notice that we can also define q as the coequalizer of R and its projection maps, and ψ_α as the pullback of $Fq \circ \alpha$ along itself.

Lemma 5. For any coalgebra (X, α): ψ_α is monotone.

The correspondence between progression and functions ψ is given by the following lemma:

Lemma 6. For any coalgebra (X, α) and for any relations $R, S \subseteq X \times X$:

$$R \subseteq \psi(S) \text{ iff } R \rightsquigarrow S\,.$$

Consequently, behavioural equivalence up to any ψ-compatible function is sound. Unfortunately, the property (†) does *not* hold for ψ, that is, in general it does not hold that $\psi(R) \circ \psi(S) \subseteq \psi(R \circ S)$. This is shown by the following example:

Example 17. Consider the identity functor $FX = X$ and the F-coalgebra with states $\{x, y\}$ and transitions $x \mapsto x$ and $y \mapsto y$. Let $R = \{(x, y)\}$. Then $\psi(R) = \{(x, x), (y, y), (x, y)\}$ and $\psi(\emptyset) = \{(x, x), (y, y)\}$. Now $\psi(R) \circ \psi(\emptyset) = \{(x, x), (y, y), (x, y)\}$, whereas $\psi(R \circ \emptyset) = \psi(\emptyset) = \{(x, x), (y, y)\}$. So $\psi(R) \circ \psi(\emptyset)$ is not included in $\psi(R \circ \emptyset)$.

This motivates to prove the compatibility of the equivalence closure e directly, which is in fact quite easy in the case of behavioural equivalence.

Theorem 6. Let (X, α) be any coalgebra. The following are ψ_α-compatible:

1 r — the reflexive closure;
2 e — the equivalence closure;
3 u_S — union with S (for a behavioural equivalence S);
4 *beh* — behavioural equivalence ($\mathit{beh}(R) = \mathit{con}_\approx \bullet \mathit{id} \bullet \mathit{con}_\approx$).

Proof. Items 1, 3 and 4 are analogous to the case of φ-compatibility in Theorem 2. We proceed with the compatibility of the equivalence closure. First, notice that $e \circ \psi = \psi$ since $\psi(R)$ is an equivalence relation for any relation R. Second, since $e(R) = e(e(R))$ for any R, the quotient maps in the definition of $\psi(R)$ and $\psi(e(R))$ are equal, so $\psi(R) = \psi(e(R))$. Thus $e \circ \psi = \psi = \psi \circ e$. □

Notice that the ψ-compatibility of the equivalence closure does not require any assumptions on the functor.

We proceed with the compatibility of behavioural equivalence up to context. As we will see, in the proof we need a certain coequalizer to be a T-algebra homomorphism; more precisely, we will need the forgetful functor U from the category of T-algebras to Set to create *reflexive* coequalizers[‡]. This does not hold for abitrary monads, see (Adámek et al., 2000, page 538). Our solution is to restrict to *finitary* monads, that is, monads where the underlying functor preserves filtered colimits. For a free monad over a signature, this means that each operation has finite arity (but there may be infinitely many operations). Other examples include the free vector space monad and the finite powerset monad.

Lemma 7. If T is a finitary (Set) monad then it preserves reflexive coequalizers.

This allows to deduce the following result, the proof of which is routine:

Lemma 8. Let (X, β) be a T-algebra, $f, g\colon R \to X$ a pair of algebra homomorphisms which is reflexive in Set, T a finitary monad, and $q\colon X \to X'$ the coequalizer of $(TR, \alpha \circ$

[‡] A reflexive coequalizer is a coequalizer of a reflexive pair, that is, a pair of functions $f, g\colon X \to Y$ such that there is a function $h\colon Y \to X$ with $f \circ h = g \circ h = id$.

$Tf, \alpha \circ Tg)$. Then there is a T-algebra on X' such that q is an algebra homomorphism from (X, β) to (X', β').

Now we are ready to prove compatibility of up-to-context. However, notice that the above results depend on the relations being reflexive; thus we will directly prove compatibility of $c \circ r$ instead of c.

