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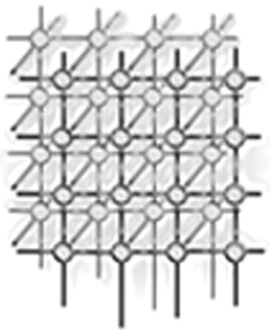
On the optimality of Feautrier's scheduling algorithm



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SUMMARY

Feautrier's scheduling algorithm is the most powerful existing algorithm for parallelism detection and extraction. But it has always been known to be suboptimal. However, the question whether it may miss some parallelism because of its design was still open. We show that this is not the case. Therefore, to find more parallelism than this algorithm does, one needs to get rid of some of the hypotheses underlying its framework.

KEY WORDS: Automatic parallelization, scheduling, Farkas' lemma, vertex method.

1. INTRODUCTION

One of the fundamental steps of automatic parallelization is the detection and extraction of parallelism. This extraction can be done in very different ways, from the try and test of *ad hoc* techniques to the use of powerful scheduling algorithms. In the field of dense matrix code parallelization, lots of algorithms have been proposed along the years. Among the main ones, we have the algorithms proposed by Lamport [12], Allen and Kennedy [2], Wolf and Lam [18], Feautrier [9, 10], and Darte and Vivien [7]. This collection of algorithms spans a large domain of techniques (loop distribution, unimodular transformations, linear programming, etc.) and a large domain of dependence representations (dependence levels, direction vectors, affine dependences, dependence polyhedra). One may wonder which algorithm to chose from such a collection. Fortunately, we have some theoretical comparative results on these algorithms, as well as some optimality results.

Allen and Kennedy's, Wolf and Lam's, and Darte and Vivien's algorithms are optimal for the representation of the dependences they respectively take as input [6]. This means that each

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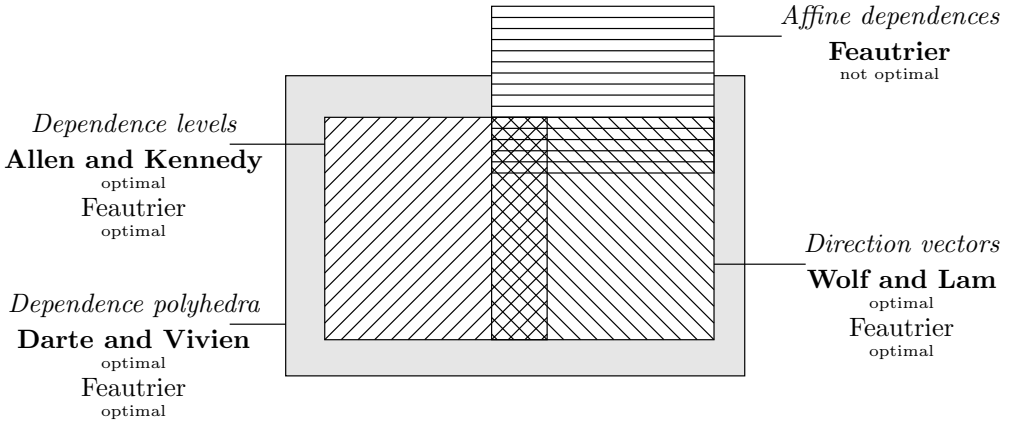

Figure 1. Schematic representation of the domains of the main dependence representations and of the main parallelism detection algorithms. Each dependence representation, named in italic, is represented by the domain of all the sets of dependences that it can describe. Each parallelism detection algorithm, named in bold, is associated with the dependence representation it was designed for. Feautrier's algorithm can in fact be used on the four dependence representations and is optimal whenever one of the other algorithms is optimal.

of these algorithms extracts all the parallelism contained in its input (some representation of the code dependences). Wolf and Lam's algorithm is a generalization of Lamport's; Darte and Vivien's algorithm is a generalization of those of Allen and Kennedy, and of Wolf and Lam, and is generalized by Feautrier's [6]. Finally, Feautrier's algorithm can handle any of the dependence representations used by the other algorithms [6]. These results are summarized on Figure 1.

It appears from these results that Feautrier's algorithm is the most powerful algorithm we have at hand. Although this algorithm has always be known to be suboptimal, its exact efficiency was so far unknown. Hence the questions we address in this paper: What are its weaknesses? Is its suboptimality only due to its framework or also to its design? What can be done to improve this algorithm? How can we build a more powerful algorithm?

These questions may seem rather theoretical. However, Feautrier's algorithm has a very interesting property: when called on a program, it does not output a single answer, but exhibits all the schedules (in its framework) that are valid for the program. Therefore, this algorithm can be used as a tool which exhibits the set of all the valid schedules, and this set can be searched for a solution to a particular problem: in automatic parallelization, it could be a schedule that respects a given mapping of the computations [5]; in program optimization it could be an affine occupancy vector that enables to shrink the memory used by the program [17]; etc. Therefore, even outside the scope of automatic parallelization, Feautrier's algorithm appears to be an

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interesting tool to build program optimizations on. However, before considering using such a tool, one would like to know how sound and powerful it is!

In Section 2 we briefly recall Feautrier's algorithm. Then we discuss its weaknesses in Section 3. In Section 4 we present what seems to be a "better" algorithm. Section 5 presents the major new result of this paper: to find "more" parallelism than Feautrier's algorithm one needs to use far more powerful techniques. This result is proved in Section 6.

2. THE ALGORITHM

Feautrier uses schedules to detect and extract parallelism. This section gives a brief overview of the problem and of Feautrier's algorithm. All the missing details can be found either in [9, 10] or [6].

2.1. Framework: static control programs

To enable an exact dependence analysis, the control-flow must be predictable at compile time. The necessary restrictions define the class of the *static control programs*. These are the programs:

- whose only data structures are integers, floats, arrays of integers, and arrays of floats, with no pointers or pointer-like mechanisms;
- whose elementary statements are assignments of scalars or array elements;
- whose only control structure are sequences and `do` loops with constant steps;
- where the array subscripts and the loop bounds are affine functions of surrounding loop indices and structural parameters.

Static control programs are mainly sets of nested loops. Figure 2 presents Example 1 which is an example of such a program. Let S be any statement. The iteration domain of S, denoted $\mathcal{D}_S$, is the set of all possible values of the vector of the indices (the *iteration vector*) of the loops surrounding S: in Example 1, $\mathcal{D}_S = \{(i,j) \mid 1 \leq i \leq N, 1 \leq j \leq i\}$. In static control programs, an iteration domain is always a polyhedron. In other words, there always exist a matrix A and a vector b such that : $\mathcal{D}_S = \{x \mid A.x \leq b\}$.

2.2. Dependence representation

In the framework of static control programs, an exact dependence analysis is feasible [8] and each exact *dependence relation* e from statement S_e to statement T_e is defined by a $\mathbb{Z}$-polyhedron $\mathcal{D}_e$, the domain of existence of the dependence relation, and a quasi-affine [†] function h_e as follows: for any value $j \in \mathcal{D}_e$, operation $T_e(j)$ depends on operation $S_e(h_e(j,N))$, which we note:

$$j \in \mathcal{D}_e \quad \Rightarrow \quad S_e(h_e(j,N)) \rightarrow T_e(j)$$

[†]See the original paper [8] for more details.

