

# Reliable and accurate solutions of linear and nonlinear systems

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# Reliable computing: Interval Arithmetic

(Moore 1966, Kulisch 1983, Neumaier 1990, Rump 1994, Alefeld and Mayer 2000. . . )

**Numbers are replaced by intervals.**

Ex:  $\pi$  replaced by  $[3.14159, 3.14160]$

**Advantages:** every result is guaranteed.

Data known up to measurement errors are representable.

**Drawbacks:** overestimation of the results.

**Variable dependency**

$$\begin{aligned} I * I &= \{x * y ; x \in I, y \in I\} \\ &\neq I^2 = \{x^2 ; x \in I\} \end{aligned}$$

**Wrapping effect**



# Hansen's algorithm for global optimization

$\mathcal{L}$  = list of boxes to process :=  $\{X_0\}$

while  $\mathcal{L} \neq \emptyset$  loop

suppress  $X$  from  $\mathcal{L}$

**reject  $X$  ?**

yes if  $F(\bar{X}) > \bar{f}$

yes if  $\text{Grad}F(\bar{X}) \neq 0$

yes if  $HF(\bar{X})$  has its diag. non  $> 0$

**reduce  $X$**

Newton applied with the gradient

solve  $Y \subset X$  such that  $F(\bar{Y}) \leq \bar{f}$

**bisect  $Y$**  into  $Y_1$  and  $Y_2$  if  $Y$  is not a result

insert  $Y_1$  and  $Y_2$  in  $\mathcal{L}$

# When is extra precision needed?

**Target application:** global optimization of a continuous function.

**Difficulties with** “very flat valley” function and with “egg-box” function.

**Solution: bisection intervals**

if  $X = X_1 \cup X_2$ , then  $F(X_1) \cup F(X_2) \subset F(X)$

and the left interval is usually tighter than the right one.

**Theorem:** Under mild assumptions on  $f$  and with a direct evaluation,

$$\text{dist}(f(X), F(X)) \leq \mathcal{O}(w(X))$$

**Moral:** in order to refine the accuracy on the results, the precision on the inputs must be refined. . . as much as required.

# Accurate computing: arbitrary precision

(Knuth, Schönhage & Strassen, Brent 1976 . . . )

**Representable numbers** are floating-point numbers with arbitrary precision, chosen by the user.

**MPFI:** Multiple Precision Floating-point Interval arithmetic library.

- **what:** a C library, based upon MPFR ([www.mpfr.org](http://www.mpfr.org)), a multiple precision floating-point library, providing directed exact rounding, based upon GMP ([www.swox.se/gmp](http://www.swox.se/gmp)) for portability and efficiency.
- **why:** Maple or Mathematica packages are incorrect.
- **who:** N. Revol and F. Rouillier.
- **how:** sources and documentation available from [www.ens-lyon.fr/~nrevol](http://www.ens-lyon.fr/~nrevol)

# Reliable and accurate computing: transforming usual algorithms

**Usual approach:** (cf. **iRRAM**)

if accuracy of result is not sufficient then  
    increase computing precision  
    restart whole computation

**Goal: avoid restarting whole computation.**

if accuracy of result is not sufficient then  
    increase computing precision  
    go on

**Stopping criterion:** arbitrary accuracy can be required **both** on results and on residuals.

# Agenda

- **Reliable and accurate computations**

- reliability: interval arithmetic
- accuracy: arbitrary precision
- MPFI: multiple precision interval arithmetic library