Theorem 7. Let (X, β, α) be a λ-bialgebra for a distributive law $\lambda\colon TF \Rightarrow FT$, where T is a finitary monad. The following are ψ_α-compatible:

1 $c_\beta \circ r$ — contextual closure;
2 $e \circ c_\beta \circ r$ — congruence;
3 $e \circ c_\beta \circ r \circ u_S$ — context, S-union, reflexivity and equivalence;
4 $b \circ c_\beta \circ r$ — context, reflexivity and bisimilarity.

Proof. We only need to prove ψ-compatibility of $c_\beta \circ r$. Suppose $R \subseteq \psi(S)$ for some relations $R, S \subseteq X \times X$. Then clearly $r(R) \subseteq \psi(r(S))$, and since $r(S) \subseteq c_\beta(r(S))$ we have $r(R) \subseteq \psi(c_\beta(r(S)))$. By Lemma 6 then the diagram

$$r(R) \underset{\pi_2}{\overset{\pi_1}{\rightrightarrows}} X \xrightarrow{\alpha} FX \xrightarrow{Fq} FX' \tag{5}$$

commutes, where q is the coequalizer of $e(c_\beta(r(S)))$. But using the definition of c_β we find that q is also the coequalizer of $(T(r(S)), \beta \circ T\pi_1, \beta \circ T\pi_2)$, which is a reflexive pair (in Set), and consequently by Lemma 8 we see that q is an algebra homomorphism, i.e., the following diagram commutes:

$$\begin{array}{ccc} TX & \xrightarrow{Tq} & TX' \\ \downarrow{\scriptstyle \beta} & & \downarrow{\scriptstyle \beta'} \\ X & \xrightarrow{q} & X' \end{array} \tag{6}$$

for some β'. Now consider the following diagram:

$$\begin{array}{ccccccc} T(r(R)) & \underset{T\pi_2}{\overset{T\pi_1}{\rightrightarrows}} & TX & \xrightarrow{T\alpha} & TFX & \xrightarrow{TFq} & TFX' \\ & & \Big\downarrow & & \downarrow{\scriptstyle \lambda_X} & & \downarrow{\scriptstyle \lambda_{X'}} \\ & & {\scriptstyle \beta} & & FTX & \xrightarrow{FTq} & FT' \\ & & \Big\downarrow & & \downarrow{\scriptstyle F\beta} & & \downarrow{\scriptstyle F\beta'} \\ & & X & \xrightarrow{\alpha} & FX & \xrightarrow{Fq} & X' \end{array}$$

The top horizontal paths commute by (5) and functoriality. The left rectangle commutes by the assumption that (X, β, α) is a λ-bialgebra. The upper square commutes by naturality of λ, and the lower square by (6) and functoriality. Thus we have $Fq \circ \alpha \circ \beta \circ T(\pi_1) = Fq \circ \alpha \circ \beta \circ T(\pi_2)$, and consequently

$$c_\beta(r(R)) \underset{\pi_2}{\overset{\pi_1}{\rightrightarrows}} X \xrightarrow{\alpha} FX \xrightarrow{Fq} F(X')$$

commutes, which means $c_\beta(r(R)) \rightsquigarrow c_\beta(r(S))$. Thus $c_\beta(r(R)) \subseteq \psi(c_\beta(r(S)))$ by Lemma 6, and by Lemma 1 now $c_\beta \circ r$ is ψ-compatible. □

For compatibility of the contextual closure function for coalgebras which are (part of) λ-bialgebras when paired with the identity function, we can perform a similar development as in Section 6.2, by defining $\psi'(R) = \psi(R) \cap R$.

