

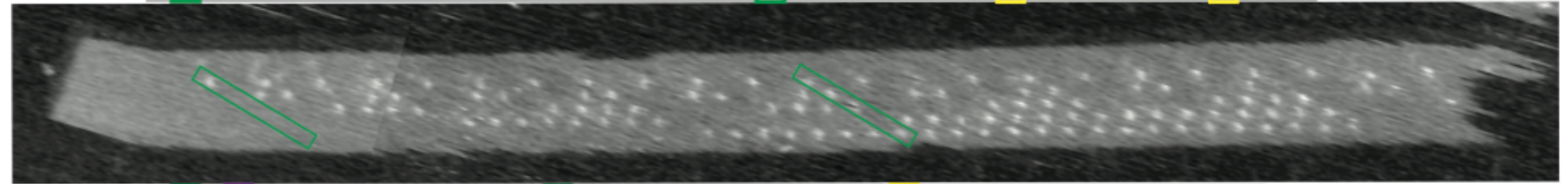
Molecules that compute and assemble nano-objects

Nicolas Schabanel

Directeur de Recherches CNRS

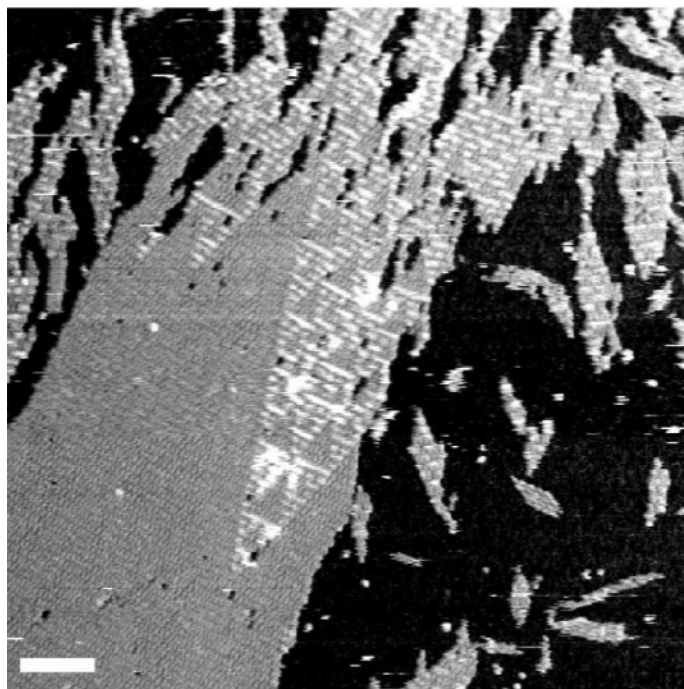
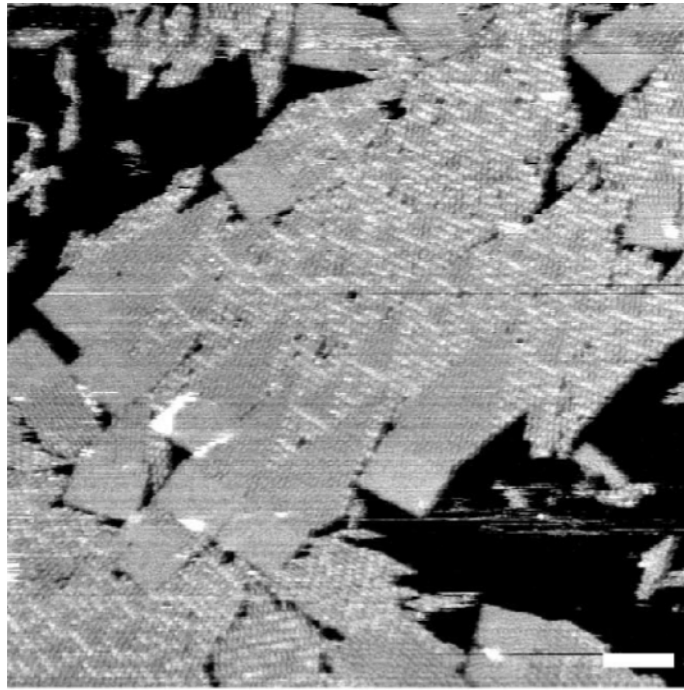
LIP & IXXI - ÉNS de Lyon

— ~100 nm

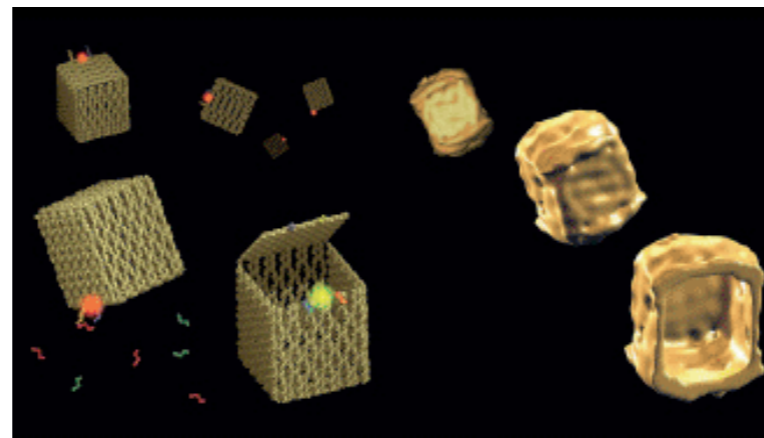


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0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0		
0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1	0	0	0	0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	
0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	0	0	0	0	0	0	1	1	1	1	1	0	0	0	0	0	0	0	1	1	1	0	0	0	0	1	1	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1

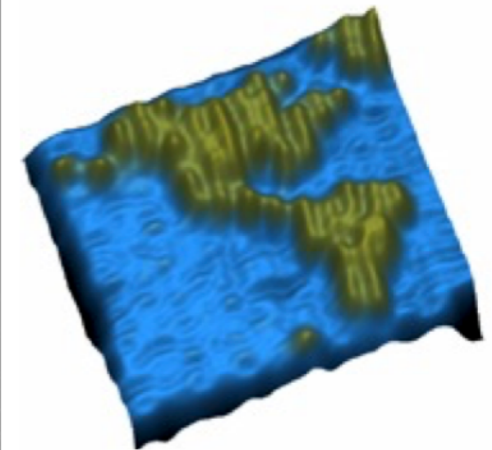
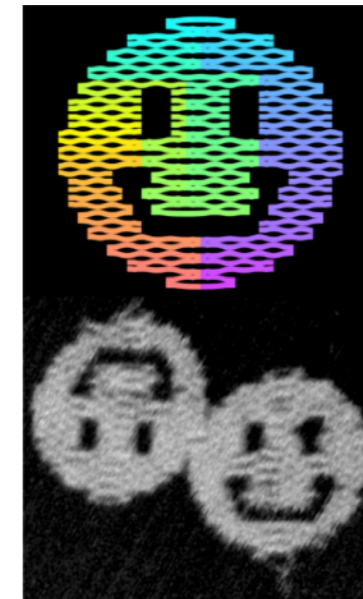
Constantine Evans, PhD Thesis, Caltech 2014



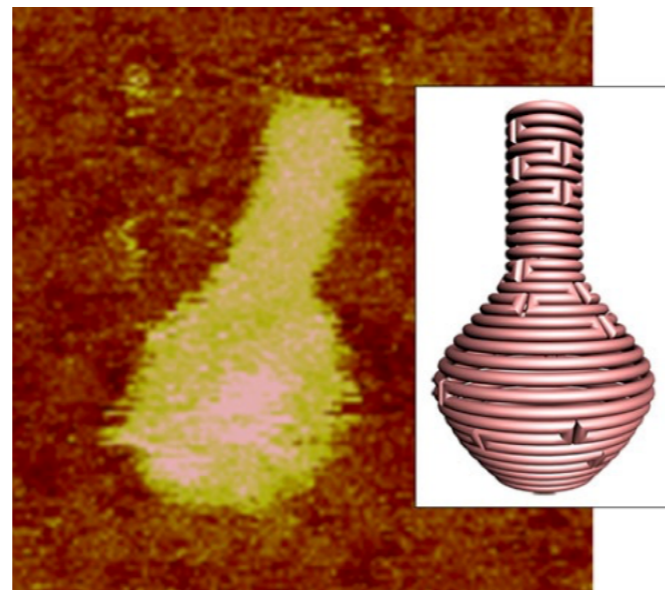
Fujibayashi et al, 2007



Andersen et al, 2009



Rothmund, Nature 2006

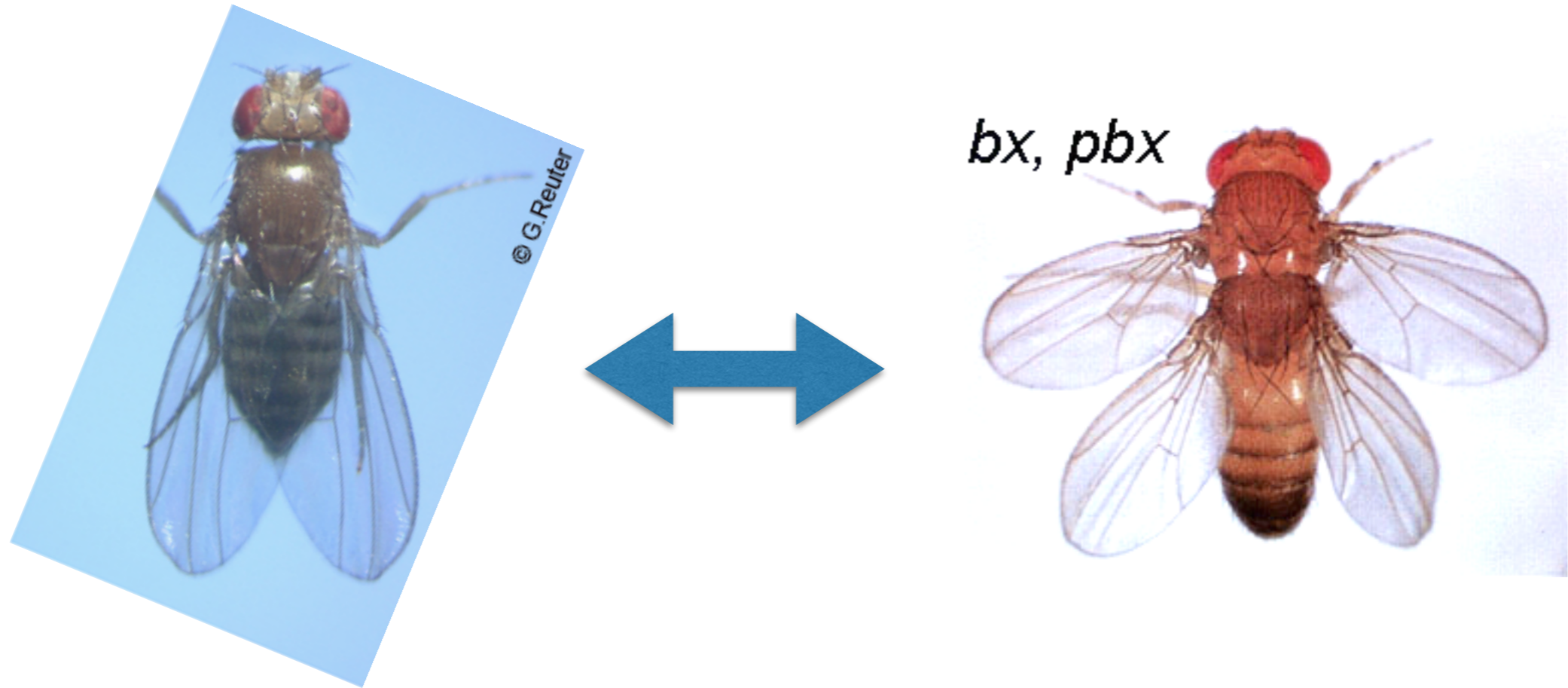


Han et al, Science 2011



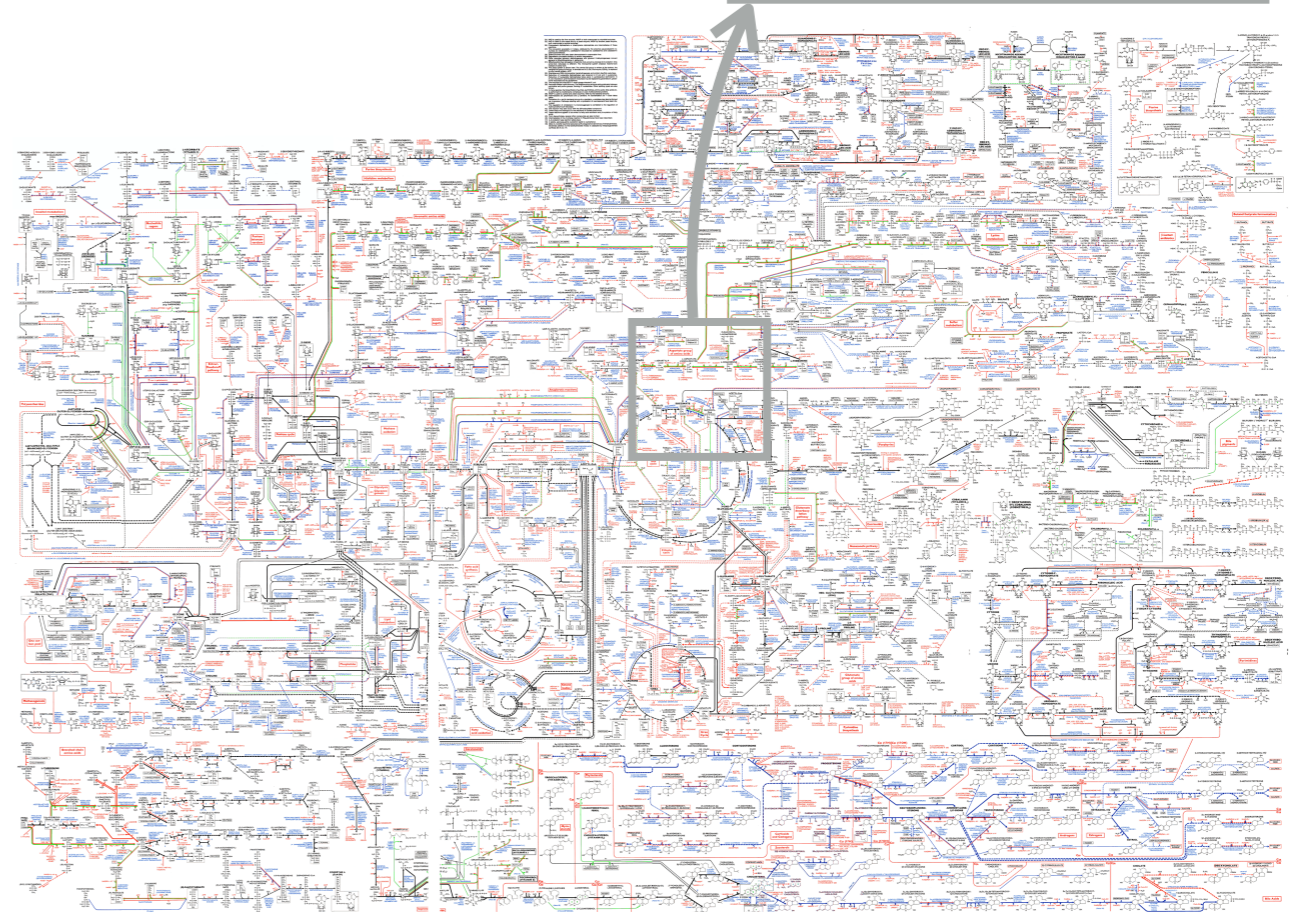
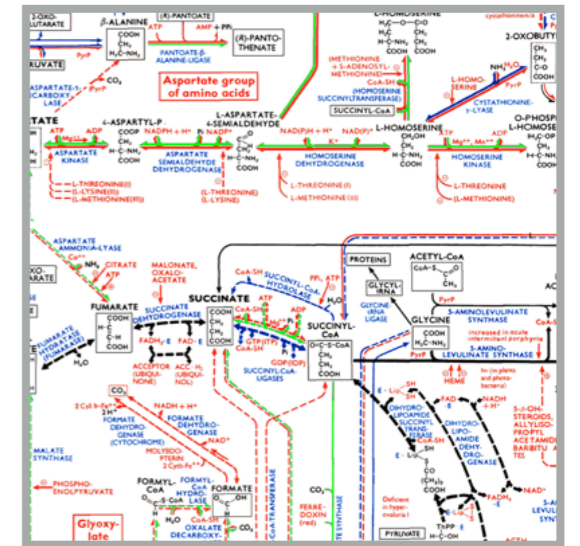
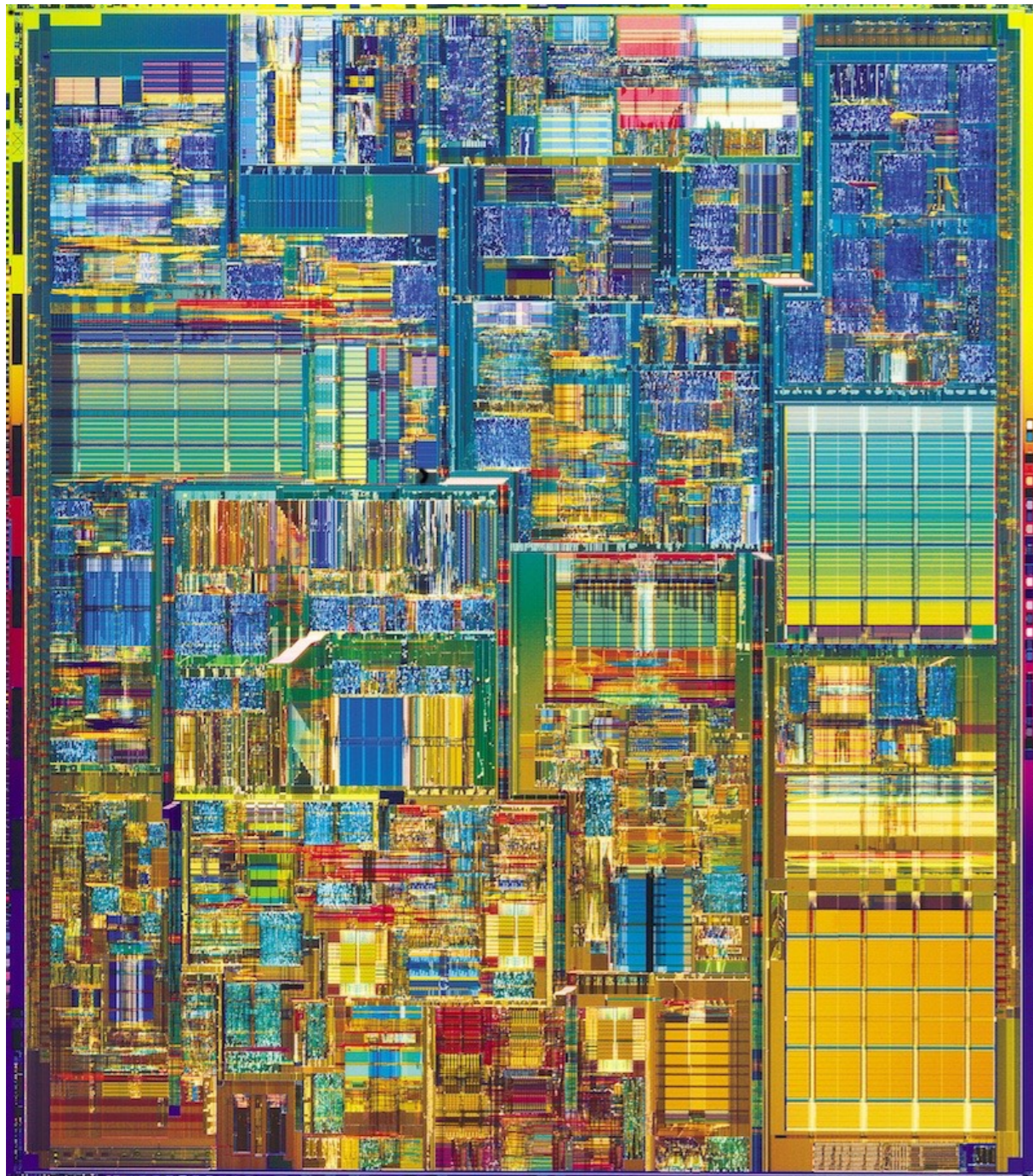
Wei, Dai, Yin, Nature 2013

Genetic code behaves as a program

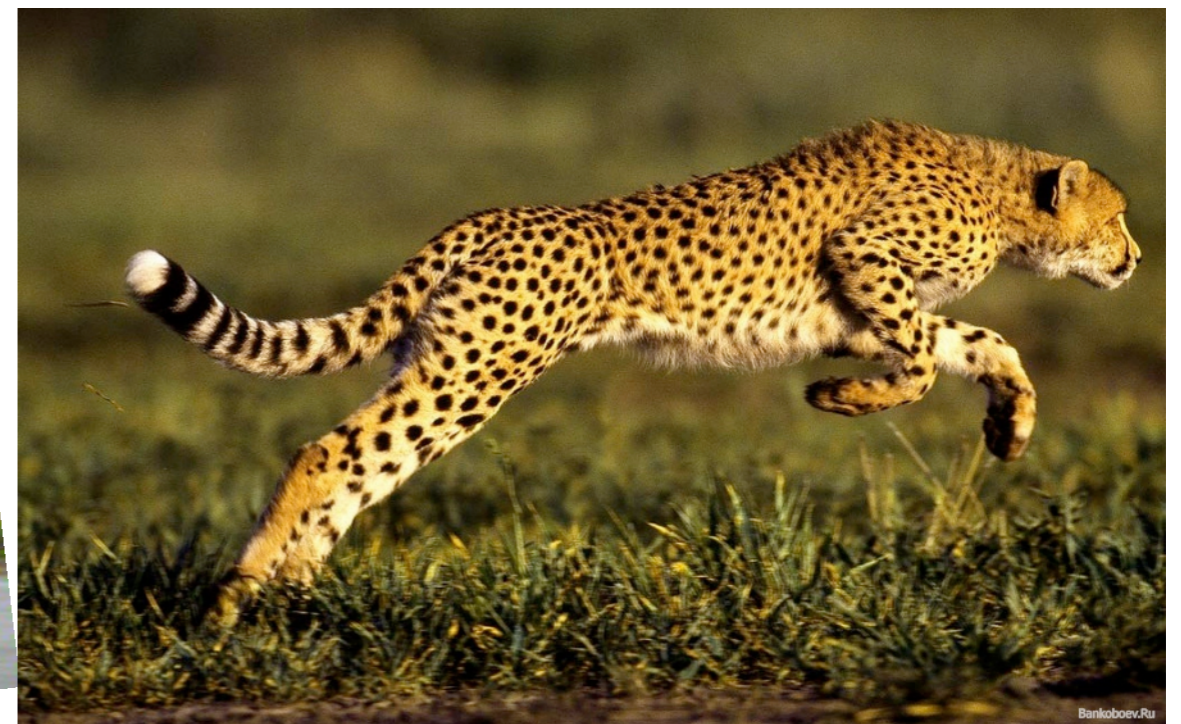


For instance, small changes in the code
imply big differences

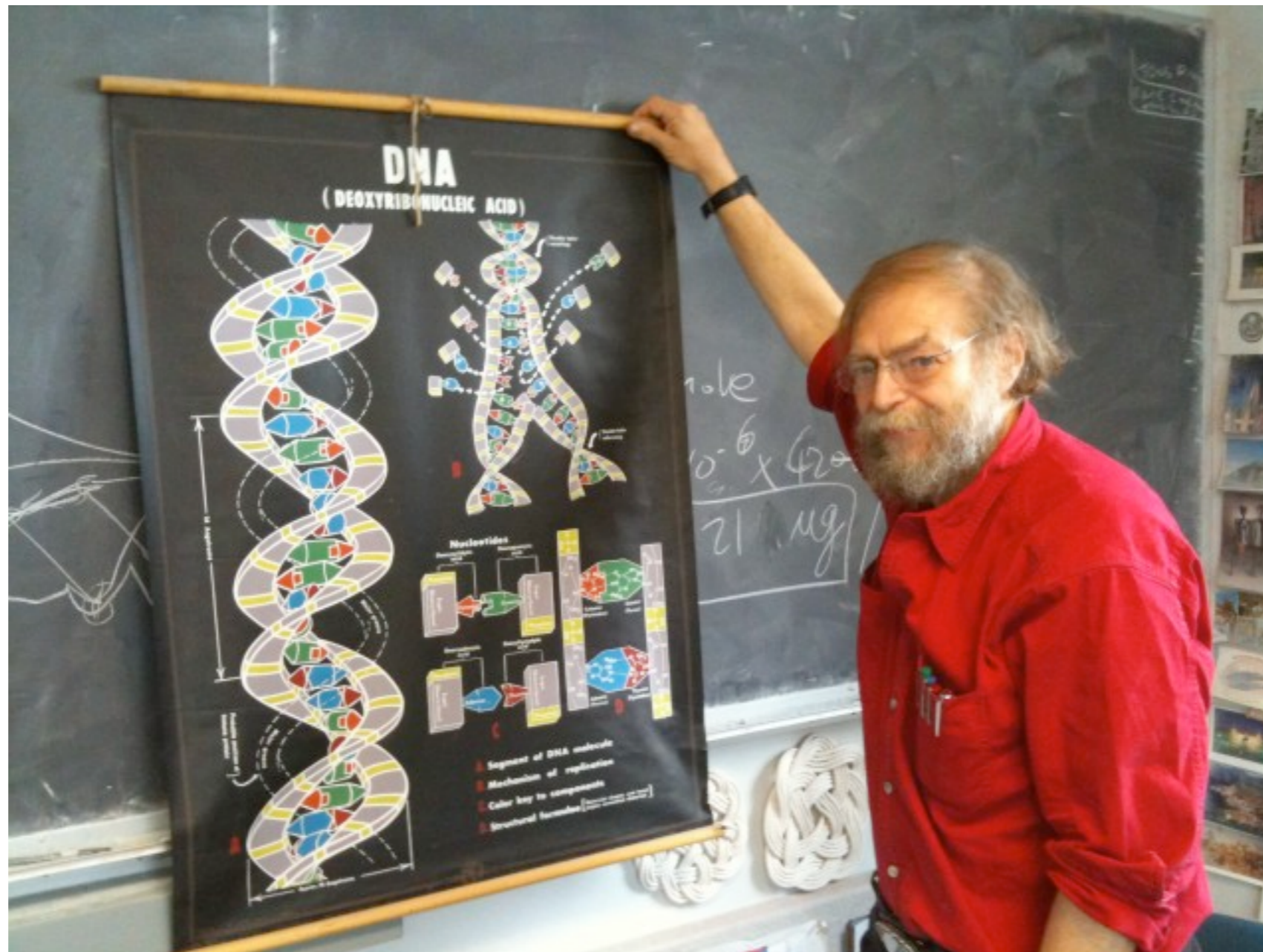
Nature is very complicated



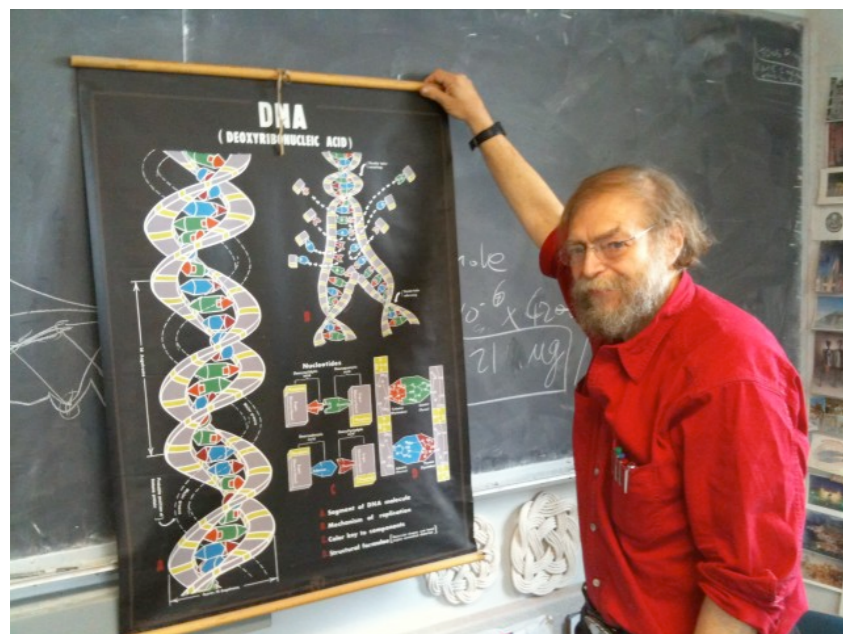
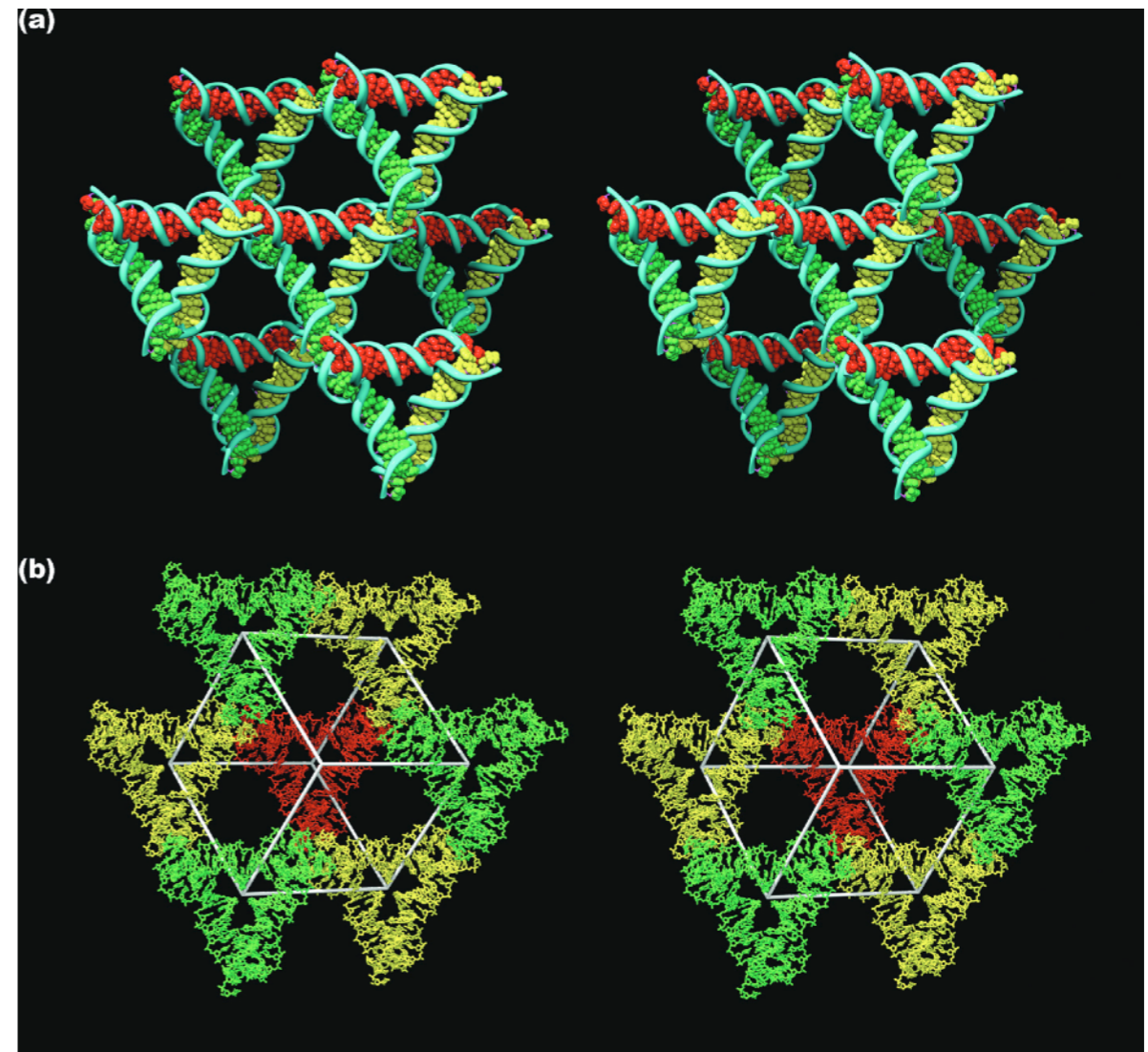
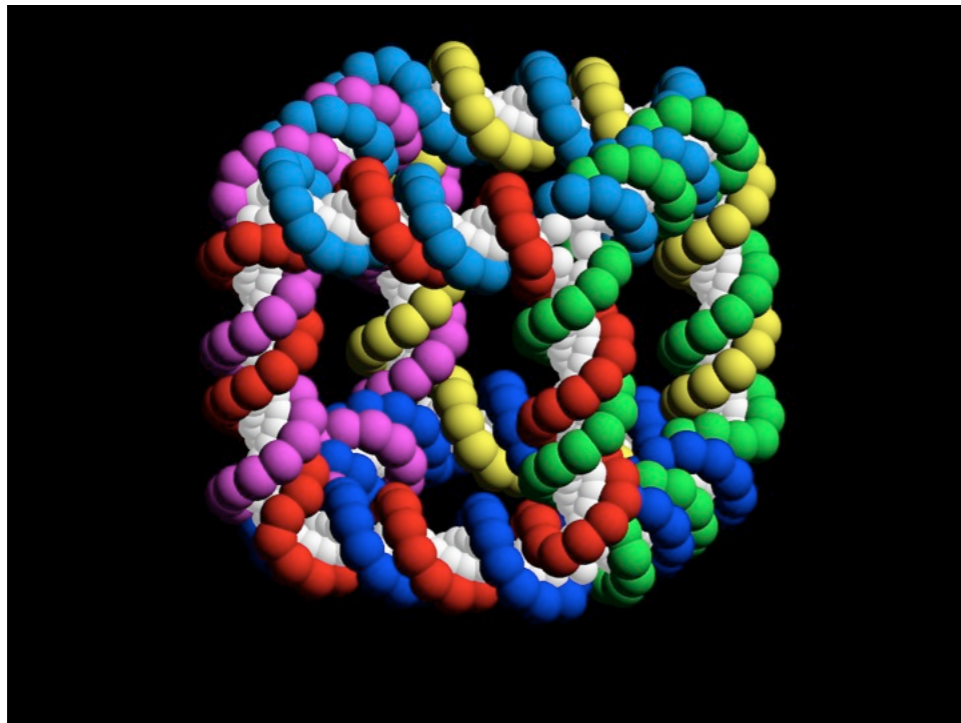
However we can try
doing differently



Using DNA to create shapes, Ned Seeman (1990-)



Using DNA to create shapes, Ned Seeman (1990-)