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```
DO i=1, N
 DO j=1, i
  S: a(i,i+j+1) = a(i-1,2*i-1) + a(j,2*j)
 ENDDO
ENDDO
```

Figure 2. Example 1.

$$e_1\text{:}\quad S(i\text{–}1, i\text{–}1) \rightarrow S(i,j), h_{e_1}(i,j)\text{=}(i\text{–}1, i\text{–}1)$$
$$\mathcal{D}_{e_1} = \{(i,j) \mid 2 \leq i \leq N, 1 \leq j \leq i\}$$

$$e_2\text{:}\quad S(j, j\text{–}1) \rightarrow S(i,j), h_{e_2}(i,j)\text{=}(j, j\text{–}1)$$
$$\mathcal{D}_{e_2} = \{(i,j) \mid 1 \leq i \leq N, 2 \leq j \leq i\}$$

Figure 3. Dependence relations for Example 1.

where N is the vector of structural parameters. Obviously, the description of the exact dependences between two statements may involve the union of many such dependence relations. A dependence relation e describes for any value $j \in \mathcal{D}_e$ a dependence between the two operations $S_e(h_e(j,N))$ and $T_e(j)$, what we call an *operation to operation dependence*. In other words, a *dependence relation* is a set of elementary *operation to operation dependences*. Figure 3 presents the dependence relations for Example 1.

Following Feautrier [9], we suppose that all the quasi-affine functions we have to handle are in fact affine functions and that all $\mathbb{Z}$-polyhedra are in fact polyhedra (at the possible cost of a conservative approximation of the dependences).

2.3. Searched schedules

Feautrier does not look for any type of functions to schedule affine dependences. He only considers nonnegative functions, with rational values, that are affine functions in the iteration vector and in the vector of structural parameters. Therefore he only handles (affine) schedules of the form:

$$\Theta(S,j,N) = X_S.j + Y_S.N + \rho_S \tag{1}$$

where X_S and Y_S are non-parameterized rational vectors and ρ_S is a rational constant. The hypothesis of nonnegativity of the schedules is not restrictive as all schedules must be lower bounded.

2.4. Problem statement

Once chosen the form of the schedules, the scheduling problem *seems* to be simple. For a schedule to be valid, it must (and only has to) satisfy the dependences. For example, if operation $T(j)$ depends on operation $S(i)$, $T(j)$ must be scheduled *after* $S(i) : \Theta(T,j,N) > \Theta(S,i,N)$. Therefore, for each statement S, we just have to find a vector X_S, a vector Y_S, and a constant

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ρ_S such that, for each dependence relation e, the schedule satisfies: [‡]

$$j \in \mathcal{D}_e \quad \Rightarrow \quad \Theta(S_e, h_e(j,N), N) + 1 \ \leq \ \Theta(T_e, j, N). \tag{2}$$

All the constraints are affine, and one can imagine using linear system solvers to find a solution. Actually, there are now two difficulties to overcome:

1. Equation (2) must be satisfied for any possible value of the structural parameters. If polyhedron $\mathcal{D}_e$ is parameterized, Equation (2) may correspond to an infinite set of constraints, which cannot be enumerated. There are two means to overcome this problem: the polyhedron vertices (cf. Section 4) and the affine form of Farkas' lemma (cf. Section 2.5). Feautrier uses the latter.
2. There does not always exist a solution for such a set of constraints. We will see, in Section 2.6, how the use of multidimensional schedules can overcome this problem.

2.5. The affine form of Farkas' lemma and its use

This lemma [9, 16] predicts the shape of certain affine forms.

Theorem 1 (Affine form of Farkas' lemma) *Let $\mathcal{D}$ be a nonempty polyhedron defined by p inequalities: $a_k x + b_k \geq 0$, for any $k \in [1,p]$. An affine form Φ is nonnegative over $\mathcal{D}$ if and only if it is a nonnegative affine combination of the affine forms used to define $\mathcal{D}$:*

$$\Phi(x) \equiv \lambda_0 + \sum_{k=1}^{p} \lambda_k (a_k x + b_k), \ \textit{with } \lambda_k \geq 0 \textit{ for any } k \in [0,p].$$

This theorem is useful as, in static control programs, all the important sets are polyhedra : iteration domains, dependence existence domains [8], etc. Feautrier uses this lemma to predict the shape of the schedules and to simplify the set of constraints.

Schedules. By hypothesis, the schedule $\Theta(S,j,N)$ is a nonnegative affine form defined on a polyhedron $\mathcal{D}_S$: the iteration domain of statement S. Therefore, the affine form of Farkas' lemma states that $\Theta(S,j,N)$ is a nonnegative affine combination of the affine forms used to define $\mathcal{D}_S$. Let $\mathcal{D}_S = \{x \mid \forall i \in [1,p_S], \ A_{S,i}.x + B_{S,i}.N + c_{S,i} \geq 0\}$ ($\mathcal{D}_S$ is thus defined by p_S inequalities). Then Theorem 1 states that there exist some nonnegative values $\mu_{S,0}$, ..., μ_{S,p_S} such that:

$$\Theta(S,j,N) \equiv \mu_{S,0} + \sum_{i=1}^{p_S} \mu_{S,i}(A_{S,i}.j + B_{S,i}.N + c_{S,i}). \tag{3}$$

Dependence constraints. Equation (2) can be rewritten as an affine function that is nonnegative over a polyhedron because the schedules and the function h_e are affine functions:

$$j \in \mathcal{D}_e \quad \Rightarrow \quad \Theta(T_e,j,N) - \Theta(S_e, h_e(j,N), N) - 1 \geq 0.$$

[‡]The transformation of the inequality, from $a > b$ to $a \geq 1 + b$, is obvious for schedules with integral values and a classical approximation for schedules with rational values [15].

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Once again we can apply the affine form of Farkas' lemma, this time the polyhedron is the existence domain of the dependence. Let $\mathcal{D}_e = \{x \mid \forall i \in [1, p_e],\ A_{e,i}.x + B_{e,i}.N + c_{e,i} \geq 0\}$ ($\mathcal{D}_e$ is thus defined by p_e inequalities). Theorem 1 states that there exist some nonnegative values $\lambda_{e,0}, ..., \lambda_{e,p_e}$ such that:

$$\Theta(T_e, j, N) - \Theta(S_e, h_e(j, N), N) - 1 \equiv \lambda_{e,0} + \sum_{i=1}^{p_e} \lambda_{e,i}(A_{e,i}.j + B_{e,i}.N + c_{e,i}).$$

Using Equation (3), we rewrite the left-hand side of this equation:

$$\begin{aligned}\left(\mu_{T_e,0} + \sum_{i=1}^{p_{T_e}} \mu_{T_e,i}(A_{T_e,i}.j + B_{T_e,i}.N + c_{T_e,i})\right) \\ - \left(\mu_{S_e,0} + \sum_{i=1}^{p_{S_e}} \mu_{S_e,i}(A_{S_e,i}.h_e(j,N) + B_{S_e,i}.N + c_{S_e,i})\right) - 1 \\ \equiv \lambda_{e,0} + \sum_{i=1}^{p_e} \lambda_{e,i}(A_{e,i}.j + B_{e,i}.N + c_{e,i}). \quad (4)\end{aligned}$$

Equation 4 is a formal equality ($\equiv$). Thus, the coefficients of a given component of either of the vectors j and N must be the same on both sides. The constant terms on both sides of this equation must also be equal. This identification process leads to a set of $(n+q+1)$ equations, equivalent to Equation (4), where n is the size of the iteration vector j, and q the size of the parameter vector N.

The way Feautrier uses the affine form of Farkas' lemma enables him to obtain a finite set of linear equations and inequations, equivalent to the original scheduling problem, and that can easily be solved using any solver of linear systems.

2.6. Extension to multidimensional scheduling

There exist some static control programs that cannot be scheduled with (monodimensional) affine schedules (e.g. Example 1, cf. Section 4). Hence the need for multidimensional schedules, i.e. schedules whose values are not rationals but rational *vectors* (ordered by lexicographic ordering). Before introducing Feautrier's solution, we introduce some vocabulary.

Definition 1 (Delay of an affine function, satisfied and respected dependences) *Let $S(i)$ and $T(j)$ be two operations of the loop nest such that $T(j)$ depends on $S(i)$: $S(i) \rightarrow T(j)$. Let $\tau(U, i, N)$ be any affine function. We call delay induced by τ on the dependence $S(i) \rightarrow T(j)$ the difference of the "execution times" of $T(j)$ and $S(i)$ defined by τ:*

$$\tau(T, j, N) - \tau(S, i, N).$$

*τ **respects** the dependence if it induces on it a nonnegative delay.*

*τ **satisfies** the dependence if it induces on it a delay greater than one.*

As we have already stated, a dependence relation *is a set of elementary operation to operation dependences. An affine function is said to **fully satisfy** (resp. **partially satisfy**) a dependence*

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relation e if and only if it satisfies all (resp. respects all and satisfies some but not all of) the operation to operation dependences in e.

The solution proposed by Feautrier is simple and greedy. For the first dimension of the schedules one looks for affine functions that 1) respect all the dependences; 2) satisfy as many dependence relations as possible. The algorithm is then recursively called on the unsatisfied dependence relations. This, plus a strongly connected component distribution[§] that reminds us of Allen and Kennedy's algorithm, defines the algorithm below. G denotes the multigraph defined by the statements and the dependence relations. The multidimensional schedules built satisfy the dependences according to the lexicographic order [6].

FEAUTRIER(G)

1. Compute the strongly connected components of G.
2. **For each** strongly connected component G_i of G **do** in topological order:
   (a) Find, using the method exposed in Section 2.5, an affine schedule that induces a nonnegative delay on all the dependences and satisfies Equation (2) for as many dependences as possible.
   Formally: find an affine function that satisfies
   $$\forall e,\ j \in \mathcal{D}_e \ \Rightarrow\ \theta(S_e, h_e(j,N), N) + z_e \leq \theta(T_e, j, N) \text{ with } 0 \leq z_e \leq 1 \qquad (5)$$
   and which maximizes the sum $\sum_e z_e$.
   (b) Build the subgraph G'_i generated by the unsatisfied dependences. If G'_i is not empty, recursively call FEAUTRIER(G'_i).

Notations: if Θ is the multidimensional schedule of a loop nest, we denote by Θ_i its i-th dimension, and by $\Theta(S)$ the multidimensional schedule it defines for the statement S.

2.7. Extension to other dependence representations

We recall here that FEAUTRIER can be extended to process a representation of the dependences by levels, direction vectors, and/or polyhedra. We focus on the latter type because it is a generalization of the two others [6, p. 266]. Let us consider a perfect loop nest[¶], two instructions S and T of same domain $\mathcal{D}$, with T depending on S by a dependence e described by the polyhedron $\mathcal{P}_e$. The schedule must respect this dependence. Therefore, it must satisfy the equation:

$$\forall j \in \mathcal{D},\ \forall d \in \mathcal{P}_e,\ (j-d) \in \mathcal{D} \quad \Rightarrow \quad \Theta(T, j, n) >_{lex} \Theta(S, j-d, n). \qquad (6)$$

The set of values of the couple (j, d) satisfying the conditions on the left-hand side of the implication is a polyhedron as $\mathcal{D}$ and $\mathcal{P}_e$ are polyhedra. Indeed, if $\mathcal{D} = \{x \mid Ax + b \geq 0\}$ and

[§]This distribution is rather aesthetic as the exact same result can be achieved without using it. This distribution is intuitive and eases the computations.
[¶]This extension is also valid for non perfect loop nests but is far more complicated to write...

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$\mathcal{P}_e = \{x \mid Cx + d \geq 0\}$ then:

$$\{(j,d) \mid j \in \mathcal{D}, d \in \mathcal{P}_e, j-d \in \mathcal{D}\} = \left\{ (j,d) \left| \begin{bmatrix} A & 0 \\ 0 & C \\ A & -A \end{bmatrix} \begin{bmatrix} j \\ d \end{bmatrix} + \begin{bmatrix} b \\ d \\ b \end{bmatrix} \geq 0 \right. \right\}$$

which is a polyhedron. Then we can apply the affine form of Farkas' lemma to Equation (6), as we did for Equation (2). Hence we can schedule a loop nest whose dependences are described by levels, direction vectors, and/or polyhedra using the affine form of Farkas' lemma and Feautrier's scheme.

3. THE ALGORITHM'S WEAKNESSES

3.1. Definitions of optimality

Depending on the definition one uses, an algorithm extracting parallelism is optimal if

1. It finds all the parallelism that can be extracted in its framework (only certain program transformations are allowed, etc.).
2. It finds all the parallelism that is contained in the representation of the dependences it handles.
3. It finds all the parallelism that is contained in the program to be parallelized (not taking into account the dependence representation used nor the transformations allowed).

For example, Allen, Callahan, and Kennedy uses the first definition [1], and Darte and Vivien the second [7] to prove that their respective algorithms are optimal. Feautrier uses the third definition to prove that its algorithm is *not* optimal [10]. We now recall that FEAUTRIER is not optimal under any of the last two definitions.

3.2. The classical counter-example to optimality

Feautrier proved in his original article [9] that his algorithm was not optimal for parallelism detection in static control programs. In his counterexample (Example 2, Figure 4) the source of any dependence is in the first half of the iteration domain and the sink in the second half. Cutting the iteration domain "in the middle" enables the trivial parallelization presented on Figure 5. However, the only loop in Example 2 contains some dependences. Thus, Feautrier's schedules must be of dimension at least one (hence there must be at least one sequential loop in the code after parallelization) to satisfy these dependences. Therefore, FEAUTRIER finds no parallelism in this example.

3.3. Weaknesses

The weaknesses in Feautrier's algorithm are either a consequence of the algorithm framework, or of the algorithm design.

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```
DO i=0, 2n
   x(i) = x(2n-i)
ENDDO
```