Example 18. For an example of behavioural equivalence up-to, we consider the "general process algebra with transitions costs" (GPA) from (Buchholz and Kemper, 2001). GPA processes are defined for a given set of labels A and a semiring $\mathbb{S}$ which, for this example, we fix to be the semiring of reals $\mathbb{R}$. The operational semantics of GPA is expressed in terms of *weighted transition systems* which are coalgebras for the functor $(\mathbb{R}^-_\omega)^A$ where $\mathbb{R}^-_\omega \colon \mathsf{Set} \to \mathsf{Set}$ is defined as follows:

— For each set X, $\mathbb{R}^X_\omega$ is the set of functions from X to $\mathbb{R}$ with finite support (see the *Notation* paragraph in the introduction).
— For each function $h \colon X \to Y$, $\mathbb{R}^h_\omega \colon \mathbb{R}^X_\omega \to \mathbb{R}^Y_\omega$ is the function mapping each $\varphi \in \mathbb{R}^X_\omega$ into $\varphi^h \in \mathbb{R}^Y_\omega$ defined, for all $y \in Y$, by

$$\varphi^h(y) = \sum_{x' \in h^{-1}(y)} \varphi(x')$$

In a nutshell, a weighted transition system is a pair (X, t) where X is a set of states and $t \colon X \to (\mathbb{R}^X_\omega)^A$ is the transition relation that associates a weight to each transition. We use the same notation of weighted automata: $x \stackrel{a,r}{\to} y$ means that $t(x)(a)(y) = r$ and $r \neq 0$.

As shown in Section 2.3 of (Bonchi et al., 2012), the functor $(\mathbb{R}^-_\omega)^A$ does not preserve weak pullbacks and therefore bisimulations up-to cannot be used in this context. However, thanks to Theorems 6 and 7 we can use behavioural equivalence up-to. First observe that, by instantiating the definition of ψ above to an $(\mathbb{R}^-_\omega)^A$-coalgebra (X, t), one obtains the function $\psi_t \colon \mathcal{P}(X \times X) \to \mathcal{P}(X \times X)$ defined for all relations $R \subseteq X \times X$ as

$$\psi_t(R) = \{(x_1, x_2) \mid \forall a \in A, y \in X, \sum_{y' \in [y]_R} t(x_1)(a)(y') = \sum_{y' \in [y]_R} t(x_2)(a)(y')\}$$

where $[y]_R$ denotes the equivalence class of y w.r.t. R. With this explicit definition, it is easy to see that our notion of behavioural equivalence coincides with the one of bisimulation in (Buchholz and Kemper, 2001).

In order to illustrate our example is enough to consider a small fragment of GPA. The set P of basic GPA processes is defined by

$$p ::= 0 \mid p + p \mid (a, r).p$$

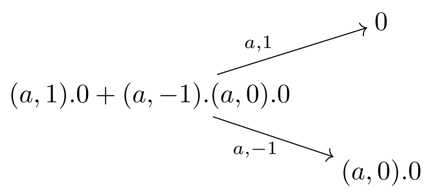
where $a \in A$, $r \in \mathbb{R}$. The operational semantics of basic GPA processes is given by the coalgebra $t \colon P \to (\mathbb{R}^P)^A$ defined for all $a' \in A$ and $p' \in P$ as follows:

$$\begin{array}{rcl} t(0)(a')(p') & = & 0 \\ t((a,r).p)(a')(p') & = & \begin{cases} r & \text{if } a = a', p = p' \\ 0 & \text{otherwise} \end{cases} \\ t(p_1 + p_2)(a')(p') & = & t(p_1)(a')(p') + t(p_2)(a')(p') \end{array}$$

Equivalently, it is described by the following rules:

$$\frac{r \neq 0}{(a,r).p \stackrel{a,r}{\to} p} \qquad \frac{p_1 \stackrel{a,r}{\to} p' \quad r \neq 0}{p_1 + p_2 \stackrel{a,s}{\to} p'} \qquad \frac{p_2 \stackrel{a,r}{\to} p' \quad r \neq 0}{p_1 + p_2 \stackrel{a,s}{\to} p'}$$

where $s = t(p_1)(a)(p') + t(p_2)(a)(p')$. For instance, the operational semantics of $(a,1).0 + (a,-1).(a,0).0$ is depicted below.



