

# HW1 Molecular Programming

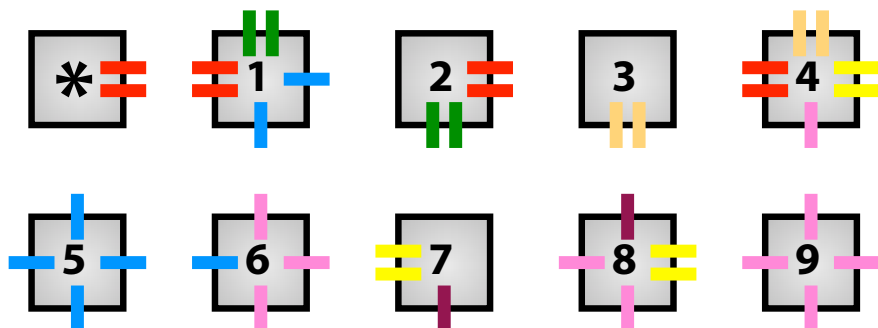
MPRI 2.11.1 9.10.2019 - Due on Wed. 16/10 before 12:45



You are asked to complete the exercise marked with a [★] and to send me your solutions to:  
[nicolas.schabanel@ens-lyon.fr](mailto:nicolas.schabanel@ens-lyon.fr)  
 as a PDF file named **HW1-Lastname.pdf** on **Wed. 16/10 before 12:45**.

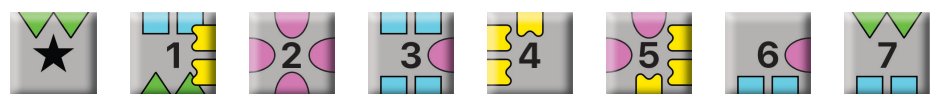
■ **Exercise 1 (Algorithmic Self-Assembly).** Recall that the self-assembly process consists in, given a finite tileset (with infinitely many tiles of each type), starting from the seed tile (marked with a ★), glueing tiles with matching colors to the current aggregate so that each new tile is attached by at least *two* links to the aggregate (either on the same border or on two borders). Recall that a shape is *final* if no tiles can be attached to it anymore.

► **Question 1.1)** *What is the exact family of final shapes self-assembled by the following tileset? (No proof nor justification is asked.) Indicate the local order of assembly by drawing arrows over the tiles of a generic final shape. Which are the two competing tiles that decide the size of the resulting final shape?*



■ **Exercise 2 (Guess the shape).** Recall that the self-assembly process consists in, given a finite tileset (with infinitely many tiles of each type), starting from the seed tile (marked with a ★), glueing tiles with matching colors to the current aggregate so that each new tile is attached by at least  $T^\circ = 2$  links to the aggregate (either on the same border or on two borders). Recall that a shape is *final* if no tiles can be attached to it anymore.

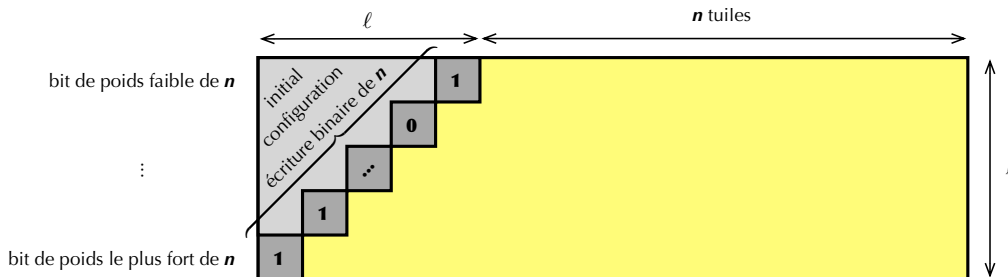
► **Question 2.1)** *What is the exact family of final shapes self-assembled by the following tileset at temperature  $T^\circ = 2$ ? (No proof nor justification is asked.)*