Figure 4. Example 2.

```
DOPAR i=0, n
   x(i) = x(2n-i)
ENDDOPAR
DOPAR i=n+1, 2n
   x(i) = x(2n-i)
ENDDOPAR
```

Figure 5. Parallelized version of Example 2.

Framework

Given a program, we extract its implicit parallelism and then we rewrite it. The new order of the computations must be rather regular to enable the code generation. Hence the restriction on the schedule shape: affine functions. The parallel version of Example 2 presented on Figure 5 can be expressed by a non affine schedule, but not by an affine schedule. The restriction on the schedule shape is thus a cause of inefficiency. Another problem with Example 2 is that Feautrier looks for a transformation conservative in the number of loops. Breaking a loop into several loops, i.e., cutting the iteration domain into several subdomains, can enable to find more parallelism (even with affine schedules). The limitation here comes from the hypothesis that all instances of a statement are scheduled the same way, i.e., with the same affine function. Note that this hypothesis is almost always made [12, 2, 18, 7], but Griebl, Feautrier, and Lengauer have already tried to get rid of it [11].

Some of the weaknesses of FEAUTRIER are thus due to its framework. Before thinking of changing this framework, we must check whether one can design a more powerful algorithm, or even improve FEAUTRIER, in Feautrier's framework.

Algorithm design

FEAUTRIER is a greedy algorithm which builds multidimensional schedules whose first dimension satisfies as many *dependence relations* as possible, and not as many *operation to operation dependences* as possible. We may wonder with Darte [4, p. 80] whether this can be the cause of a loss of parallelism. We illustrate this possible problem with Example 1.

The first dimension of the schedule must satisfy Equation (5), i.e. the constraints due to the dependences, for both dependence relations e_1 and e_2. This gives us respectively Equations (7) and (8):

$$X_S \begin{vmatrix} i\text{–}1 \\ i\text{–}1 \end{vmatrix} + z_{e_1} \leq X_S \begin{vmatrix} i \\ j \end{vmatrix} \Leftrightarrow z_{e_1} \leq X_S \begin{vmatrix} 1 \\ j\text{–}i\text{+}1 \end{vmatrix} \Leftrightarrow z_{e_1} \leq \alpha + \beta(j - i + 1) \text{ with } \begin{matrix} 2 \leq i \leq N \\ 1 \leq j \leq i \end{matrix} \quad (7)$$

$$X_S \begin{vmatrix} j \\ j\text{–}1 \end{vmatrix} + z_{e_2} \leq X_S \begin{vmatrix} i \\ j \end{vmatrix} \Leftrightarrow z_{e_2} \leq X_S \begin{vmatrix} i\text{–}j \\ 1 \end{vmatrix} \Leftrightarrow z_{e_2} \leq \alpha(i - j) + \beta \text{ with } \begin{matrix} 1 \leq i \leq N \\ 2 \leq j \leq i \end{matrix} \quad (8)$$

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if we note $X_S = (\alpha, \beta)$ [‖]. Equation (7) with $i = N$ and $j = 1$ is equivalent to $z_{e_1} \leq \alpha+\beta(2-N)$. The schedule must be valid for any (nonnegative) value of the structural parameter N, which implies $\beta \leq 0$. Equation (8) with $i = j$ is equivalent to $z_{e_2} \leq \beta$. Hence $z_{e_2} \leq 0$ as $\beta \leq 0$. As z_{e_2} must be nonnegative (cf. Equation (5)), this leads to $z_{e_2} = 0$. This means that the first dimension of any affine schedule cannot satisfy the dependence relation e_2.

The dependence relation e_1 can be satisfied, a solution being $X_S = (1, 0)$ ($\alpha = 1$, $\beta = 0$). Therefore, FEAUTRIER **is called recursively on the whole dependence relation** e_2. However, most of the dependences described by e_2 are satisfied by the schedule $\Theta(S, (i, j), N) = i$ (defined by $X_S = (1, 0)$). Indeed, Equation (7) is then satisfied for any value $(i, j) \in \mathcal{D}_{e_2}$ except when $i{=}j$. Thus, **one only needed to call recursively** FEAUTRIER **on the dependence relation** e_2':

$$S(j, j{-}1) \rightarrow S(i, j), h_{e_2}(i, j) = (j, j{-}1), \mathcal{D}_{e_2'} = \{(i, j) \mid 2{\leq}i{\leq}N, i = j\}.$$

Therefore, this example shows that the search for the schedules in FEAUTRIER is overconstrained by design.