Since $0 \approx (a,0).0$, we have that $(a,1).0 + (a,-1).(a,0).0 \approx 0$. More generally, it holds that for all $a \in A$, $r \in R$, p_1 and $p_2 \in P$,

$$\text{if } p_1 \approx p_2 \text{ then } 0 \approx (a,r).p_1 + (a,-r).p_2. \tag{7}$$

We are going to prove (7) by means of behavioural equivalence up to *beh* (Theorem 6). To this end, consider the relation

$$R = \{(0, (a,r).p_1 + (a,-r).p_2) \mid p_1 \approx p_2\}$$

and note that R is *not* a behavioural equivalence by taking $p_1 = 0$ and $p_2 = (a,0).0$ (namely, $R \not\subseteq \psi_t(R)$). However R is a behavioural equivalence up to *beh*: to see that $R \subseteq \psi_t(beh(R))$, fix $p = (a,r).p_1 + (a,-r).p_2$ and observe that for all processes $q \in P$

$$\sum_{y' \in [q]_{beh(R)}} t(0)(a)(y') = 0 = \sum_{y' \in [q]_{beh(R)}} t(p)(a)(y').$$

The leftmost equality comes from the semantics of the process 0. For the rightmost, we have that either $q \not\approx p_1$ or $q \approx p_1$. In the first case, the above rightmost equivalence is obvious. In the second case,

$$\sum_{y' \in [q]_{beh(R)}} t(p)(a)(y') = t(p)(a)(p_1) + t(p)(a)(p_2) = r - r = 0$$

since $p_1 \in [q]_{beh(R)}$ and $p_2 \in [q]_{beh(R)}$.

8. Conclusions

Coalgebraic bisimulation-up-to enhances the proof method for bisimilarity, allowing for smaller proofs and equational reasoning on bisimulation equivalence for a large class of state-based systems and calculi. We presented a compositional framework for up-to-techniques and showed the compatibility (and thus the soundness) of the more common techniques: any novel compatible enhancements could be combined with these as well, without the necessity of re-proving soundness.

While showing this we also obtained interesting side results, such as Proposition 3, which provides a novel characterization of weak pullback preservation. This result is based on a definition of relational lifting for weak pullback preserving functors. We leave as future work the study of either a more broadly applicable notion of bisimulation, as in (Gorín and Schröder, 2013), or a more general definition of relation lifting, which applies to arbitrary functors on Set, as in (Marti and Venema, 2012). While in our work relators must preserve binary composition, in (Levy, 2011) a framework has been developed which only laxly preserves composition. The combination of this theory with our framework allows for a study of up-to-techniques for simulations, rather than bisimulations.

Bisimulation up to context for proving weighted language equivalence of weighted automata is, to the best of our knowledge, an original contribution. Finally, the soundness results generalize those of (Rot et al., 2013a) in that, for bisimulation up to context, the monad T does not need to be finitary. Future work includes the study of up-to techniques for particular types of systems, for example for name-passing languages such as the π-calculus.

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Appendix A. Bisimulations on different systems

The validity of the implication (3) $\Rightarrow$ (1) of Theorem 3 is shown in (Gumm and Schröder, 2000), based on the standard notion of bisimulation on different systems: given F-coalgebras (X, α_X) and (Y, α_Y) a relation $R \subseteq X \times Y$ is a bisimulation if there exists a transition structure $\gamma\colon R \to FR$ such that the following diagram commutes:

$$
\begin{array}{ccccc}
X & \xleftarrow{\pi_1} & R & \xrightarrow{\pi_2} & Y \\
{\scriptstyle \alpha_X}\downarrow & & \downarrow{\scriptstyle \gamma} & & \downarrow{\scriptstyle \alpha_Y} \\
FX & \xleftarrow[F\pi_1]{} & FR & \xrightarrow[F\pi_1]{} & FY
\end{array}
$$

The notion of bisimulation which we adopted in this paper is based on single systems, i.e., where $(X, \alpha_X) = (Y, \alpha_Y)$. We proceed to show that in Set, if bisimulations on single systems are closed under composition, then bisimulations on different systems are closed under composition as well; this proves that the implication from (3) to (1) of Theorem 3 holds in our setting as well.

