Kernel Level Speculative DSM

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Motivation

- Main interest is performance, fault-tolerance, and correctness of distributed systems
- Present our ideas in the context of a DSM system
- We are developing tools that
 - Improve performance
 - Address reliability
 - Simplify programming of distributed applications





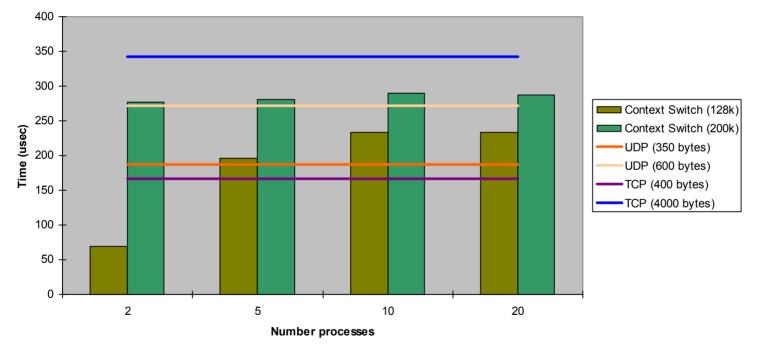
KDSM System Design

- System V IPC is a useful, popular model
- When generalizing to distributed IPC, semantics must be preserved
- Programs expect *sequential consistency*
- Total order communication protocol
 - Obvious choice
 - Usually expensive





KDSM Implementation issues: kernel level module



Benchmarks collected using LMBench tool on a cluster of RedHat Linux 8.0 (kernel v.2.4.19) machines, connected through 100Mb LAN, each a dual processor PIII 700Mhz

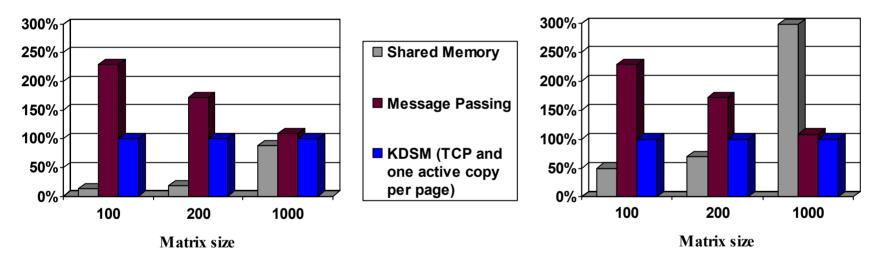




Preliminary Results: matrix multiplication

Normalized running times per thread

Total normalized running times



Results obtained on a cluster of RedHat Linux 8.0 (kernel v.2.4.19) machines, connected through 100Mb LAN, each a dual processor PIII 700Mhz



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Speculations and fault tolerance

```
Transfer (file1, file2, k) {
   //we want operation to be atomic
   // read and write operations are atomic
   if (read (file1, buf1, k) != k)
      return failure:
   if (read (file2, buf2, k) != k)
      return failure;
   if (write(file1, buf2, k) != k)
      return failure;
   if (write(file2, buf1, k) != k)
     // Undo first write
     write(file1, buf1, k);
     // Unrecoverable error on write failure
     // Inconsistent state
     return failure;
   return success;
```

Transfer (file1, file2, k) {
 speculate { // Speculation entry
 // read_spec and write_spec throw
 // an exception in case of failure
 read_spec (file1, buf1, k);
 read_spec (file2, buf2, k);
 write_spec (file1, buf2, k);
 write_spec (file2, buf1, k);
 return success;
 // Speculation implicitly commited
 } catch { // Abort speculation
 return failure;
 }
}



}

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}



Speculations

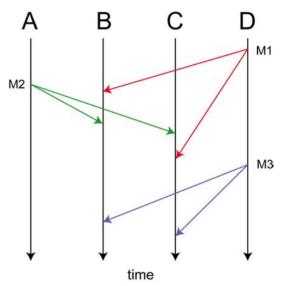
- Speculation
 - Computation based on an assumption that has not been verified
- Provide a simple programming paradigm
- Represented by three operations
 - Speculate
 - Commit
 - Abort





Speculations and total order communication

- Total order communication:
 - N machines communicating with each other
 - Each machine has a unique id: m
 - Each message is uniquely stamped
- Goal
 - Message delivery is based on global total order relation on stamps



Correct message delivery order: M1 < M2 < M3

- Problem
 - Delivering a message can have higher latency than desired







Speculations and total order communication

- Enter speculation when message arrives
- If speculation is valid at safe delivery time then commit
- Otherwise abort and start over