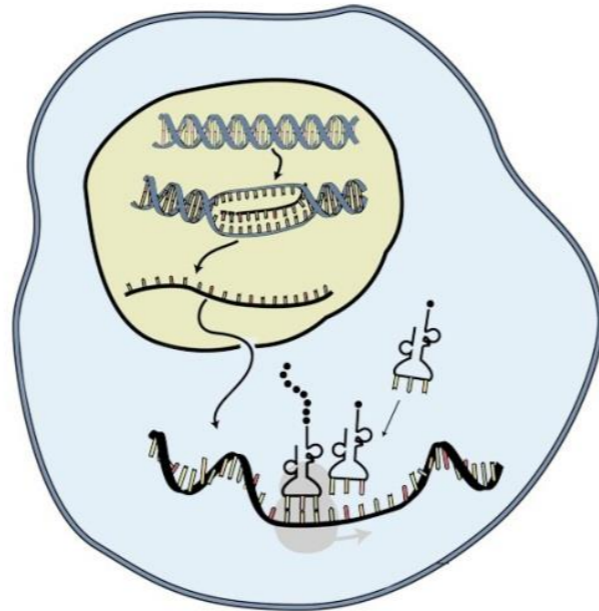


Create complementary strands
inducing particular shapes

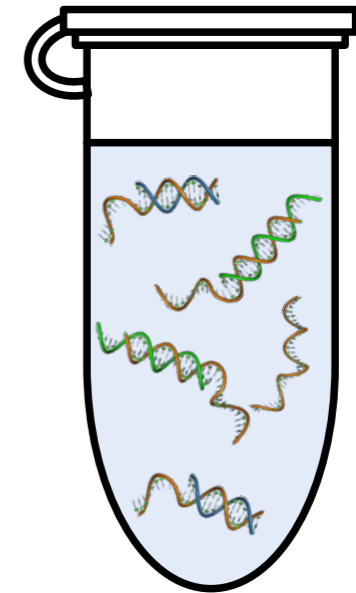
A tour of achievements

Why do we want to use nucleic acids to build structures, motors and circuits?

natural
biological
interface



easy
chemical
synthesis



combinatorial
design space

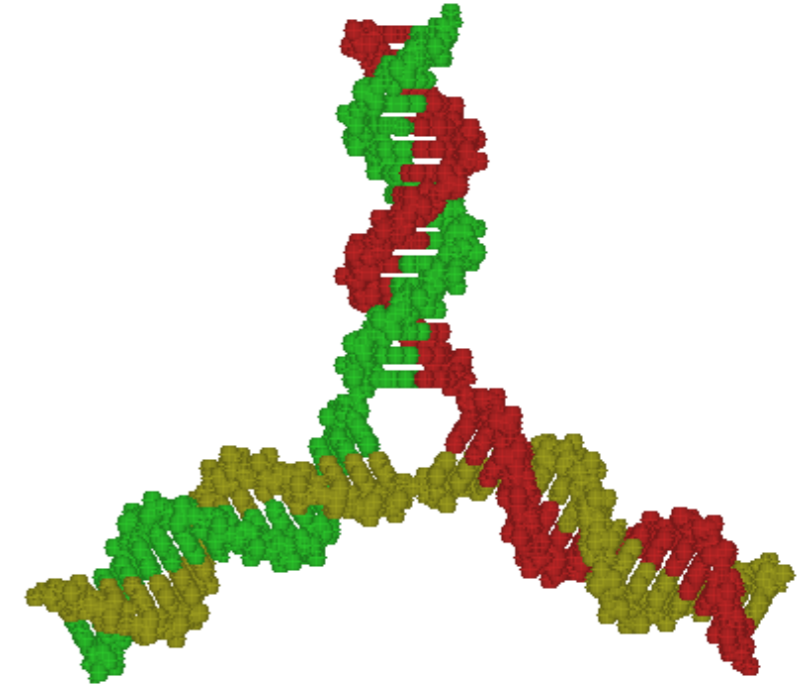


S1 = ATCGAATTCCGTAGGCC
S2 = CCCGATCGTTACGTCAT
S3 = GGCATTTTGTGGAACCA
S4 = TTAGAATCCACAGTTAG

predictable
behavior



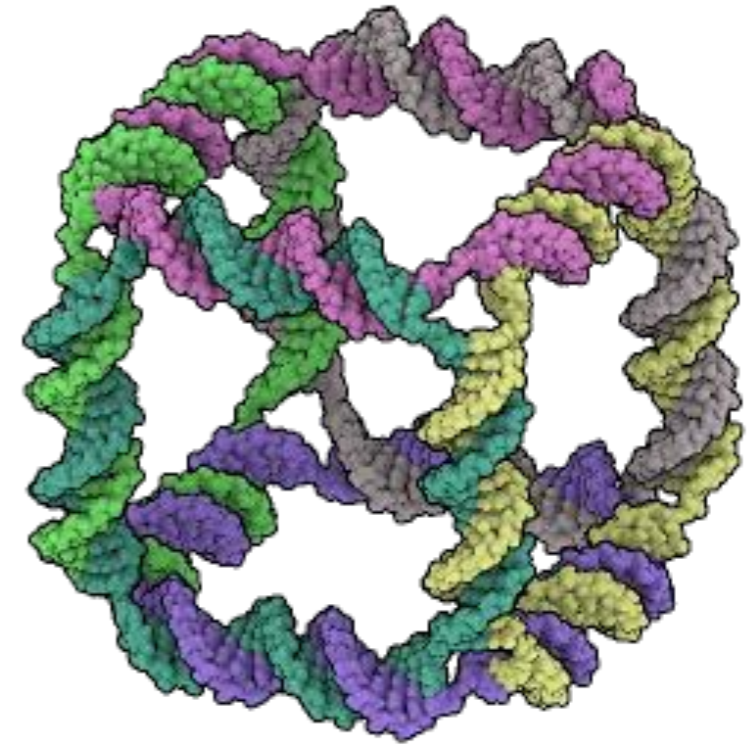
First structures



Seeman, *J. Theor. Biol.* 1982

self-assembly of
nanostructures

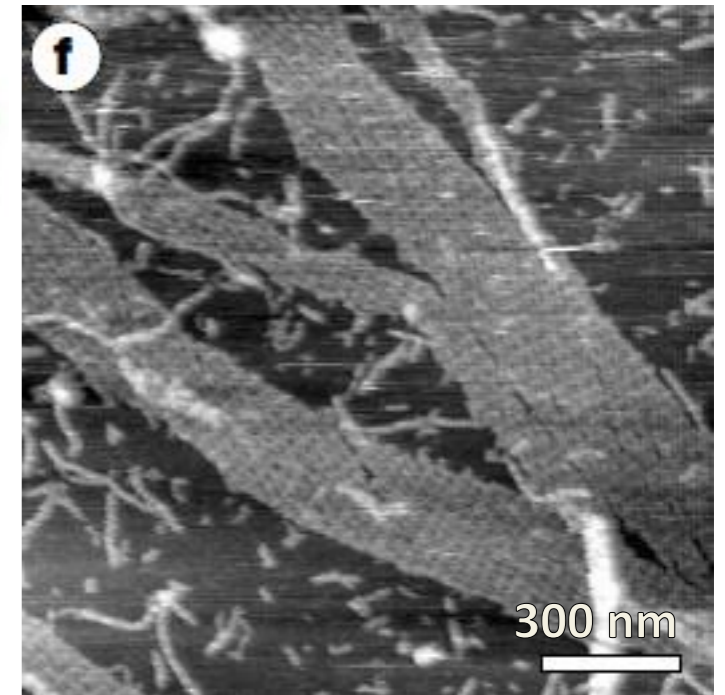
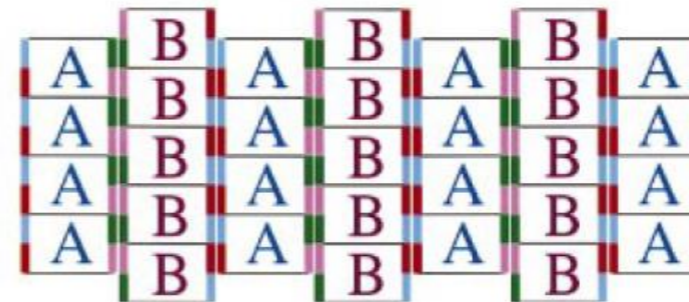
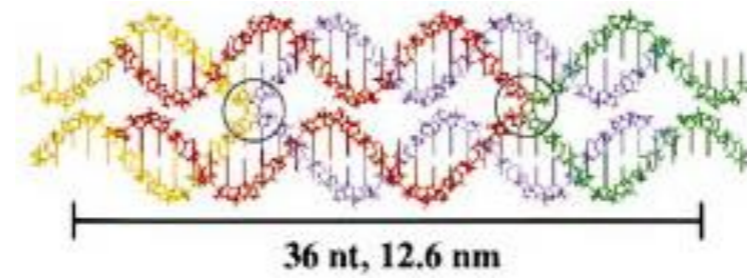
First structures



Chen & Seeman, *Nature* 1991

self-assembly of
nanostructures

Meshes



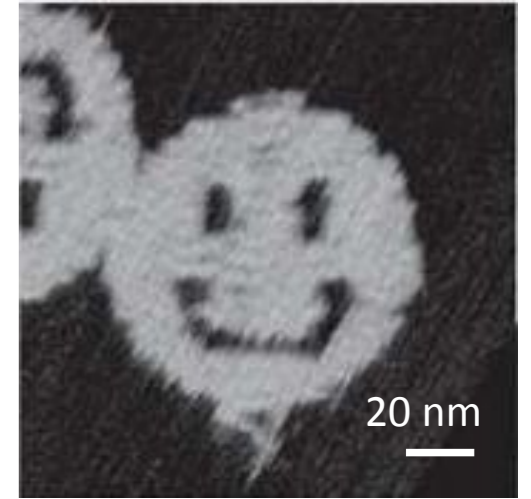
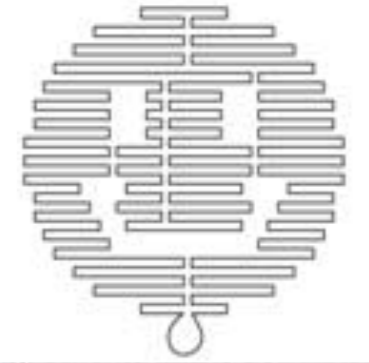
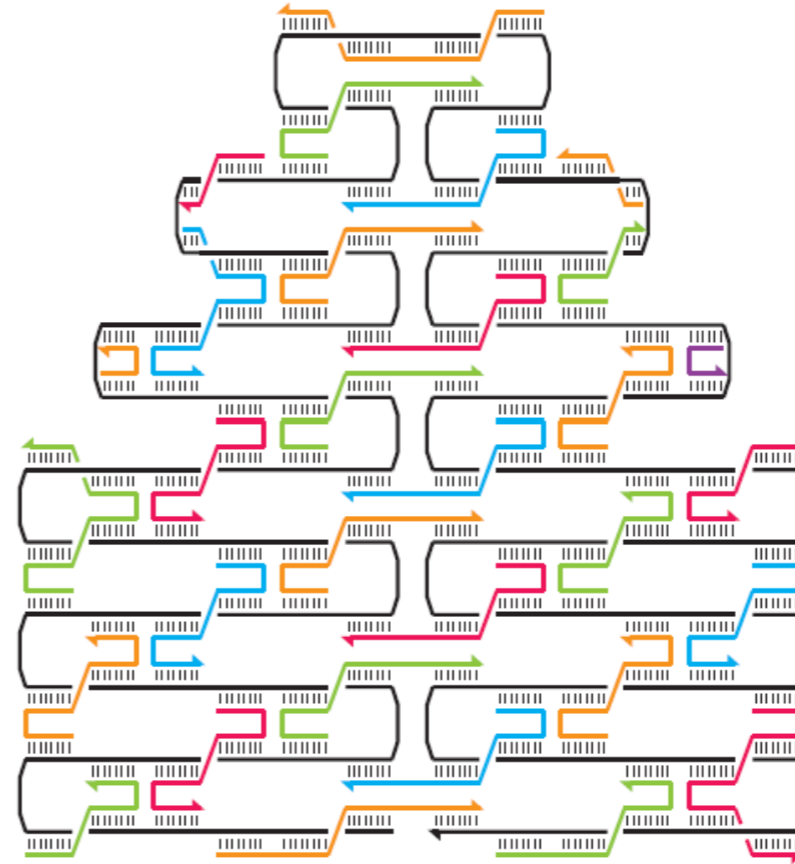
Winfrey et al, *Nature* 1998

self-assembly of
nanostructures

Larger elementary structures

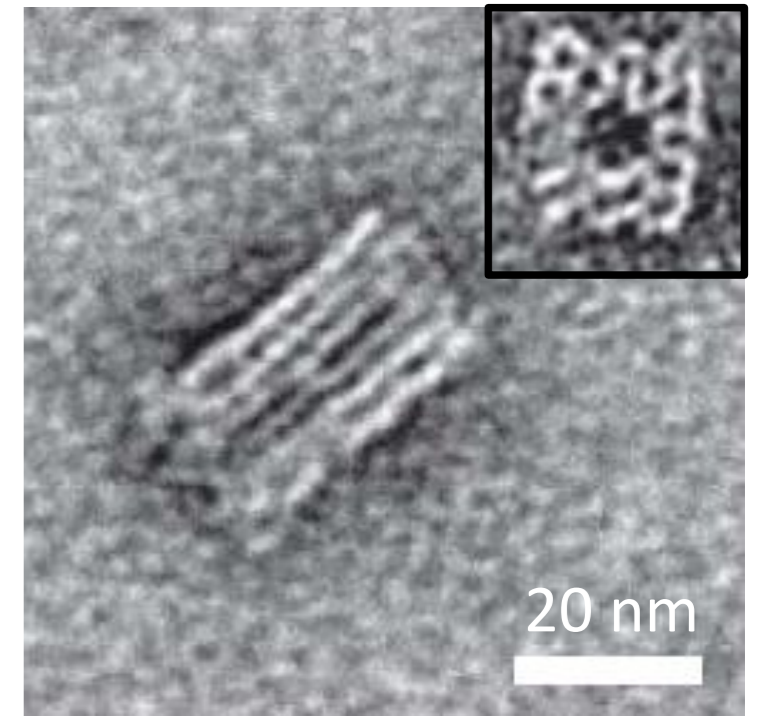
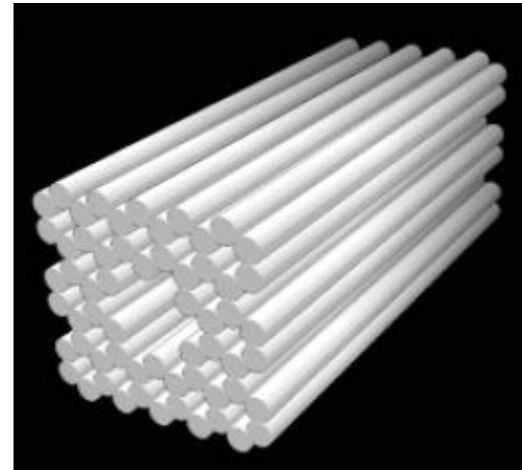


self-assembly of nanostructures



Rothemund, *Nature* 2006

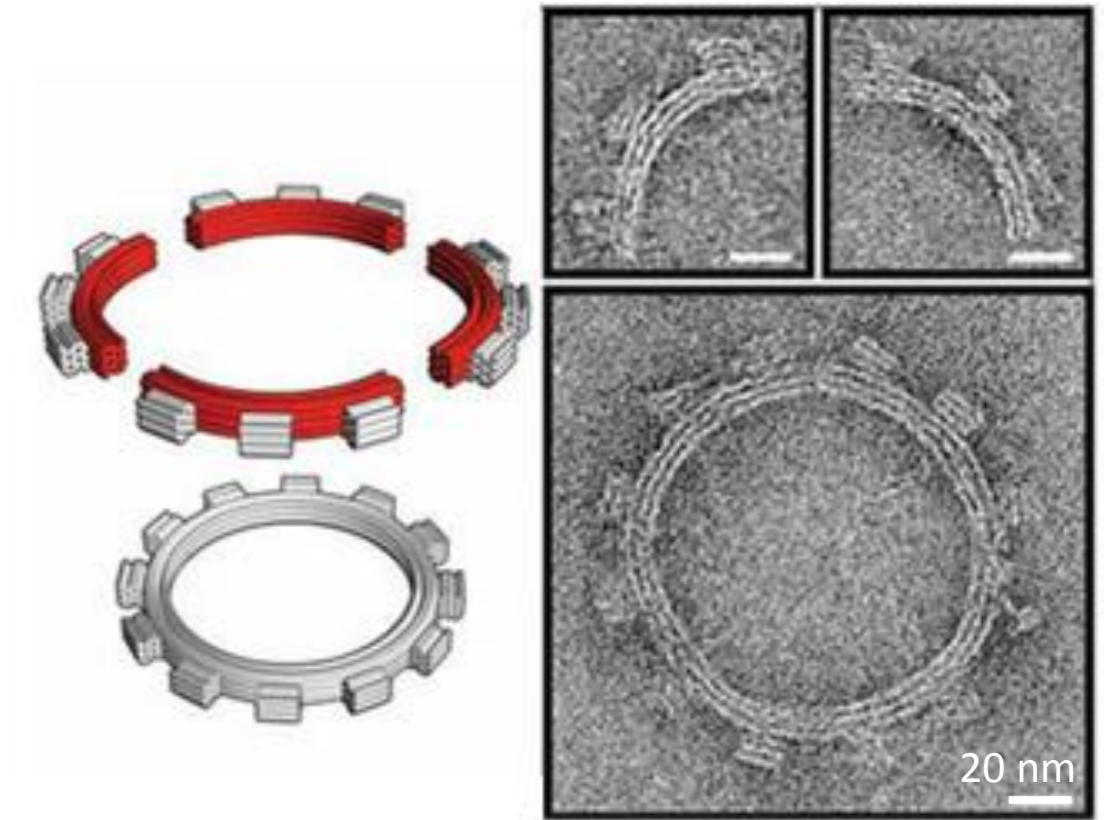
Going 3D



Douglas et al, *Nature* 2009

self-assembly of
nanostructures

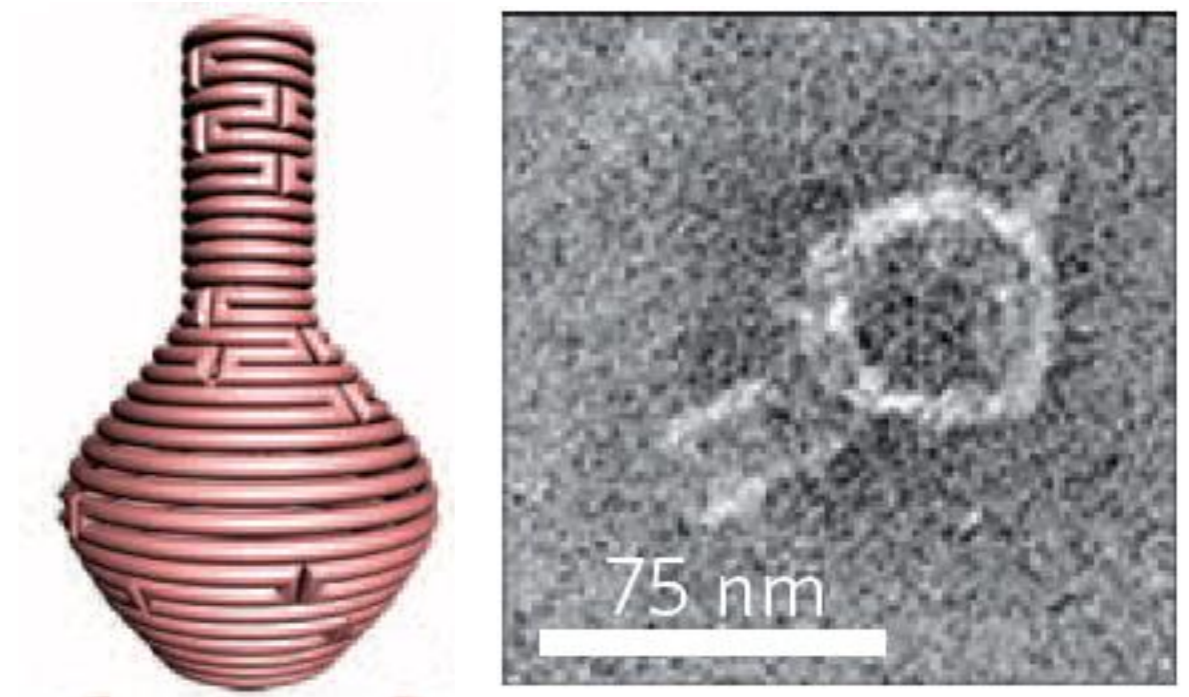
Mastering angles



Dietz et al, *Science* 2009

self-assembly of
nanostructures

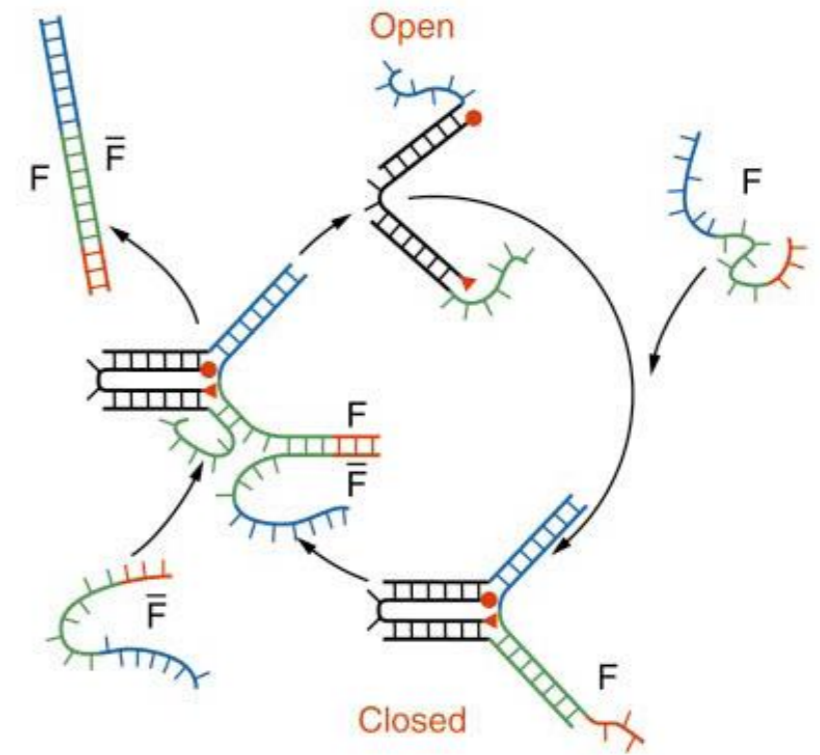
Complex 3D structures



Han et al, *Science* 2011

self-assembly of
nanostructures

Building logic

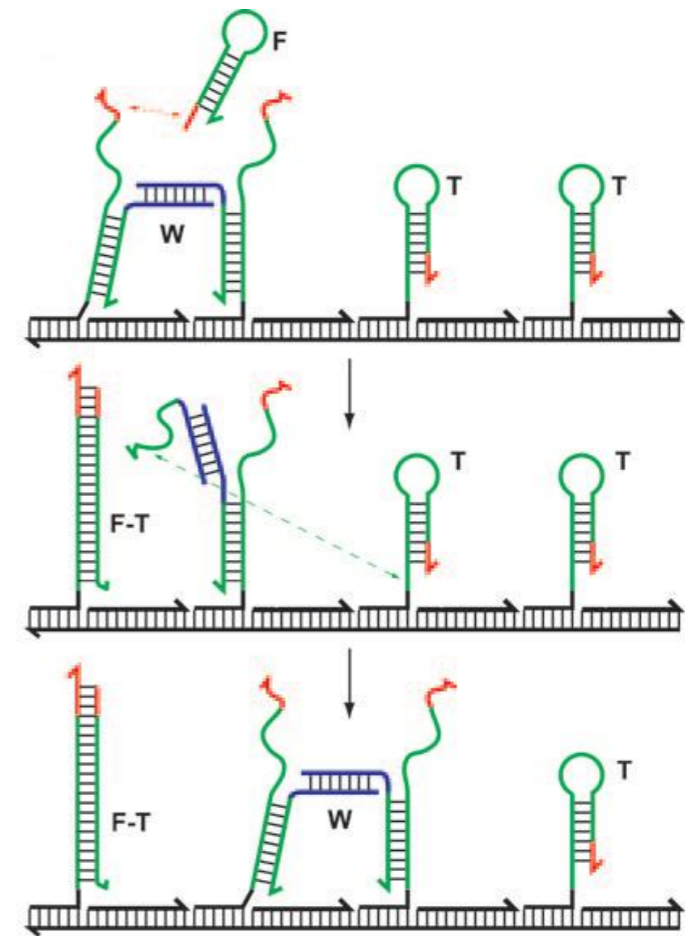


Yurke et al, *Nature* 2000

nanomechanical
devices

self-assembly of
nanostructures

Active devices

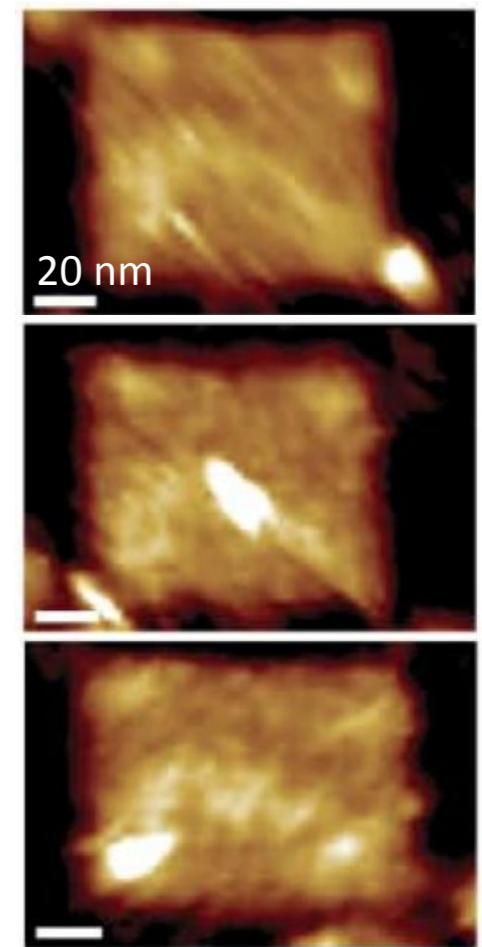
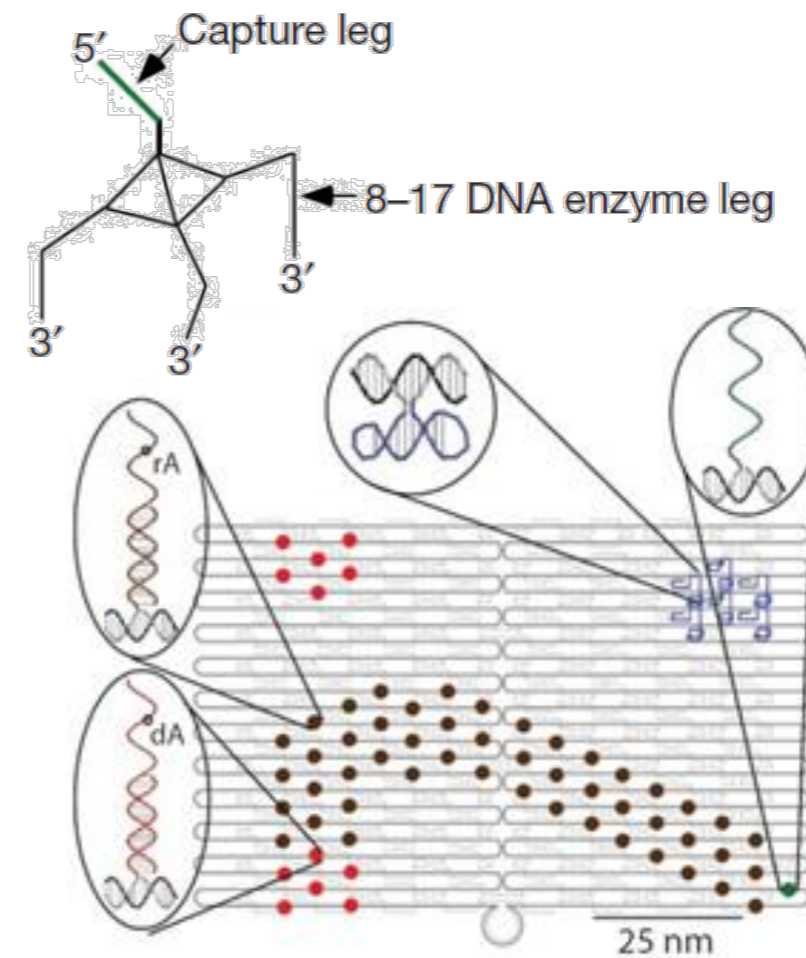


