

# MPFI: a library for arbitrary precision interval arithmetic

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# Reliable computing: Interval Arithmetic

(Moore 1966, Kulisch 1983, Neumaier 1990, Rump 1994, Alefeld and Mayer 2000. . . )

**Numbers are replaced by intervals.**

Ex:  $\pi$  replaced by  $[3.14159, 3.14160]$

**Computations:** can be done with interval operands.

**Advantages:**

- Every result is guaranteed (including rounding errors).
- Data known up to measurement errors are representable.
- Global information can be reached (“computing in the large”).

**Drawbacks:** overestimation of the results.

# Hansen's algorithm for global optimization

Hansen 1992

$\mathcal{L}$  = list of boxes to process :=  $\{X_0\}$        $F$ : function to minimize

while  $\mathcal{L} \neq \emptyset$  loop

  suppress  $X$  from  $\mathcal{L}$

**reject  $X$  ?**

  yes if  $F(\bar{X}) > \bar{f}$

  yes if  $\text{Grad}F(\bar{X}) \neq 0$

  yes if  $HF(\bar{X})$  has its diag. non  $> 0$

**reduce  $X$**

  Newton applied with the gradient

  solve  $Y \subset X$  such that  $F(\bar{Y}) \leq \bar{f}$

**bisect  $Y$**  into  $Y_1$  and  $Y_2$  if  $Y$  is not a result

  insert  $Y_1$  and  $Y_2$  in  $\mathcal{L}$

# Solving linear systems preconditioned Gauss-Seidel algorithm

Hansen & Sengupta 1981

**Linear system:**  $Ax = b$  with  $A$  and  $b$  given.

**Problem:** compute an enclosure of

$\text{Hull}(\Sigma_{\exists\exists}(A, b)) = \text{Hull}(\{x : \exists A \in A, \exists b \in b, Ax = b\})$ .

**Hansen & Sengupta's algorithm**

compute  $C$  an approximation of  $\text{mid}(A)^{-1}$

apply Gauss-Seidel to  $CAx = Cb$  until convergence.

**Idea (hope):**

$CA$  contains the identity matrix, is diagonally dominant, thus the iteration matrix has a spectral radius  $< 1$  and this iteration is a **contraction**.

# Agenda

- **Introduction: when is more precision wanted?**
  - global optimization (Hansen's algorithm)
  - linear system solving (Hansen & Sengupta's algorithm)
- **MPFI: a library for arbitrary precision interval arithmetic**
  - general view of MPFI
  - MPFI in details
  - other libraries
- **Applications**
  - robot kinematics
  - isolation of real roots of a polynomial
  - approximation of zeros of a function

# MPFI Multiple Precision Floating-point Interval arithmetic library

## general view

- **what:** C library for arbitrary precision interval arithmetic
- **why:** to compute reliably and accurately
- **where:** freely available, including source code and documentation (last release: April 2002)
- **who:** N. Revol and F. Rouillier
- **how:** based on MPFR for arbitrary precision floating-point arithmetic (based on GMP: efficient and portable)

# MPFI: in details

- **Requirements:** directed roundings are required, exact is better. Even for elementary functions! (provided by MPFR)
- **Interval:** connected closed subset of  $\mathbb{R}$
- **Operation:**  $op(\mathbf{X}_1, \dots, \mathbf{X}_n) \supset \{op(x_1, \dots, x_n) : x_i \in \mathbf{X}_i\}$
- **Undefined operation:** NaN if  $op(x_1, \dots, x_n)$  is undefined (no *a priori* intersection with the domain)
- **Precision:** intervals are represented by their endpoints, each endpoint carries its own precision.

