

# Measures of maximal entropy of SFT on lattices and trees

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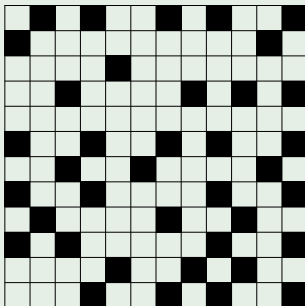
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### Example: two-dimensional Fibonacci SFT

Set of configurations without two consecutive black squares, vertically or horizontally.



# Outline of the talk

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# One-dimensional subshift of finite type

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## Subshift of finite type

The **subshift of finite type** associated to  $A$  is the set  $\Sigma_A$  of words  $w \in \mathcal{A}^{\mathbb{Z}}$  such that if  $A_{i,j} = 0$ ,  $w$  does not contain the pattern  $ij$ .

$$A_{i,j} = \begin{cases} 1 & \text{if } ij \text{ is an allowed pattern,} \\ 0 & \text{if } ij \text{ is a forbidden pattern.} \end{cases}$$

$$\Sigma_A = \{w \in \mathcal{A}^{\mathbb{Z}}; \forall k \in \mathbb{Z}, A_{w_k, w_{k+1}} = 1\}.$$

In what follows, we assume that the matrix  $A$  irreducible and aperiodic.

Let  $\Sigma_A$  be a SFT, and let  $\mathcal{W}_k$  be the set of allowed words of length  $k$ .

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### Questions:

- 1 What is the growth rate of  $|\mathcal{W}_k|$ ?  
Precisely, we would like to be able to compute the **topological entropy** of the SFT:

$$h(\Sigma_A) = \lim_{k \rightarrow \infty} \frac{\log |\mathcal{W}_k|}{k}.$$

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- 2 What do “typical” configurations look like?  
How to generate “uniformly” patterns of  $\Sigma_A$ ?

# Topological entropy

From Perron-Frobenius theorem, the matrix  $A$  has an eigenvalue  $\lambda > 0$  such that  $|\mu| \leq \lambda$  for any other eigenvalue  $\mu$ .

## Proposition

$$h(\Sigma_A) = \lim_{k \rightarrow \infty} \frac{\log |\mathcal{W}_k|}{k} = \log \lambda.$$

Furthermore, there is a unique choice of  $r_1, \dots, r_n \geq 0$  such that  $\sum_{i=1}^n r_i = 1$  and

$$A \begin{pmatrix} r_1 \\ \vdots \\ r_n \end{pmatrix} = \lambda \begin{pmatrix} r_1 \\ \vdots \\ r_n \end{pmatrix}.$$

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### Definition of the Parry measure

The **Parry measure** is the Markov measure  $\pi$  of transition matrix  $P$  defined, for any  $i, j \in \mathcal{A}$ , by

$$P_{i,j} = A_{i,j} \frac{r_j}{\lambda r_i}.$$

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*Remark.*

$$P_{i,j} = A_{i,j} \frac{r_j}{\sum_{k \in \mathcal{A}} A_{i,k} r_k} = A_{i,j} \frac{r_j}{\sum_{k \in \mathcal{S}(i)} r_k},$$

where  $\mathcal{S}(i) = \{k \in \mathcal{A}; A_{i,k} = 1\}$ .



# Markov-uniform property of the Parry measure

## Proposition

The Parry measure is **Markov-uniform**: for given  $k \geq 1$  and  $a, b \in \mathcal{A}$ , the value

$$\pi(awb)$$

does not depend on the word  $w \in \mathcal{A}^k$  such that  $awb$  is allowed.

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*Proof.* By definition,  $P_{i,j} = A_{i,j} \frac{r_j}{\lambda r_i}$ . If  $awb \in \mathcal{W}_{k+2}$ , then:

$$\pi(awb) = \pi_a P_{a,w_1} P_{w_1,w_2} \dots P_{w_{k-1},w_k} P_{w_k,b}$$

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# Measure of maximal entropy

## Theorem

Let  $\mathcal{M}_{\Sigma_A}$  be the set of translation invariant measures on the SFT  $\Sigma_A$ , and let  $\pi \in \mathcal{M}_{\Sigma_A}$ . The following properties are equivalent.

- (i)  $\pi$  is the **Parry measure** associated to  $\Sigma_A$ ,
- (ii)  $\pi$  is a **Markov-uniform** measure on  $\Sigma_A$ ,
- (iii)  $\pi$  is the measure of **maximal entropy** of  $\Sigma_A$ ,
- (iv) the entropy of  $\pi$  is equal to the **topological entropy**  $h(\Sigma_A)$ .