Yin et al, *Nature* 2008

nanomechanical
devices

self-assembly of
nanostructures

Active devices

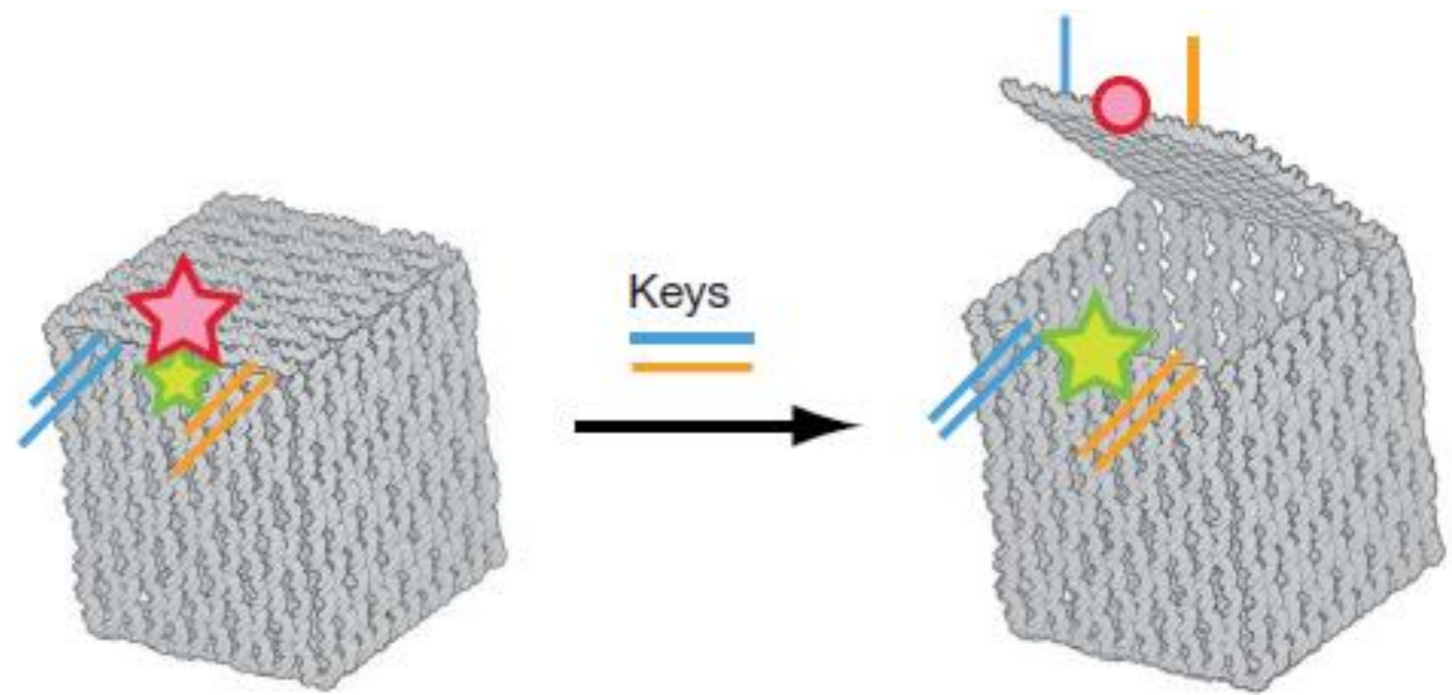


Lund et al, *Nature* 2010

nanomechanical devices

self-assembly of nanostructures

Reactive devices

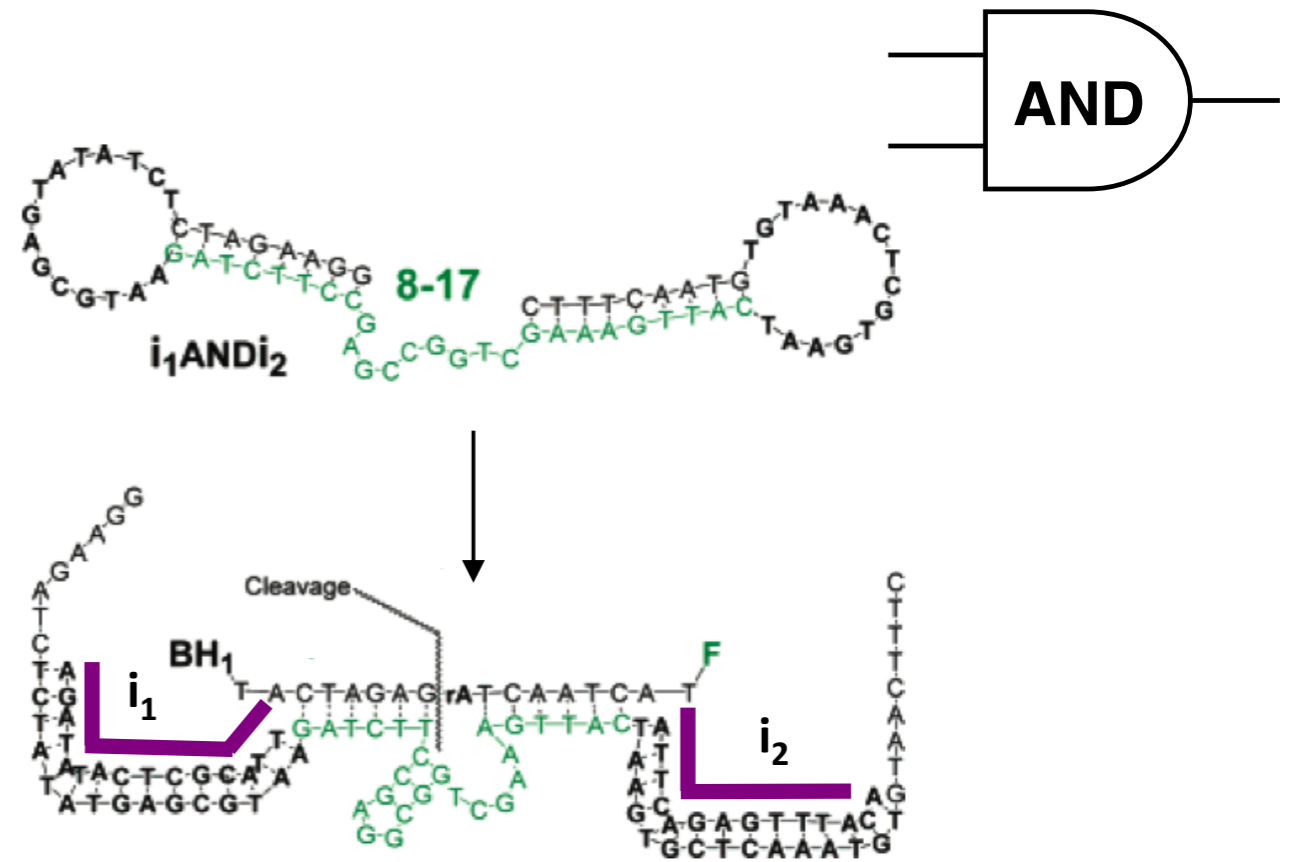


Andersen et al, *Nature* 2009

nanomechanical
devices

self-assembly of
nanostructures

Logical gates



Stojanovic & Stefanovic, *JACS* 2002

biochemical
circuits

nanomechanical
devices

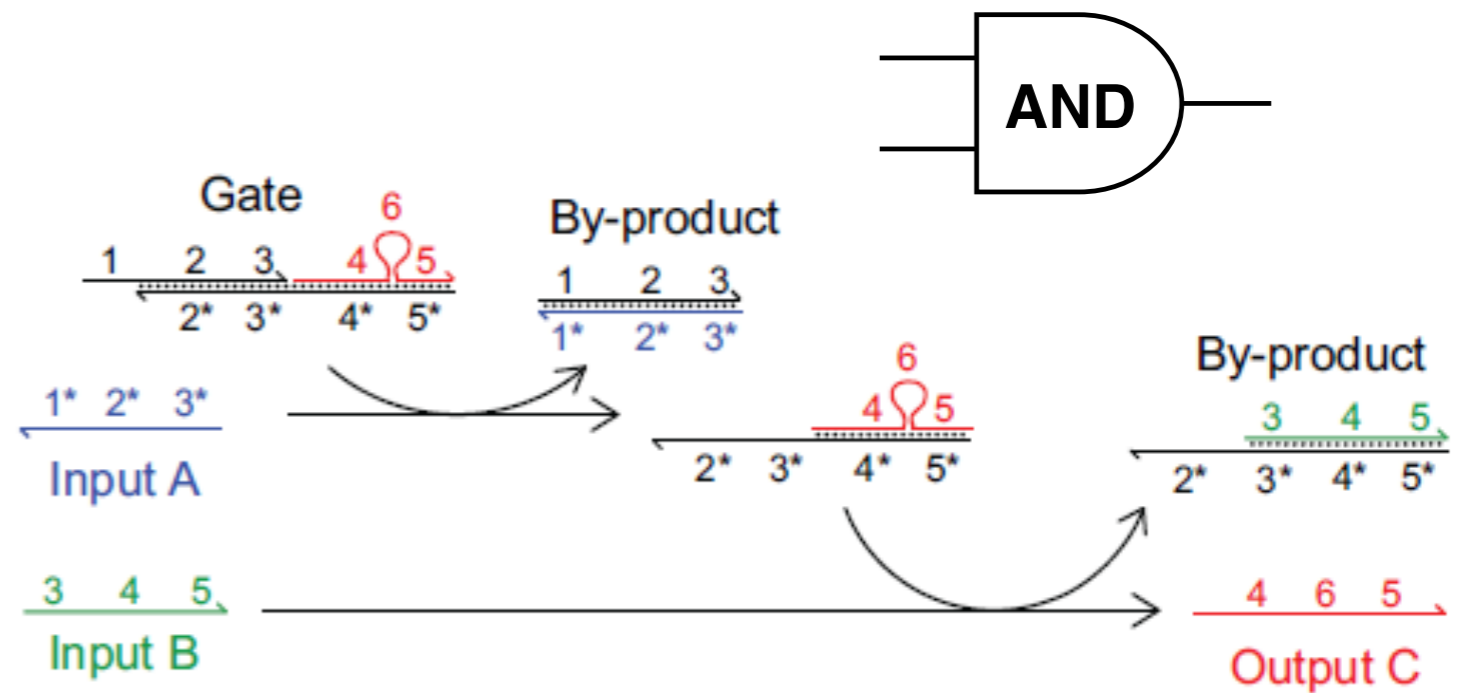
self-assembly of
nanostructures

Logical gates

biochemical
circuits

nanomechanical
devices

self-assembly of
nanostructures



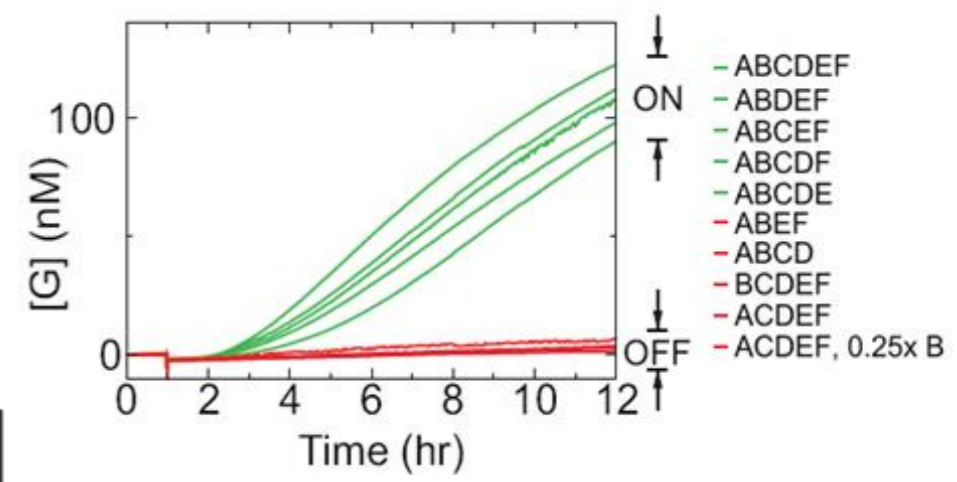
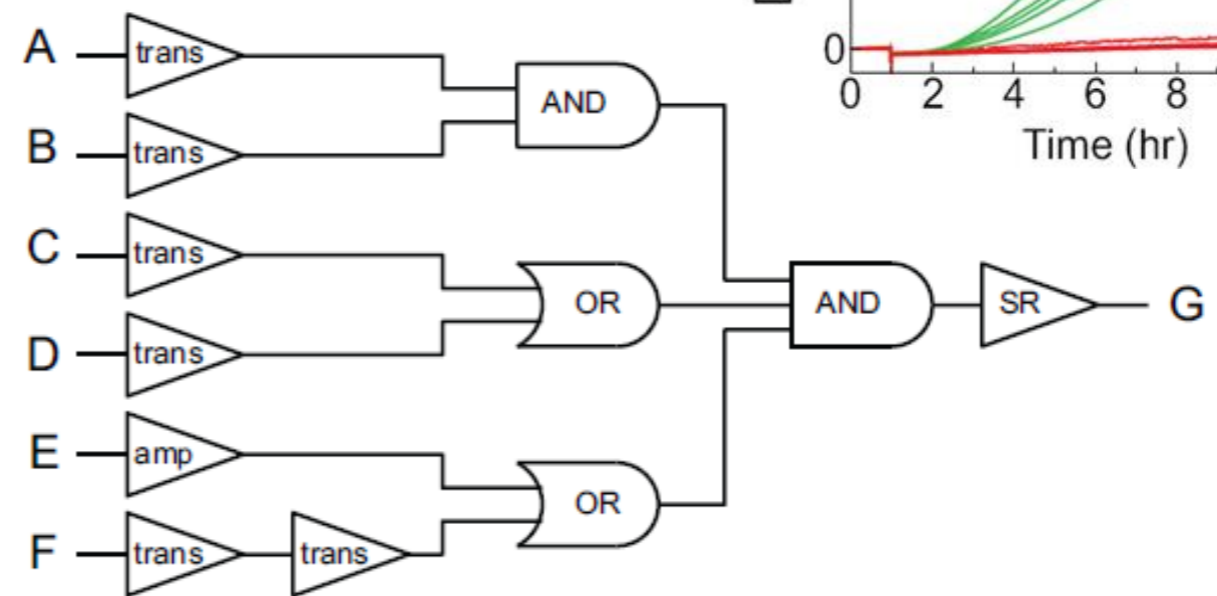
Seelig et al, *Science* 2006

Boolean circuits

biochemical
circuits

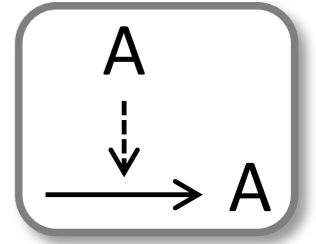
nanomechanical
devices

self-assembly of
nanostructures



Seelig et al, *Science* 2006

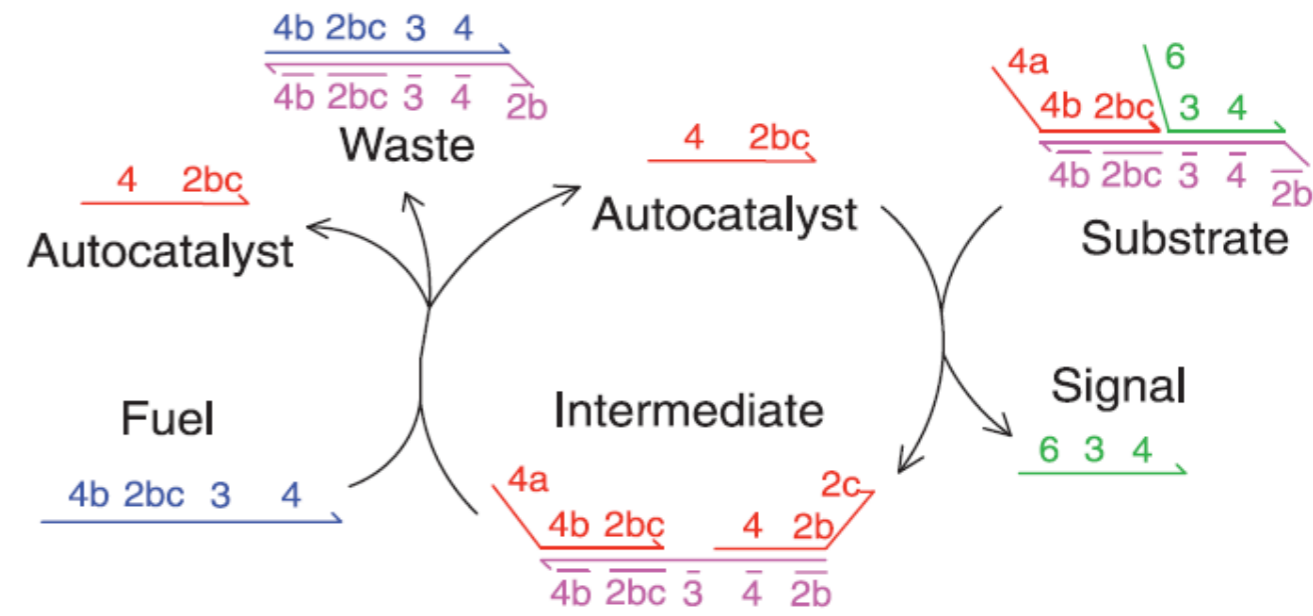
Boolean circuits



biochemical
circuits

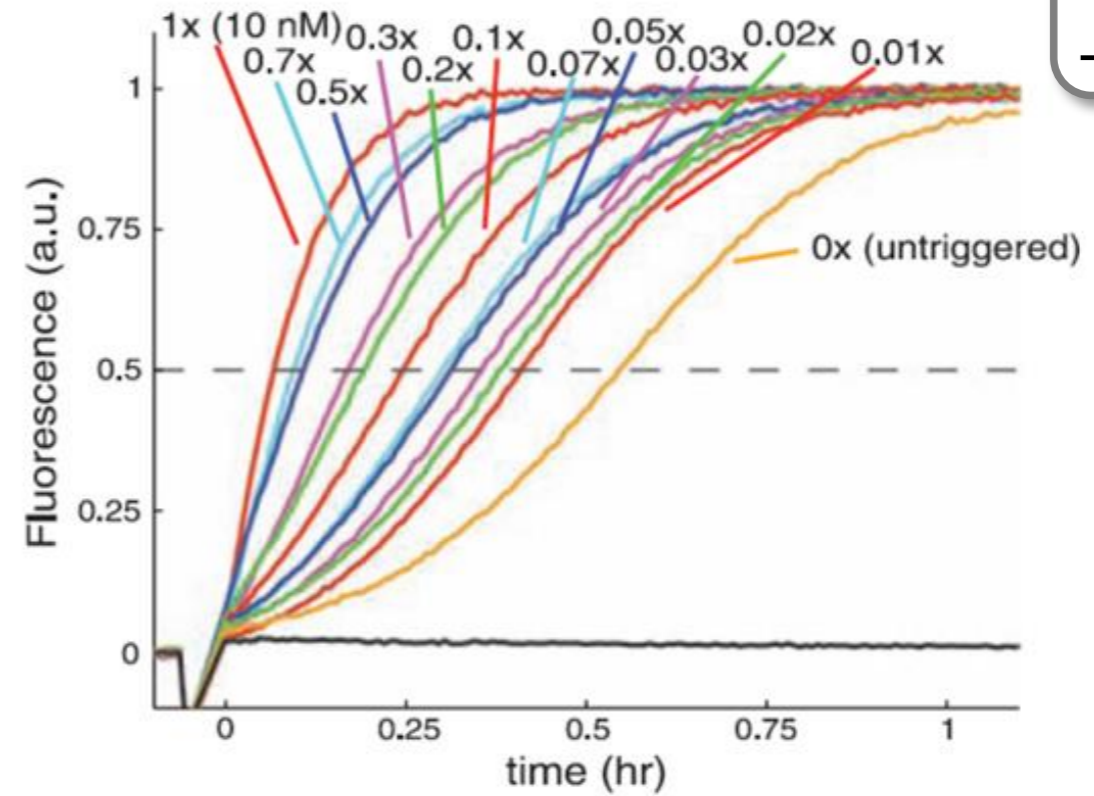
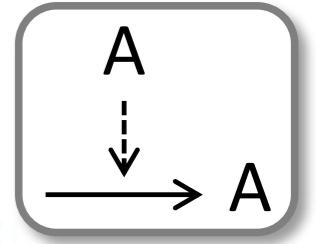
nanomechanical
devices

self-assembly of
nanostructures



Zhang et al, *Science* 2007

Boolean circuits



Zhang et al, *Science* 2007

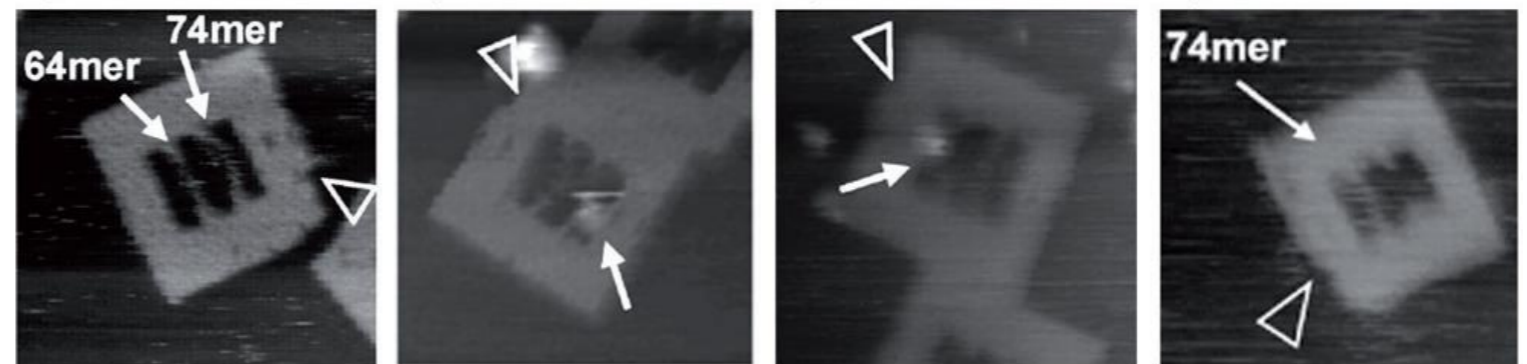
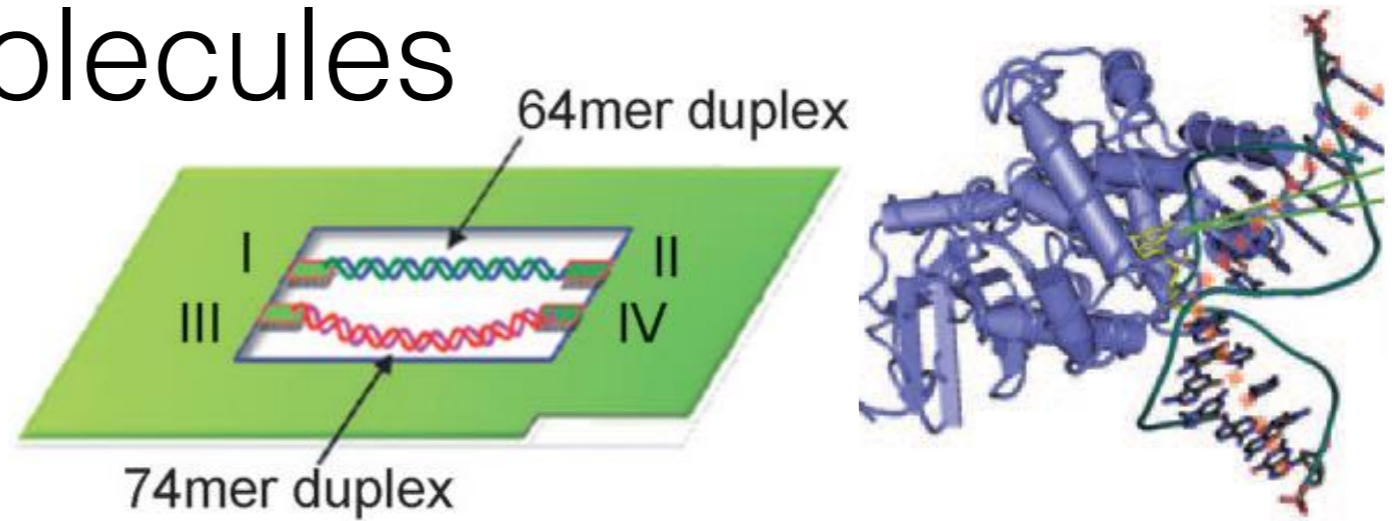
biochemical
circuits

nanomechanical
devices

self-assembly of
nanostructures

Interacting with biomolecules

fundamental architectures



Endo et al, *JACS* 2010

biochemical
circuits

nanomechanical
devices

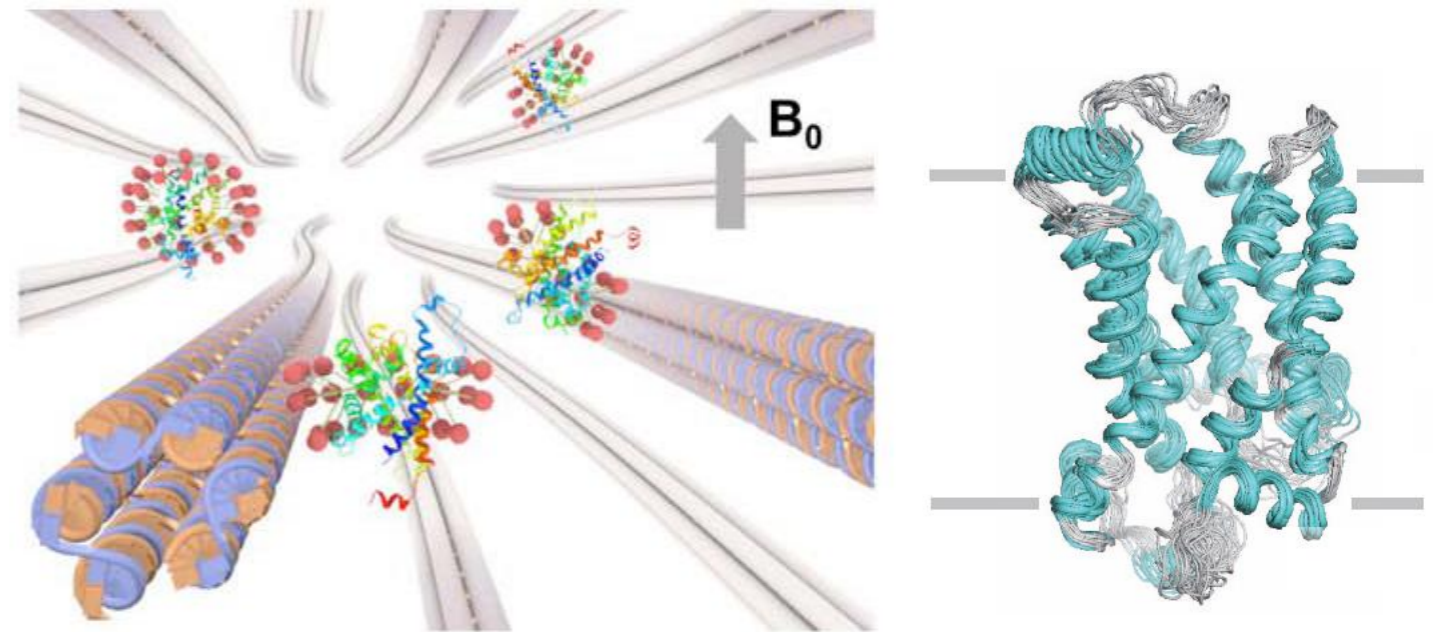
self-assembly of
nanostructures

study
DNA-protein
interactions

real-world applications

Interacting with biomolecules

fundamental architectures



Berardi et al, *Nature* 2011

biochemical
circuits

nanomechanical
devices

self-assembly of
nanostructures

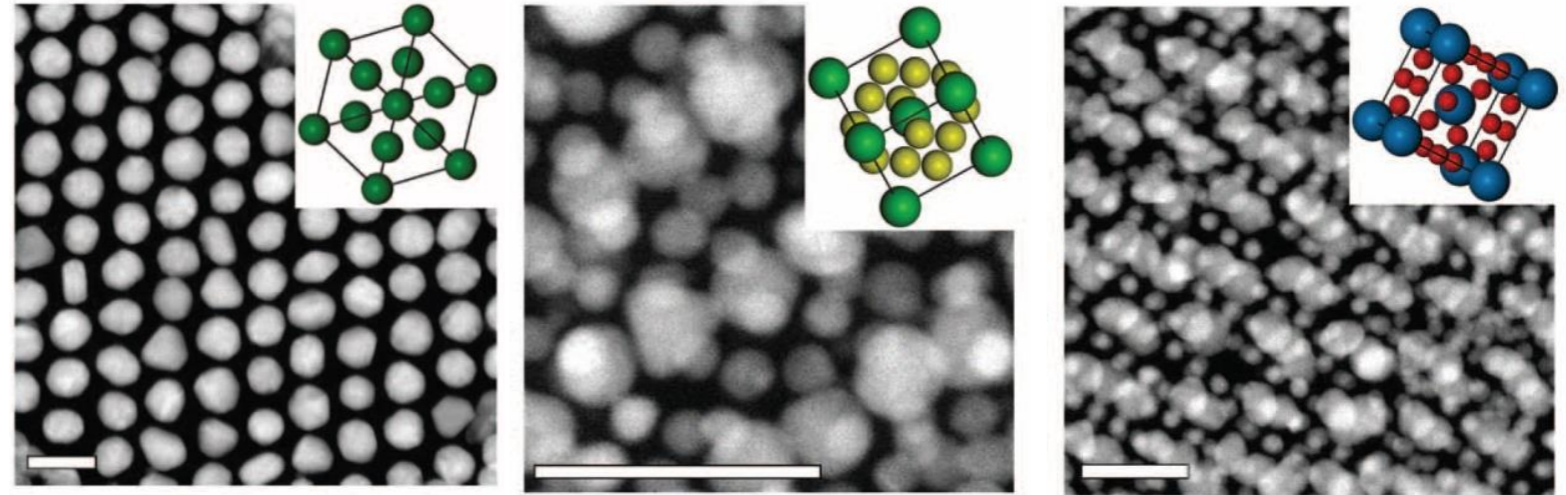
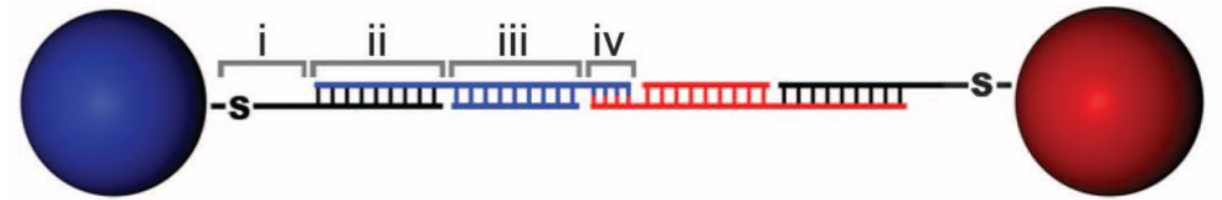
study
DNA-protein
interactions

determine
protein
structures

real-world applications

Organize nanoparticles

fundamental architectures



Macfarlane et al, *Science* 2011

biochemical
circuits

nanomechanical
devices

self-assembly of
nanostructures

study
DNA-protein
interactions

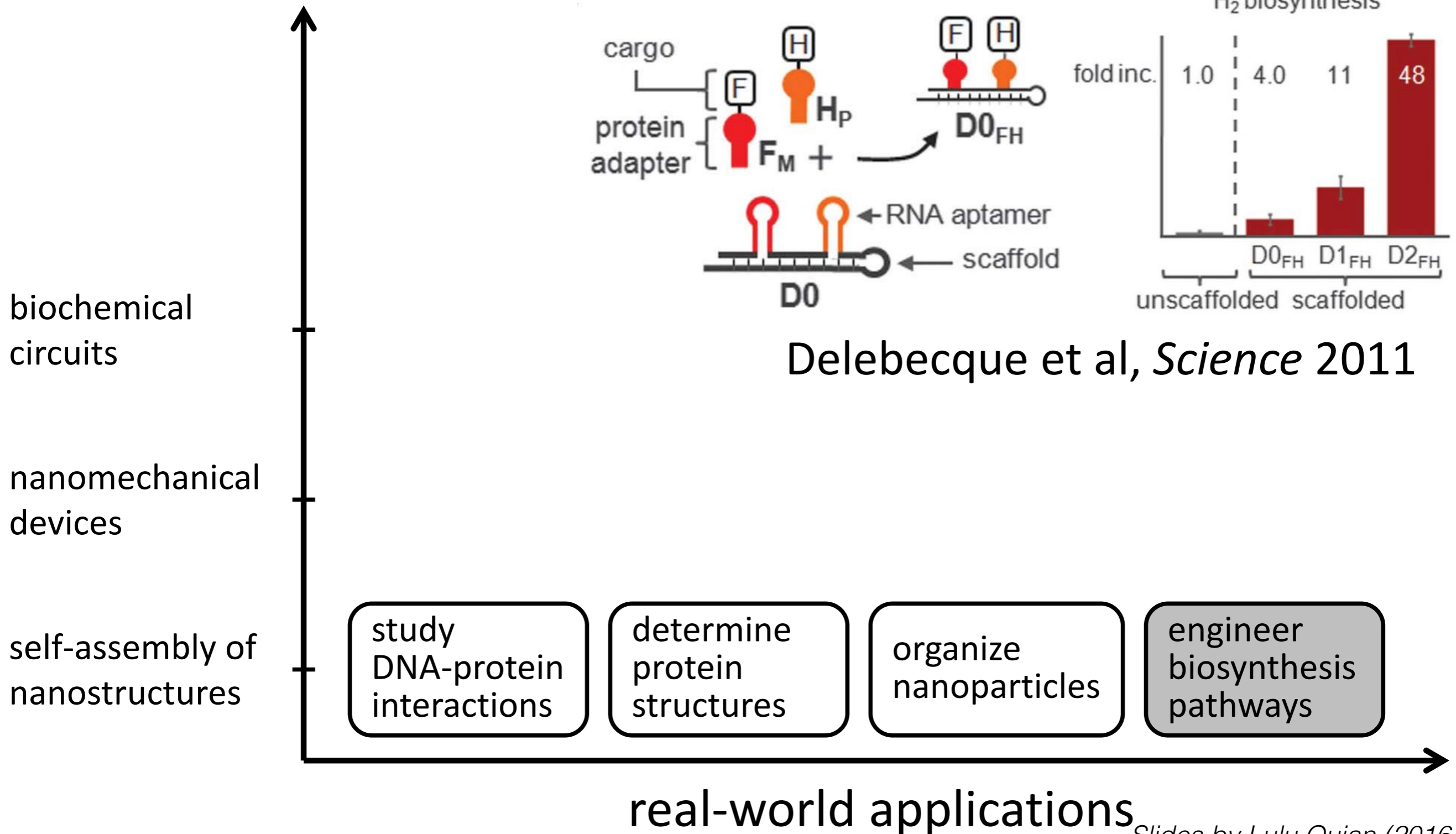
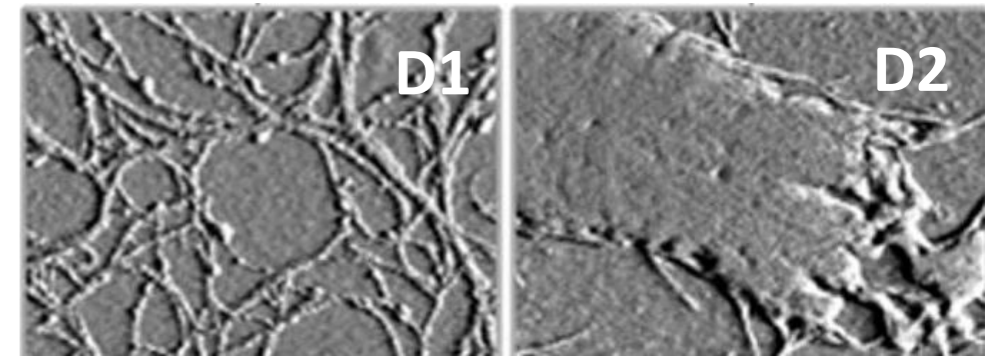
determine
protein
structures

organize
nanoparticles

real-world applications

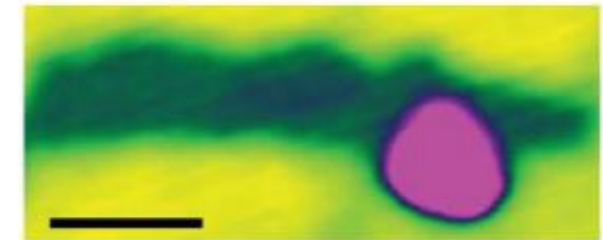
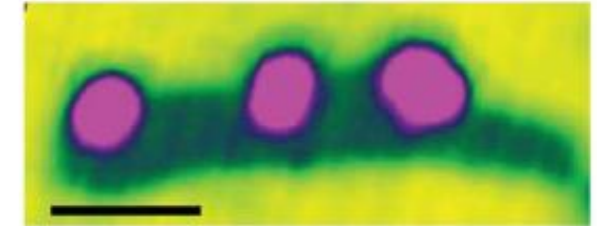
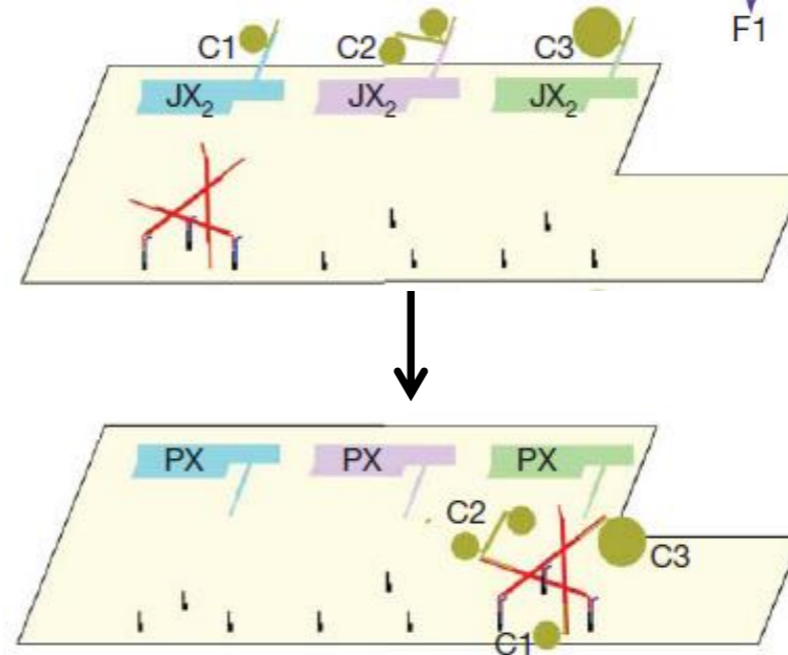
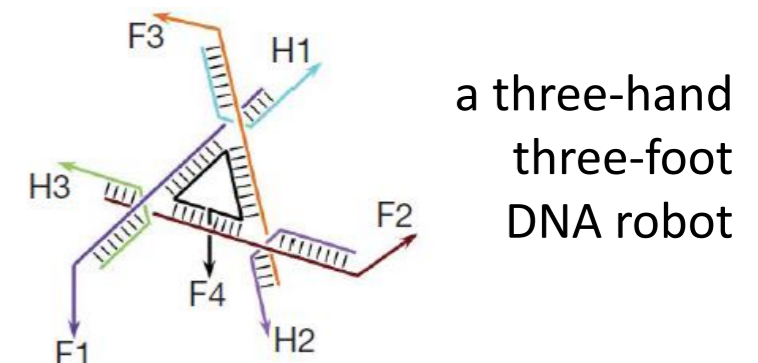
Engineer synthesis

fundamental architectures



Engineer chemical assembly lines

fundamental architectures



Gu et al, *Nature* 2010

biochemical circuits

nanomechanical devices

self-assembly of nanostructures

create an assembly line of molecules

study DNA-protein interactions

determine protein structures

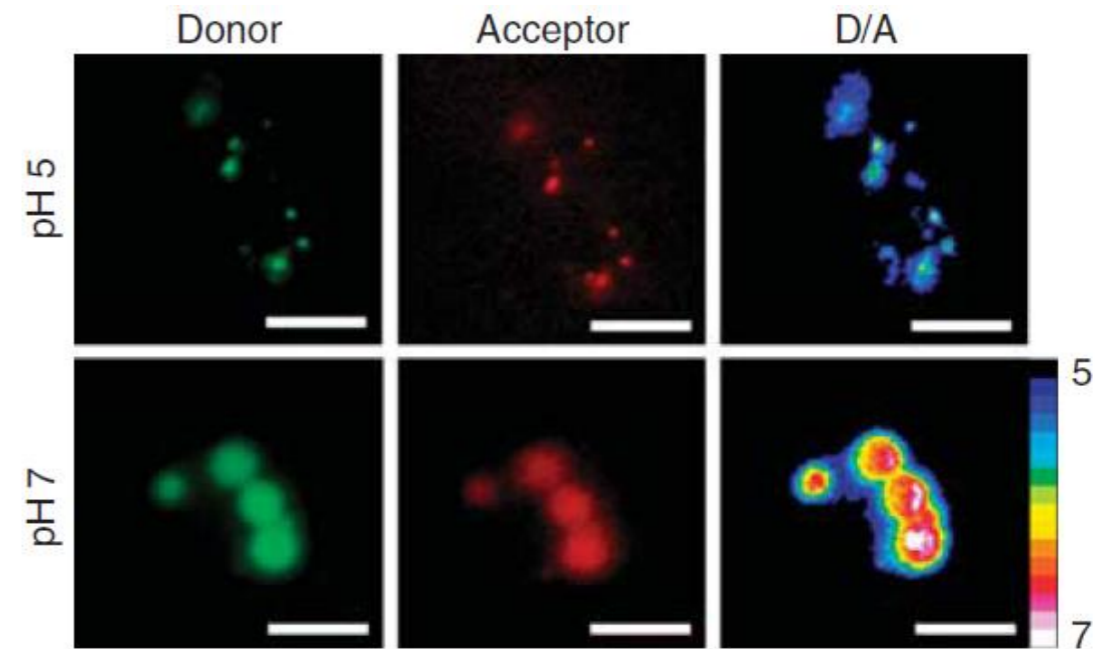
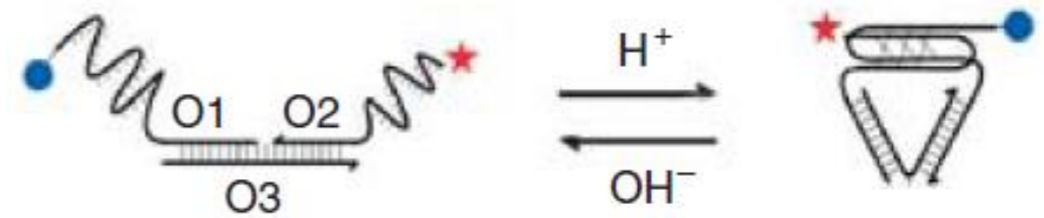
organize nanoparticles