We may now wonder whether this overconstraining may lead FEAUTRIER to build some affine schedules of non minimal dimensions and thus to miss some parallelism. We first present in Section 4 an algorithm which gets rid of this potential problem. Then, we show in Section 5 that no parallelism is lost because of this design particularity.

4. A GREEDIER ALGORITHM

4.1. The Vertex method

A polyhedron can always be decomposed as the sum of a polytope (i.e. a convex bounded polyhedron) and a polyhedral cone, called the characteristic cone (see [16] for details). A polytope is defined by its vertices, and any point of the polytope is a nonnegative convex combination of the polytope vertices. A polyhedral cone is finitely generated and is defined by its rays and lines. Any point of a polyhedral cone is the sum of a nonnegative combination of its rays and any combination of its lines. Therefore, a polyhedron $\mathcal{D}$ can be equivalently defined by a set of *vertices*, $\{v_1, \ldots, v_\omega\}$, a set of *rays*, $\{r_1, \ldots, r_\rho\}$, and a set of *lines*, $\{l_1, \ldots, l_\lambda\}$. Then $\mathcal{D}$ is the set of all vectors p such that

$$p = \sum_{i=1}^{\omega} \mu_i v_i + \sum_{i=1}^{\rho} \nu_i r_i + \sum_{i=1}^{\lambda} \xi_i l_i \tag{9}$$

with $\mu_i \in \mathbb{Q}^+$, $\nu_i \in \mathbb{Q}^+$, $\xi_i \in \mathbb{Q}$, and $\sum_{i=1}^{\omega} \mu_i = 1$. As we have already stated, all the important sets in static control programs are polyhedra. As we just recalled, any nonempty polyhedron

[‖]Example 1 contains a single statement S. Therefore, the components Y_S and ρ_S of Θ (cf. Equation (1)) have no influence here on Equation (5) which is equivalent to: $(X_S.h_e(j, N)+Y_S.N+\rho_S)+z_e \leq (X_S.j+Y_S.N+\rho_S) \Leftrightarrow X_S.h_e(j, N) + z_e \leq X_S.j$.

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is fully defined by its vertices, rays, and lines, which can be computed even for parameterized polyhedra [13]. The vertex method, introduced by Quinton [15], uses the vertices, rays, and lines to reduce the size of the set of constraints described by Equation (2).

Theorem 2 (The vertex method) *Let $\mathcal{D}$ be a nonempty polyhedron defined by a set of vertices (denoted by $\{v_1, \ldots, v_\omega\}$), a set of rays (denoted by $\{r_1, \ldots, r_\rho\}$), and a set of lines (denoted by $\{l_1, \ldots, l_\lambda\}$). Let Φ be an affine form of linear part A and constant part b ($\Phi(x) = A.x + b$). Then the affine form Φ is nonnegative over $\mathcal{D}$ if and only if 1) Φ is nonnegative on each of the vertices of $\mathcal{D}$ and 2) the linear part of Φ is nonnegative (respectively null) on the rays (resp. lines) of $\mathcal{D}$. This can be written :*

$$\forall p \in \mathcal{D},\ A.p + b \geq 0 \quad \Leftrightarrow$$
$$\forall i \in [1, \omega],\ A.v_i + b \geq 0,\ \forall i \in [1, \rho],\ A.r_i \geq 0,\ \textit{and}\ \forall i \in [1, \lambda],\ A.l_i = 0.$$

The polyhedra produced by the dependence analysis of programs (e.g. existence domain of dependences) are in fact polytopes. Then, according to Theorem 2, an affine form is nonnegative on a polytope if and only if it is nonnegative on the vertices of this polytope. We use this property to simplify Equation (2) and define a new scheduling algorithm.

4.2. The greediest algorithm

Feautrier's algorithm is a greedy heuristic which maximizes the number of *dependence relations* satisfied by the first dimension of the schedule, and then proceeds recursively. The algorithm below is a greedy heuristic which maximizes the number of *operation to operation dependences* satisfied by the first dimension of the schedule, and then proceeds recursively. To achieve this goal, this algorithm greedily considers the vertices of the existence domains of the dependence relations.

Let e_1, ..., e_n be the dependence relations in the studied program. For any $i \in [1, n]$, let $v_{i,1}, ..., v_{i,m_i}$ be the vertices of $\mathcal{D}_{e_i}$, and let, for any $j \in [1, m_i]$, $e_{i,j}$ be the operation to operation dependence from $S_{e_i}(h_{e_i}(v_{i,j}, N), N)$ to $T_{e_i}(v_{i,j})$: $e_{i,j} : S_{e_i}(h_{e_i}(v_{i,j}, N), N) \rightarrow T_{e_i}(v_{i,j})$. G denotes here the multigraph generated by the dependences $e_{i,j}$.

GREEDY(G)

1. Compute the strongly connected components of G.
2. **For each** strongly connected component G_k of G **do** in topological order:
   (a) Find an *integral* affine function Θ that satisfies
   $$\forall e_{i,j},\quad \Theta(S_{e_i}, h_{e_i}(v_{i,j}, N), N) + z_{i,j} \ \leq\ \Theta(T_{e_i}, v_{i,j}, N) \ \text{ with } \ 0 \leq z_{i,j} \leq 1 \qquad (10)$$
   and which maximizes the sum $\sum_{e_{i,j}} z_{i,j}$.
   (b) Build the subgraph G'_k generated by the unsatisfied dependences. If G'_k is not empty, recursively call GREEDY(G'_k).

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Lemma 1 (Correctness and maximum greediness) *The output of algorithm* GREEDY *is a schedule and the first dimension of this schedule satisfies all the operation to operation dependences that can be satisfied by the first dimension of an affine schedule (of the form defined in Section 2).*

Proof. Let $S(i)$ and $T(j)$ be any two operations of the loop nest such that $T(j)$ depends on $S(i)$: $S(i) \rightarrow T(j)$. In the graph G there exists a dependence relation e which includes the operation to operation dependence $S(i) \rightarrow T(j)$. Let v_1, ..., v_p be the vertices of $\mathcal{D}_e$. By definition of these vertices, there exist some nonnegative rationals μ_1, ..., μ_p such that : $j = \sum_{k=1}^{p} \mu_k v_k$ and $\sum_{k=1}^{p} \mu_k = 1$.

We first show that the multidimensional affine functions built by GREEDY respect all the dependences (by showing that the dependence from $S(i)$ to $T(j)$ is respected). Then we show the maximum greediness: the first dimension of Θ satisfies all the operation to operation dependences that can be satisfied by an affine schedule (if there exists an affine schedule whose first dimension satisfies the dependence from $S(i)$ to $T(j)$ then so does the first dimension of Θ). As a consequence, if there exists one multidimensional affine schedule for the studied set of dependences, we show that GREEDY builds such a schedule.

All dependences are respected. Before all, we must show that this algorithm effectively builds a schedule. Equation (10) insures us that, for any $k \in [1, p]$, $\Theta(S, h_e(v_k, N), N) \leq \Theta(T, v_k, N)$. Therefore:

$$\begin{aligned}
\sum_{k=1}^{p} \mu_k \Theta\left(S, h_e\left(v_k, N\right), N\right) &\leq \sum_{k=1}^{p} \mu_k \Theta\left(T, v_k, N\right) &\Leftrightarrow \\
\Theta\left(S, \sum_{k=1}^{p} \mu_k h_e\left(v_k, N\right), N\right) &\leq \Theta\left(T, \sum_{k=1}^{p} \mu_k v_k, N\right) &\Leftrightarrow \\
\Theta\left(S, h_e\left(\sum_{k=1}^{p} \mu_k v_k, N\right), N\right) &\leq \Theta\left(T, \sum_{k=1}^{p} \mu_k v_k, N\right) &\Leftrightarrow \\
\Theta\left(S, h_e\left(j, N\right), N\right) &\leq \Theta\left(T, j, N\right) &\Leftrightarrow \\
\Theta\left(S, i, N\right) &\leq \Theta\left(T, j, N\right)
\end{aligned}$$

The equivalence between the first and the second inequations (resp. between the second and the third) comes from the facts that Θ (resp. h_e) is an affine function and that $\sum_{k=1}^{p} \mu_k = 1$.

From what precedes, $\Theta\left(S, i, N\right) \leq \Theta\left(T, j, N\right)$, and Θ (the first dimension of the schedule) respects the dependence between $S(i)$ and $T(j)$. Therefore, the schedule built by GREEDY respects all the dependences.

We still have to show that GREEDY *satisfies* any operation to operation dependence. This is a consequence of the next point.