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- (iv) the entropy of  $\pi$  is equal to the **topological entropy**  $h(\Sigma_A)$ .

*Remark.* On  $\mathbb{Z}^d$ ,  $d \geq 2$ , there can be several measures of maximal entropy. But the equivalence between (ii), (iii), and (iv) can be extended to some multi-dimensional SFT.

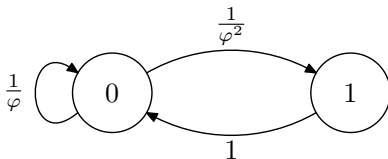
## Example: one-dimensional Fibonacci SFT

Let  $\mathcal{A} = \{0, 1\}$ . The **one-dimensional Fibonacci SFT** is the set of words that do not contain two consecutive 1's. It is given by:

$$A = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}.$$

Its topological entropy is equal to  $\log \varphi$ , where  $\varphi = \frac{1+\sqrt{5}}{2}$ .

The Parry measure is the Markov measure given by



with  $\pi_0 = \frac{\varphi^2}{1+\varphi^2}$  and  $\pi_1 = \frac{1}{1+\varphi^2}$ .

## First way to generate the Parry measure

The Parry measure of the Fibonacci SFT can be generated by:

- choosing independently to write a 0 with probability  $r_0 = \frac{1}{\phi}$  and a 1 with probability  $r_1 = \frac{1}{\phi^2}$ ,
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*Proof.*

$$P_{i,j} = A_{i,j} \frac{r_j}{\sum_{k \in \mathcal{S}(i)} r_k}.$$

## Second way to generate the Parry measure

The Parry measure of the Fibonacci SFT can be generated by:

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### Confluent SFT

A SFT is **confluent** if for any  $i, j, k \in \mathcal{A}$  such that both  $ij$  and  $jk$  are forbidden, then  $i = k$ .

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### Proposition

For **confluent** SFT, the Parry measure can be generated by independent draws of letters and deletion of forbidden patterns.

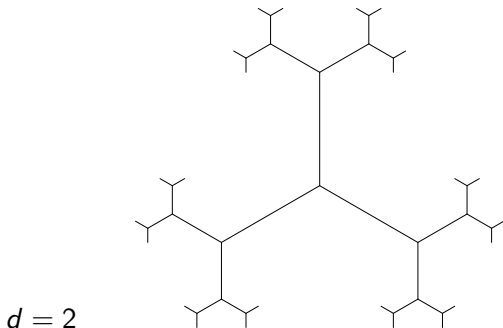
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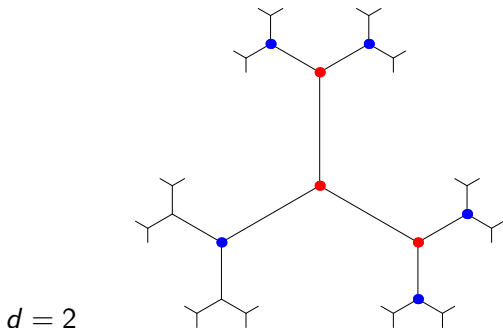
**Question:** how to construct **Markov-uniform** measures on  $\Sigma_A^d$ ?





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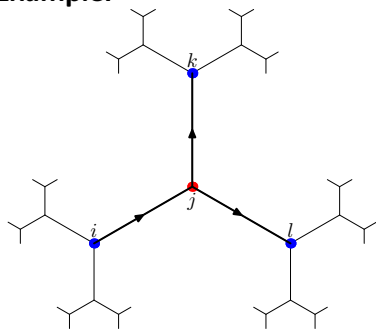


**Idea 1:** consider a (reversible) Markov chain  $P$  on the alphabet  $\mathcal{A}$ , of stationary distribution  $\pi$ .  
Choose the letter at one given vertice according to  $\pi$  and then label the other vertices using  $P$ .

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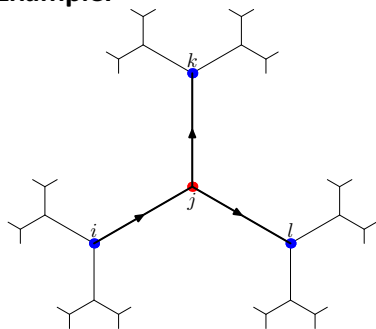


Probability of this pattern:  
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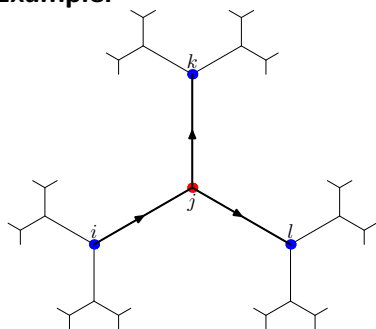
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For given  $i, k, l$ , we want this value to be independent of the letter  $j$  such that the pattern is allowed.