engineer biosynthesis pathways

real-world applications

Sensors

fundamental architectures



Surana et al, *Nat. Commun.* 2011

biochemical
circuits

nanomechanical
devices

self-assembly of
nanostructures

create
an assembly line
of molecules

sense
pH conditions
in living organisms

study
DNA-protein
interactions

determine
protein
structures

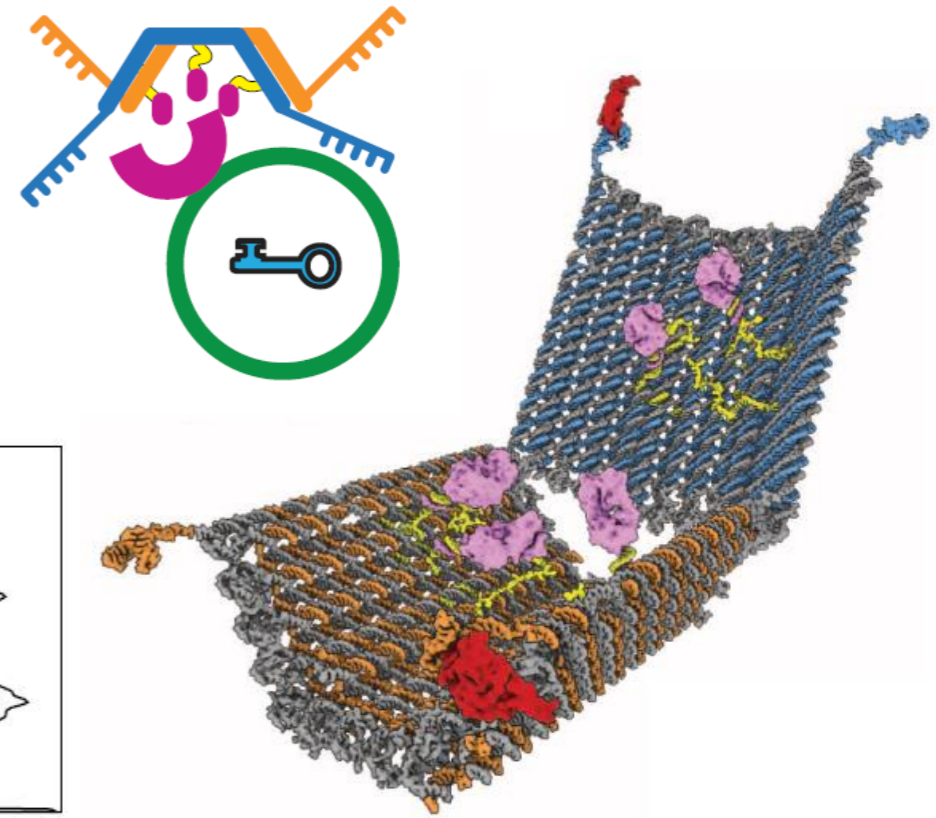
organize
nanoparticles

engineer
biosynthesis
pathways

real-world applications

Delivery devices

fundamental architectures



biochemical
circuits

Douglas et al, *Science* 2012

nanomechanical
devices

create
an assembly line
of molecules

sense
pH conditions
in living organisms

target
delivery
to cells

self-assembly of
nanostructures

study
DNA-protein
interactions

determine
protein
structures

organize
nanoparticles

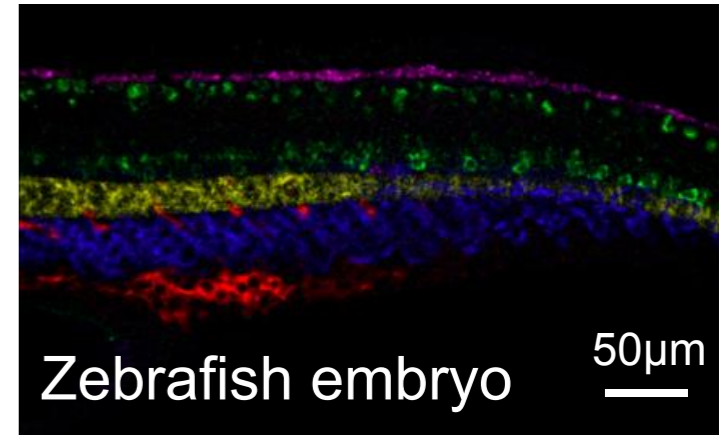
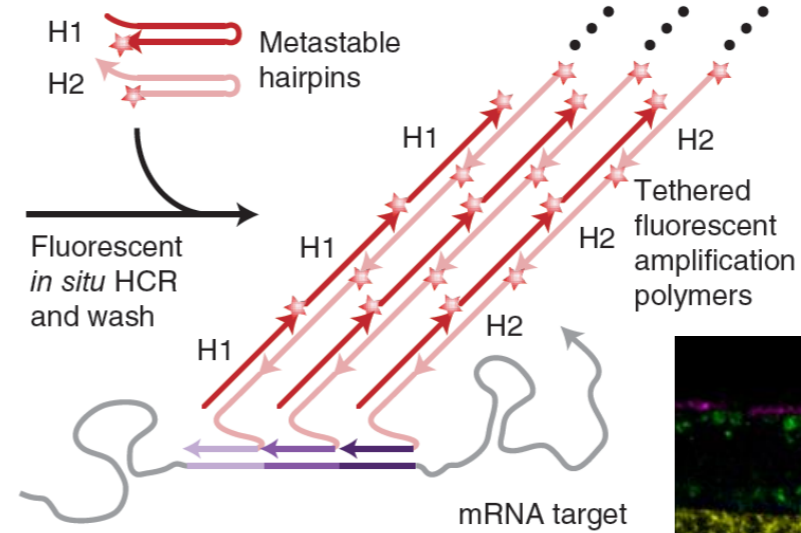
engineer
biosynthesis
pathways

real-world applications

Slides by Lulu Quian (2016)

Imaging devices

fundamental architectures



Choi et al, *Nat. Biotech.* 2010

biochemical circuits

image mRNA expression

nanomechanical devices

create an assembly line of molecules

sense pH conditions in living organisms

target delivery to cells

self-assembly of nanostructures

study DNA-protein interactions

determine protein structures

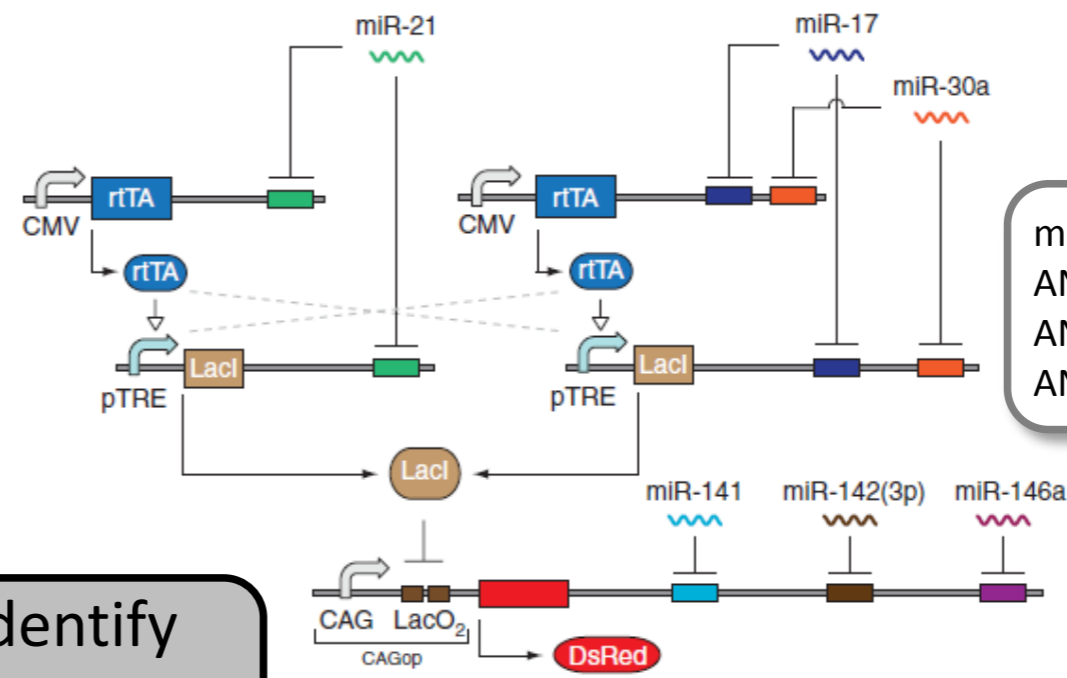
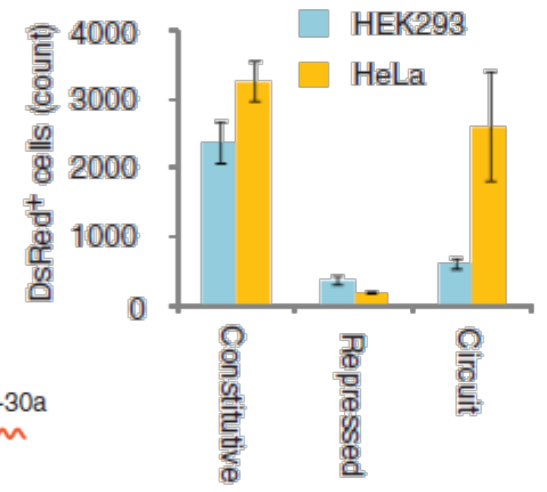
organize nanoparticles

engineer biosynthesis pathways

real-world applications

Logic-based sensor

fundamental architectures



miR-21 AND miR-17-30a
AND NOT (miR-141)
AND NOT (miR-142(3p))
AND NOT (miR-146a)

Xie et al, *Science* 2011

biochemical circuits

- image mRNA expression
- identify cancer cells

nanomechanical devices

- create an assembly line of molecules
- sense pH conditions in living organisms
- target delivery to cells

self-assembly of nanostructures

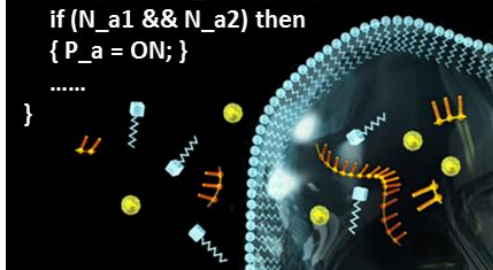
- study DNA-protein interactions
- determine protein structures
- organize nanoparticles
- engineer biosynthesis pathways

real-world applications

Towards programming languages

fundamental architectures

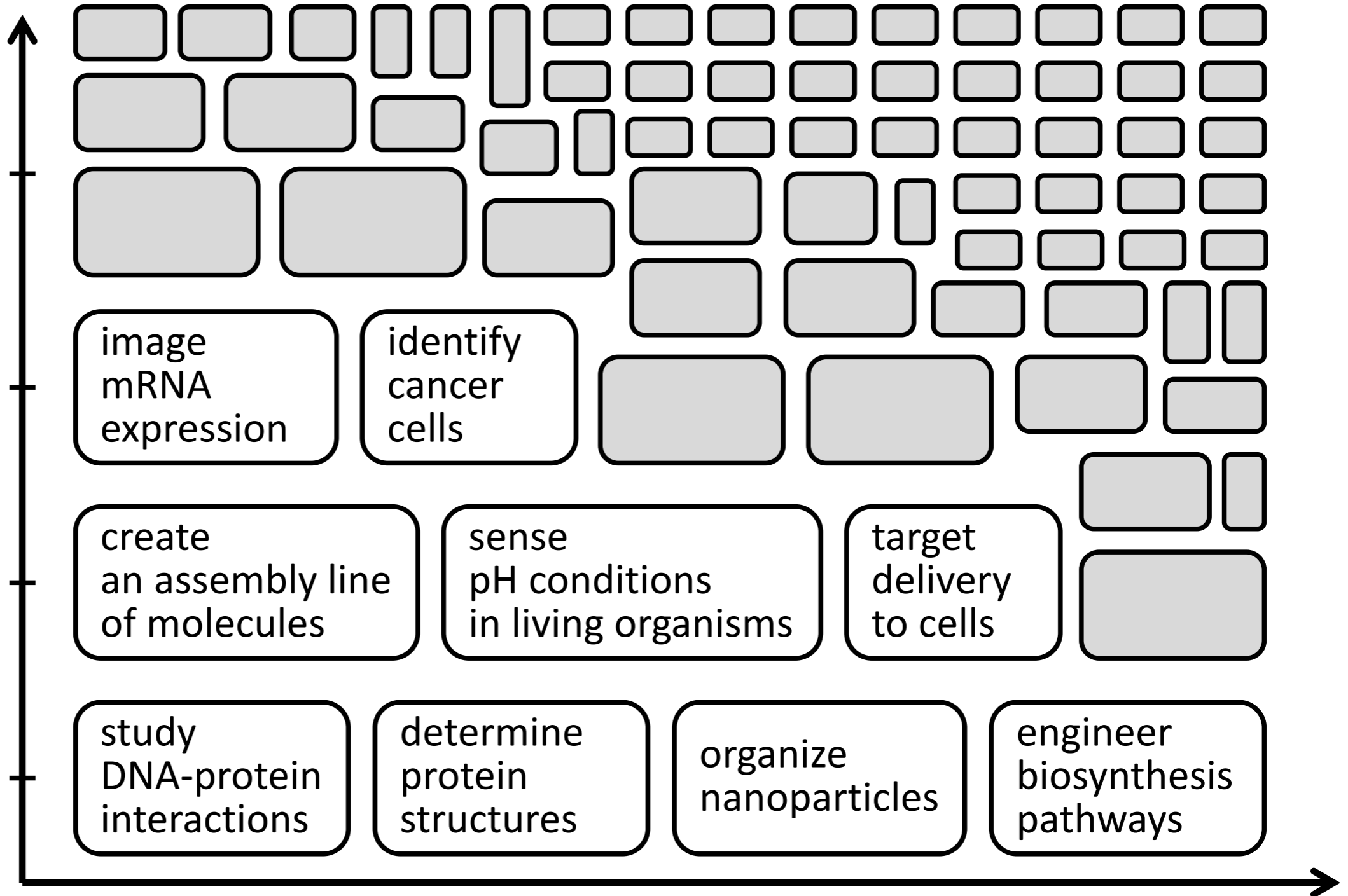
```
type P = ●, N = ↗, L = ⚡;  
declare system Bacterium Q13_7 {  
  component P_a = P(1, N_a1, N_a2);  
  component N_a1 = N(7), N_a2 = N(9);  
  component L_m = L(100);  
  if (N_a1 && N_a2) then  
  { P_a = ON; }  
  .....  
}
```



biochemical
circuits

nanomechanical
devices

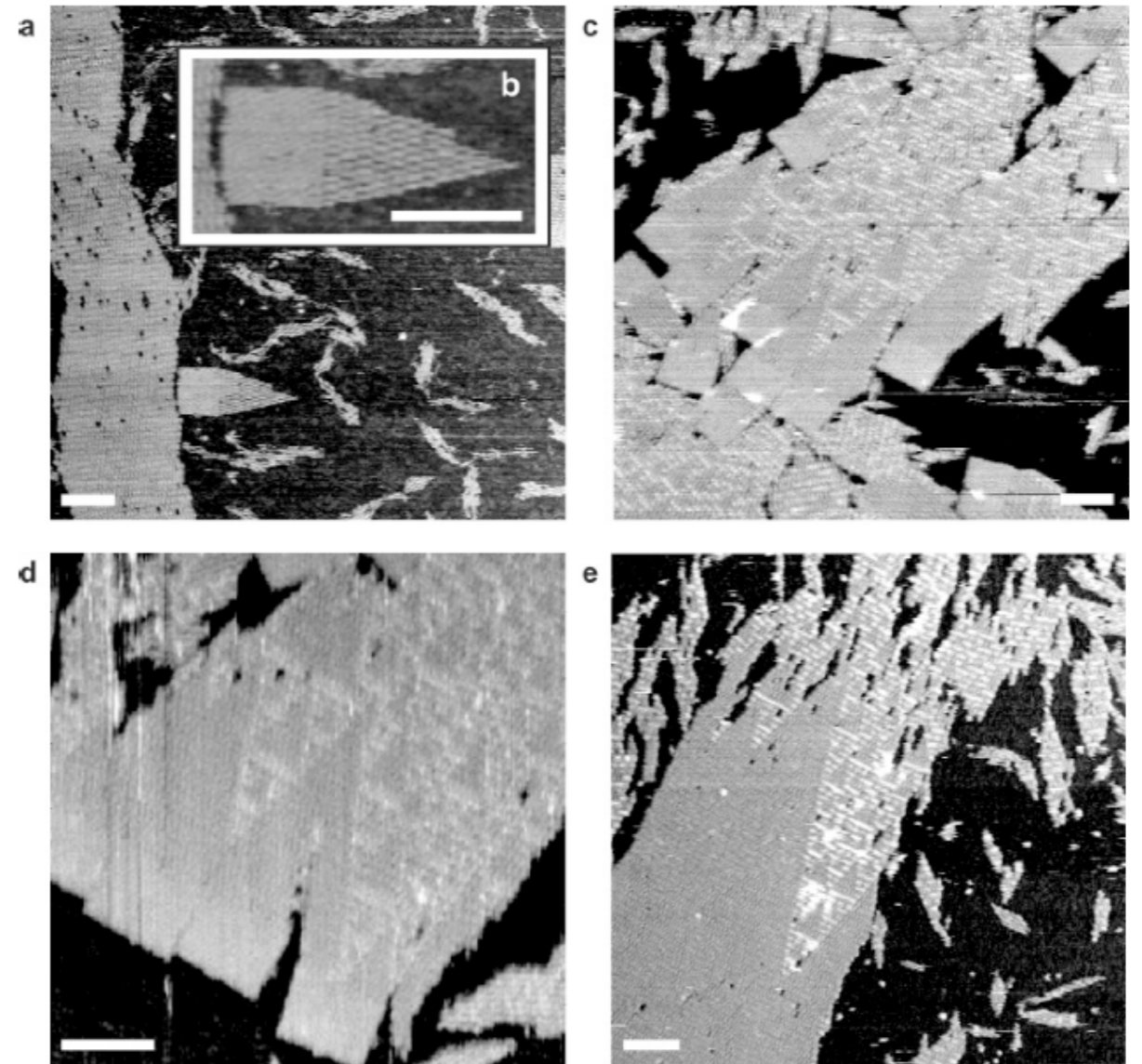
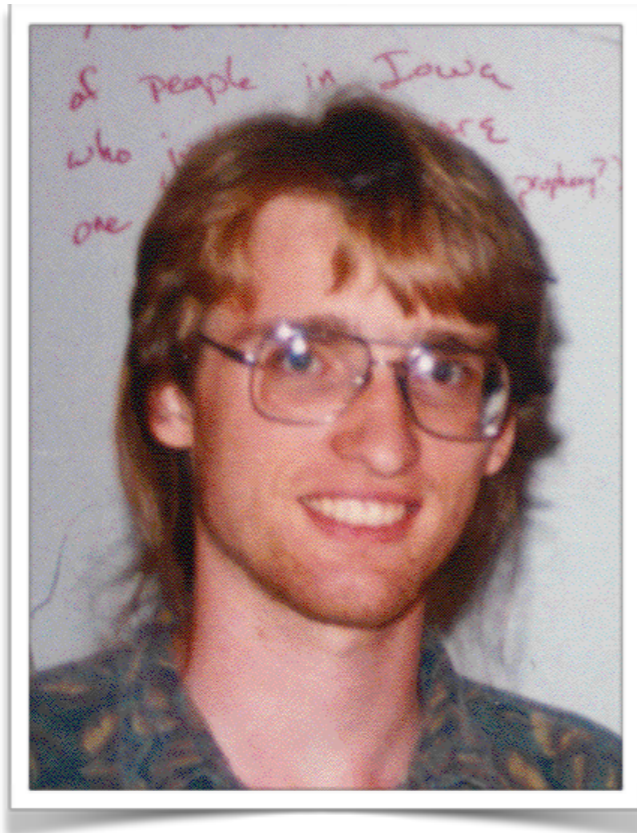
self-assembly of
nanostructures



real-world applications

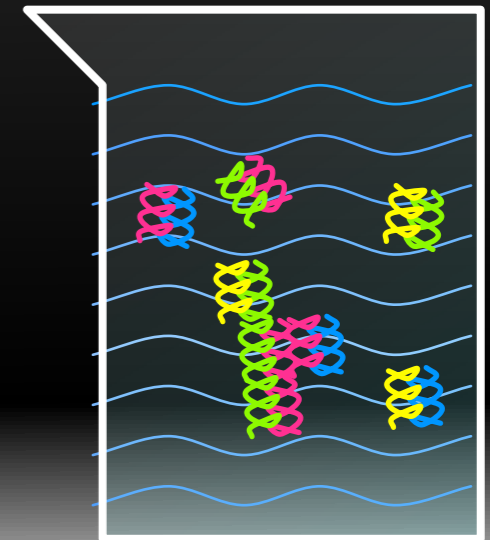
Part I: Tile assembly

Erik Winfree (1998-): DNA algorithmic self-assembly



Principle of DNA algorithmic self-assembly

FedEx



1. Thermal cycler

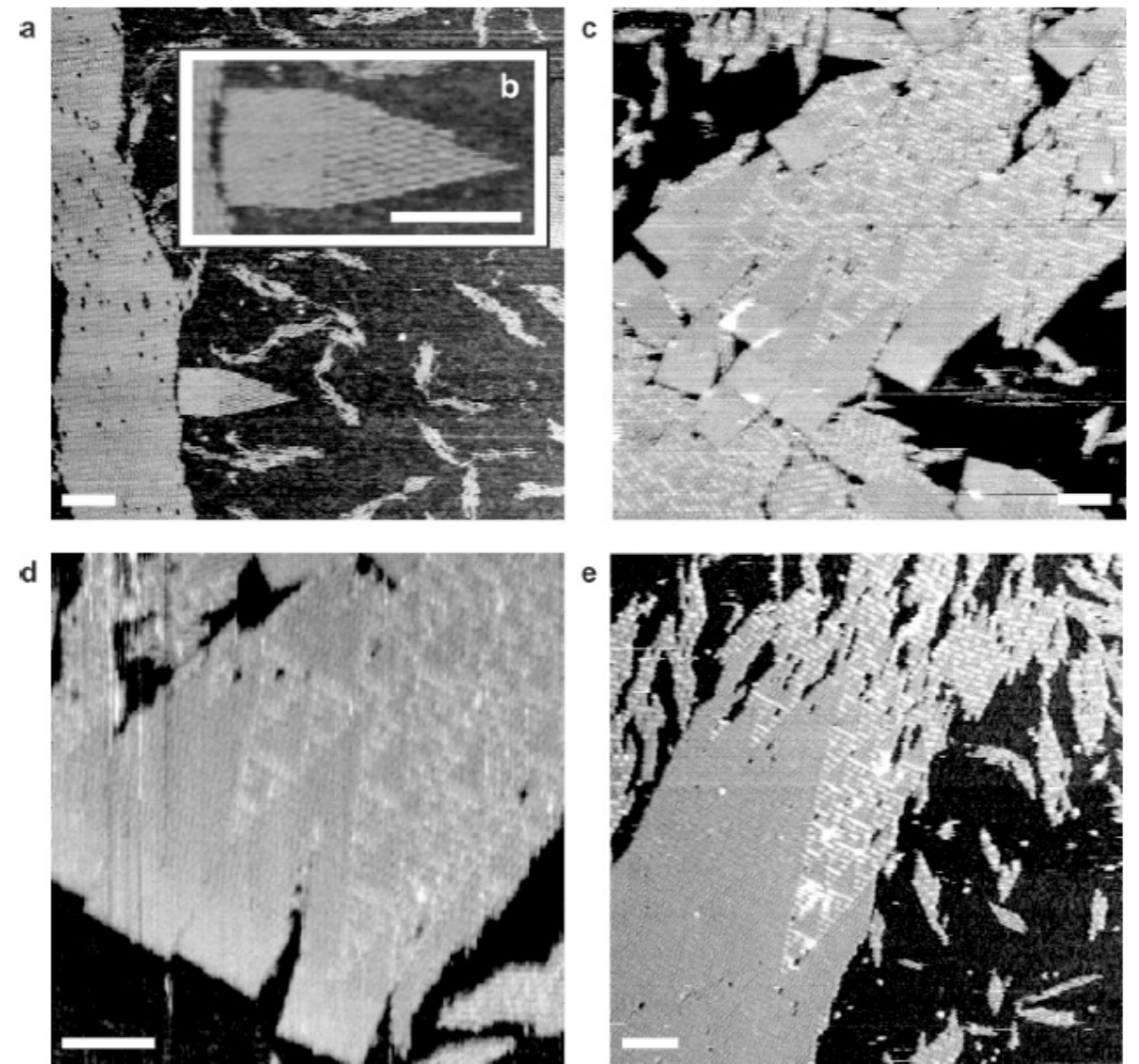
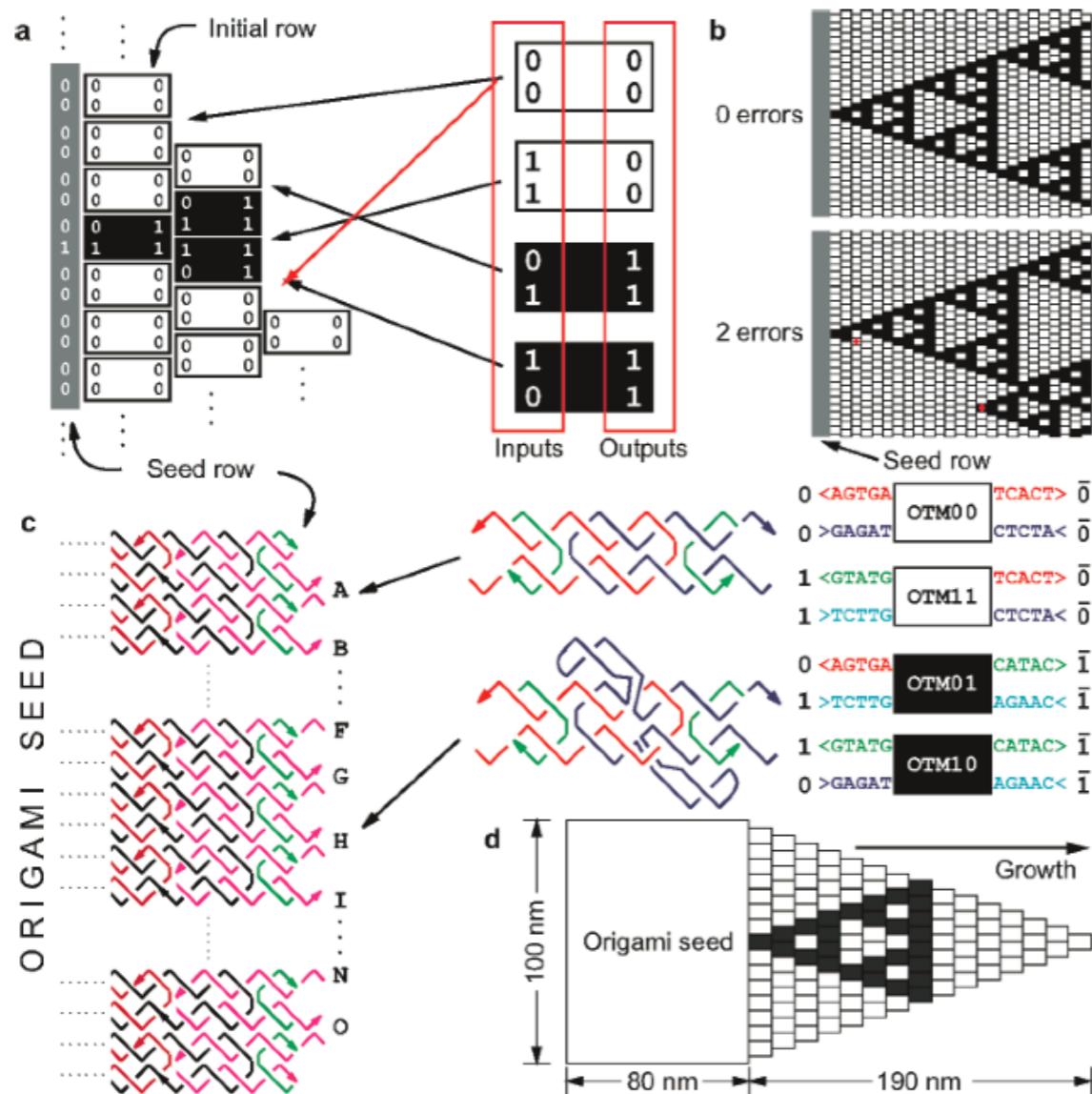
Hot / less hot cycles for exponential duplication of the DNA strands by Polymerase Chain Reaction (PCR)

⇒ the tiles 

2. One-pot reaction

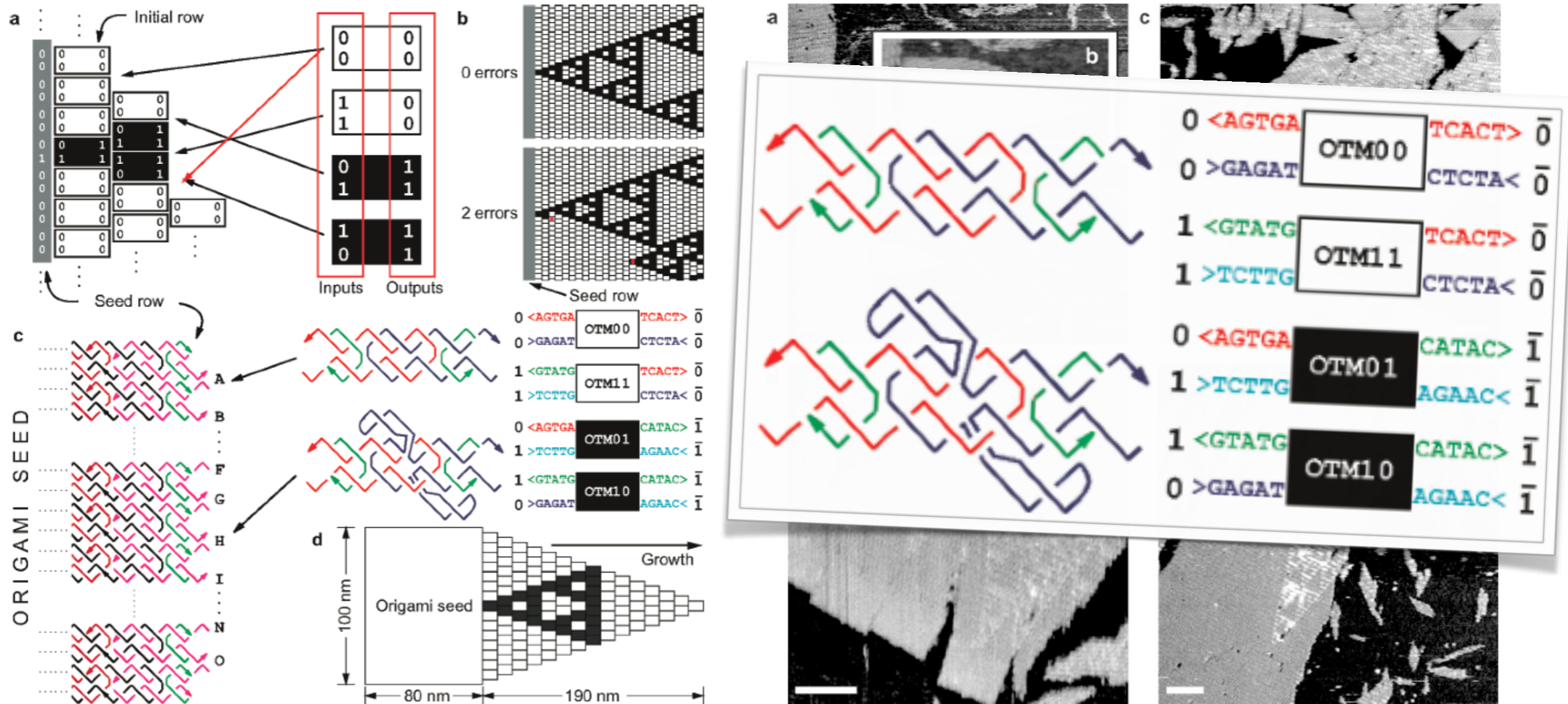
Mixing the tiles at 90°C and letting the solution cool down to room temperature for few hours