- **Solving nonlinear equations**

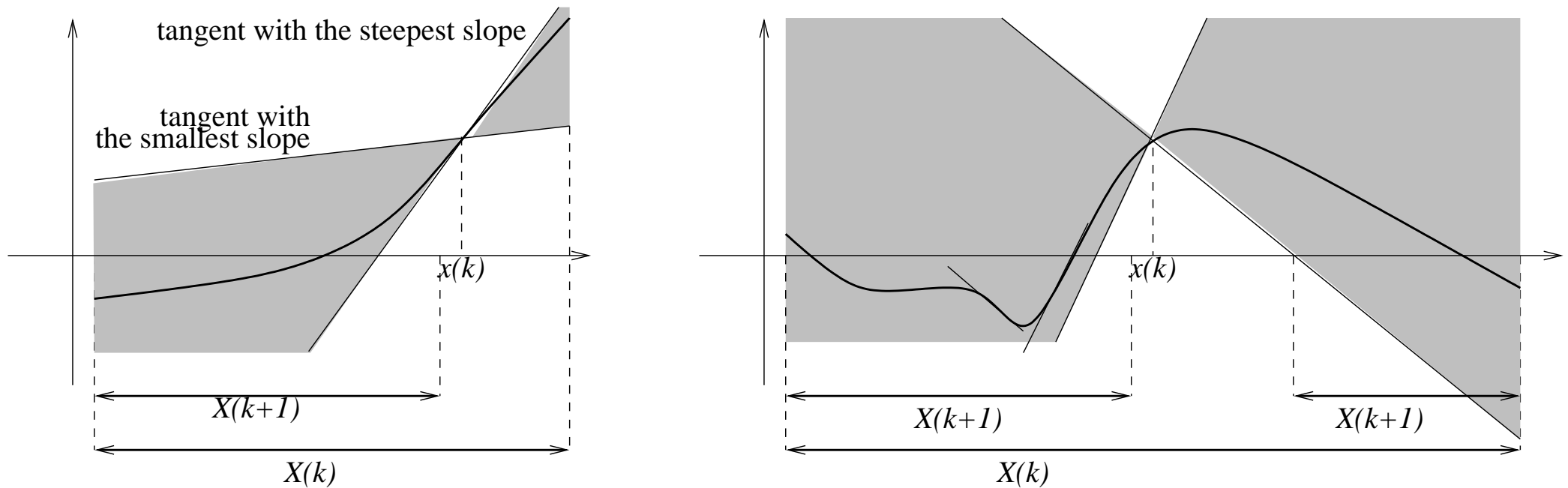
- interval Newton algorithm
- some properties
- experimental results

- **Solving linear systems**

- Hansen & Sengupta's algorithm
- adaptation to multiple precision

# Interval Newton algorithm principle of one iteration

(Greenberg & Hansen 1983, Kearfott 1995, Mayer 1995, van Hentenryck et al. 1997. . . )



The result will be a list of intervals.

# Interval Newton algorithm

**Input:**  $F, F', X_0$  //  $X_0$  initial search interval

**Initialization:**  $\mathcal{L} = \{X_0\}, \alpha = 0.75$  // any value in  $]0.5, 1[$  is suitable

**Loop:** while  $\mathcal{L} \neq \emptyset$

    Suppress  $(X, \mathcal{L})$

    Increase the working precision if needed

$x := \text{mid}(X)$

$(X_1, X_2) := \left(x - \frac{F(\{x\})}{F'(X)}\right) \cap X$  //  $X_1$  and  $X_2$  can be empty

    if  $w(X_1) > \alpha w(X)$  or  $w(X_2) > \alpha w(X)$  then  $(X_1, X_2) := \text{bisect}(X)$

    if  $X_1 \neq \emptyset$  and  $F(X_1) \ni 0$  then

        if  $w(X_1)/|\text{mid}(X_1)| \leq \varepsilon_X$  and  $w(F(X_1)) \leq \varepsilon_Y$  then Insert  $X_1$  in Res

        else Insert  $X_1$  in  $\mathcal{L}$

    same handling of  $X_2$

**Output:**  $Res$ , a list of intervals that may contain the roots.

# Interval Newton algorithm stopping criterion and termination proof

(Baker Kearfott and Walster 2000)

## Stopping criterion

**Relative Root Accuracy:**  $w(X_1)/|x_1| \leq \varepsilon_X$

**Absolute Residual Accuracy:**  $w(F(X_1)) \leq \varepsilon_Y$

( $\varepsilon_X$  and  $\varepsilon_Y$  given by the user)

## Termination proof

One step of the algorithm =  $\left\{ \begin{array}{l} 1 \text{ step of Newton} \\ \text{OR} \\ 1 \text{ step of dichotomy} \end{array} \right.$

This ensures that  $w(X_{k+1}) \leq \alpha w(X_k)$ .