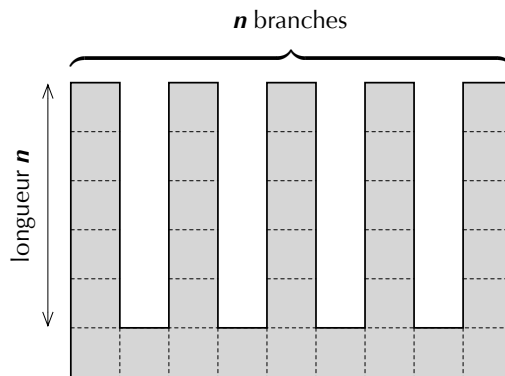
*Indicate the local order of assembly by drawing arrows over the tiles of a generic final shape. Which are the competing tiles that decide the size of the resulting final shape?*

■ **Exercise 3 (Counter at  $T^\circ = 2$  (★)).** Given an integer  $n$ , and an seed configuration consisting of an isosceles rectangle triangle isocèle of side  $\ell = \lceil \log_2 n \rceil$  where the bits of  $n$  are encoded on the diagonal as shown in grey below.

Propose a well-ordered (finite) tileset which assembles the yellow at  $T^\circ = 2$  to realise a rectangle of size  $\ell \times (n + \ell)$ . Carefully indicate the position of the glue of strength 1 and 2 on the diagonal of the seed configuration. Indicate the assembly order<sup>(1)</sup>. What does the tiles encode?



■ **Exercise 4.** Propose a staged assembly scheme at temperature  $T^\circ = 1$  of the shape family  $E$  of candelabums with  $n$  branches of length  $n$ .



Describe the tiles, glues, their number, the number of stages and the number of different bechers needed. Give an illustration of the stages to build a generic production.

■ **Exercise 5.** Assume a random Poisson model where the random time  $X$  between two consecutive appearances of a tile of a given type  $\tau$  at a given empty location follows an exponential law:  $p(x) = c \cdot e^{-cx}$  where  $c > 1$  is the concentration of the tiles of type  $\tau$ . We want to prove the following theorem:

**Theorem 1 (Adleman et al, 2001).** Consider an ordered tile system  $\mathcal{T}$  that assembles deterministically a single shape  $S$ . Let  $\prec$  be the partial order of the assembly, i.e. such that  $(i, j) \prec (k, l)$  if the tile at position  $(i, j)$  is attached before the tile at  $(k, l)$  by  $\mathcal{T}$ . With very high probably, the assembly time of a shape  $S$  by  $\mathcal{T}$  is:

$$O(\gamma \times \text{rank}(S))$$

where  $\gamma$  only depends on the concentrations and  $\text{rank}(S)$  is the highest rank in the shape  $S$  (i.e. the length of the longest path in  $\prec$ ).

► **Question 5.1)** Let  $X$  be an exponential random variable such that  $p(X = x) = ce^{-cx}$  for all real  $x \geq 0$ , for some  $c > 0$ . Show that  $X$  is memoryless, i.e. for all  $u, t \geq 0$ ,

$$p(X = t + u | X \geq u) = p(X = t)$$

Let  $T$  be the assembly time of the shape  $S$ , i.e. the time at which the last tile of shape  $S$  is attached. We denote by  $w(P)$  the random variable for the weight of a  $\prec$ -path  $P$ , defined as:  $w(P) = \sum_{(i,j) \in P} X_{i,j}$ .

► **Question 5.2)** Let  $X_{i,j}$  be the independent exponential random variable for the time between two consecutive appearances of the tile to be attached at position  $(i, j)$  in  $S$ . Show that:

$$T = \max_{\prec\text{-path } P} w(P)$$

▷ Hint. Proceed by recurrence on the rank of the tiles and show that for all tile  $(i, j)$ , its assembly time is the random variable  $T_{ij} = \max_{\prec\text{-path } P \text{ from } (0,0) \text{ to } (i,j)} w(P)$ .

► **Question 5.3)** Let  $X_1, \dots, X_\ell$  be  $\ell$  independent exponential variables s.t.  $p(X_i = x) = c_i e^{-c_i x}$  with  $c_i > 1$ . Show that there is  $\gamma$  which depends only of  $\min_i c_i$  such that: for all  $n \geq \ell$ ,

$$\Pr\{X_1 + \dots + X_\ell \geq \gamma \cdot n\} \leq 1/4^\ell \cdot e^{-\gamma(n-\ell)}$$

▷ Hint. Note that  $\mathbb{E}[e^{X_i}] < \infty$  and apply Markov inequality to  $Z = e^{X_1 + \dots + X_\ell}$ .

► **Question 5.4)** Conclude.