Erik Winfree (1998-): DNA algorithmic self-assembly



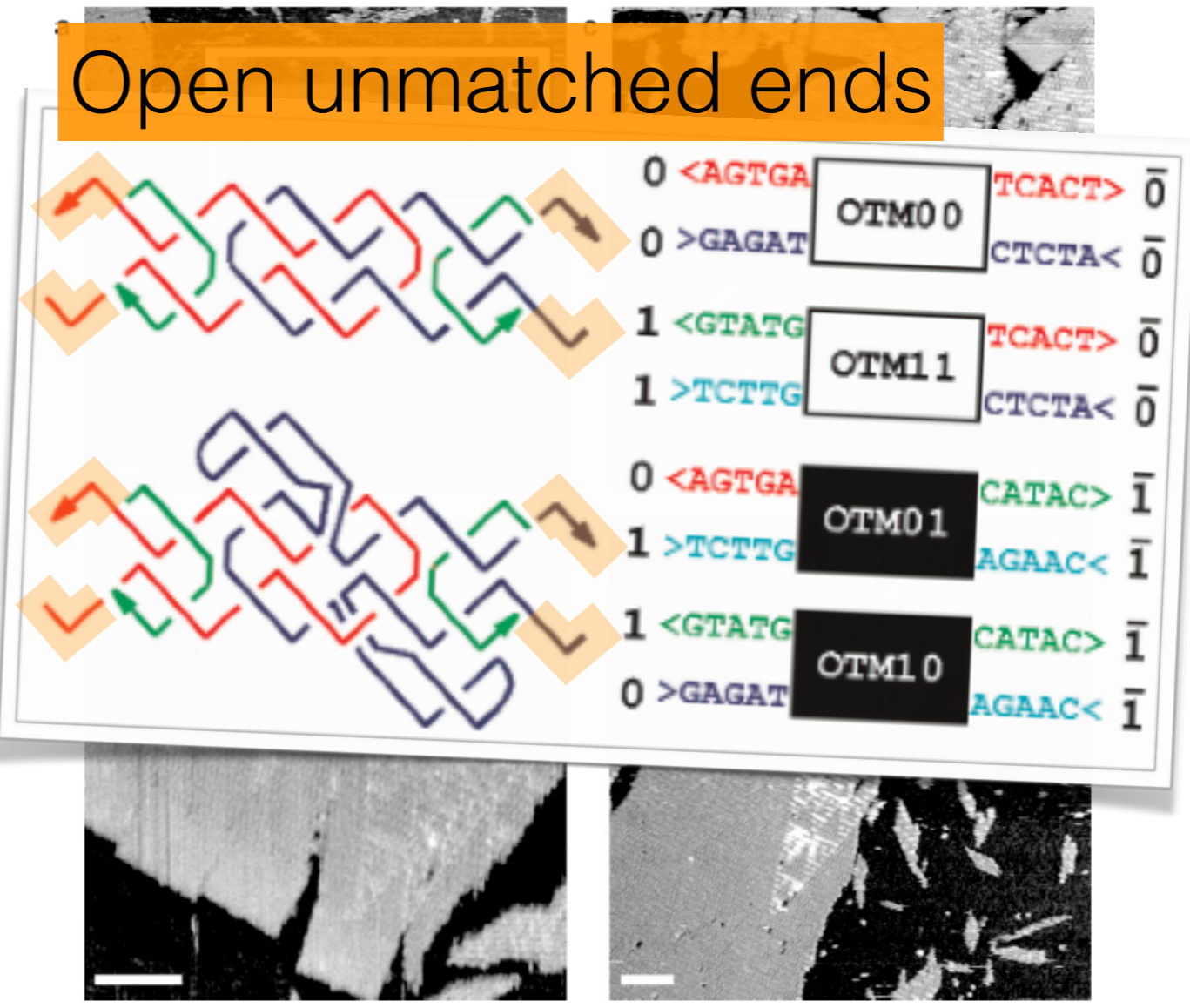
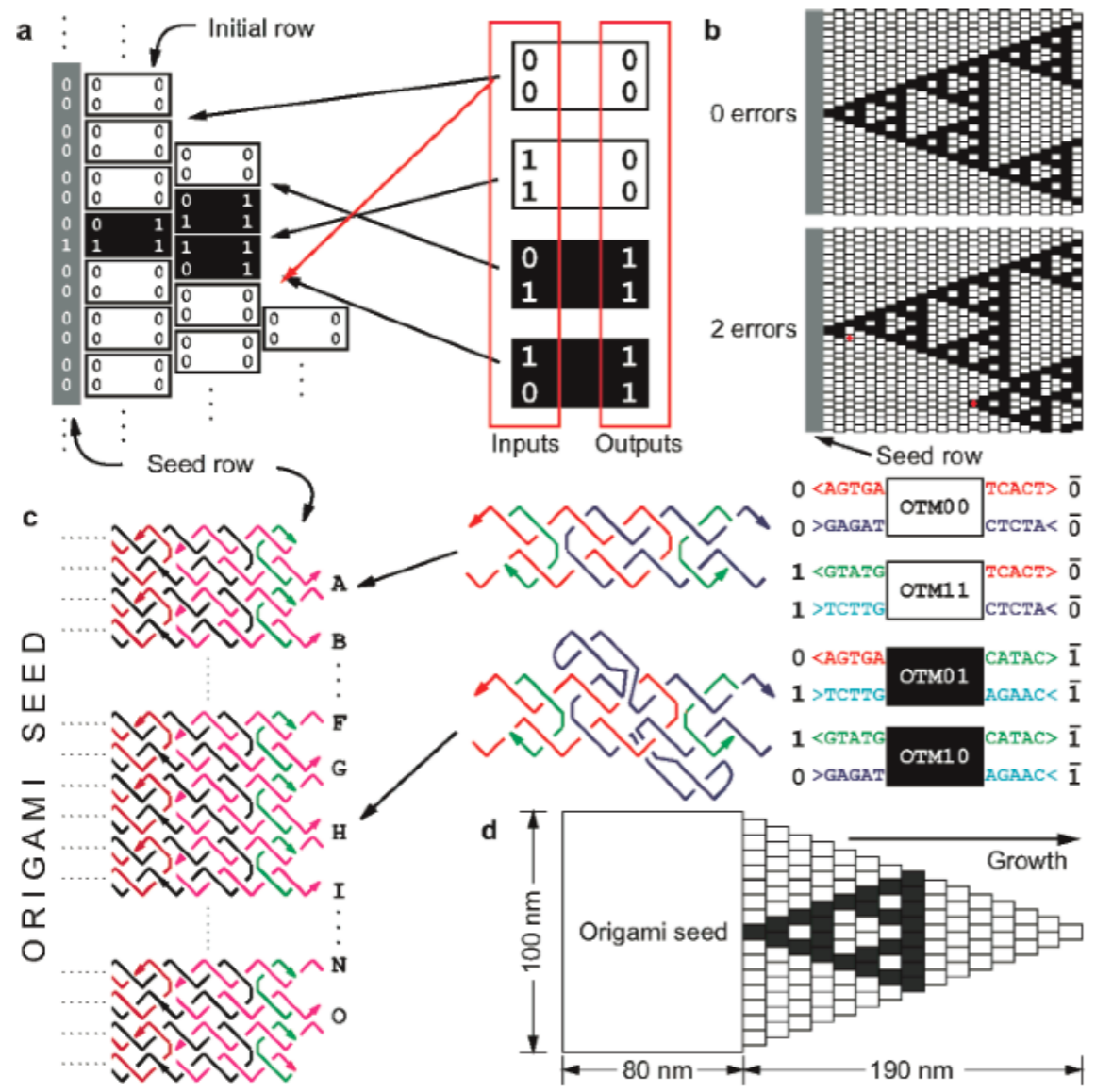
Credits: K. Fujibayashi, R. Hariadi, S.H. Park, E. Winfree & S. Murata

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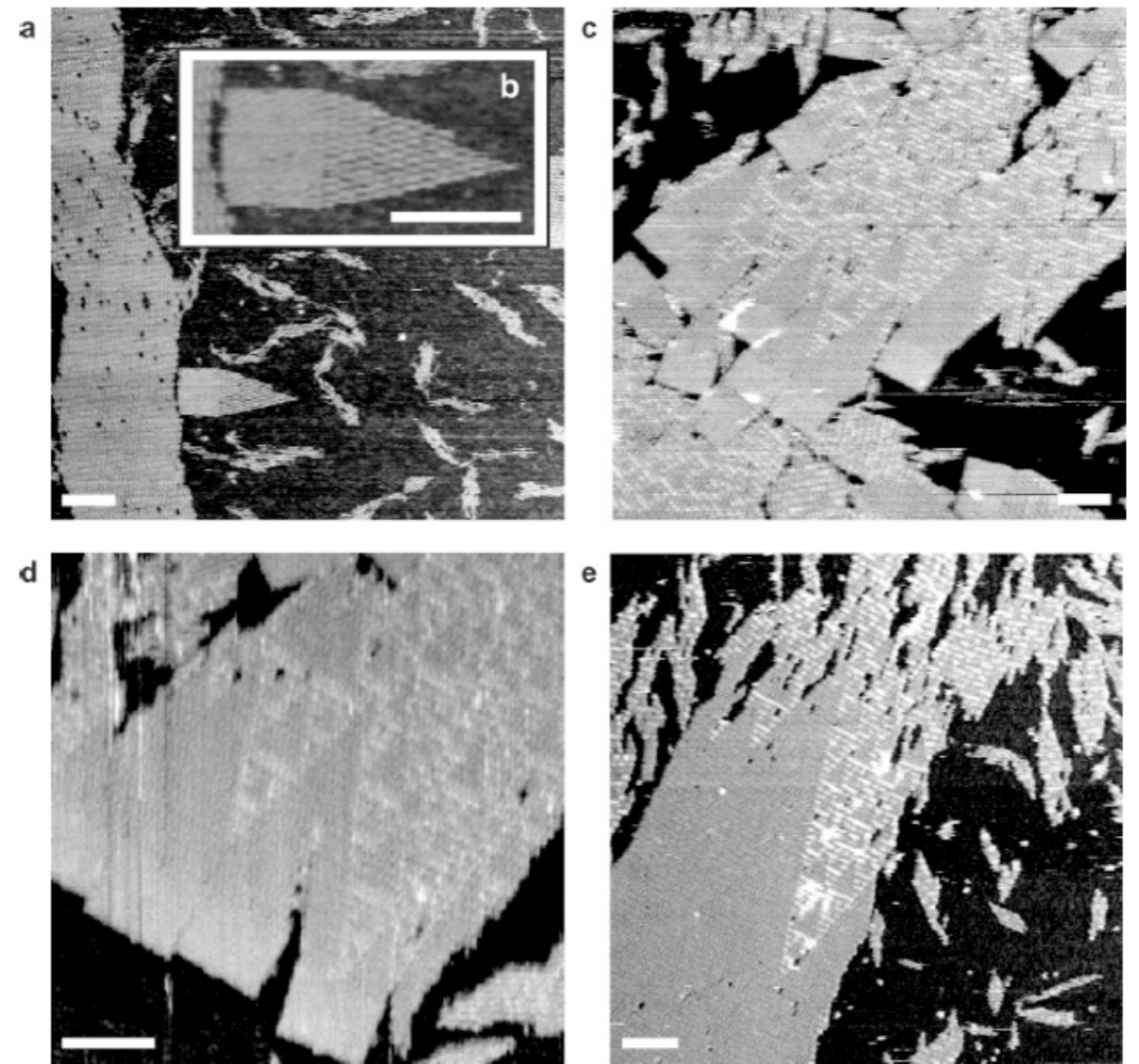
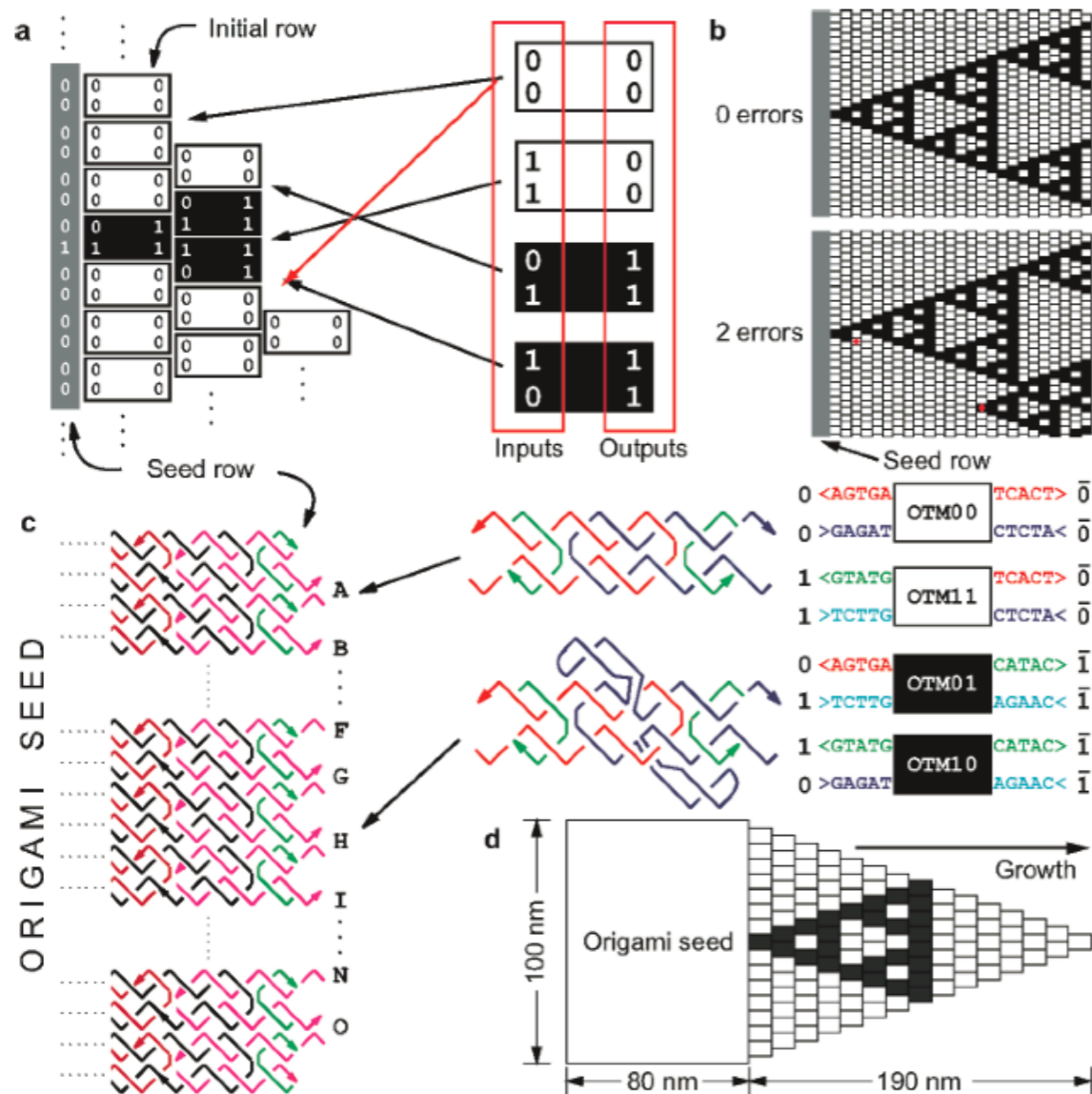
Credits: K. Fujibayashi, R. Hariadi, S.H. Park, E. Winfree & S. Murata

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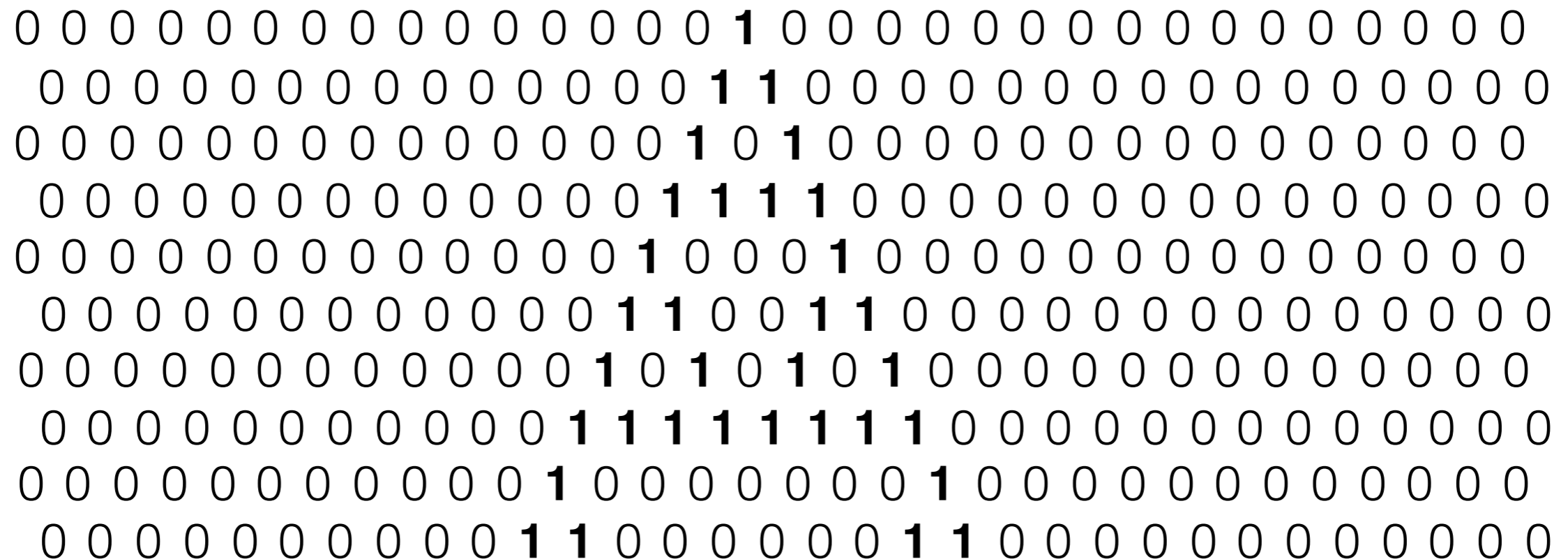
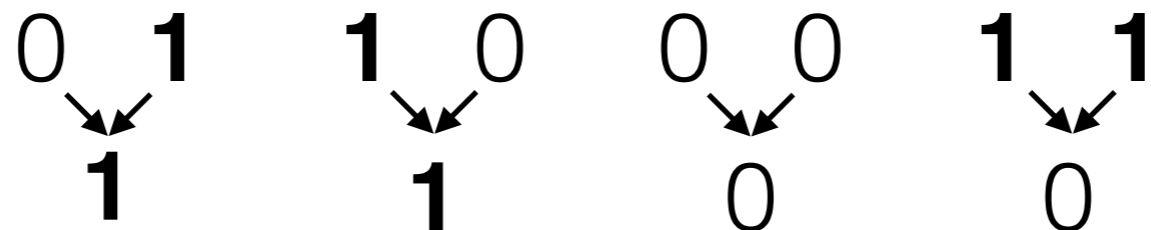
Credits: K. Fujibayashi, R. Hariadi, S.H. Park, E. Winfree & S. Murata

Calculer =

Transformer l'information

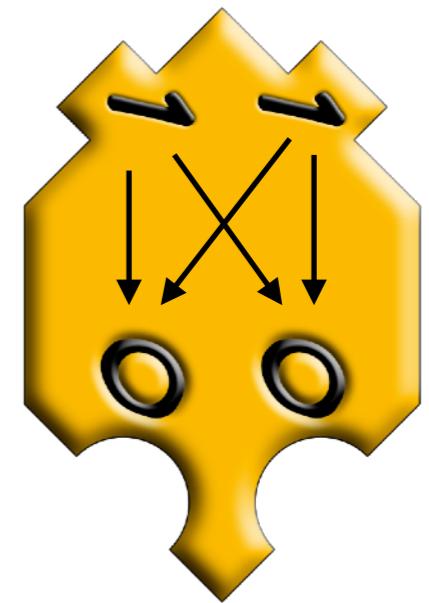
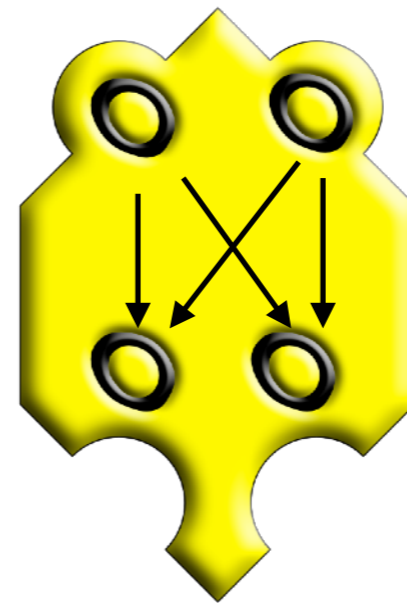
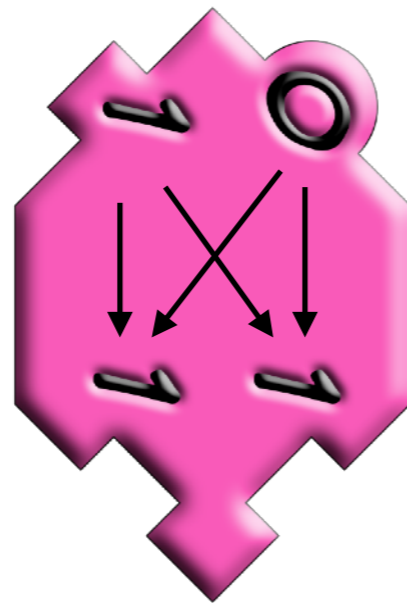
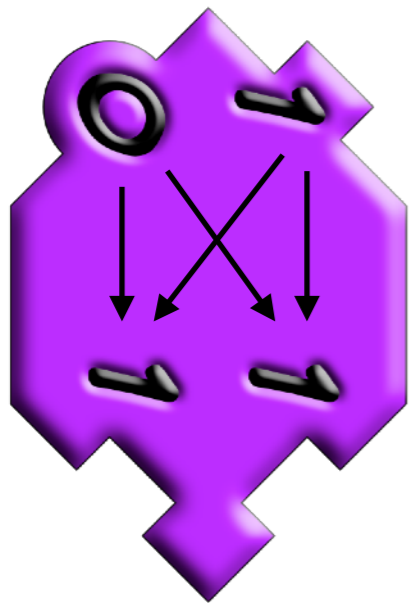
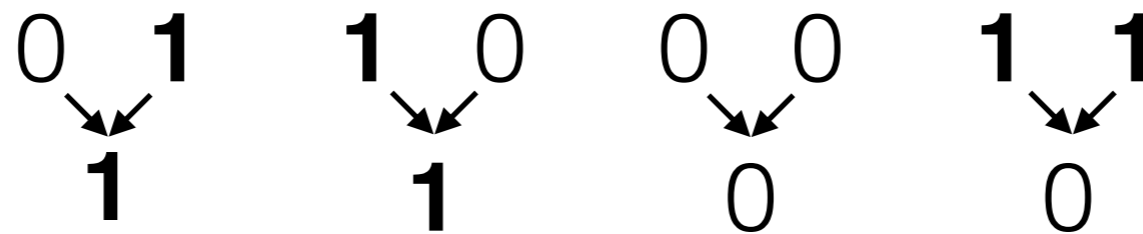
Triangles de Sierpiński

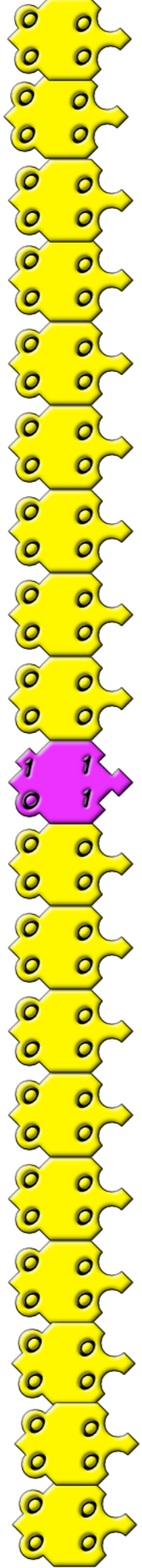
Règles de substitutions:



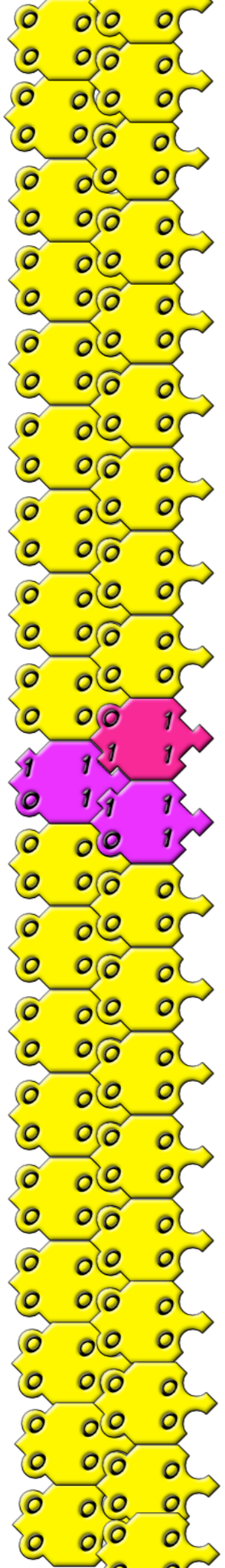
Triangles de Sierpiński

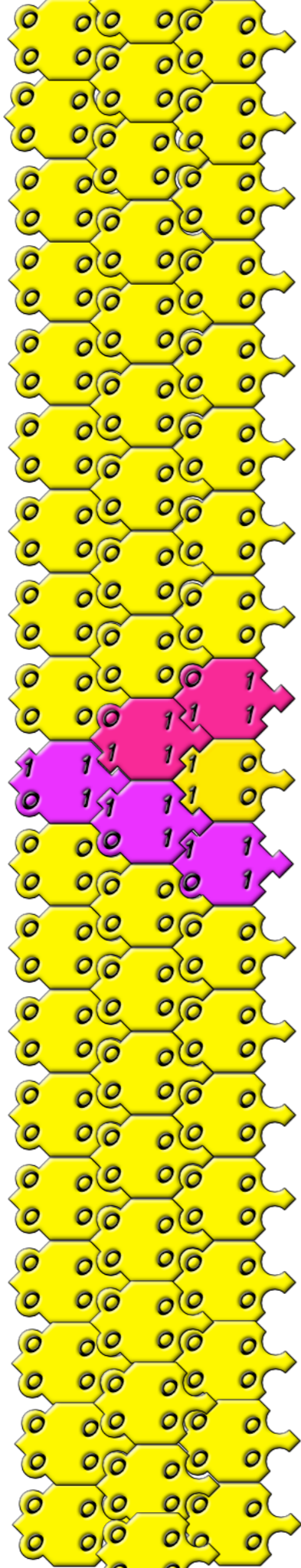
Règles de substitutions:

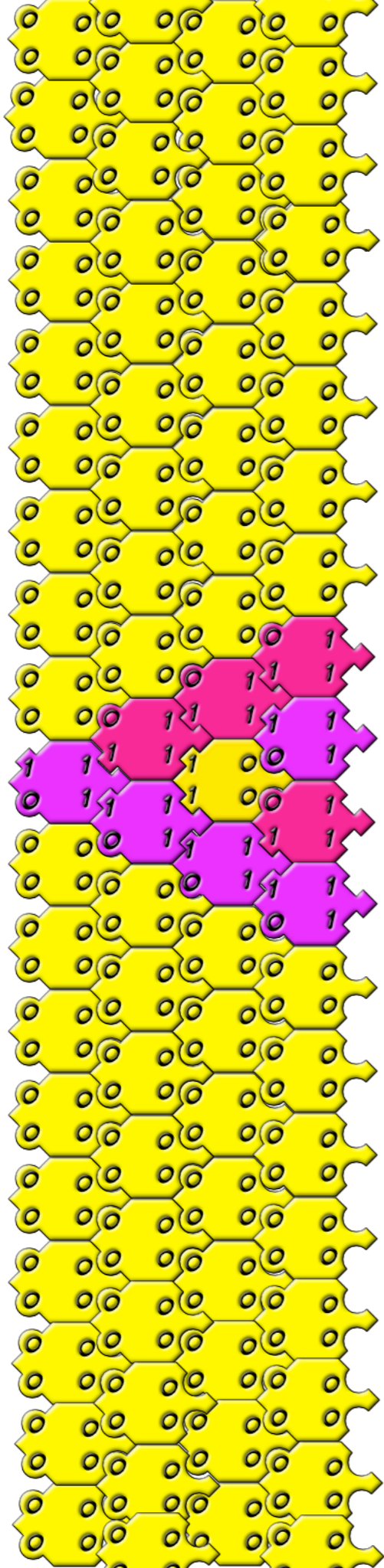


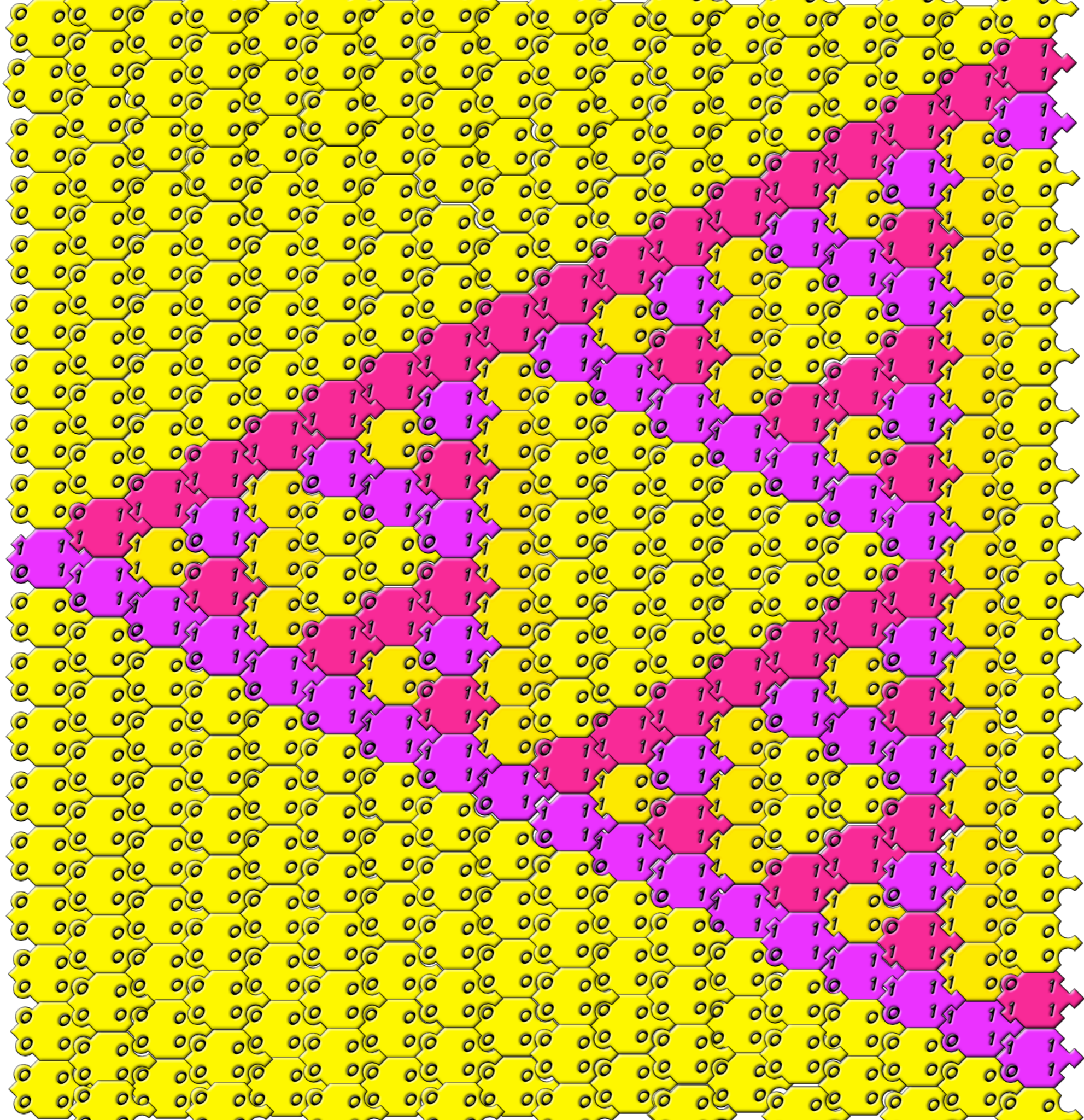


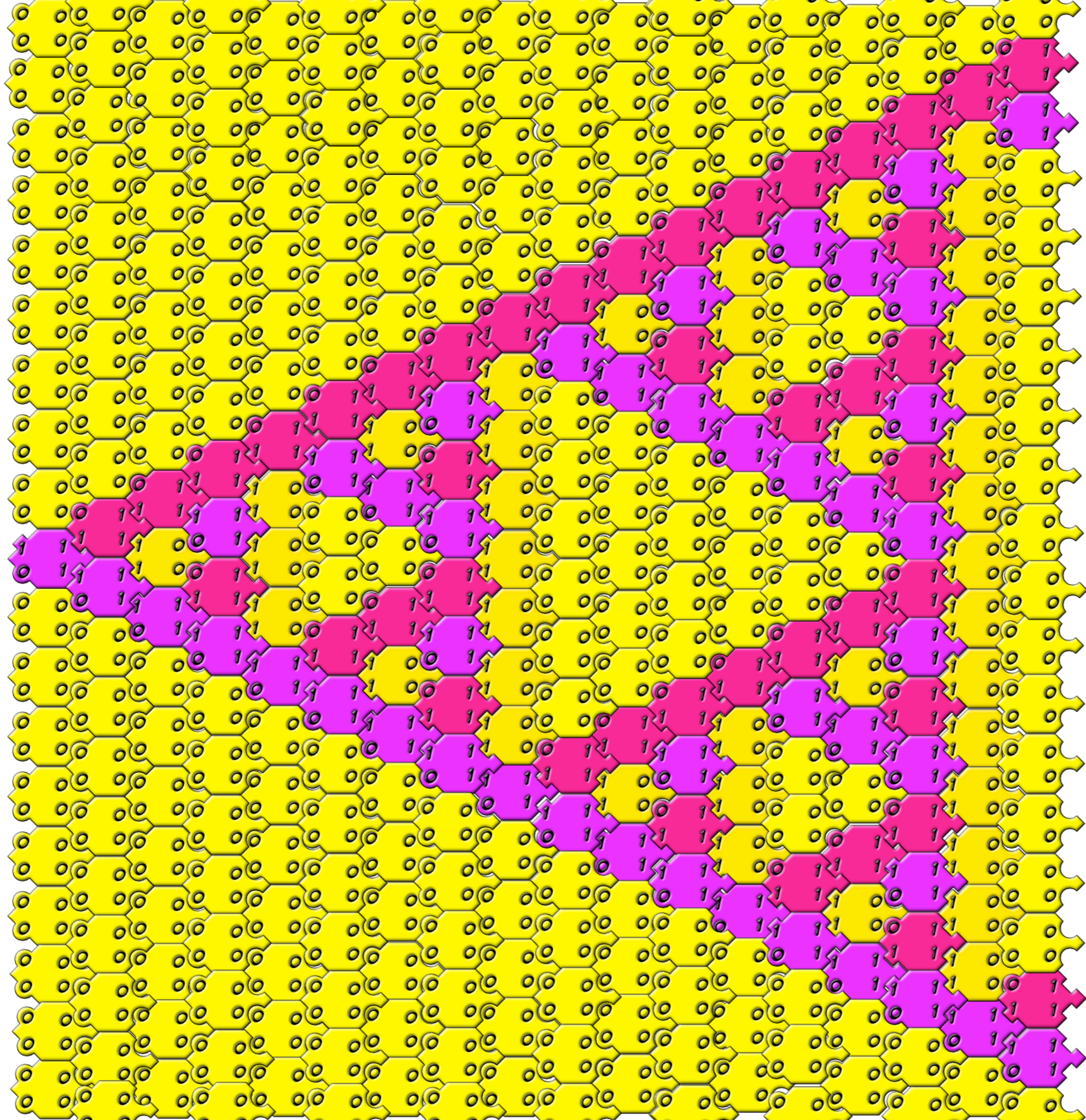
1 1
1



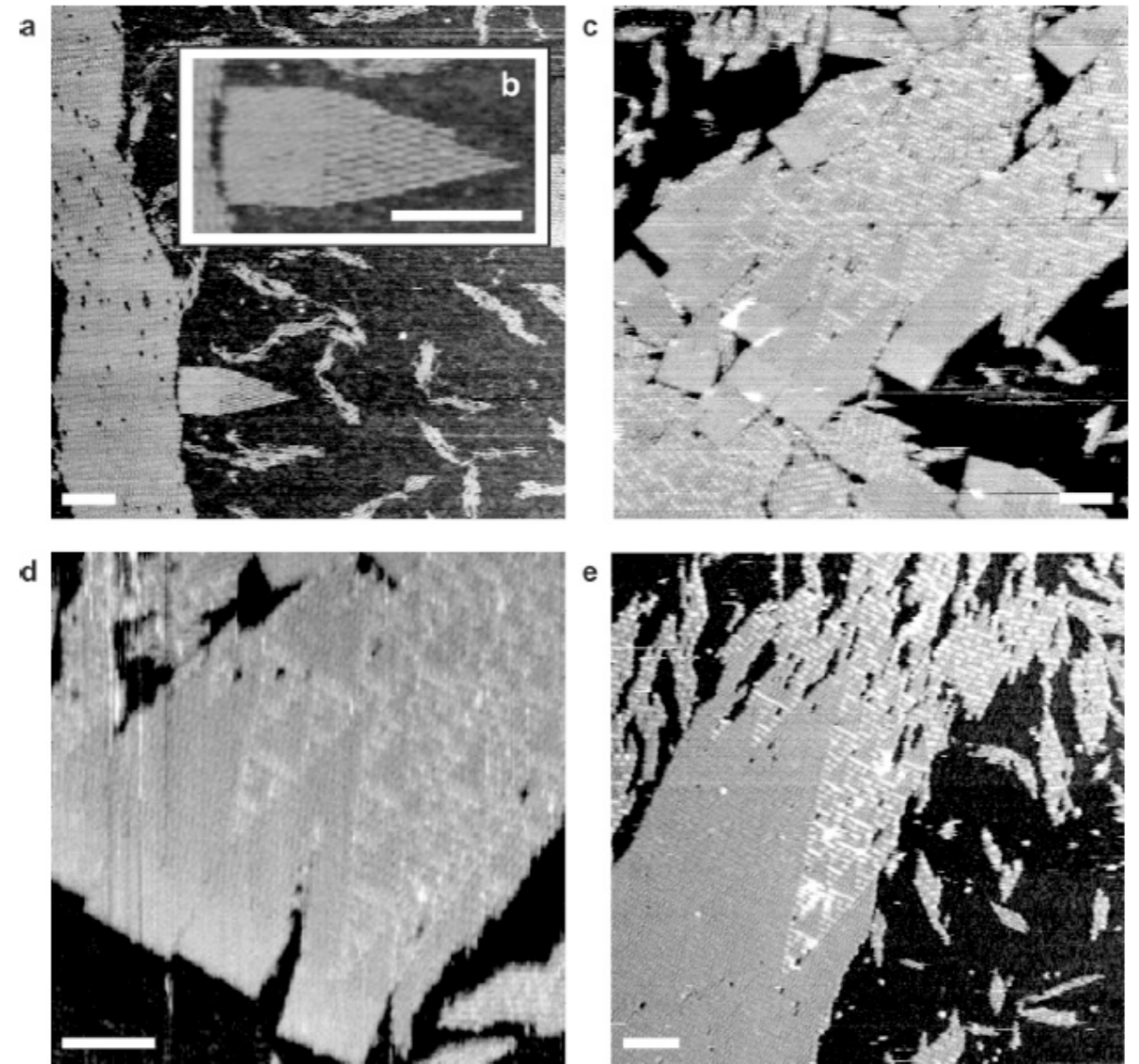
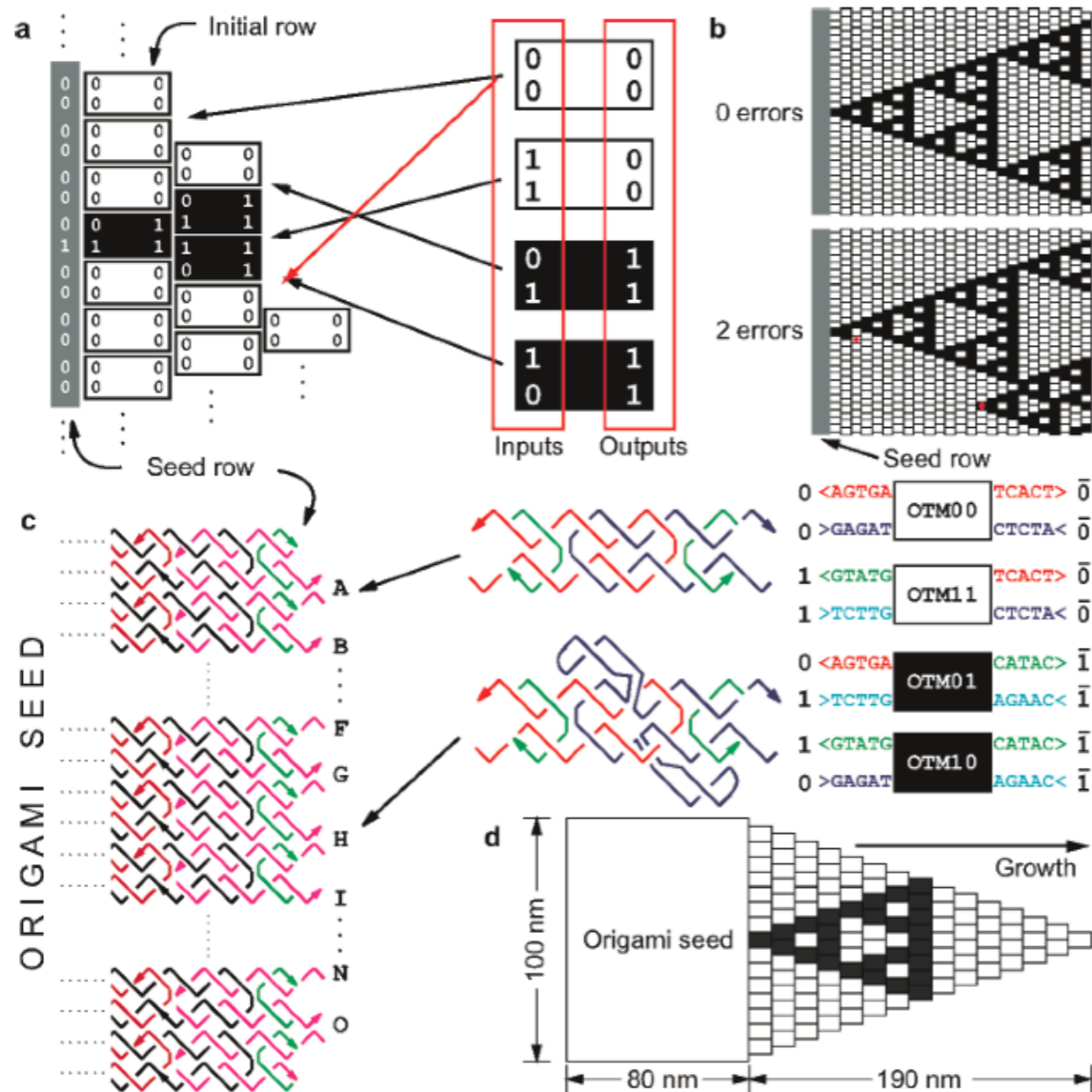