We denote a coproduct by $X + Y$ and the associated injections by i_X and i_Y.

Proposition 5. Let (X, α_X) and (Y, α_Y) be F-coalgebras (F is a Set endofunctor) and $R \subseteq X \times Y$ a relation. Then R is a bisimulation on X and Y iff $(i_X \times i_Y)(R)$ is a bisimulation on $X + Y$.

Proof. Let $R \subseteq X \times Y$. The cases $X = \emptyset$ or $Y = \emptyset$ are trivial, so we may assume $X \neq \emptyset$ and $Y \neq \emptyset$. Let $(X + Y, \alpha_{X+Y})$ be the coproduct coalgebra (Rutten, 2000). So in the diagram below, the outer two squares commute. Suppose R is a bisimulation on X and Y; then there exists a γ such that the middle squares of the diagram below commute:

$$
\begin{array}{ccccccccc}
X+Y & \xleftarrow{i_X} & X & \xleftarrow{\pi_1} & R & \xrightarrow{\pi_2} & Y & \xrightarrow{i_Y} & X+Y \\
\downarrow{\scriptstyle \alpha_{X+Y}} & & \downarrow{\scriptstyle \alpha_X} & & \downarrow{\scriptstyle \gamma} & & {\scriptstyle \alpha_Y}\downarrow & & \downarrow{\scriptstyle \alpha_{X+Y}} \\
F(X+Y) & \xleftarrow{Fi_X} & X & \xleftarrow{F\pi_1} & FR & \xrightarrow{F\pi_2} & FY & \xrightarrow{Fi_Y} & FX+Y
\end{array}
$$

which means the entire diagram commutes and $(i_X \times i_Y)(R)$ is a bisimulation on $X + Y$.

Conversely suppose $((i_X, i_Y)(R), \gamma)$ is a bisimulation on $X + Y$ for some relation $R \subseteq X \times Y$; so in the above diagram, the outer rectangles commute (i.e., $\alpha_{X+Y} \circ i_X \circ \pi_1 =$

$Fi_X \circ F\pi_1 \circ \gamma$ and similarly for Y). Now since i_X is mono, X is nonempty, and F is a Set functor, Fi_X is mono as well (see, e.g., (Rutten, 2000)). Further $Fi_X \circ \alpha_X \circ \pi_1 = \alpha_{X+Y} \circ i_X \circ \pi_1 = Fi_X \circ F\pi_1 \circ \gamma$, and since Fi_X is mono we may conclude $\alpha_X \circ \pi_1 = F\pi_1 \circ \gamma$. Similarly we derive $\alpha_Y \circ \pi_2 = F\pi_2 \circ \gamma$. So R is a bisimulation on X and Y. □

Proposition 6. Let (X, α_X) and (Y, α_Y) be F-coalgebras, and $R \subseteq X \times X$ a relation. Then R is a bisimulation on X iff $(i_X \times i_Y)(R)$ is a bisimulation on $X + Y$.

Proof. Similar to that of Proposition 5. □

From the above two propositions, one can deduce the following:

Corollary 3. Let F be a Set endofunctor. Suppose F-bisimulations on single systems (i.e., of type $R \subseteq X \times X$) are closed under composition. Then F-bisimulations on different coalgebras (i.e., of type $R \subseteq X \times Y$) are closed under composition as well.

Proof. The outline of the proof is as follows. Let $R \subseteq X \times Y$ and $S \subseteq Y \times Z$ be bisimulations. Then by applying first Proposition 5 and then Proposition 6 we can turn both into bisimulations on $X + Y + Z$. The composition of these two then is a bisimulation by assumption; and applying Proposition 6 and Proposition 5 in the other direction again we obtain that $R \circ S$ is a bisimulation on X and Z. □