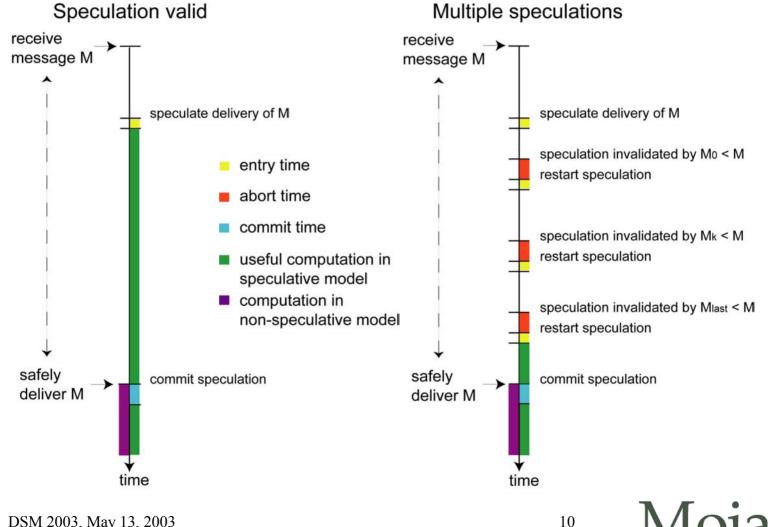
```
T – a waiting window
time – current local time
last(P) – last message from machine P
```

```
1: read message M from network; t<sub>0</sub> <- time
 2: while (time < t_0+T) \Lambda (3P.last(P) < M) do
        process messages from network
 3:
    end while
 4 •
 5: if (\exists P.last(P) < M) then
        enter speculation; deliver M
 6:
        abort if receive message M<sub>o</sub> s.t. M<sub>o</sub> < M
 7:
      commit when \forall P.last(P) > M
 8:
 9: else
10:
        deliver M
11: end if
```





Mathematical Model





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Speculations and distributed locking

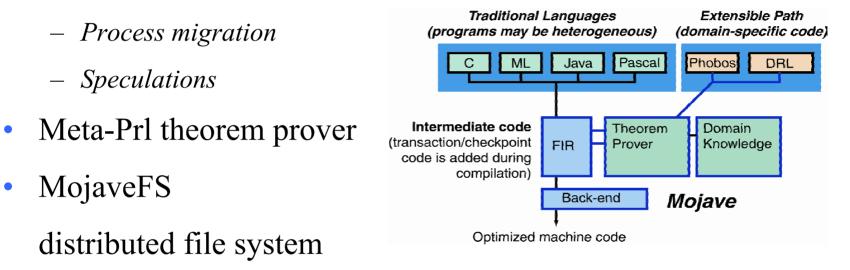
- Shared memory needs complementary synchronization mechanisms
- Semaphores can be easily implemented using speculations
- Speculations provide optimistic lock acquisition
 - Enter a new speculations when request for lock is issued
 - Commit speculation when lock is granted
 - Abort otherwise and try again.





Mojave System

• Compiler for ML, C, Java with support for



• KDIPC kernel level distributed inter-process communication



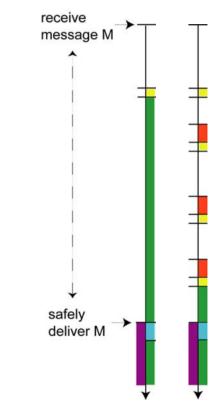


Experimental Results – Speculation overhead

	Operation (and mutation percentile)				
	(time in μ sec)				
Heap	Entry	Abort		Commit	
Size		10%	100%	10%	100%
100k	27	65	84	57	54
200k	40	120	135	81	87
1000k	63	131	466	111	109



- abort time
- commit time
- useful computation in speculative model
- computation in non-speculative model



Results obtained on a cluster of RedHat Linux 8.0 (kernel v.2.4.19) machines, connected through 100Mb LAN, each a dual processor PIII 700Mhz





Conclusions

- Speculations address fault-tolerance and performance
 - Speculations can be used to implicitly provide fault-tolerance thus a simpler programming model
 - Improve performance by speculating instead of waiting for a condition to be verified
- Used speculations to show how we can improve a DSM system which supports sequential consistency





Future work

- Support full System V IPC API
 - Semaphores
- Use speculations to design faster communication protocols
- Distributed file system design
 - Support for speculations
 - Support for totally transparent process migration







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