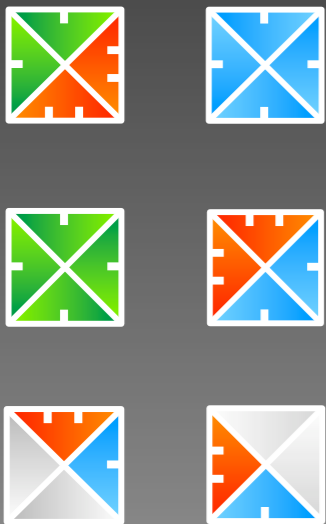


Erik Winfree (1998-): DNA algorithmic self-assembly



Credits: K. Fujibayashi, R. Hariadi, S.H. Park, E. Winfree & S. Murata

Algorithmic model



- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



*White glue with
strength 0*



*Red glue with
strength 2*



*Blue glue with
strength 1*

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



*White glue with
strength 0*

*Red glue with
strength 2*



*Blue glue with
strength 1*



The seed

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



*White glue with
strength 0*

*Red glue with
strength 2*



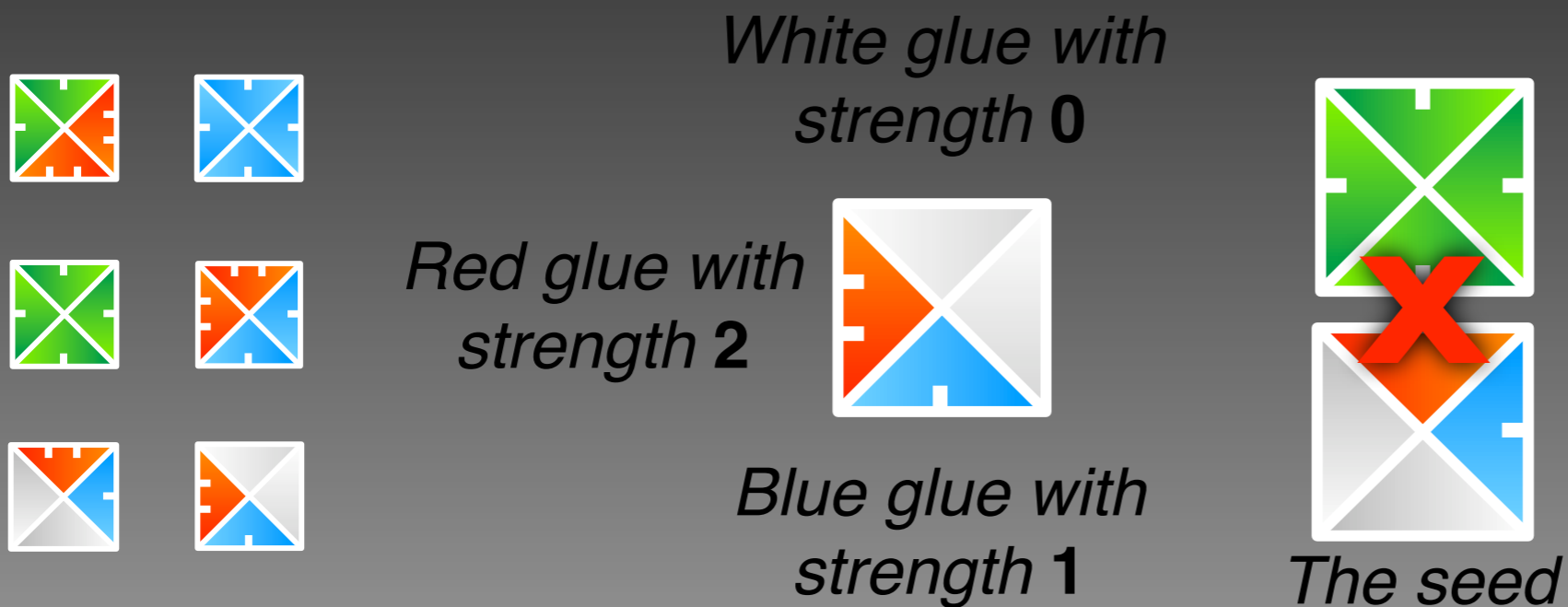
*Blue glue with
strength 1*



The seed

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



*White glue with
strength 0*

*Red glue with
strength 2*



*Blue glue with
strength 1*



The seed

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model

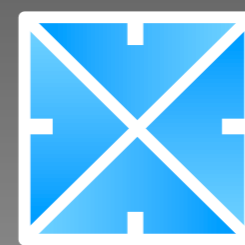


*Red glue with
strength 2*



*Blue glue with
strength 1*

*White glue with
strength 0*



The seed

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



White glue with strength 0



Red glue with strength 2



Blue glue with strength 1



The seed

#bonds $<$ $T^\circ = 2$

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



*White glue with
strength 0*

*Red glue with
strength 2*



*Blue glue with
strength 1*



The seed

- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



*White glue with
strength 0*

*Red glue with
strength 2*



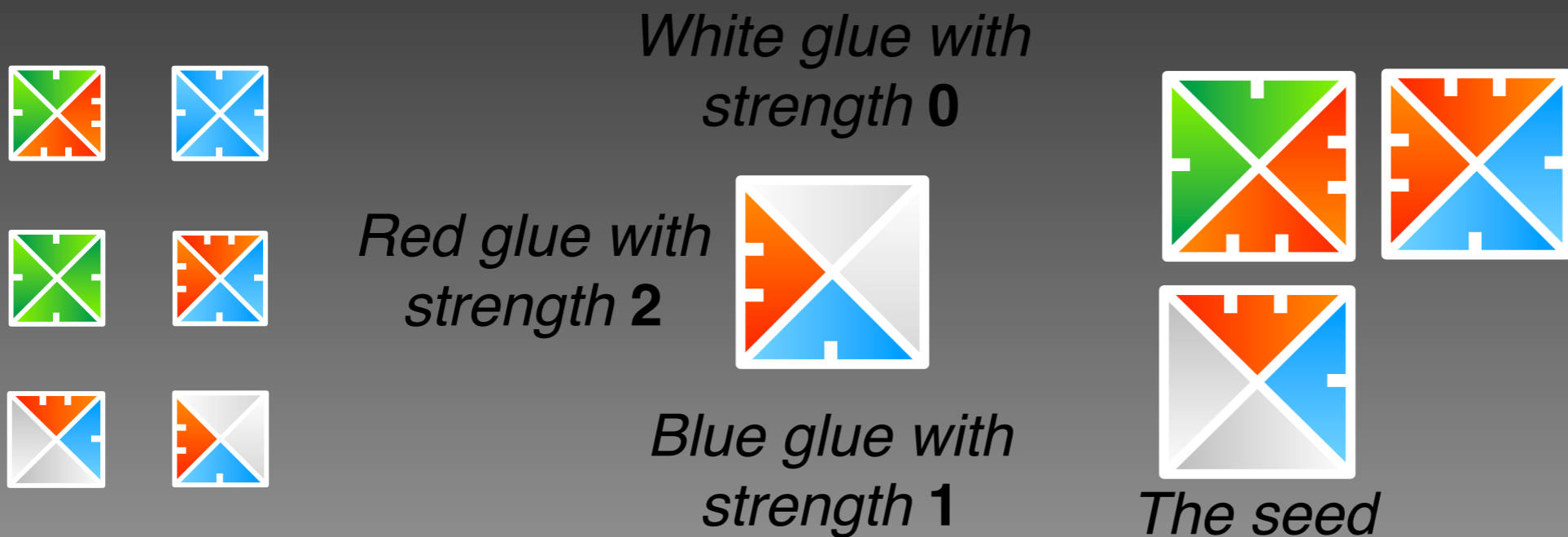
*Blue glue with
strength 1*



The seed

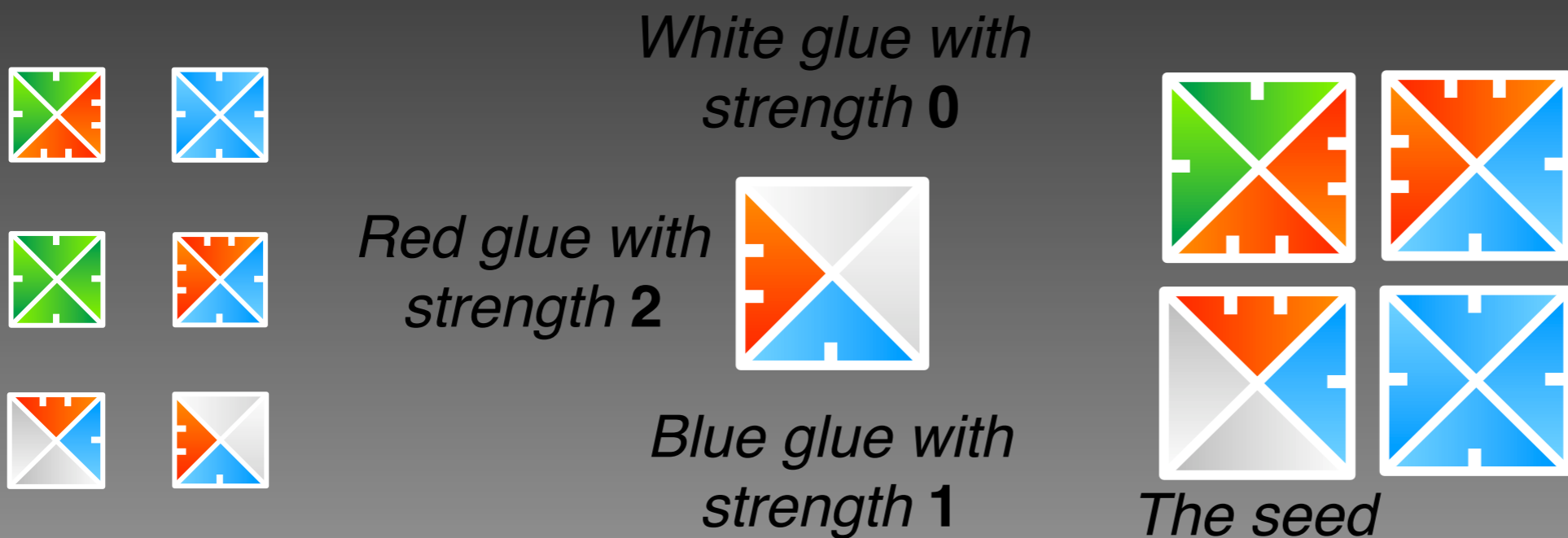
- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



- A seed tile
- A temperature (= **2** in practice)

Algorithmic model



- A seed tile
- A temperature (= **2** in practice)

Assembling a square

The tiles
at $T^{\circ} = 2$



The seed

Assembling a square

The tiles
at $T^{\circ} = 2$



The seed

Assembling a square

The tiles
at $T^{\circ} = 2$



The seed

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at $T^{\circ} = 2$



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at $T^{\circ} = 2$



The seed

Assembling a square

The tiles
at $T^{\circ} = 2$



The seed

Assembling a square

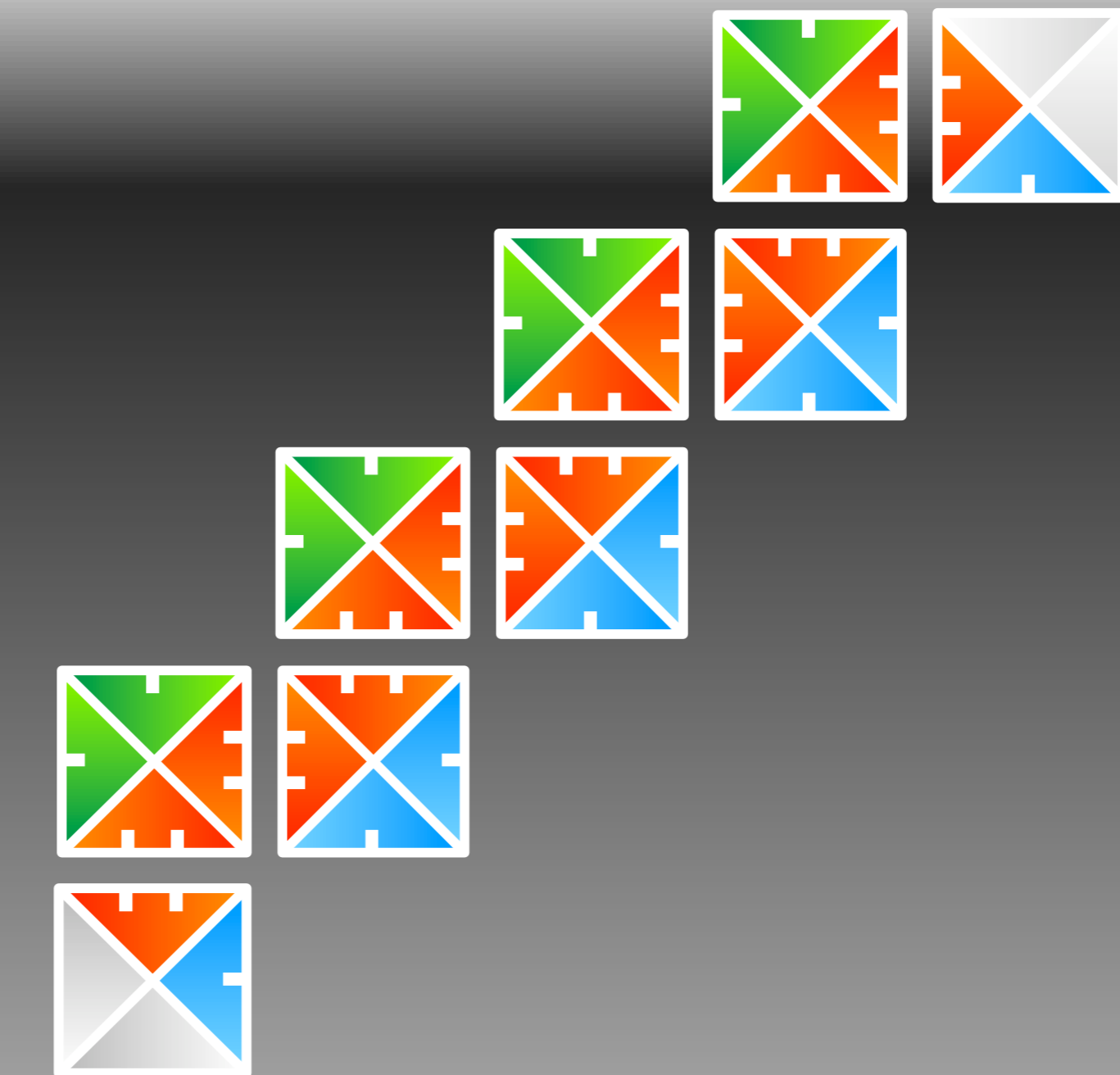
The tiles
at $T^{\circ} = 2$



The seed

Assembling a square

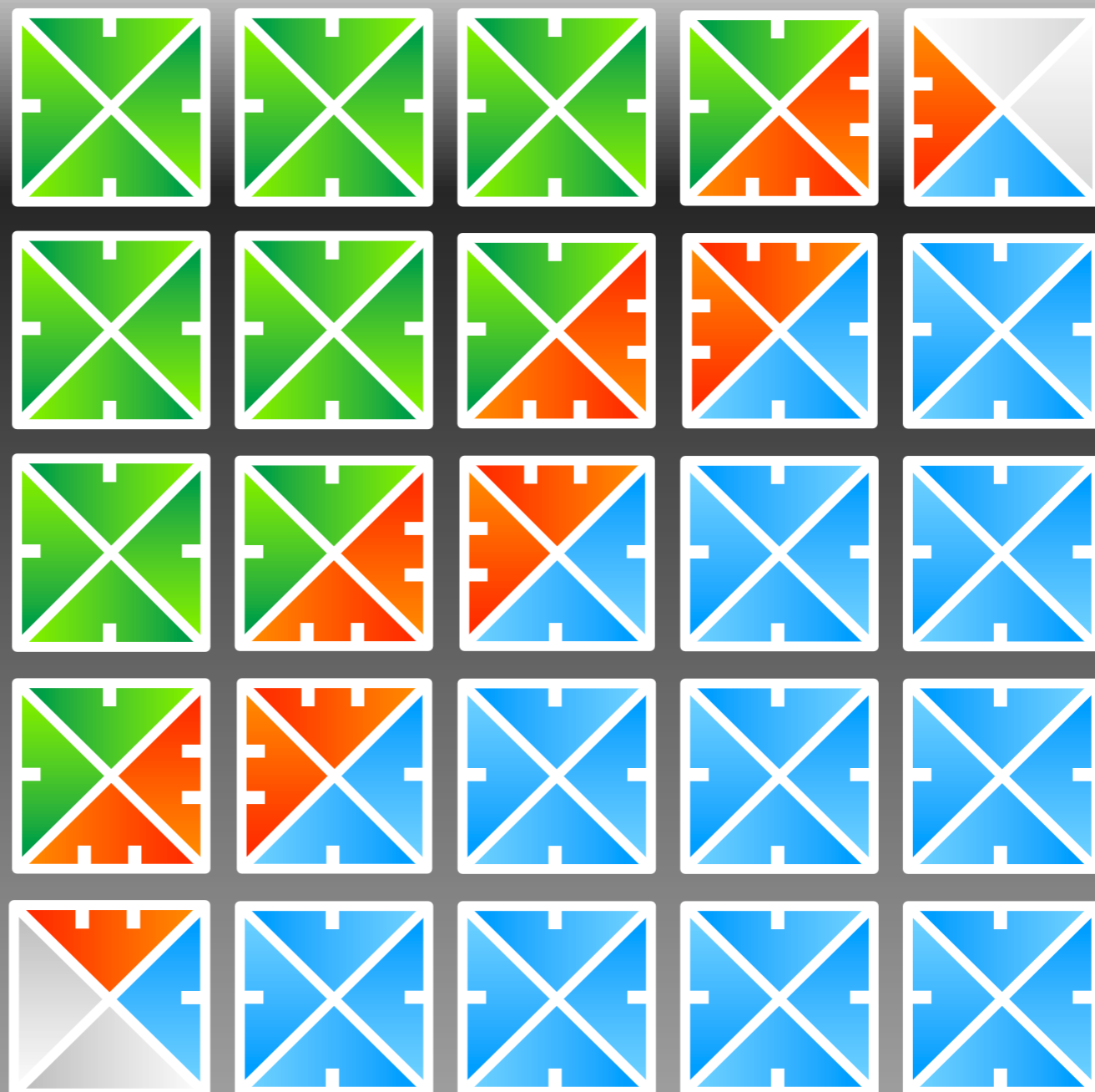
The tiles
at $T^{\circ} = 2$



The seed

Assembling a square

The tiles
at $T^{\circ} = 2$

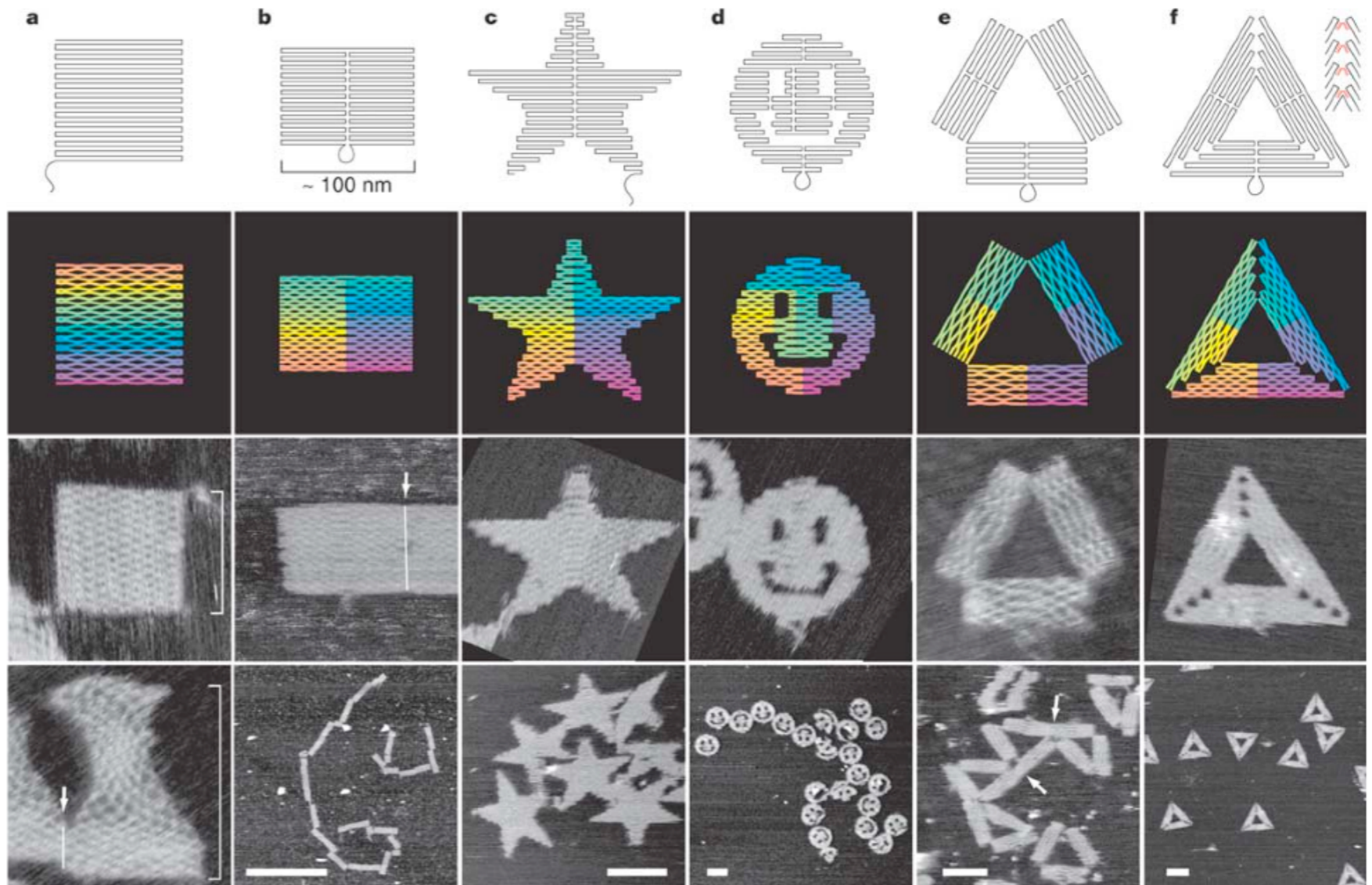


The seed

No more tile can be added

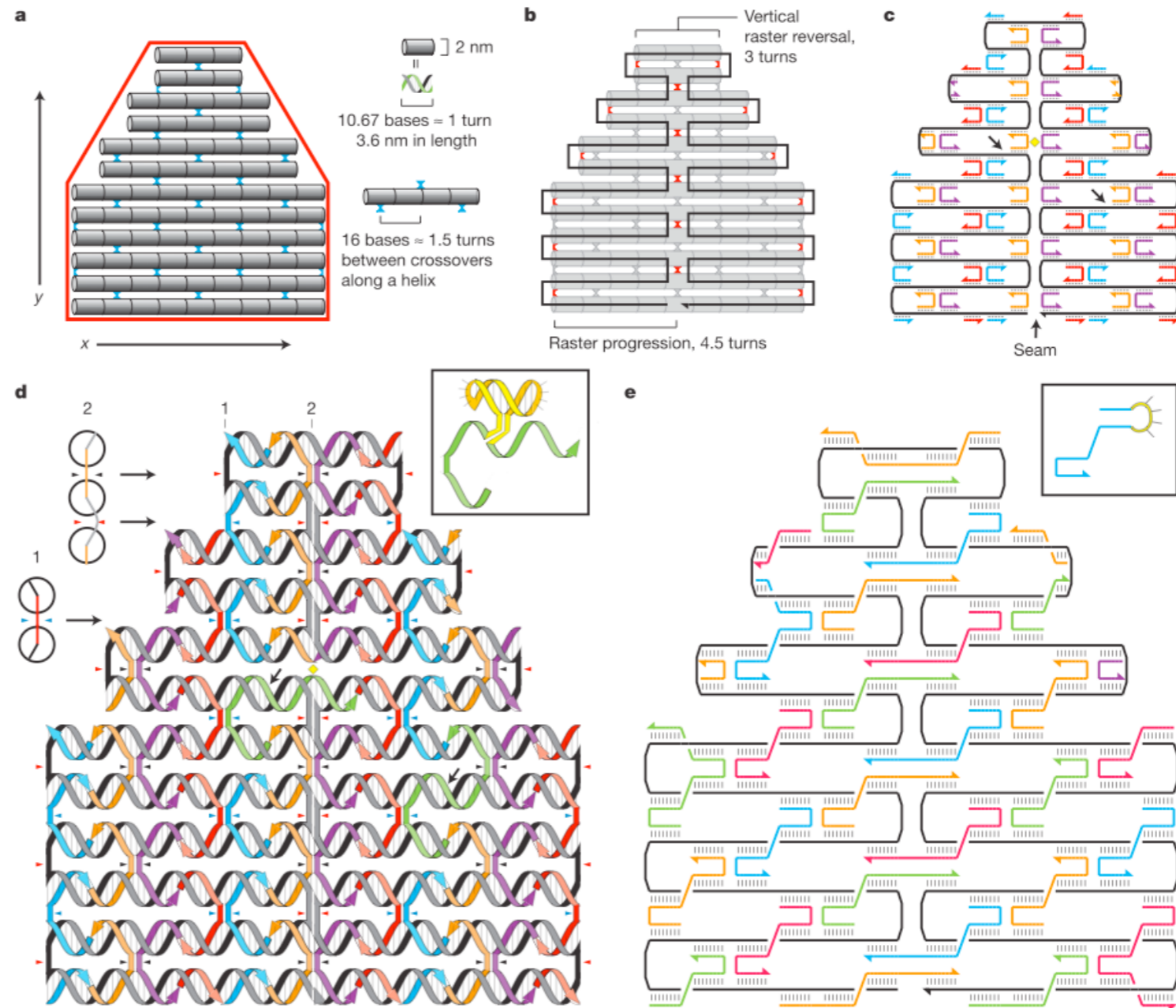
Building the Seed

Paul **Rothemund** (2001-): DNA Origami



Building the Seed

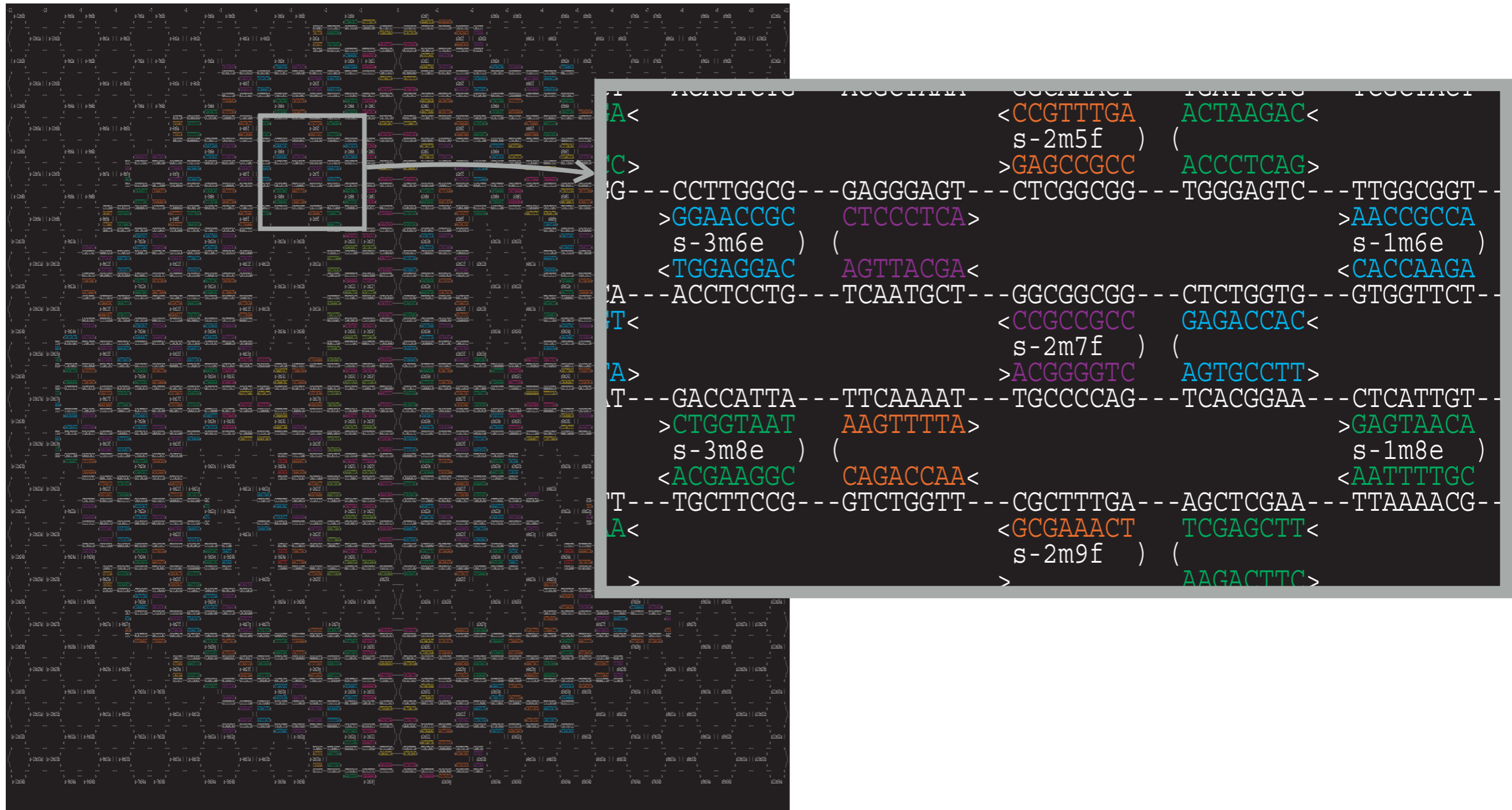
Paul **Rothemund** (2001-): DNA Origami



The DNA of a virus is folded by stapples!