As many dependences as possible are satisfied. We suppose that there exists at least one affine schedule whose first dimension, denoted σ, satisfies the operation to operation dependence

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$S(i) \rightarrow T(j)$: $\sigma(S,i,N)+1 \leq \sigma(T,j,N)$.

$$\begin{aligned}
\sigma(S,i,N)+1 &\leq \sigma(T,j,N) \quad \Leftrightarrow \\
\sigma(S,h_e(j,N),N)+1 &\leq \sigma(T,j,N) \quad \Leftrightarrow \\
\sigma\left(S,h_e\left(\sum_{k=1}^{p}\mu_k v_k,N\right),N\right)+1 &\leq \sigma\left(T,\sum_{k=1}^{p}\mu_k v_k,N\right) \quad \Leftrightarrow \\
\sigma\left(S,\sum_{k=1}^{p}\mu_k h_e(v_k,N),N\right)+1 &\leq \sigma\left(T,\sum_{k=1}^{p}\mu_k v_k,N\right) \quad \Leftrightarrow \\
1+\sum_{k=1}^{p}\mu_k\sigma(S,h_e(v_k,N),N) &\leq \sum_{k=1}^{p}\mu_k\sigma(T,v_k,N) \qquad (11)
\end{aligned}$$

As $\sum_{k=1}^{p}\mu_k = 1$ with $\mu_k \geq 0$ for any $k \in [1,p]$, Inequation (11) implies that there exists at least one vertex v_q such that $1+\sigma(S,h_e(v_q,N),N) \leq \sigma(T,v_q,N)$ with $\mu_q > 0$. Therefore, the first dimension, Θ, of the schedule built by GREEDY satisfies the dependence corresponding to v_q. Otherwise $(\Theta+\sigma)$ would be an affine function inducing:

1. a nonnegative delay on any vertex (as the sum of two such affine functions);
2. a delay greater than or equal to one on all the vertices on which either Θ or σ induces a delay greater than or equal to one; therefore $(\Theta+\sigma)$ would induce a delay greater than or equal to one on v_q and therefore on at least one more vertex than Θ, which is impossible by definition of Θ.

As Θ induces a delay greater than or equal to one on v_q, it induces a strictly positive delay on $S(i) \rightarrow T(j)$. Indeed, from the inequations:

$$1+\sigma(S,h_e(v_q,N),N) \leq \sigma(T,v_q,N)$$

and

$$\forall k \in [1,p],\ \sigma(S,h_e(v_k,N),N) \leq \sigma(T,v_k,N)$$

we can obtain (using the same transformations than previously):

$$\Theta(S,i,N)+\mu_q \leq \Theta(T,j,N) \quad \Rightarrow \quad \Theta(S,i,N) < \Theta(T,j,N).$$

As Θ is by construction an integral schedule, the strictly positive delay $\Theta(T,j,N)-\Theta(S,i,N)$ is greater than or equal to one, and Θ satisfies the operation to operation dependence $S(i) \rightarrow T(j)$. ■

The following lemma provides a property which is crucial for code generation.

Lemma 2 (Independence of the linear parts) *Among all the schedules that algorithm* GREEDY *can build for any given loop nest, there exists at least one schedule Θ such that, for any statement S in the loop nest, the set of the linear parts for statement S is independent. In other words, if d_S is the dimension of the schedule of statement S under Θ (d_S is the dimension of $\Theta(S)$), and if, for $i \in [1,d_S]$, $\Theta_i(S,j,N) = X_S^i.j + Y_S^i.N + \rho_S^i$, then $\{X_S^1,...,X_S^{d_S}\}$ is an independent set of vectors.*

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Proof. Let σ be any schedule built by GREEDY. We note: $\sigma_i(S,j,N) = x_S^i.j + y_S^i.N + \nu_S^i$. We build the desired schedule Θ from σ by induction on the dimension. The induction hypothesis at rank i is that, for any statement S, the linear parts of the first i dimensions of $\Theta(S)$ (the only dimensions built so far) are independent and the first i dimensions of Θ satisfy all the dependences of level less than or equal to i.

We denote by $\mathcal{F}(S,i)$ the set of vectors $\{X_S^1, ..., X_S^i\}$. We call b_j the j-th canonical vector (whose components are all null except the j-th one which is equal to one). Finally we suppose that all the loops have positive steps: the theorem obviously holds without this property but the proof is far more painful to write.

Initialization of the induction. For each statement S, we let $\Theta_1(S,j,N) = \sigma_1(S,j,N) + b_1.j$ (thus $X_S^1 = x_S^1 + b_1$). This is a valid first dimension as we have supposed that all the steps are positive (thus all dependence distances are lexicographically nonnegative [6]). Furthermore, this affine function satisfies all the dependences of level one (in fact, both terms of the sum satisfy them!).

Building dimension $i+1$. We suppose that the induction hypothesis is satisfied up to dimension i included. We look for some nonnegative values $\lambda_1, ..., \lambda_{i+1}$ as we want to define Θ_{i+1} as follows for any statement S:

$$\Theta_{i+1}(S,j,N) = \sigma_{i+1}(S,j,N) + \sum_{k=1}^{i+1} \lambda_k.b_k.j \quad (\text{then } X_S^{i+1} = x_S^{i+1} + \sum_{k=1}^{i+1} \lambda_k.b_k).$$

First, remark that whatever the values of the λ, $\Theta_{i+1}(S)$ will be a valid $(i+1)$-th dimension satisfying all the dependences satisfied by σ_{i+1} as the λ are all nonnegative, as all the steps are positive (and thus all the dependence distance are lexicographically nonnegative [6]), and as, by induction hypothesis, all the dependences up to level i are already satisfied. Furthermore, the affine function $\tau(S,j,N) = b_{i+1}.j$ defines a valid $(i+1)$-th dimension of the schedule satisfying all the dependences at level $i+1$. By maximum greediness of GREEDY (Lemma 1), σ_{i+1} satisfies at least all the dependences satisfied by this affine function, and so does Θ_{i+1}. Thus, whatever the values of the λ, the first $i+1$ dimensions of Θ satisfy all the dependences of level up to $i+1$ (included).

For each statement S there exists a (possibly empty) set of values of the λ, denoted $\mathcal{P}(S,i+1)$, such that $\mathcal{F}(S,i+1) = \{X_S^1, ..., X_S^{i+1}\}$ is not an independent set of vectors. As, by induction hypothesis, $\mathcal{F}(S,i)$ is an independent set of vectors, $\mathcal{F}(S,i+1)$ is not independent if and only if X_S^{i+1} is a combination of $X_S^1, ..., X_S^i$:

$$\mathcal{P}(S,i+1) = \left\{ (\lambda_1, ..., \lambda_{i+1}) \;\middle|\; \begin{array}{l} \lambda_1 \geq 0, ..., \lambda_{i+1} \geq 0, \\ \exists \mu_1, ..., \exists \mu_i, \\ x_S^{i+1} + \sum_{j=1}^{i+1} \lambda_j.b_j = \sum_{k=1}^{i} \mu_k.X_S^k \end{array} \right\}.$$

$\mathcal{P}(S,i+1)$ is obviously a polyhedron. Furthermore, $\mathcal{P}(S,i+1)$ is not equal to the whole vector space: $\mathcal{F}(S,i)$ is an independent set of vectors of dimension i and, thus, it cannot contain the vector space of dimension $i+1$ generated by $b_1, ..., b_{i+1}$. Thus, for any statement

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S, $\mathcal{P}(S, i+1) \subsetneq \mathbb{Q}^{i+1}$. Therefore, $\bigcup_S \mathcal{P}(S, i+1) \subsetneq \mathbb{Q}^{i+1}$, and we can take for $(\lambda_1, ..., \lambda_{i+1})$ any value in $\mathbb{Q}^{i+1} \setminus (\bigcup_S \mathcal{P}(S, i+1)) = \bigcap_S \left(\mathbb{Q}^{i+1} \setminus \mathcal{P}(S, i+1)\right)$. ■

The above proof is constructive as one can explicit the difference of two polyhedra (for example using the polylib library [14]).

5. SCHEDULES OF MINIMAL DIMENSION

As GREEDY is greedier than FEAUTRIER, one could imagine that the former may sometimes build schedules of smaller dimension than the latter and thus may find more parallelism. The following theorem shows that this never happens.