# MPFI: functionalities

- arbitrary precision interval data type
- arithmetic operations:  $+$ ,  $-$ ,  $\times$ ,  $/$ ,  $\sqrt{\quad}$
- constants:  $\pi$ ,  $\log 2$ , Euler constant
- elementary functions:  $\exp$ ,  $\log$ ,  $\text{atan}$ ,  $\text{cosh}$ ,  $\text{asinh}$ . . .
- IEEE-754 special values:  $\infty$ , signed zeros, NaN
- conversions to and from integer, double, exact naturals, exact integers, rational, “reals” (MPFR numbers)
- Input/Output (to be improved)



# Performances

## Development version

Gaussian elimination on a random M-matrix ( $300 \times 300$ )

ARITHM	double	fi_lib	MPFI	MPFI	MPFI	MPFI
PREC	53 bits	53 bits	53 bits	63 bits	127 bits	255 bits
TIME (s)	0.42	4.39	21.75	22.35	24.11	34.09
RATIO	0.10	1.00	4.95	5.09	5.49	7.76

ARITHM			MPFI	MPFI	MPFI	MPFI
PREC			511 bits	1023 bits	2047 bits	4095 bits
TIME (s)			51.61	112.30	285.57	817.05
RATIO			11.75	25.58	65.05	186.11

(Athlon 1GHz)

## Other libraries

Most of the interval arithmetic libraries are based on machine floating-point precision (Profil/BIAS, Intlib, Sun Forte. . . )

Some libraries offer multiple precision:

- **Maple, Mathematica:** floating-point arithmetic and roundings bugged
- **range:** “the interval is only crudely represented” (Aberth)
- **XSC:** limited use (in long accumulators)
- **IntLab:** limited set of operations and functions

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- **MPFI: a library for arbitrary precision interval arithmetic**
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# Parallel robot kinematics

**Formulation:** algebraic system depending on 7 to 12 variables.

**Goal:** computing all the real solutions (up to 40).

**Usual numerical computational strategies** failed on this example using hardware doubles.

## **Solution obtained by**

- a few exact methods
- MPFI using 128 bits of precision

**Remark:** some preliminary computations were performed using exact computations (Gröbner bases).

# Isolation of polynomial real roots

Rouillier and Zimmermann 2002

## Descartes' rule of sign

Let  $P = \sum_{i=0}^d p_i x^i$  be a polynomial,

$V(P)$  = number of sign changes in  $(p_0, \dots, p_d)$

and  $pos(P)$  = number of  $> 0$  real roots of  $P$ , counted with mult.

Then  $pos(P) \leq V(P)$  and  $V(P) - pos(P)$  is even.

# Isolation of polynomial real roots algorithm

## “Uspensky’s” algo.

Assumption:  $P$  is a square-free polynomial with roots in  $[0, 1]$ .

Transform  $P$  into  $Q = \sum_i q_i x^i$  such that  $V(Q) = 0$  or  $1$

*i.e.*  $P$  has 0 or 1 root on a given interval:

a root of  $P$  is **isolated**.

## Modified algorithm using MPFI

replace exact computations in *Transform  $P$  into  $Q$*

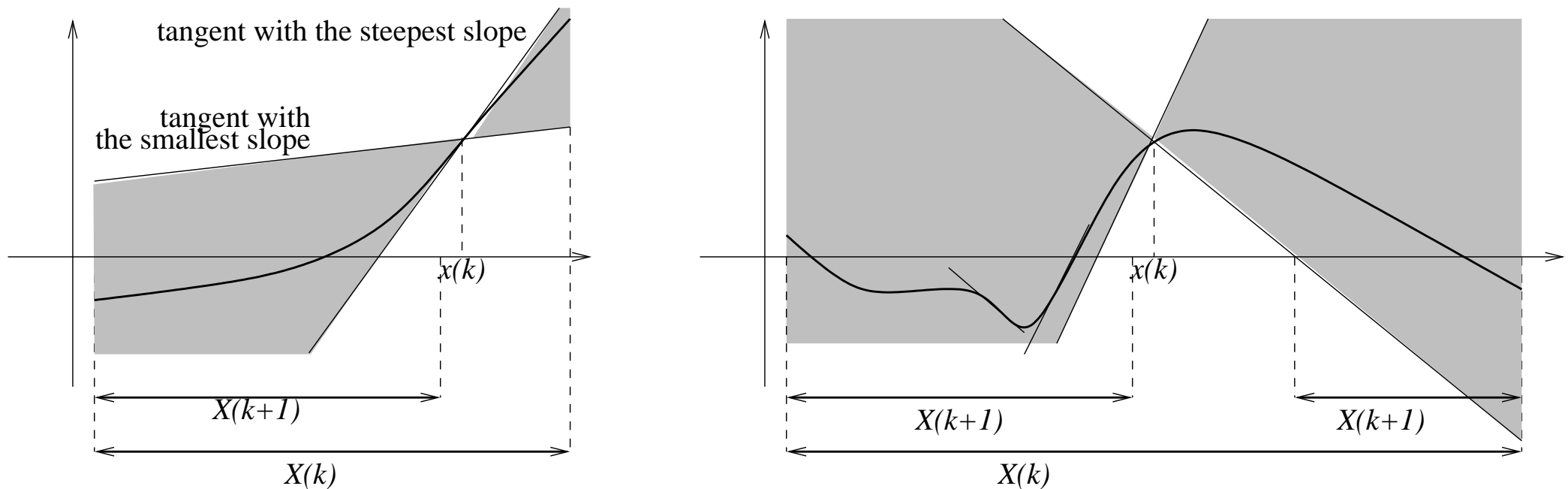
by MPFI computations:

$q_i$  is now an interval.

If  $q_i \ni 0$  then stop interval computation, increase MPFI precision or go on with exact computation.

# Approximation of zeros: interval Newton algorithm principle of one iteration

(Greenberg & Hansen 1983, Kearfott 1995, Mayer 1995, van Hentenryck et al. 1997. . . )



The result will be a list of intervals.

# Interval Newton algorithm

**Input:**  $F, F', X_0$  //  $X_0$  initial search interval    **Initialization:**  $\mathcal{L} = \{X_0\}$ ,

$\alpha = 0.75$  // any value in  $]0.5, 1[$  is suitable

**Loop:** while  $\mathcal{L} \neq \emptyset$

    Suppress  $(X, \mathcal{L})$

    Increase the working precision if needed

$x := \text{mid}(X)$

$(X_1, X_2) := \left(x - \frac{F(\{x\})}{F'(X)}\right) \cap X$  //  $X_1$  and  $X_2$  can be empty

    if  $w(X_1) > \alpha w(X)$  or  $w(X_2) > \alpha w(X)$  then  $(X_1, X_2) := \text{bisect}(X)$

    if  $X_1 \neq \emptyset$  and  $F(X_1) \ni 0$  then

        if  $w(X_1)/|\text{mid}(X_1)| \leq \varepsilon_X$  and  $w(F(X_1)) \leq \varepsilon_Y$  then Insert  $X_1$  in Res

        else Insert  $X_1$  in  $\mathcal{L}$

    same handling of  $X_2$

**Output:** *Res*, a list of intervals that may contain the roots.



## Interval Newton algorithm: experiments

**Chebyshev polynomials:**  $C_n(\cos(\theta)) = \cos(n\theta)$ .

**Results:** very precise roots for degrees up to 40, with proof of existence and uniqueness.

**Wilkinson polynomial:**  $\prod_{i=1}^{20} (X - i)$ .

**With enough precision to be able to store exactly the coefficients:** roots found with a precision  $5 \cdot 10^{-2}$  and a proof of existence (but not uniqueness). A lot of intervals are not eliminated:  $[0.96, 1.02]$  et  $[1.62, 20.984]$ .

**With enough precision and a perturbation  $\pm 2^{-19}$  on the coefficient of  $X^{19}$ :** roots (with proof of existence but not uniqueness):  $1 \pm 4 \cdot 10^{-2}$ ,  $2 \pm 5 \cdot 10^{-2}$ ,  $3 \pm 4 \cdot 10^{-2}$ ,  $4 \pm 4 \cdot 10^{-2}$ ,  $5 \pm 4 \cdot 10^{-2}$ ,  $6 \pm 5 \cdot 10^{-2}$ ,  $7 \pm 6 \cdot 10^{-2}$  and  $[7.91, 22.11]$ . A lot of intervals are not eliminated:  $[0.96, 22.64]$ .

# Conclusion

## MPFI:

- library for arbitrary precision interval computation
- at a reasonable cost.

## Applications

- guaranteed results with arbitrary accuracy
- algorithms: increase precision when needed and continue

## Coming soon

- MPFI++
- linear algebra
- more algorithms.