Building the Seed

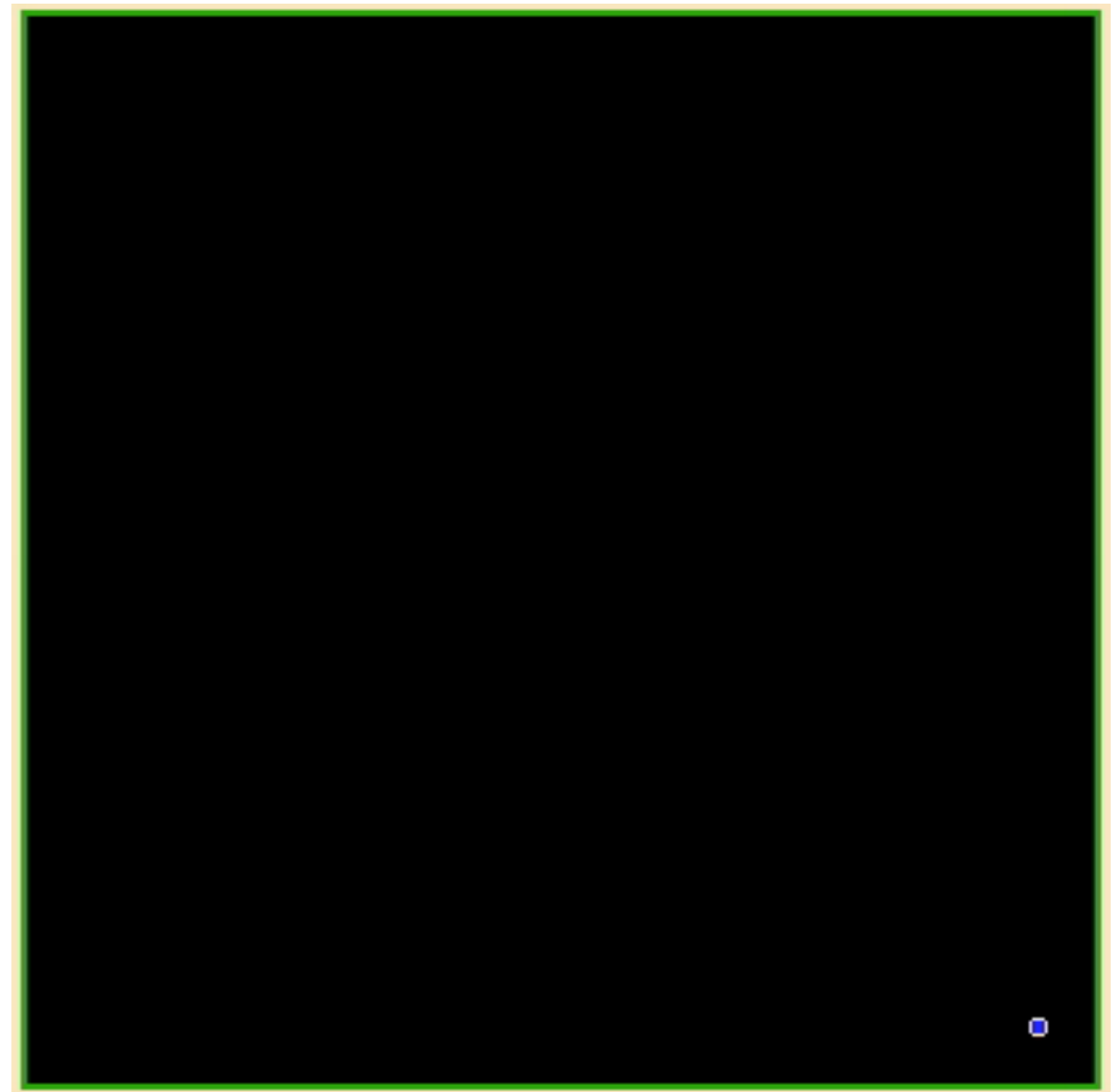
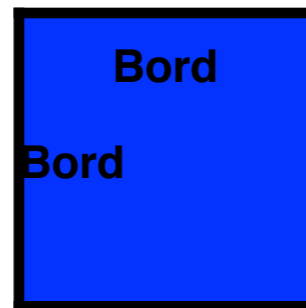
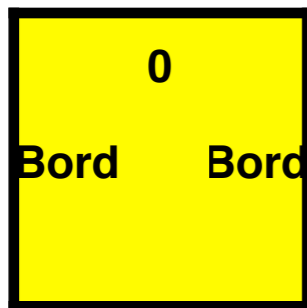
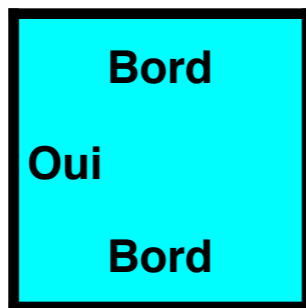
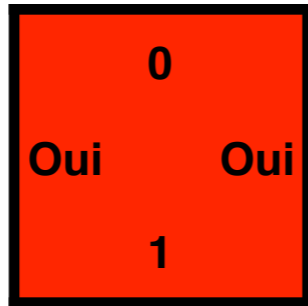
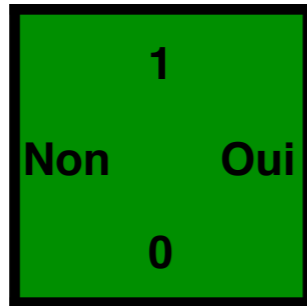
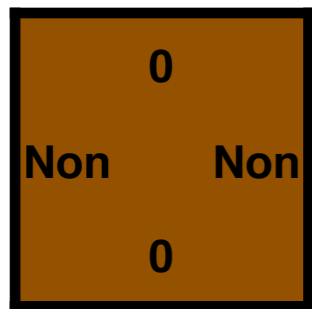
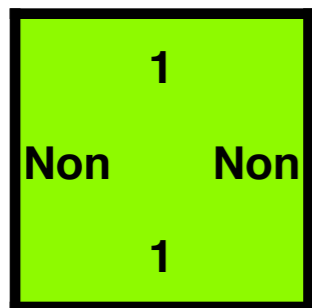
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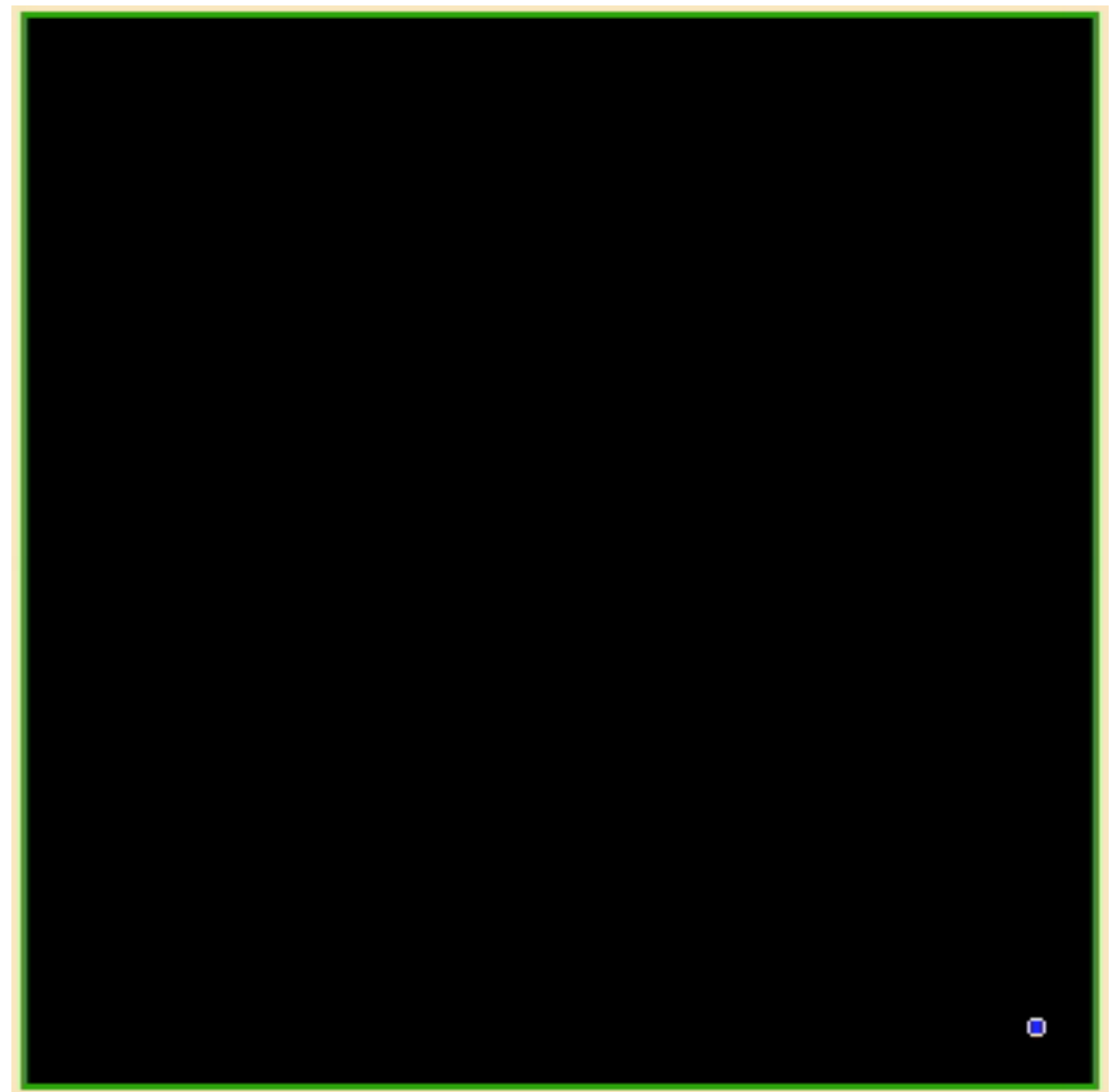
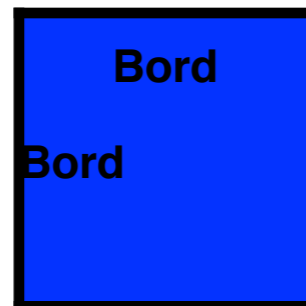
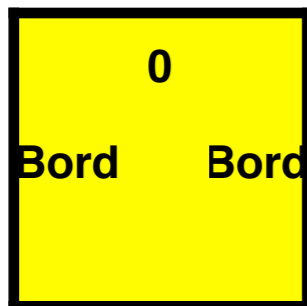
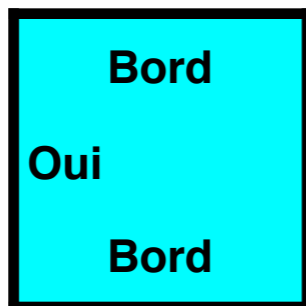
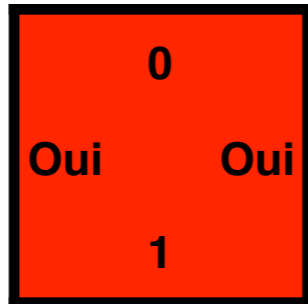
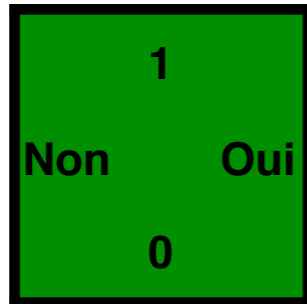
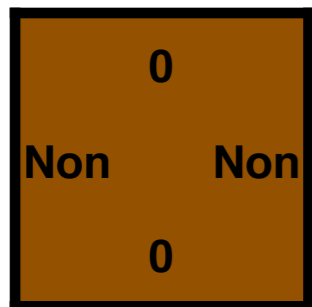
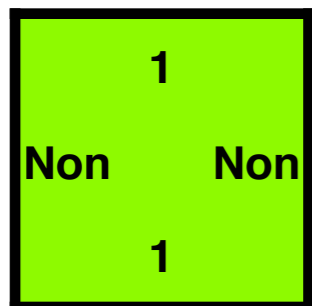
How to control the growth ?



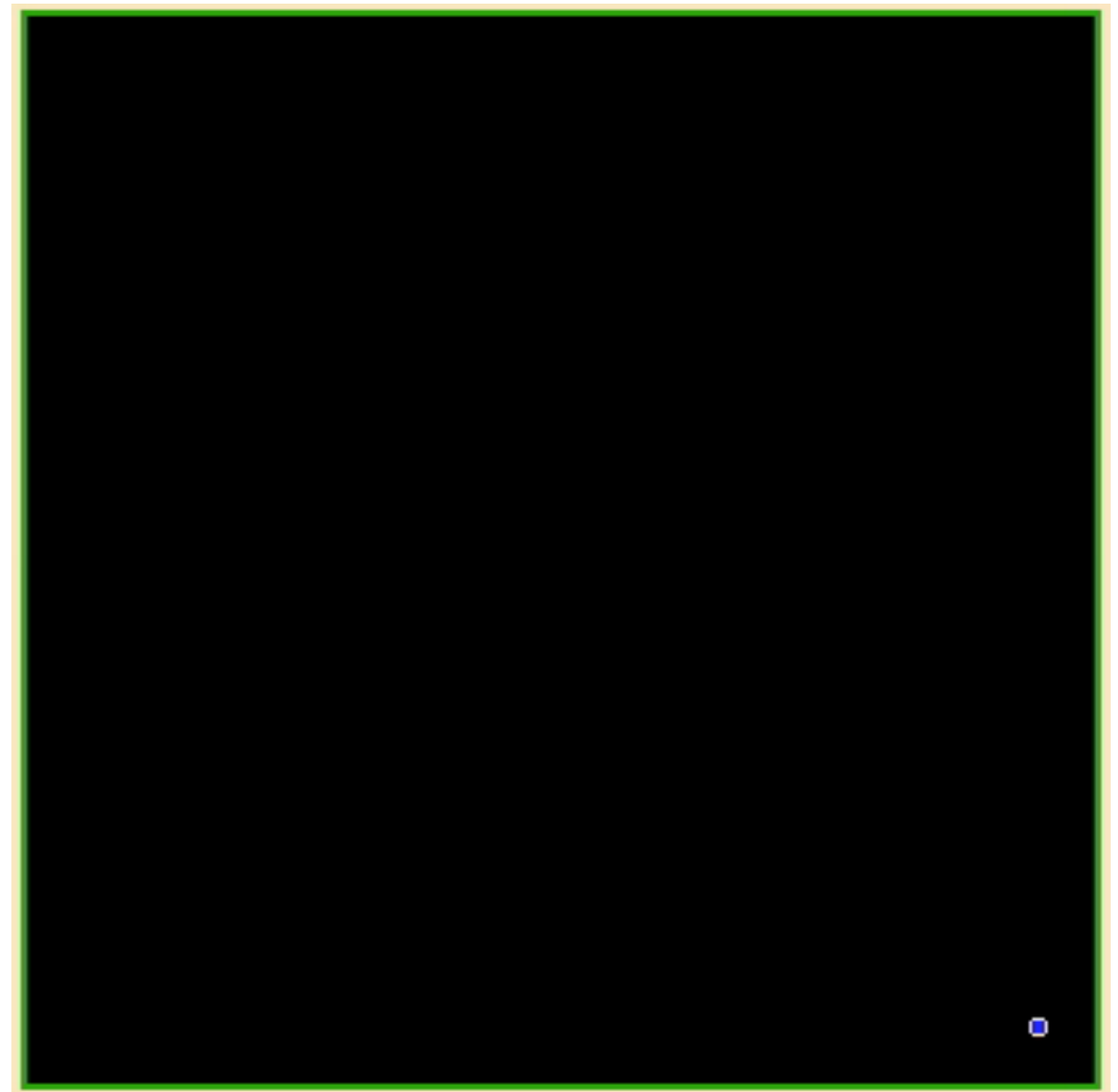
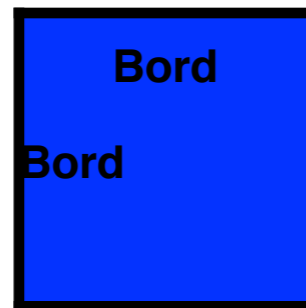
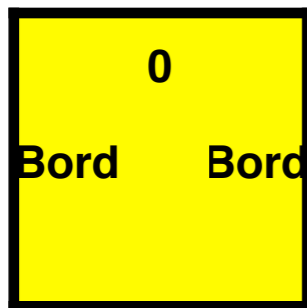
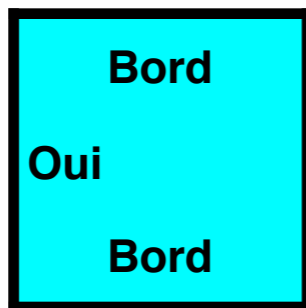
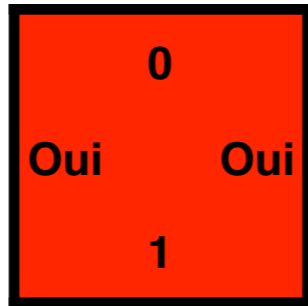
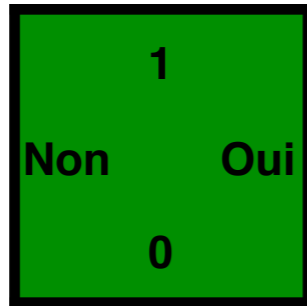
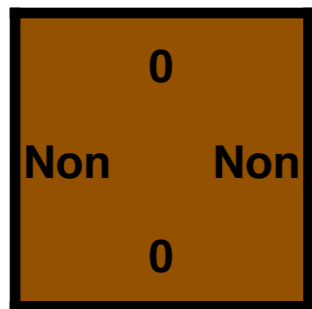
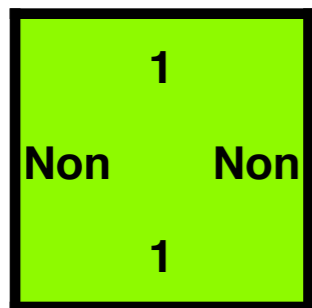
With a binary counter!



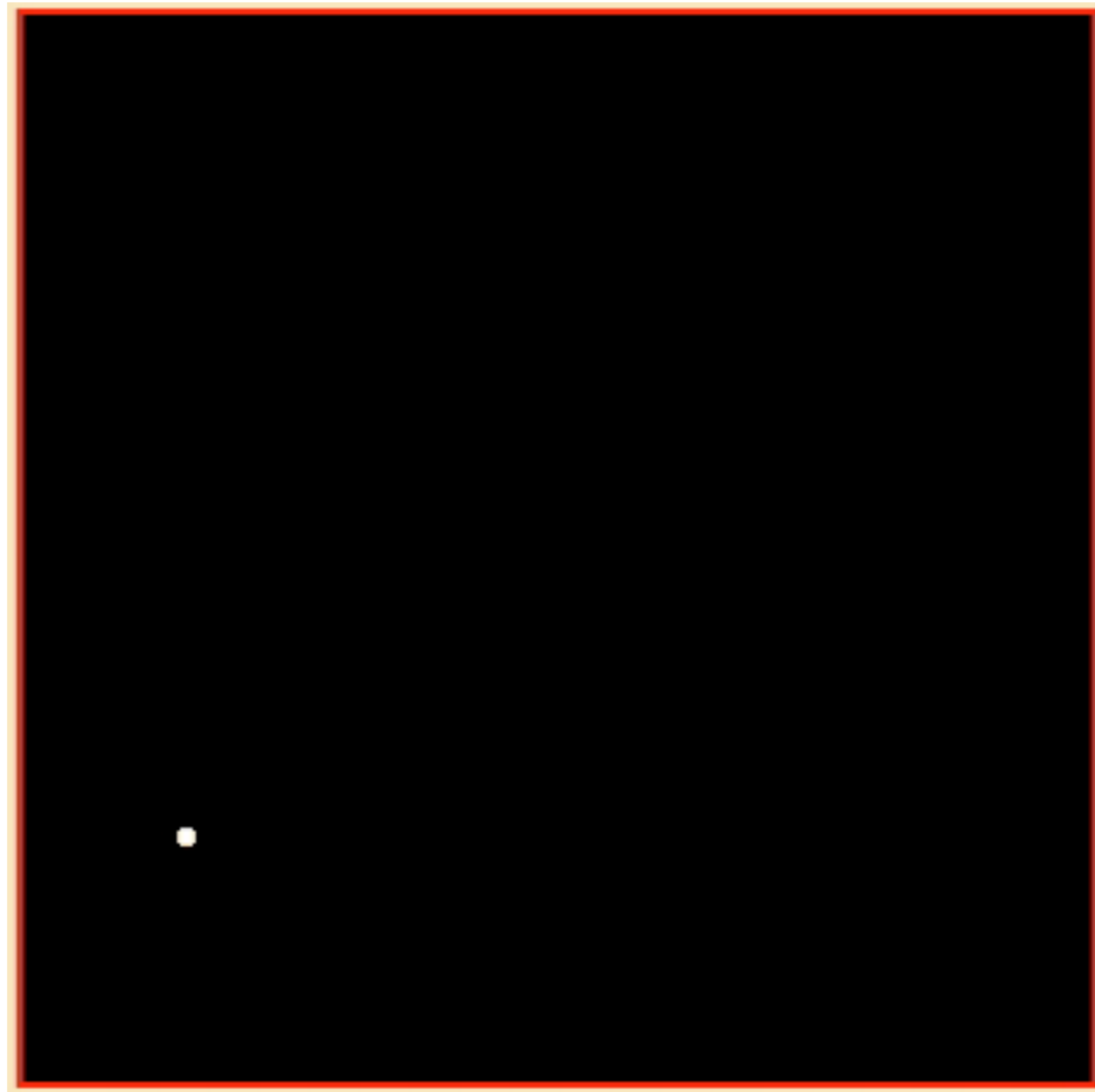
With a binary counter!



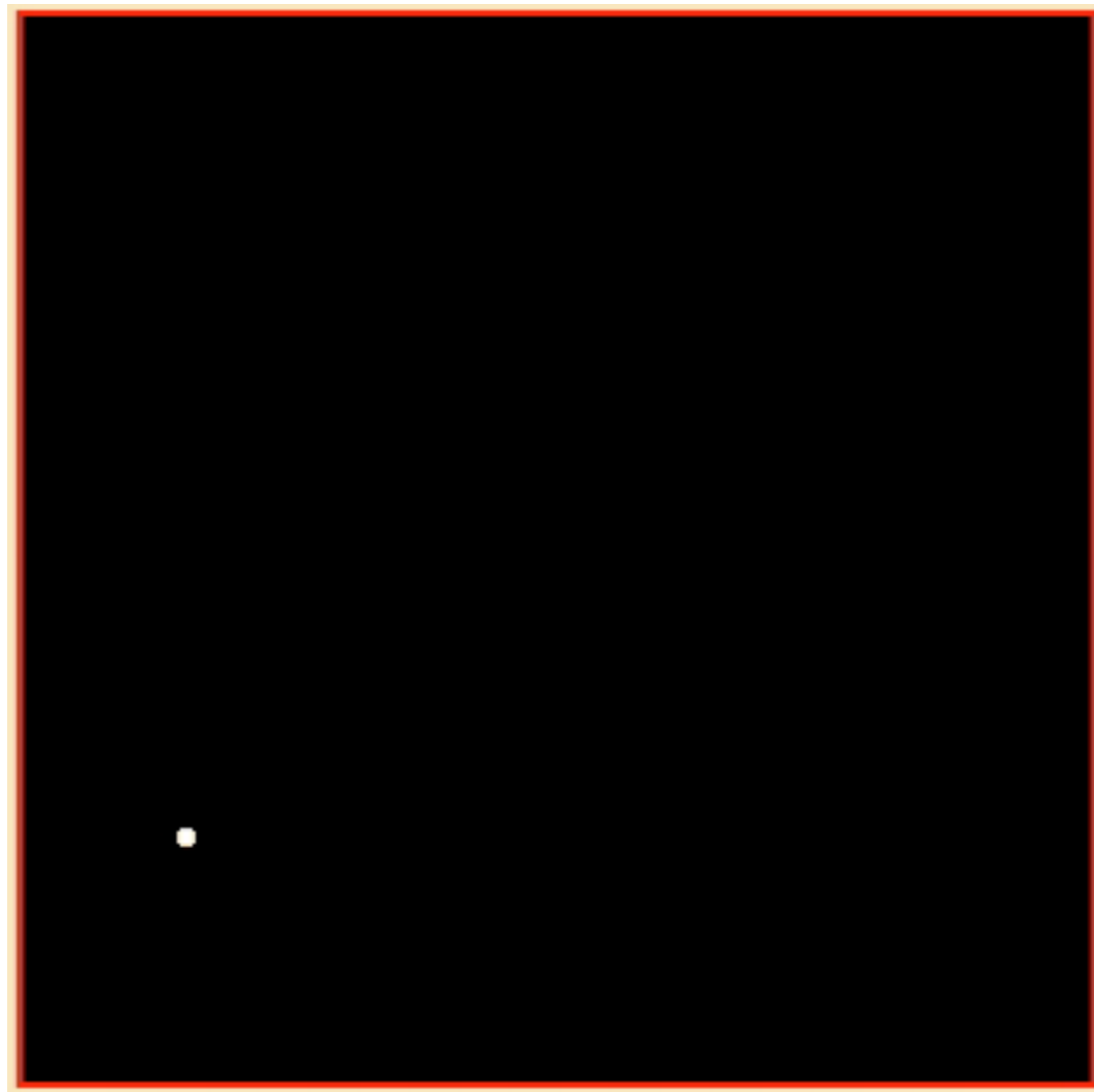
With a binary counter!



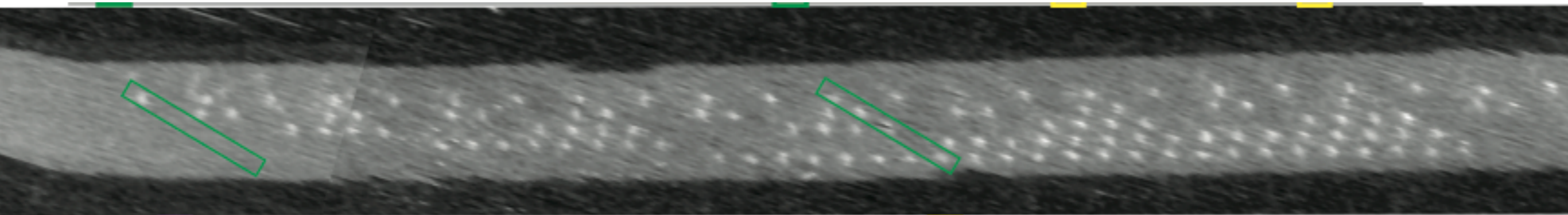
Stopping the growth of a square



Stopping the growth of a square

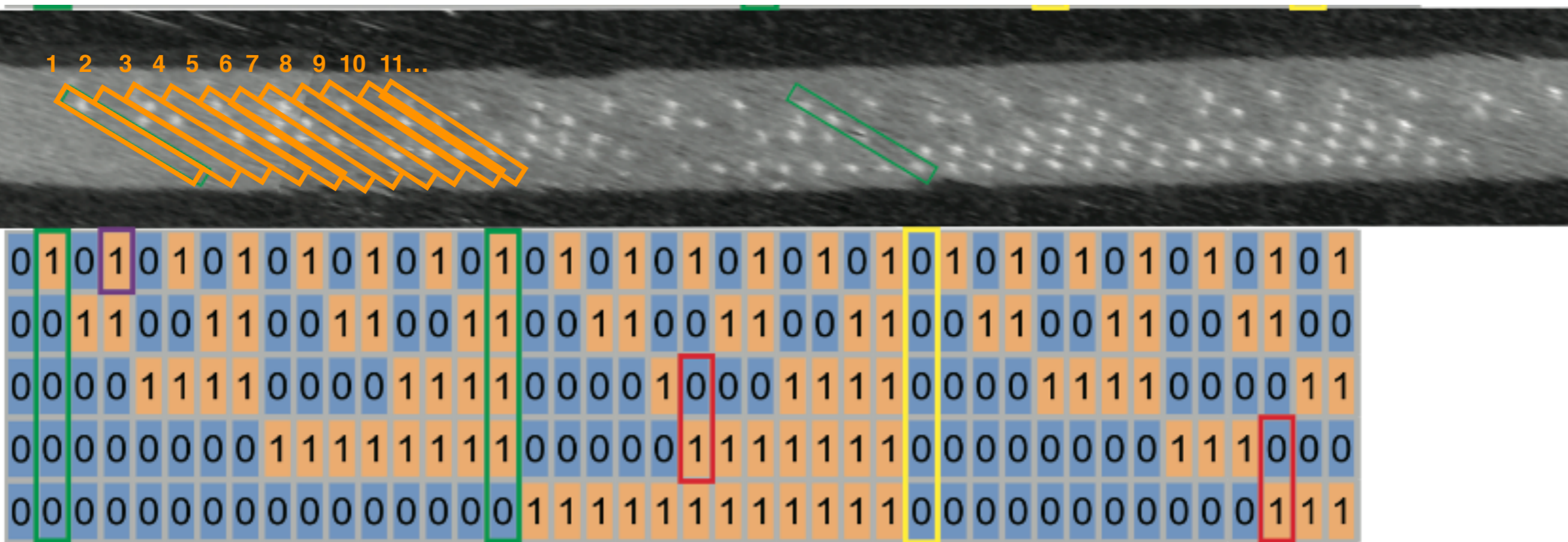


State of the art in experiments



Constantine Evans, PhD Thesis, Caltech 2014

State of the art in experiments



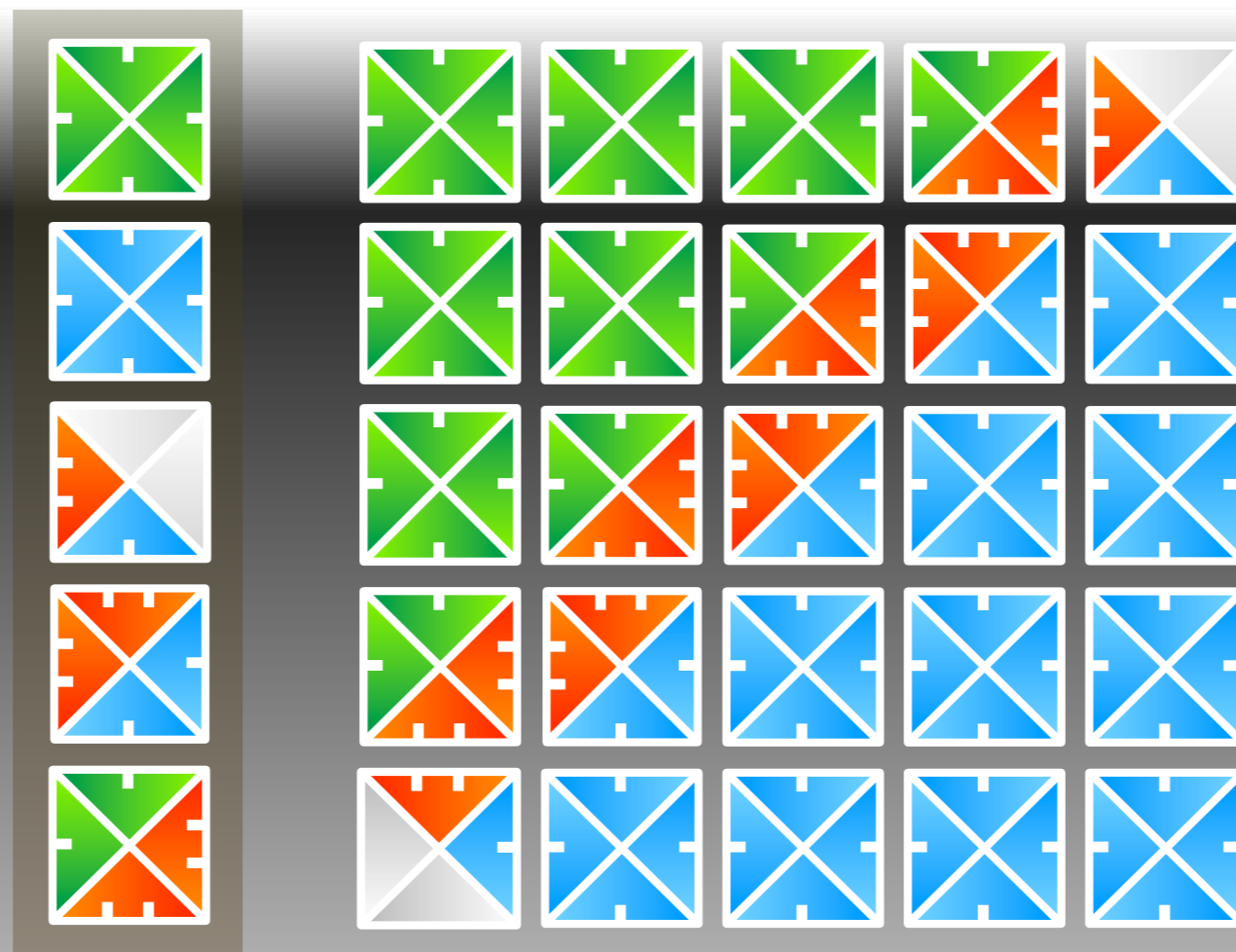
Constantine Evans, PhD Thesis, Caltech 2014

Algorithmic questions:

- Minimize the number of tiles?
- Minimize assembly time?
- Is there a universal set of tiles? (i.e. a programming language)

Minimize the number of tiles

- **Undecidable** problem in general
- For the square, **5 tiles are enough and required:**



*Becker,
Rapaport,
Rémila, 2006*

Minimizing Assembly time

- Undécidable problem in general
- Squares can be assembled in optimal time: $2n-2$