Theorem 3 (The dimension of Feautrier's schedules is minimal) *Let us consider a loop nest whose dependences are all affine, or are represented by affine functions. If we are only looking for one affine schedule per statement of the loop nest, then the dimension of the schedules built by* FEAUTRIER *is minimal, for each statement of the loop nest.*

Note that this theorem cannot be improved, as the study of Example 2 shows. The proof is direct (not using algorithm GREEDY) and is presented in Section 6.

Principle of the proof of Theorem 3. Let σ be an affine schedule whose dimension is minimal for each statement in the studied loop nest (Lemma 3 proves that such a schedule exists).

From σ, we are going to build a schedule *a la* FEAUTRIER of same dimension, by combining the dimensions of σ. Let us consider one dependence relation, say e. Suppose that no affine schedule can fully satisfy e with its first dimension. The possible inefficiency comes from the fact that the first dimension of σ can (perhaps) satisfy some of the operation to operation dependences in e. Therefore the dimensions of σ of rank greater than one will no more have to satisfy these operation to operation dependences, and thus the dimensions of σ of rank greater than one will not be constrained by these operation to operation dependences. On the opposite, the dimensions of any schedule *a la* FEAUTRIER will still have to satisfy *all* the operation to operation dependences in e, even if some of them are satisfied by the first dimensions of the schedule.

Let e be a dependence relation, of existence domain $\mathcal{D}_e$. We suppose that e is not fully, but partially, satisfied by the first dimension of σ (otherwise there is no problem with e). The operation to operation dependences in e not satisfied by the first dimension of the schedule σ define a subpolyhedron $\mathcal{D}_e^1$ of $\mathcal{D}_e$: this is the subset of $\mathcal{D}_e$ on which the first dimension of σ induces a null delay. $\mathcal{D}_e^1$ is thus defined by the equations defining $\mathcal{D}_e$ and by the null delay equation involving the first dimension of σ:

$$\mathcal{D}_e^1 = \{j \in \mathcal{D}_e \mid \sigma_1(T_e, j, N) - \sigma_1(S_e, h_e(j, N), N) = 0\}.$$

The second dimension of σ must respect the dependences in $\mathcal{D}_e^1$, i.e., must induce a nonnegative delay over $\mathcal{D}_e^1$. Therefore, the second dimension of σ is an affine form nonnegative over the polyhedron $\mathcal{D}_e^1$. Using the affine form of Farkas' lemma (Theorem 1), we obtain that the second dimension of σ is defined from the (null delay equation on the) first dimension of

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σ and from the equations defining $\mathcal{D}_e$. From the equations obtained using Farkas' lemma, we build a nonnegative linear combination of the first two dimensions of σ which induces a nonnegative delay over $\mathcal{D}_e$ (and not only on $\mathcal{D}_e^1$), and which satisfies all the operation to operation dependences in e satisfied by any of the first two dimensions of σ. This way we build a schedule *a la* Feautrier of same dimension than σ: a whole dependence relation is kept as long as all its operation to operation dependences are not satisfied by the same dimension of the schedule.

Consequences. First, a simple and important corollary of the previous theorem:

Corollary 1. FEAUTRIER *is well-defined: it always outputs a valid schedule when its input is the exact dependences of an existing program.*

The original proof relied on an assumption on the dependence relations that can be easily enforced but which is not always satisfied: all operation to operation dependences in a dependence relation are of the same dependence level (are satisfied by the same dimension of the original schedule). For example, dependence relation e_2 in Example 1 does not satisfy this property.

More important, Theorem 3 shows that Feautrier's algorithm can only miss some (significant amount of) parallelism because of the limitations of its framework, but not because of its design: as the dimension of the schedule is minimal, the magnitude of the schedule's makespan is minimal, for any statement.

Even if GREEDY is a greedier algorithm than FEAUTRIER, both outputs schedules of the same dimension. On one hand, GREEDY constrains more the first dimension of its schedules than FEAUTRIER. On the other hand, FEAUTRIER sometimes overconstrains the remaining dimensions of its schedules. Therefore, no algorithm is better than the other. They are slightly different but we believe their main properties are the same.

Finally, a careful study of the proof of Theorem 3 shows that Lemma 2, established for GREEDY, also holds for FEAUTRIER (to prove it, we just have to run the constructive proof of Theorem 3 on a schedule of GREEDY satisfying Lemma 2). In other words, among all the schedules that FEAUTRIER can build, there are some schedules whose linear parts, for any statement, is an independent set of vectors. Thanks to this property we can soundly link the magnitude of the schedule's makespan to the schedule's dimension (the schedule's makespan is an (Ehrhart [3]) polynomial whose degree is the schedule's dimension).

6. THE PROOF OF THEOREM 3

Theorem 3 is proved by induction on the dimensions of the schedule. We formally state our hypotheses and notations in Paragraph 6.2 and our induction hypotheses in Paragraph 6.3. Then we verify the base case in Paragraph 6.4 and prove the inductive step in Paragraph 6.5. Before that, we prove a necessary preliminary result in Paragraph 6.1.

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6.1. Minimal dimension simultaneously for all statements

Lemma 3. *If there exists at least one valid affine schedule for the studied system, then there exists an affine schedule which is simultaneously of minimal dimension for all statements.*

Proof. Let Θ_1 and Θ_2 be two valid schedules. For each statement S, we define d_S as the minimum of the dimensions of $\Theta_1(S)$ and $\Theta_2(S)$. Then, we define a new (multidimensional) affine function Θ'_1 as follows: for each statement S, $\Theta'_1(S) = \Theta_1(S)_{|_{d_S}}$, i.e., $\Theta'_1(S)$ is equal to the first d_S dimensions of $\Theta_1(S)$. As Θ_1 is a schedule, and as $\Theta_1(S)$ is of dimension greater than or equal to d_S, Θ'_1 respects all the dependences, i.e. induces a lexicographically nonnegative delay on all the dependences. We symmetrically define Θ'_2. Then $\Theta = \Theta'_1 + \Theta'_2$ is a schedule. Indeed, let us consider an operation to operation dependence between $S(i)$ and $T(j)$. Without any loss of generality, we suppose that $d_S \leq d_T$, and that d_S is the dimension of $\Theta_1(S)$. As Θ_1 is a schedule, it induces a lexicographically (strictly) positive delay on the dependence between $S(i)$ and $T(j)$. Furthermore, as d_S is the dimension of $\Theta_1(S)$ and as $d_S \leq d_T$, Θ'_1 induces a lexicographically (strictly) positive delay on the dependence between $S(i)$ and $T(j)$. As Θ'_2 induces a lexicographically nonnegative delay on all the dependences, $\Theta'_1 + \Theta'_2$ induces a lexicographically (strictly) positive delay on the considered dependence.

We can take for each statement S a schedule of minimal dimension for S. If we apply the previous scheme to this set of schedules we end up with a schedule which is simultaneously of minimal dimension for all statements. ■

6.2. Hypotheses and notations

Hypotheses

Let Θ be an affine schedule whose dimension is minimal for each statement. The existence of such a schedule is guaranteed by Lemma 3 on the hypothesis that there exists at least one (multidimensional) affine schedule satisfying all the dependences in the studied system. If there is no schedule, Theorem 3 obviously holds! From Θ, we build by induction an affine schedule *a la* FEAUTRIER, denoted by P, whose dimension is minimal for each statement.

We suppose that Θ is an integral function. If this is not the case, we just scale up Θ by the least common multiple of the denominator of its rational coefficients.

Notations

Let e_1, ..., e_p be the p dependence relations of the considered loop nest. Each of these dependence relations is a 4-tuple: $e_j = (S_j, T_j, \mathcal{D}_j, h_j)$, where S_j and T_j are two statements, $\mathcal{D}_j$ is a polyhedral domain, and h_j an affine function such that e_j denotes the following set of dependences:

$$\forall i \in \mathcal{D}_j, T_j(i) \text{ depends on } S_j(h_j(i, N)). \tag{12}$$

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6.3. Induction hypotheses for dimension i

1. The first i dimensions of P fully satisfy the dependence relations e_1, ..., e_{m_i}, but none of the dependence relations e_{1+m_i}, ..., e_p. Therefore, the latter must be satisfied by the dimensions of P of rank greater than or equal to $i+1$.
2. Dimension i of P fully satisfies as many of the dependence relations $e_{1+m_{i-1}}$, ..., e_p as possible.
3. The dimensions of Θ of rank greater than or equal to $i+1$ satisfy the dependence relations $e^i_{1+m_i}$, ..., e^i_p where, if $j \in [1+m_i, p]$, $e^i_j = (S_j, T_j, \mathcal{D}^i_j, h_j)$, with
$$\mathcal{D}^i_j = \{x \in \mathcal{D}_j \mid \forall k \in [1,i],\ \Theta_k(T_j, x, N) - \Theta_k(S_j, h_j(x,N), N) = 0\}.$$
The operation to operation dependences described by $e_1, ..., e_p$, but not in $e^i_{1+m_i}$, ..., e^i_p, are satisfied by the first i dimensions of Θ. (Therefore, the first i dimensions of Θ fully satisfy the dependence relations e_1, ..., e_{m_i}.)
4. For any j in $[1,i]$, Θ_j is a linear combination of $P_1, ..., P_j$.
5. For any j in $[1,i]$, P_j satisfies all the operation to operation dependences which are satisfied by Θ_j.