**Idea 2:** like for the Parry measure, choose  $P$  under the form:

$$P_{i,j} = A_{i,j} \frac{r_j}{\sum_{s \in \mathcal{A}} A_{i,s} r_s} = A_{i,j} \frac{r_j}{\sum_{s \in \mathcal{S}(i)} r_s},$$

for some probability vector  $(r_i)_{i \in \mathcal{A}}$ .

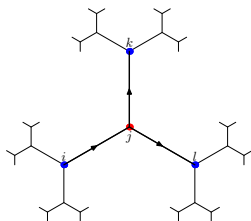
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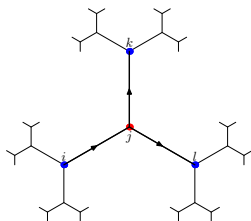
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Let us try to choose  $(r_i)_{i \in \mathcal{A}}$  such that:

$$\sum_{s \in \mathcal{A}} A_{j,s} r_s = \lambda r_j^{1/2},$$

for any  $j \in \mathcal{A}$  !





For a tree of degree  $d + 1$ , the problem is to find a probability distribution  $(r_i)_{i \in \mathcal{A}}$  such that for some  $\lambda > 0$ ,

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### Proposition

Let  $A$  be an irreducible non-negative matrix, and let  $d \geq 1$ . There exist  $\lambda > 0$  and  $r_1, \dots, r_n > 0$  satisfying  $\sum_{i=1}^n r_i = 1$  and:

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*Proof.* Fixed point theorem.

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*Proof.* Fixed point theorem.

*Remark.*  $\lambda$  and  $(r_i)_{i \in \mathcal{A}}$  may not be unique.

## Proposition

If the distribution of probability  $(r_i)_{i \in \mathcal{A}}$  satisfies

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We call such a measure a  **$d$ -Parry measure**.

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*Remark.* The reversible invariant measure is given by  $\pi_i = \gamma r_i^{1+1/d}$ .

## Example: Fibonacci SFT on trees

We search  $P = \begin{pmatrix} \alpha & 1 - \alpha \\ 1 & 0 \end{pmatrix}$ , such that

$$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \alpha \\ 1 - \alpha \end{pmatrix} = \lambda \begin{pmatrix} \alpha^{1/d} \\ (1 - \alpha)^{1/d} \end{pmatrix}.$$

For any  $d \geq 1$ , there exists a unique  $d$ -Parry measure, which is given by  $r_0 = \alpha$  and  $r_1 = 1 - \alpha$ , where  $\alpha$  is the unique positive solution of the equation

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For  $d = 1$ , we recover  $r_0 = \frac{1}{\varphi}$  and  $r_1 = \frac{1}{\varphi^2}$ .



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But we also have examples of SFT having several  $d$ -Parry measures...

**Question:** is there a relation between these  $d$ -Parry measures and some analogue of entropy on trees?

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### $f$ -invariant of Bowen

The  $f$ -**invariant** is a measure-conjugacy invariant, introduced by L. Bowen to generalize the theory of entropy to free group actions.

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Markov measures maximise the  $f$ -invariant.

## Proposition

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## Questions:

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- Is there a relation between all this and some growth rate of the number of allowed patterns?

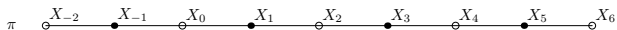


# Outline

- 1 One-dimensional SFT and the Parry measure
- 2 SFT defined on trees and  $d$ -Parry measures
- 3 SFT and probabilistic cellular automata

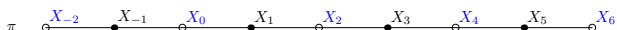
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Consider a configuration distributed according to the Parry measure  $\pi$  of the Fibonacci SFT.



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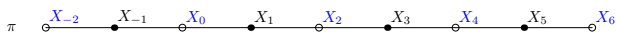


For all  $i \in \mathbb{Z}$ , if  $X_{2i} = X_{2i+2} = 0$ , we flip the value of  $X_{2i+1}$  with probability  $1/2$ .

By the [Markov-uniform](#) property, the new sequence is still distributed according to  $\pi$ .