# Interval Newton algorithm

## automatic adaptation of the working precision

**First need:** being able to bisect the current interval  $X$   
 $\Rightarrow$  increase the working precision when  $w(X) = 1$  “ulp”.

**Second need:** being able to refine the function evaluation  $F(X)$   
 $\Rightarrow$  increase the working precision when  $w(F(X)) \leq w(F(X_1))$ .

**How to increase the working precision :** many possible choices, for us it is doubled since the number of correct digits is roughly doubled at each iteration (Newton).

# Interval Newton algorithm experiments: Chebychev polynomials

**Chebychev polynomials:**  $C_n(\cos(\theta)) = \cos(n\theta)$ .

They are difficult to evaluate with a good precision: even if their values belong to  $[-1, 1]$ , their coefficients are very large.

**Results:** very precise roots for degrees up to 40, with proof of existence and uniqueness.

## Interval Newton algorithm

experiments: Wilkinson polynomial  $\prod_{i=1}^{20}(X - i)$ .

**With enough precision to be able to store exactly the coefficients:** roots found with a precision  $5 \cdot 10^{-2}$  and with a proof of existence (but not uniqueness).

A lot of intervals are not eliminated:  $[0.96, 1.02]$  et  $[1.62, 20.984]$ .

Similar behaviour (but the computing time is twice as much) with the starting intervals  $[-10, 40]$  and  $[-100, 400]$ .

**With enough precision and a perturbation  $\pm 2^{-19}$  on the coefficient of  $X^{19}$ :** roots (with proof of existence but not uniqueness):

$1 \pm 4 \cdot 10^{-2}$ ,  $2 \pm 5 \cdot 10^{-2}$ ,  $3 \pm 4 \cdot 10^{-2}$ ,  $4 \pm 4 \cdot 10^{-2}$ ,  $5 \pm 4 \cdot 10^{-2}$ ,  $6 \pm 5 \cdot 10^{-2}$ ,  $7 \pm 6 \cdot 10^{-2}$  and  $[7.91, 22.11]$ .

A lot of intervals are not eliminated:  $[0.96, 22.64]$ .

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- **Solving nonlinear equations**

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- some properties
- experimental results

- **Solving linear systems**

- Hansen & Sengupta's algorithm
- adaptation to multiple precision

# Solving linear systems

## Hansen & Sengupta's algorithm

**Linear system:**  $Ax = b$  with  $A$  and  $b$  given.

**Problem:** compute an enclosure of

$\text{Hull}(\Sigma_{\exists\exists}(A, b)) = \text{Hull}(\{x : \exists A \in A, \exists b \in b, Ax = b\})$ .

### Hansen & Sengupta's algorithm

compute  $C$  an approximation of  $\text{mid}(A)^{-1}$

apply Gauss-Seidel to  $CAx = Cb$  until convergence.

### Idea:

$CA$  contains the identity matrix, the iteration matrix has a spectral radius close to 1 (and even  $< 1$ ?) and this iteration is a contraction.

# Hansen & Sengupta's algorithm

## automatic adaptation of the working precision

**First need:** being able to improve the current interval  $x$   
 $\Rightarrow$  increase the working precision when  $w(x) = 1$  “ulp”.

**Second need:** being able to refine the Gauss-Seidel iteration evaluation.

Notation:  $[x_1, x_2] = \text{bisect}(x)$

one iteration applied to  $x$  (resp. to  $x_1$ , resp. to  $x_2$ ) gives  $y$  (resp.  $y_1$ ,  
(resp.  $y_2$ ))

increase the working precision if  $w(y_1) \geq w(y)$  or  $w(y_2) \geq w(y)$ .

**How to increase the working precision :** many possible choices, for us  
it is doubled.

# Conclusion and future work

## Conclusion

- **Reliable and accurate computing is possible**, thanks to arbitrary precision interval arithmetic.
- **Algorithms need to be adapted:** automatic adaptation of the precision, more stringent stopping criteria.

## Future work

- **Testing and improving the library:** MPFI and MPFI++ (interface à la Profil/BIAS).
- **Development of numerical algorithms** with MP interval arithmetic: multivariate Newton  $\Rightarrow$  global optim.  $\Rightarrow$  constrained global optim.
- **Applications: automatics and robotics**  
parameters estimation and robust control.