6.4. Initialization of the induction: $i = 1$

We first prove that the different induction hypotheses hold when $i = 1$.

The first dimension of Θ fully satisfies some of the dependence relations (possibly none). In other words, for some values of $j \in [1,p]$, the first dimension of Θ induces a delay greater than or equal to one on all the operation to operation dependences described by Equation (12). Without any loss of generality, we suppose that the first dimension of Θ fully satisfies the dependence relations e_1, ..., e_{m_1}, and no other dependence relations.

Hypotheses 1 and 2. By definition, the first dimension of any schedule built by Feautrier fully satisfies as many dependence relations as possible. Therefore, any schedule built by Feautrier fully satisfies at least the dependence relations e_1, ..., e_{m_1}. Indeed, suppose this is not the case. Then take a schedule built by Feautrier and not fully satisfying the dependence relations e_1, ..., e_{m_1}. We add to the first dimension of this schedule the first dimension of Θ. This way we obtain an affine schedule whose first dimension fully satisfies at least one more dependence relation than Feautrier's schedule. This is impossible as it contradicts the maximization of the number of dependence relations fully satisfied (maximization of the sum $\sum_e z_e$, cf. Section 2.6).

We suppose that any schedule built by Feautrier fully satisfies only the dependence relations e_1, ..., e_{m_1}. Is this is not the case, we take any integral schedule $\mathcal{S}$ built by Feautrier. We then define a new schedule $\mathcal{T}$ equal to Θ except for its first dimension which is the sum of the first dimension of Θ and the first dimension of $\mathcal{S}$. We replace Θ by $\mathcal{T}$ and we restart our current construction from the beginning. This is valid as Θ and $\mathcal{T}$ have the same dimension for any statement.

The dimensions of P of rank strictly greater than one must satisfy the dependence relations e_{1+m_1}, ..., e_p.

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Hypotheses 4 and 5. We define the first dimension of P to be equal to the first dimension of Θ: $P_1 = \Theta_1$. Then, Θ_1 is a linear combination of P_1 and P_1 satisfies all the operation to operation dependences which are satisfied by Θ_1.

Hypothesis 3. Let $j \in [1 + m_1, p]$. Then e_j is a dependence not fully satisfied by the first dimension of Θ. As Θ is a schedule it induces a lexicographically nonnegative delay on all dependences. Thus its first dimension induces a nonnegative delay on all dependences. Furthermore, as Θ is integral by hypothesis, its first dimension induces on all dependences a delay either null or greater than one. Let $\mathcal{D}_j^1$ denote the subset of $\mathcal{D}_j$ corresponding to operation to operation dependences of e_j not satisfied by the first dimension of Θ. Then:

$$\mathcal{D}_j^1 = \{i \in \mathcal{D}_j \mid \Theta_1(T_j, i, N) - \Theta_1(S_j, h_j(i, N), N) = 0\}.$$

Then the dimensions of Θ of rank strictly greater than one must satisfy the dependence relations $e^1_{1+m_1}$, ..., e^1_p where, if $j \in [1 + m_1, p]$, $e^1_j = (S_j, T_j, \mathcal{D}^1_j, h_j)$.

6.5. Induction: from dimension i to dimension $i + 1$

We suppose that the induction hypotheses for i are satisfied. Then we build the $(i + 1)$-th dimension of P while satisfying the induction hypotheses for $i + 1$.

The delay induced by Θ_{i+1} is a nonnegative affine function over $\mathcal{D}^i_j$. By induction hypothesis 3, the dimensions of Θ of rank greater than or equal to $i + 1$ must satisfy the dependence relations $e^i_{1+m_i}$, ..., e^i_p. Therefore, Θ_{i+1} induces a nonnegative delay on all the dependence relations $e^i_{1+m_i}$, ..., e^i_p. Let us take any value of j in $[1 + m_i, p]$. As $\Theta_{i+1}(T_j)$, $\Theta_{i+1}(S_j)$, and h_j are affine functions by hypothesis, $\Theta_{i+1}(T_j, x, N) - \Theta_{i+1}(S_j, h_j(x, N), N)$ is also an affine function. Furthermore this affine function is nonnegative over $\mathcal{D}^i_j$, the domain of the dependence relation e^i_j. $\mathcal{D}^i_j$ is obviously a polyhedron (cf. Equation (13)). Therefore, we can apply the affine form of Farkas' lemma (Theorem 1) to $\Theta_{i+1}(T_j, x, N) - \Theta_{i+1}(S_j, h_j(x, N), N)$ and $\mathcal{D}^i_j$.

Applying the affine form of Farkas' lemma. Let $f_1, ..., f_q$ be the affine functions defining polyhedron $\mathcal{D}_j$: $\mathcal{D}_j = \{x \mid \forall k \in [1, q],\ f_k(x, N) \geq 0\}$. Then, by induction hypothesis 3:

$$\mathcal{D}^i_j = \left\{ x \;\middle|\; \begin{array}{ll} \forall k \in [1, q], & f_k(x, N) \geq 0 \\ \forall k \in [1, i], & \Theta_k(T_j, x, N) - \Theta_k(S_j, h_j(x, N), N) \geq 0 \\ \forall k \in [1, i], & -\Theta_k(T_j, x, N) + \Theta_k(S_j, h_j(x, N), N) \geq 0 \end{array} \right\}. \tag{13}$$

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Then the affine form of Farkas' lemma claims the existence of some nonnegative values α, $\beta_1, ..., \beta_q$, $\gamma_1, ..., \gamma_i$, and $\delta_1, ..., \delta_i$, such that:

$$\begin{aligned}\Theta_{i+1}(T_j, x, N) - \Theta_{i+1}(S_j, h_j(x, N), N) = \alpha &+ \sum_{k=1}^{q} \beta_k f_k(x, N) \\ &+ \sum_{k=1}^{i} \gamma_k \Big(\Theta_k(T_j, x, N) - \Theta_k(S_j, h_j(x, N), N) \Big) \\ &- \sum_{k=1}^{i} \delta_k \Big(\Theta_k(T_j, x, N) - \Theta_k(S_j, h_j(x, N), N) \Big).\end{aligned}$$

By induction hypothesis 4, for any value of j in $[1, i]$, Θ_j is a linear combination of $P_1, ..., P_j$. Thus, there exist some nonnegative values $\lambda_{j,1}, ..., \lambda_{j,i}$, and $\mu_{j,1}, ..., \mu_{j,i}$ such that:

$$\begin{aligned}&\Theta_{i+1}(T_j, x, N) - \Theta_{i+1}(S_j, h_j(x, N), N) = \\ &\qquad \alpha + \sum_{k=1}^{q} \beta_k f_k(x, N) + \sum_{k=1}^{i} (\lambda_{j,k} - \mu_{j,k}) \Big(P_k(T_j, x, N) - P_k(S_j, h_j(x, N), N) \Big)\end{aligned}$$

which is equivalent to:

$$\begin{aligned}&\left(\Theta_{i+1}(T_j) + \sum_{k=1}^{i} \mu_{j,k} P_k(T_j) \right)(x, N) - \left(\Theta_{i+1}(S_j) + \sum_{k=1}^{i} \mu_{j,k} P_k(S_j) \right)(h_j(x, N), N) = \\ &\qquad \alpha + \sum_{k=1}^{q} \beta_k f_k(x, N) + \sum_{k=1}^{i} \lambda_{j,k} \Big(P_k(T_j, x, N) - P_k(S_j, h_j(x, N), N) \Big). \quad (14)\end{aligned}$$

The right-hand side of Equation (14) is nonnegative over $\mathcal{D}_j$. By induction hypothesis 1, none of the first i dimensions of P fully satisfies a dependence relation e_l if $l \in [1 + m_i, p]$. Thus, for any value of k in $[1, i]$, P_k induces a nonnegative delay on the dependence relation e_l, if $l \in [1 + m_i, p]$. In other words, with $l = j$:

$$\forall x \in \mathcal{D}_j, \;\; P_k(T_j, x, N) - P_k(S_j, h_j(x, N), N) \geq 0. \quad (15)$$

Furthermore, the functions $f_1, ..., f_q$ are by definition nonnegative on $\mathcal{D}_j$, and the constants α, β_k, and $\lambda_{j,k}$ are nonnegative. Therefore, the right-hand side of Equation (14) is a nonnegative affine combination of functions which take nonnegative values on $\mathcal{D}_j$. Then, this right-hand side takes nonnegative values on $\mathcal{D}_j$, and thus on $\mathcal{D}_j^i$ ($\mathcal{D}_j^i \subset \mathcal{D}_j$).