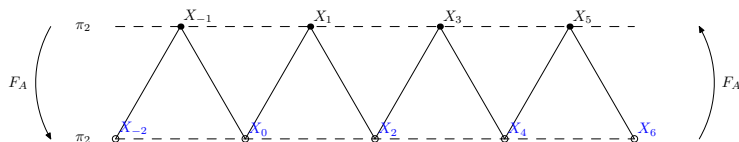
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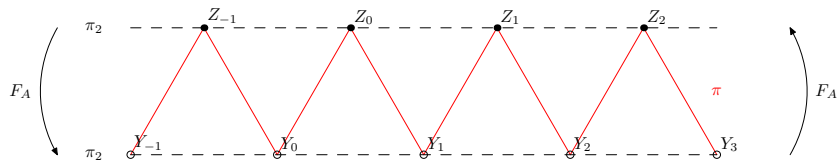
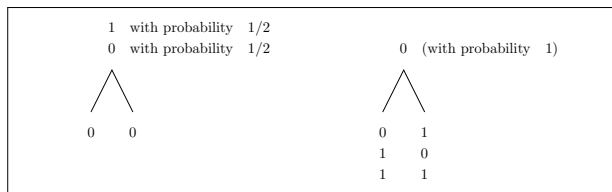


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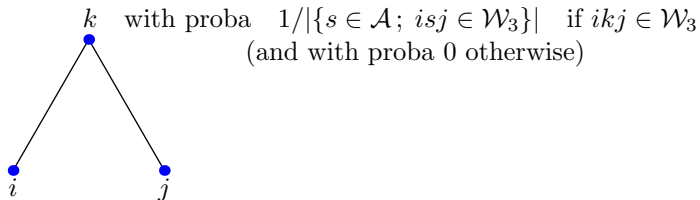
# Example: one-dimensional Fibonacci SFT



The projection  $\pi_2$  of the Parry measure on odd (resp. even) sites is an invariant measure of the **probabilistic cellular automaton**.

# General one-dimensional SFT

For a general one-dimensional SFT  $\Sigma_{\mathcal{A}}$ , let us consider the PCA  $F_{\mathcal{A}}$  defined by:

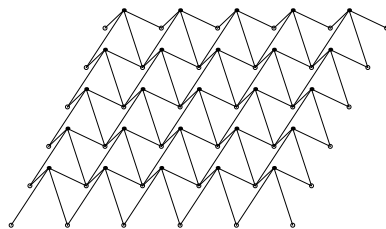
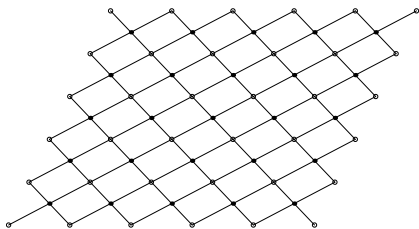




The same result holds on  $\mathbb{Z}^d$ ,  $d \geq 2$  and on infinite trees (bipartite graphs).

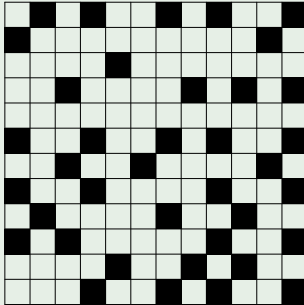


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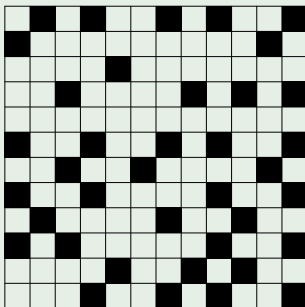


Two-dimensional case.

## Back to the two-dimensional Fibonacci SFT

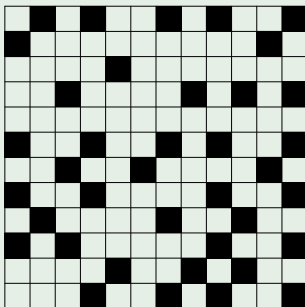


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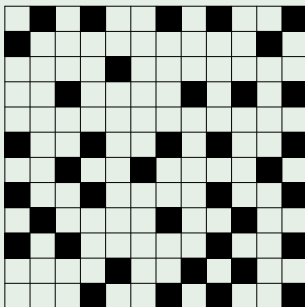
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## Back to the two-dimensional Fibonacci SFT



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- Using PCA, new ideas for **perfect sampling** of patterns according to the measure of maximal entropy of the SFT.
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- Extension to sofic SFT?