The left-hand side of Equation (14) is a delay formula. This left-hand side is obviously the formula of the delay induced by the affine function $\Theta_{i+1} + \sum_{k=1}^{i} \mu_{j,k} P_k$ on a dependence defined by the affine function h_j. As the right-hand side of Equation (14) takes nonnegative values on $\mathcal{D}_j$, $\Theta_{i+1} + \sum_{k=1}^{i} \mu_{j,k} P_k$ is an affine function which induces a nonnegative delay on e_j. Furthermore, this function induces a strictly positive delay on any points of $\mathcal{D}_j^i$ on which Θ_{i+1} induces a strictly positive delay (as $\sum_{k=1}^{i} \mu_{j,k} P_k$ induces a nonnegative delay on these points, cf. Equation (15)).

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Looking at dependence relations other than e_j. By induction hypothesis 1, none of the first i dimensions of P fully satisfies a dependence relation e_l if $l \in [1+m_i, p]$. Thus, for any value of k in $[1,i]$, P_k induces a nonnegative delay on $\mathcal{D}_l$, and thus on $\mathcal{D}_l^i$, if $l \in [1+m_i, p]$. By induction hypothesis 3, the components of Θ of rank greater than or equal to $i+1$ satisfy the dependence relations e_l^i, for $l \in [1+m_i, p]$. Thus, Θ_{i+1} induces a nonnegative delay on e_l^i, for $l \in [1+m_i, p]$. Therefore, $\Theta_{i+1} + \sum_{k=1}^{i} \mu_{j,k} P_k$ induces a nonnegative delay on e_l^i, for $l \in [1+m_i, p]$, as a nonnegative linear combination of functions inducing nonnegative delays.

Satisfying the conditions on all the e_k's. In summary, $\Theta_{i+1} + \sum_{k=1}^{i} \mu_{j,k} P_k$ is an affine function which induces:

1. A nonnegative delay on the dependence relation e_j;
2. A nonnegative delay on the dependence relations $e_{1+m_i}^i, ..., e_p^i$;
3. A strictly positive delay on the operation to operation dependences of e_j^i on which Θ_{i+1} induces a strictly positive delay.

Furthermore, a quick look to the above arguments shows that these properties also hold for $\Theta_{i+1} + \sum_{k=1}^{i} \nu_k P_k$, if for any k in $[1,i]$, $\nu_k \geq \mu_{j,k}$. Thus, for any $k \in [1,i]$, let $\mu_k = \max_{j \in [1+m_i, p]} \mu_{j,k}$. Then, $\Theta_{i+1} + \sum_{k=1}^{i} \mu_k P_k$ is an affine function which induces:

1. A nonnegative delay on the dependence relations $e_{1+m_i}, ..., e_p$;
2. A strictly positive delay on the operation to operation dependences of e_l^i ($l \in [1+m_i, p]$) on which Θ_{i+1} induces a strictly positive delay.

Furthermore, a quick look to the above arguments shows that these properties also hold for $\Theta_{i+1} + \sum_{k=1}^{i} \nu_k P_k$, if for any k in $[1,i]$, $\nu_k \geq \mu_k$.

P_{i+1} and induction hypotheses 4 and 5 for dimension $i+1$. We are now ready to define the $(i+1)$-th component of P: $P_{i+1} = \Theta_{i+1} + \sum_{k=1}^{i} (1+\mu_k) P_k$ (thus induction hypothesis 4 is satisfied for $i+1$). From what precedes, P_{i+1} induces a nonnegative delay on all the dependence relations $e_{1+m_i}, ...e_p$, as required. Furthermore, P_{i+1} satisfies all the operation to operation dependences of e_k^i ($k \in [1+m_i, p]$) satisfied by Θ_{i+1}. Let us consider any k in $[1+m_i, p]$, and any operation to operation dependence f in e_k. We have two cases to consider.

1. f belongs to e_k^i. We have already established that P_{i+1} induces a strictly positive delay on f, if Θ_{i+1} does so.
2. f does not belong to e_k^i. Then, by definition of e_k^i, this dependence is satisfied by one of the first $(i-1)$ dimensions of Θ, say the l-th dimension. Thus, Θ_l induces a strictly positive delay on f. Thus, because of induction hypothesis 5 at depth l, P_l induces a strictly positive delay on f. Therefore, P_{i+1} satisfies the operation to operation dependence f, as the coefficient of P_l in P_{i+1} is $(1+\mu_l)$ and not just μ_l.

Therefore, P_{i+1} satisfies all the operation to operation dependences satisfied by Θ_{i+1}. Furthermore, P_{i+1} fully satisfies any dependence relation e_j, $j \in [1+m_i, p]$, if e_j^i is fully satisfied by Θ_{i+1}.

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Induction hypothesis 2 for dimension $i+1$. For P_{i+1} to be the $(i+1)$-th dimension of a schedule *a la* FEAUTRIER, we still have to check that P_{i+1} fully satisfies as many dependence relations as possible, among $e^i_{1+m_i}, ..., e^i_p$. If this is not the case, Θ_{i+1} does not fully satisfy all the fully satisfiable e^i_j. Then, we add to the $(i+1)$-th dimension of Θ an affine function that induces a nonnegative delay on all the dependence relations $e^i_{1+m_i}, ..., e^i_p$, and which fully satisfies as many of them as possible. This way we obtain a new schedule of minimal dimension. We restart our recursive construction (directly at the $(i+1)$-th rank). This time P_{i+1} maximizes the number of the dependence relations it fully satisfies.

Induction hypothesis 1 for dimension $i+1$. Without any loss of generality, let the dependence relations fully satisfied by P_{i+1} be $e_{1+m_i}, ..., e_{m_{i+1}}$. Then, the first $(i+1)$ dimensions of P satisfy the dependence relations e_1, ..., $e_{m_{i+1}}$, but none of the dependence relations $e_{1+m_{i+1}}$, ..., e_p.

Induction hypothesis 3 for dimension $i+1$. As P_{i+1} only fully satisfies the dependence relations $e_{1+m_i}, ..., e_{m_{i+1}}$, then Θ_{i+1} only fully satisfies $e^i_{1+m_i}, ..., e^i_{m_{i+1}}$. Θ, and thus Θ_{i+1}, is integral by hypothesis. Then the dimensions of Θ of rank greater than or equal to $i+2$ must satisfy the dependence relations $e^{i+1}_{1+m_{i+1}}$, ..., e^{i+1}_p where, if $j \in [1+m_{i+1}, p]$, $e^{i+1}_j = (S_j, T_j, \mathcal{D}^{i+1}_j, h_j)$, with

$$\mathcal{D}^{i+1}_j = \{x \in \mathcal{D}^i_j \mid \Theta_k(T_{i+1}, x, N) - \Theta_{i+1}(S_j, h_j(x, N), N) = 0\}$$

which, by induction hypothesis 3, is equivalent to:

$$\mathcal{D}^{i+1}_j = \{x \in \mathcal{D}_j \mid \forall k \in [1, i+1],\ \Theta_k(T_j, x, N) - \Theta_k(S_j, h_j(x, N), N) = 0\}.$$

7. CONCLUSION

Feautrier's scheduling algorithm is the most powerful existing algorithm for parallelism detection and extraction. But it has always been known to be suboptimal. We have shown that Feautrier's algorithm does not miss any significant amount of parallelism *because* of its design, even if one can design a greedier algorithm. Therefore, to improve Feautrier's algorithm or to build a more powerful algorithm, one must get rid of some of the restrictive hypotheses underlying its framework: affine schedules — but more general schedules will cause great problems for code generation — and one scheduling function by statement — Feautrier, Griebl, and Lengauer have already begun to get rid of this hypothesis by splitting the iteration domains [11].

What Feautrier historically introduced as a "greedy heuristic" is nothing but the most powerful algorithm in its class! This is a sound and powerful tool that may safely be used to build other program optimization algorithms.

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