An example: *Assembling a square*

the tiles - Temperature = 2



the seed

An example: *Assembling a square*

the tiles - Temperature = 2



the seed

An example: *Assembling a square*

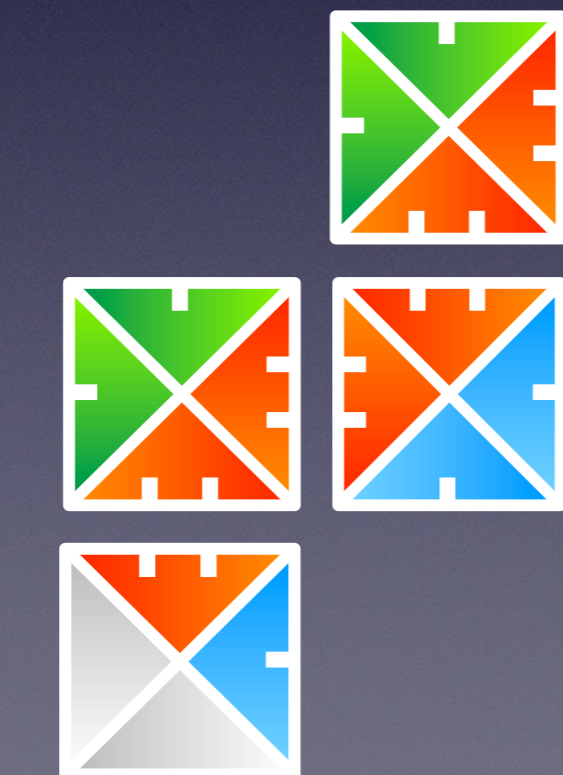
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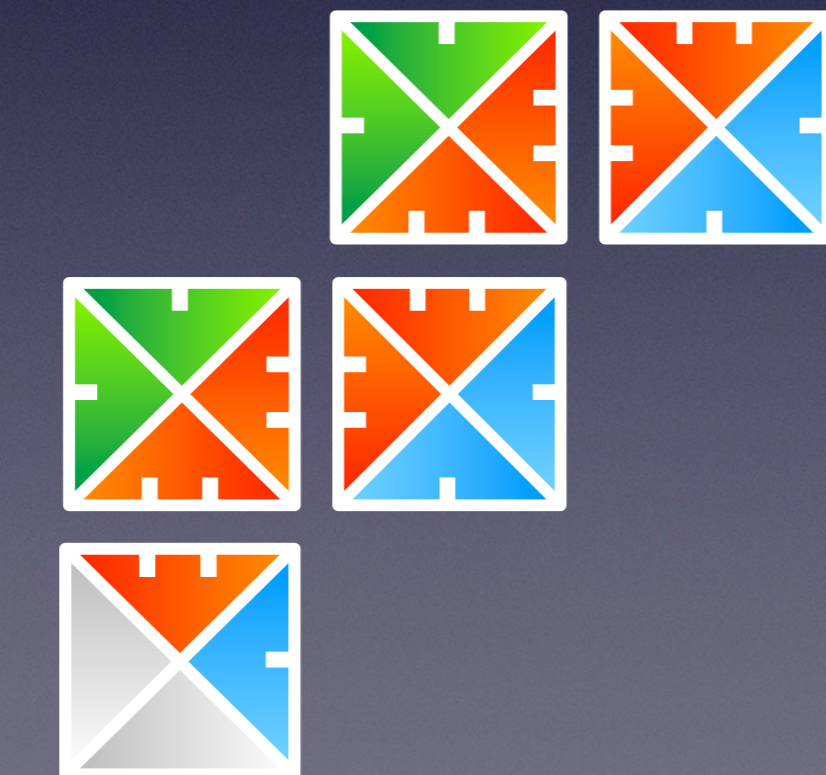
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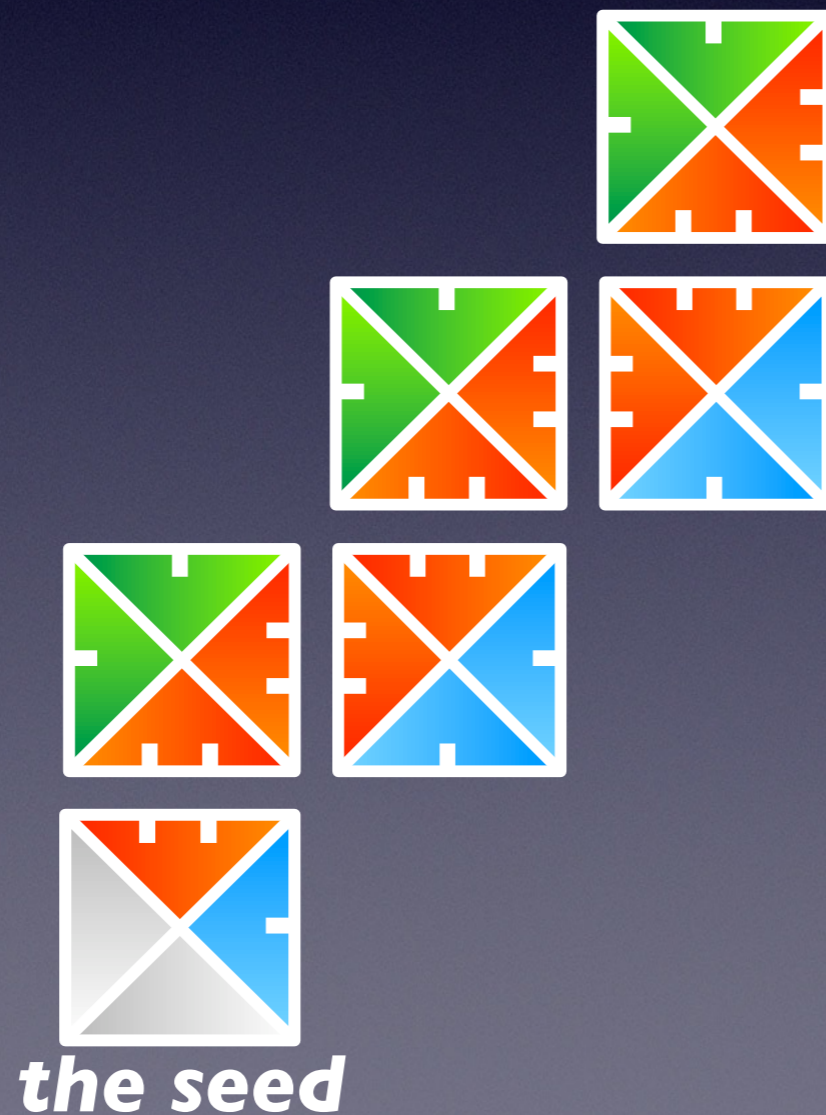
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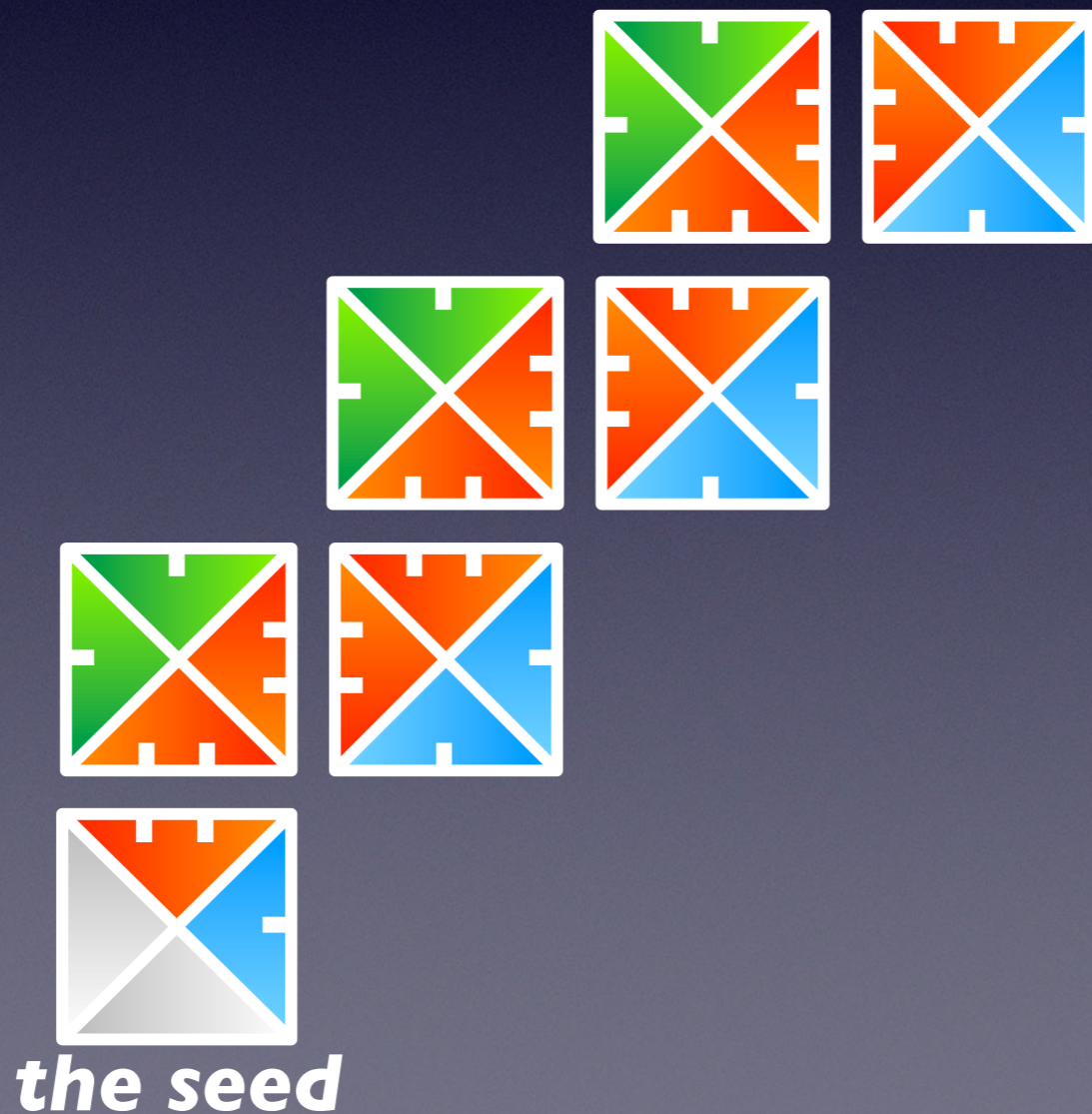
the tiles - Temperature = 2



the seed

An example: *Assembling a square*

the tiles - Temperature = 2



An example: *Assembling a square*

the tiles - Temperature = 2



the seed

An example: *Assembling a square*

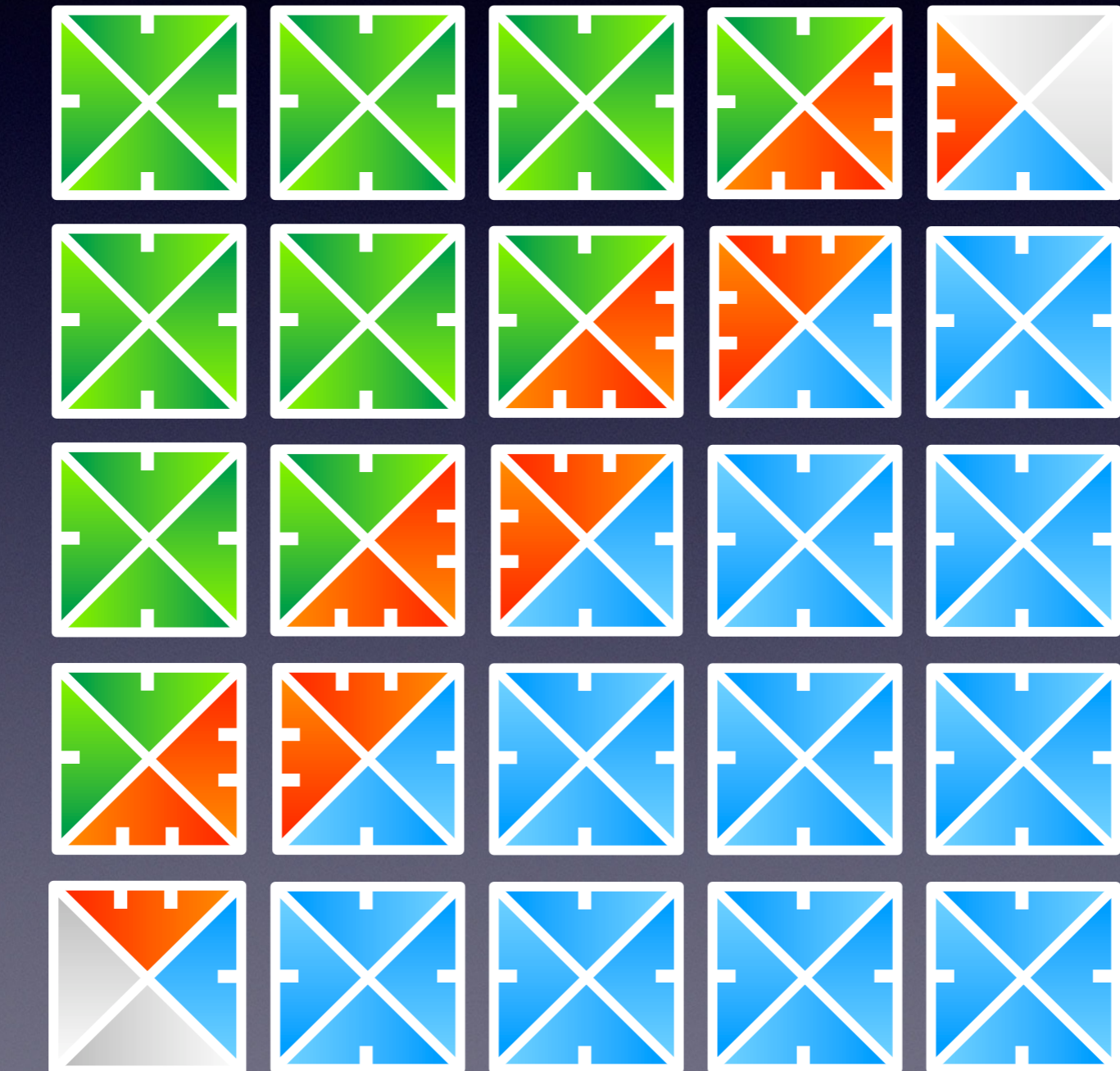
the tiles - Temperature = 2



the seed

An example: *Assembling a square*

the tiles - Temperature = 2

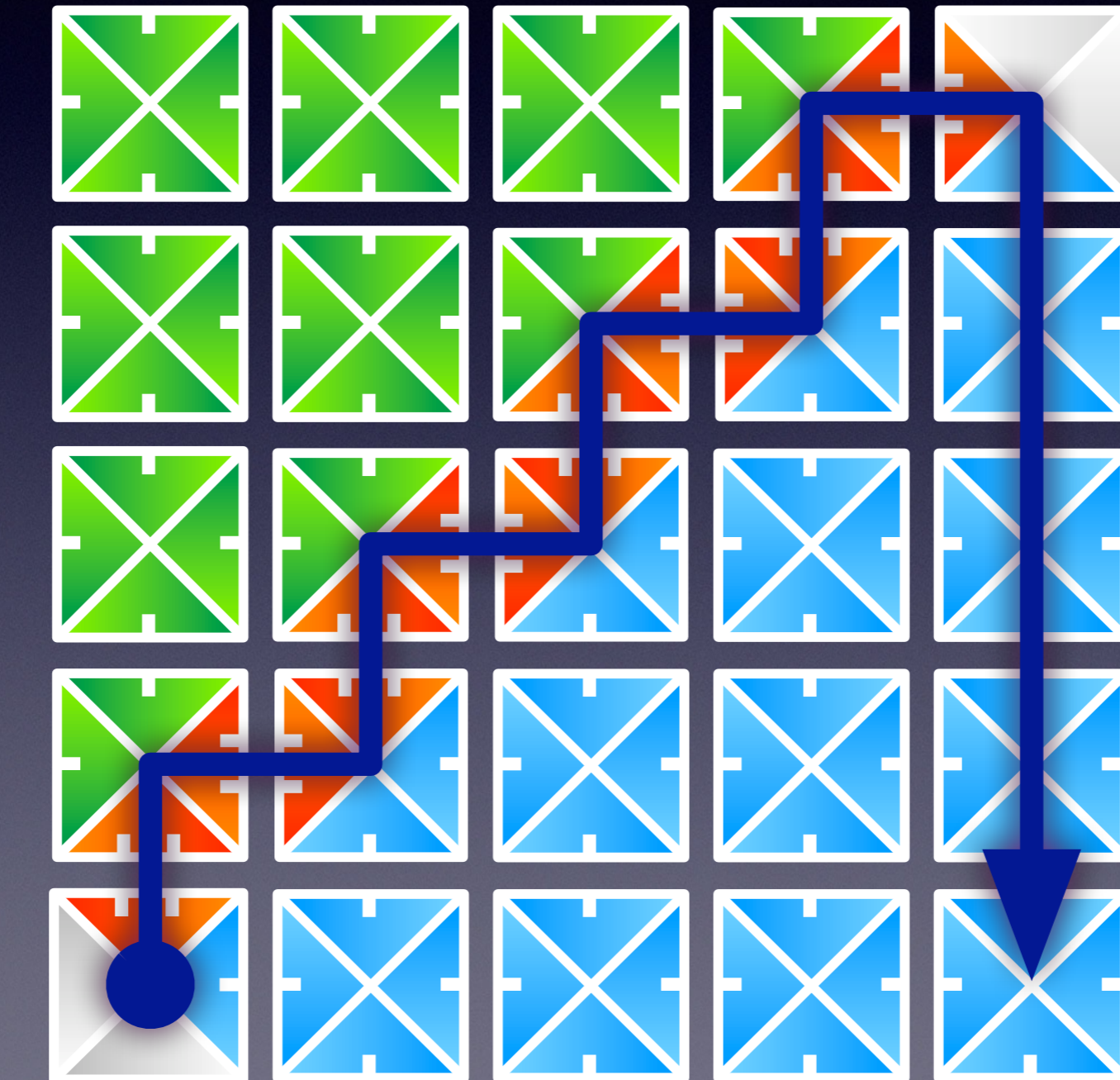


the seed

No more tiles can be attached

An example: Assembling a square

the tiles - Temperature = 2



the seed

No more tiles can be attached

Longest chain
of dependencies
=
 $3n-3$

Assembling vs Tiling & Cellular Automata

Facts

Irrevocability.

As opposed to CA, the “state/tile” of a cell cannot change

Everything that can bind, may bind so be careful with signals self-ignition.

Guaranteeing an order.

Being sure of the predecessors of each cell to guarantee the global behavior

Only one “main signal” per coordinate.

Otherwise it is not possible to guarantee the predecessors

Some consequences

Filling tiles carry information and are not interchangeable.

As opposed to quiescent states in CA or “blank tiles” in tilings

There exists flows of information within the shape.

Signals cannot go against the flows but still have to intersect predictably

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Everything that can bind, may bind so be careful with signals self-ignition.

Guaranteeing an order.

Being sure of the predecessors of each cell to guarantee the global behavior

Only one “main signal” per coordinate.

Otherwise it is not possible to guarantee the predecessors

Some consequences

Filling tiles carry information and are not interchangeable.

As opposed to quiescent states in CA or “blank tiles” in tilings

There exists flows of information within the shape.

Signals cannot go against the flows but still have to intersect predictably

Assembling vs Tiling & Cellular Automata

Facts

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Ordered Tilesystem

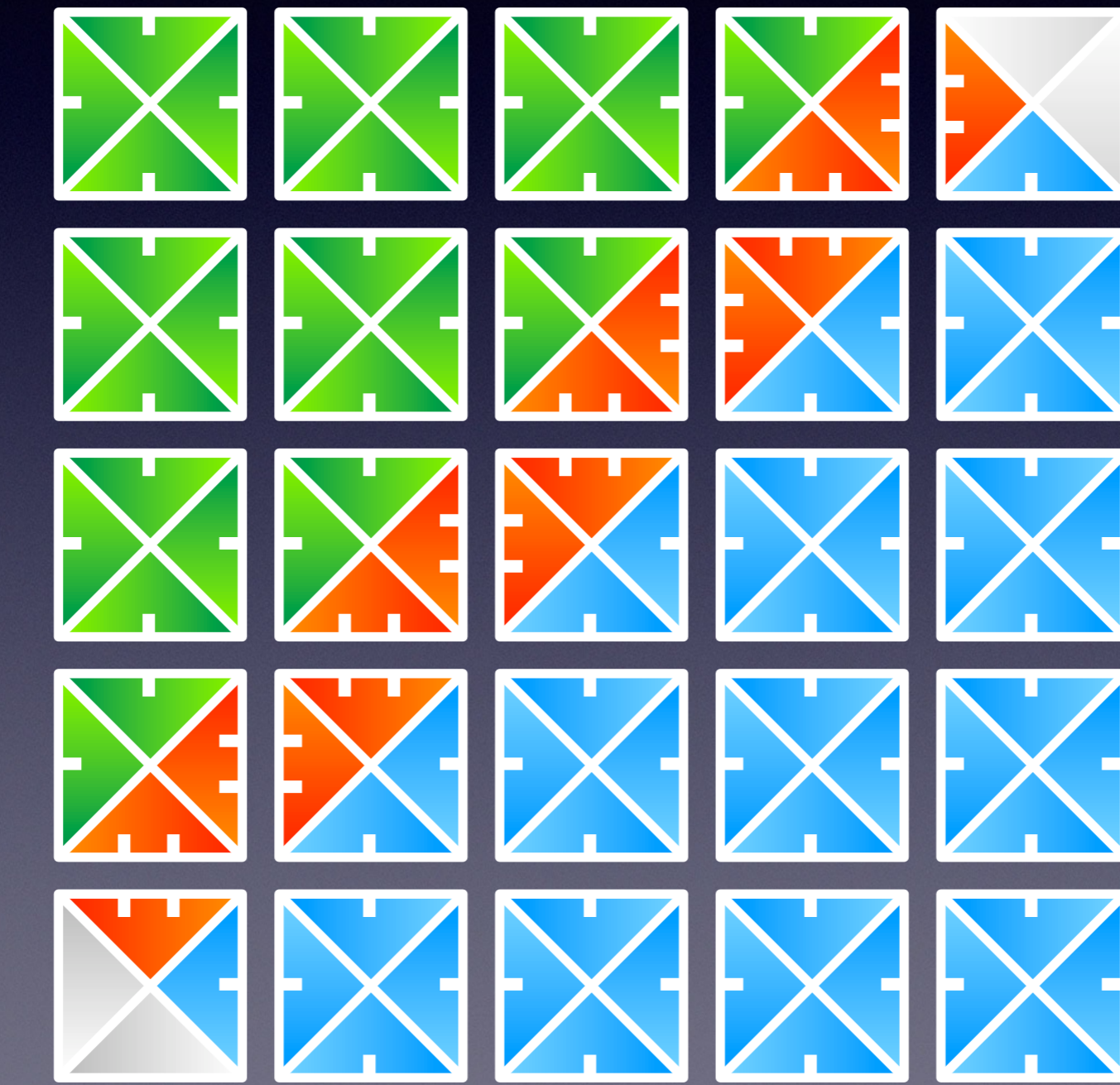
A tilesystem is **ordered** if for each production, the predecessors cells of each cell are **independent** of the construction path.

Consequences

- No one has more than T° predecessors.
- The only **non-determinism** relies in the choice of the tile to attach, which determines which shape is assembled.

Example of an ordered tile system

the tiles - Temperature = 2



the seed

Example of an ordered tile system

the tiles - Temperature = 2



the seed

Example of an ordered tile system

the tiles - Temperature = 2



the seed

Rank

Rank.

The **rank** of a site (i,j) in a given shape is the length of its longest chain of dependancies from the seed.



Time model

Poisson Markov Chain Model.

Each tile appears at each unoccupied site according to some Poisson process at a **rate** proportional to its **concentration**.

Only **matching** tile with **enough bonds** remains attached to the current aggregate.

Time & order

Theorem. [Adleman et al, 2001]

The expected time to build a shape P is:

$$O(c \cdot \text{rank}(P))$$

where c only depends on the concentrations and $\text{rank}(P)$ is a highest rank in the shape P .

⇒ we focus on **minimizing the highest rank**

Real time

Lower bounding the construction time.

Given that tiles are placed one next to the others, $\|P\|_1$ is a lower bound on the highest rank of a site of a shape P .

Real time construction.

A shape P is built in real time if

$$\text{the highest rank of a site} = \|P\|_1$$

For the $n \times n$ square S_n : $\|S_n\|_1 = 2n - 2$

Skeleton

understanding the flow of information

Skeleton. *The skeleton of a shape in an ordered tile system is the set of the sites with at most one predecessor.*

the **y-skeleton** in orange
opens the rows



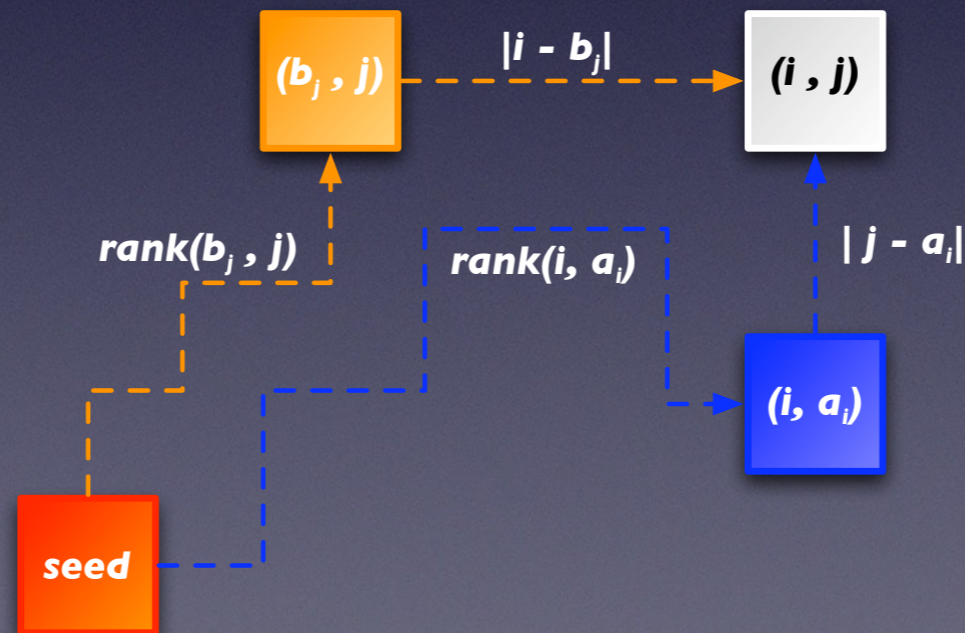
the **x-skeleton** in blue
opens the columns

Assembling a Square in Real Time

Lower bounding the rank

Let $(i, a_i)_{i \geq 0}$ and $(b_j, j)_{j \geq 0}$ be the x - and y -skeletons

Since the x - and y -skeletons sites are the **first** tiled on each column and each row respectively, then for each site (i, j) :

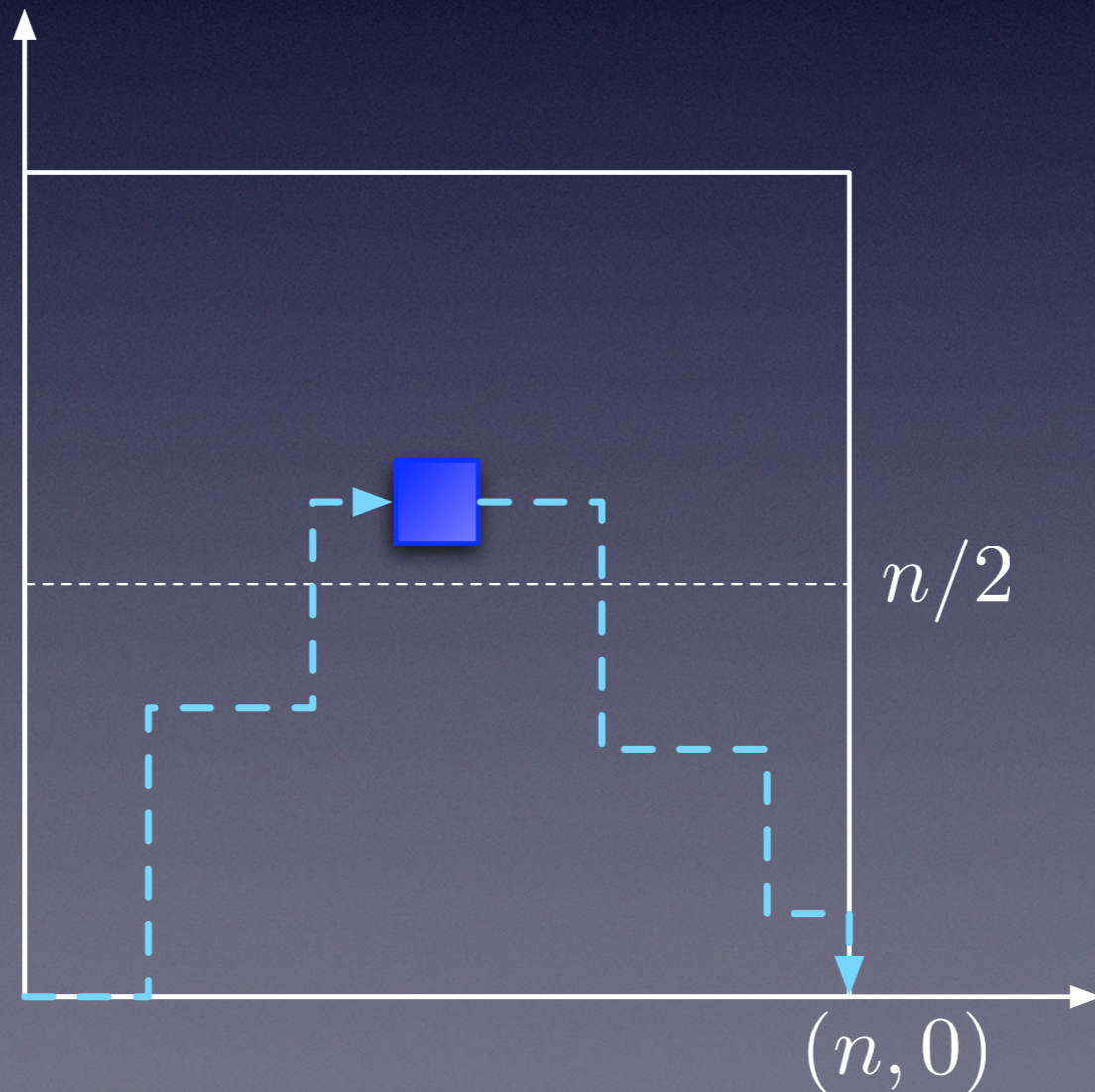


$$\text{rank}(i, j) \geq \max\{\text{rank}(i, a_i) + |j - a_i|, \text{rank}(b_j, j) + |i - b_j|\}$$

Where should be the skeleton?

the x -skeleton **cannot** go above $n/2$

the y -skeleton **cannot** go to the right of $n/2$

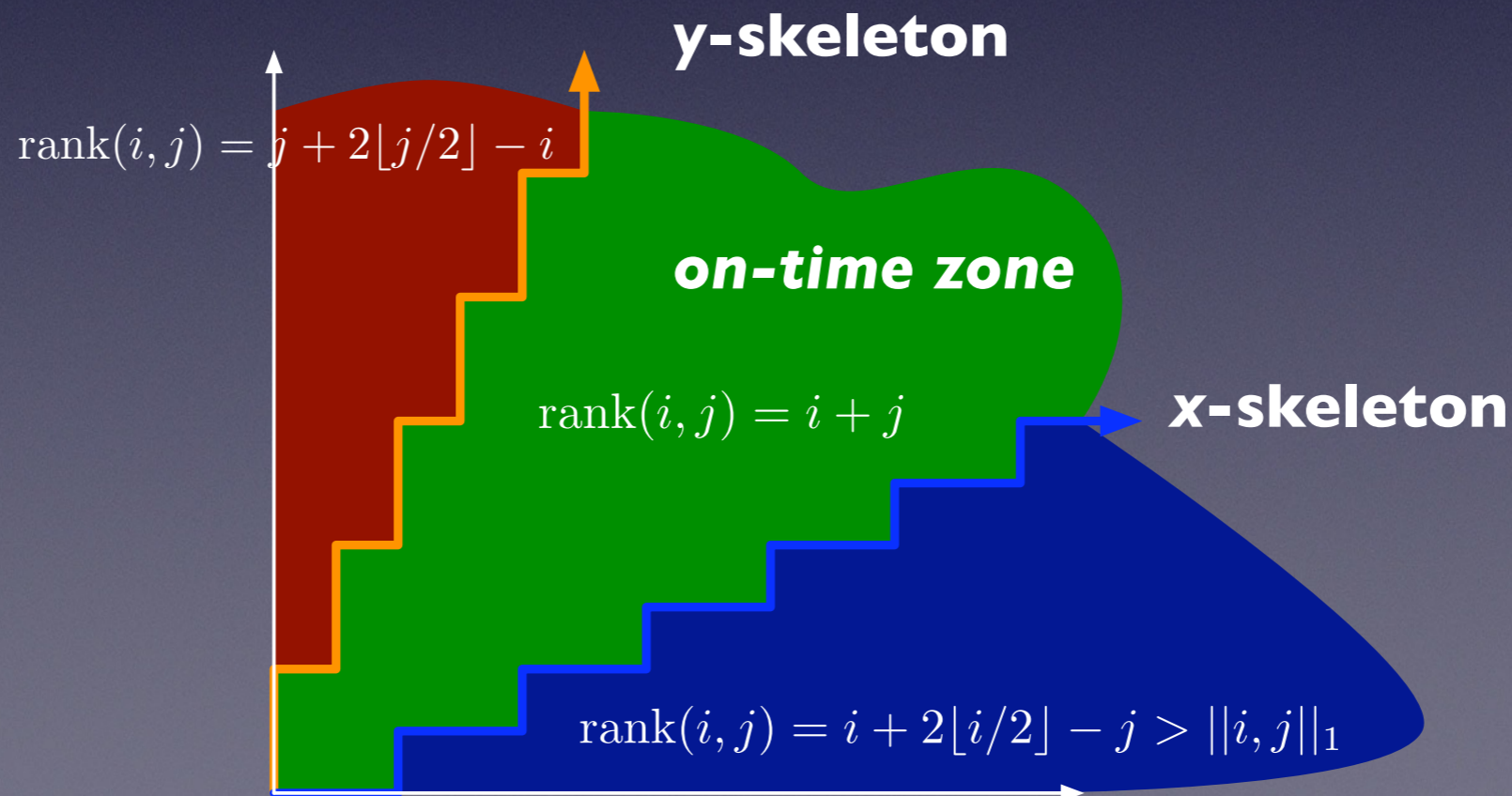


Rank function induced by the skeleton

The skeleton: $a_i = (i, \lfloor i/2 \rfloor)$ and $b_j = (\lfloor j/2 \rfloor, j)$

The rank induced:

$$\text{rank}(u) = \max\{ \|a_i\|_1 + \|u - a_i\|_1, \|b_j\|_1 + \|u - b_j\|_1 \}$$



Order induced by the skeleton

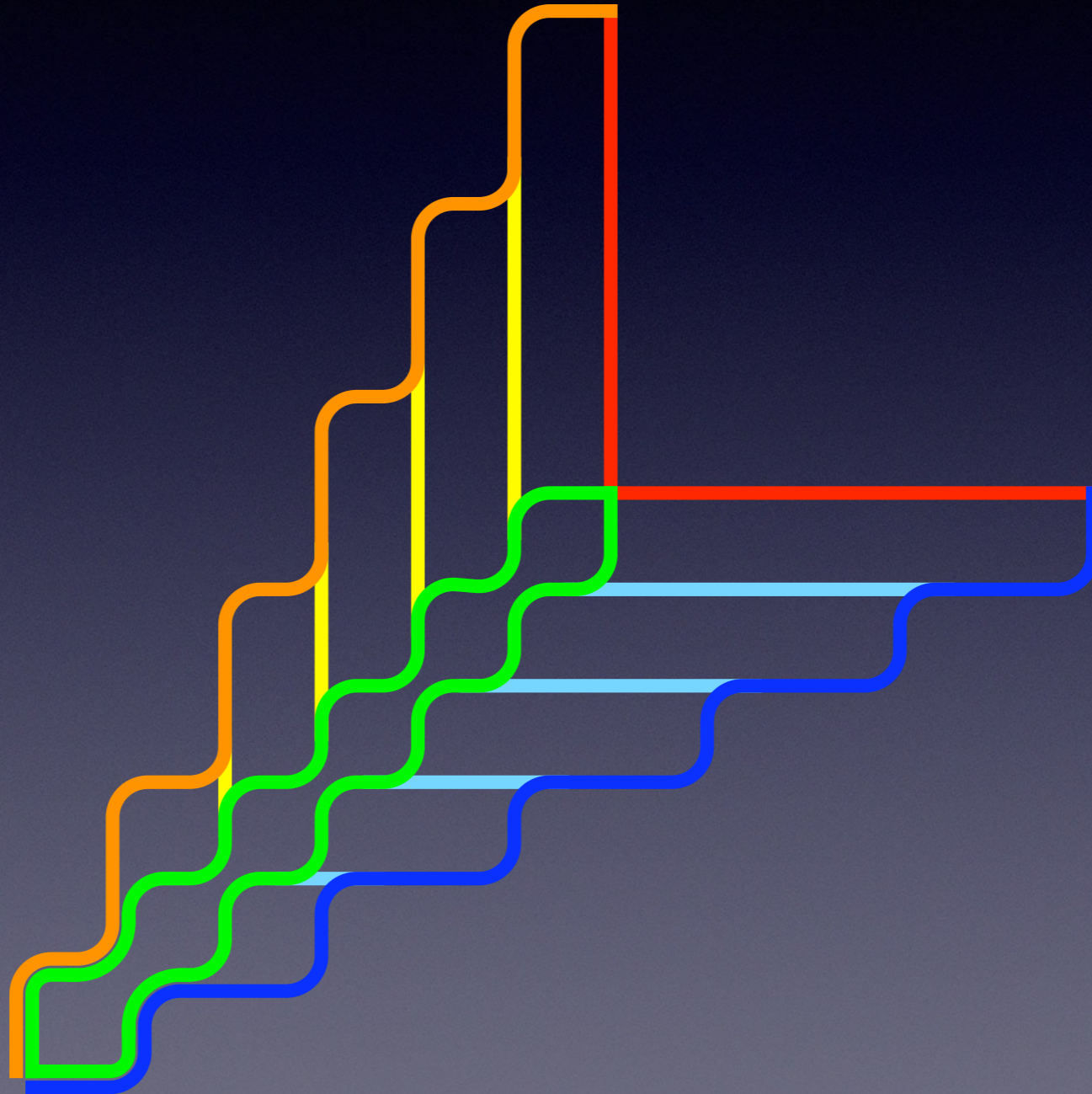


Key to construct the tileset:

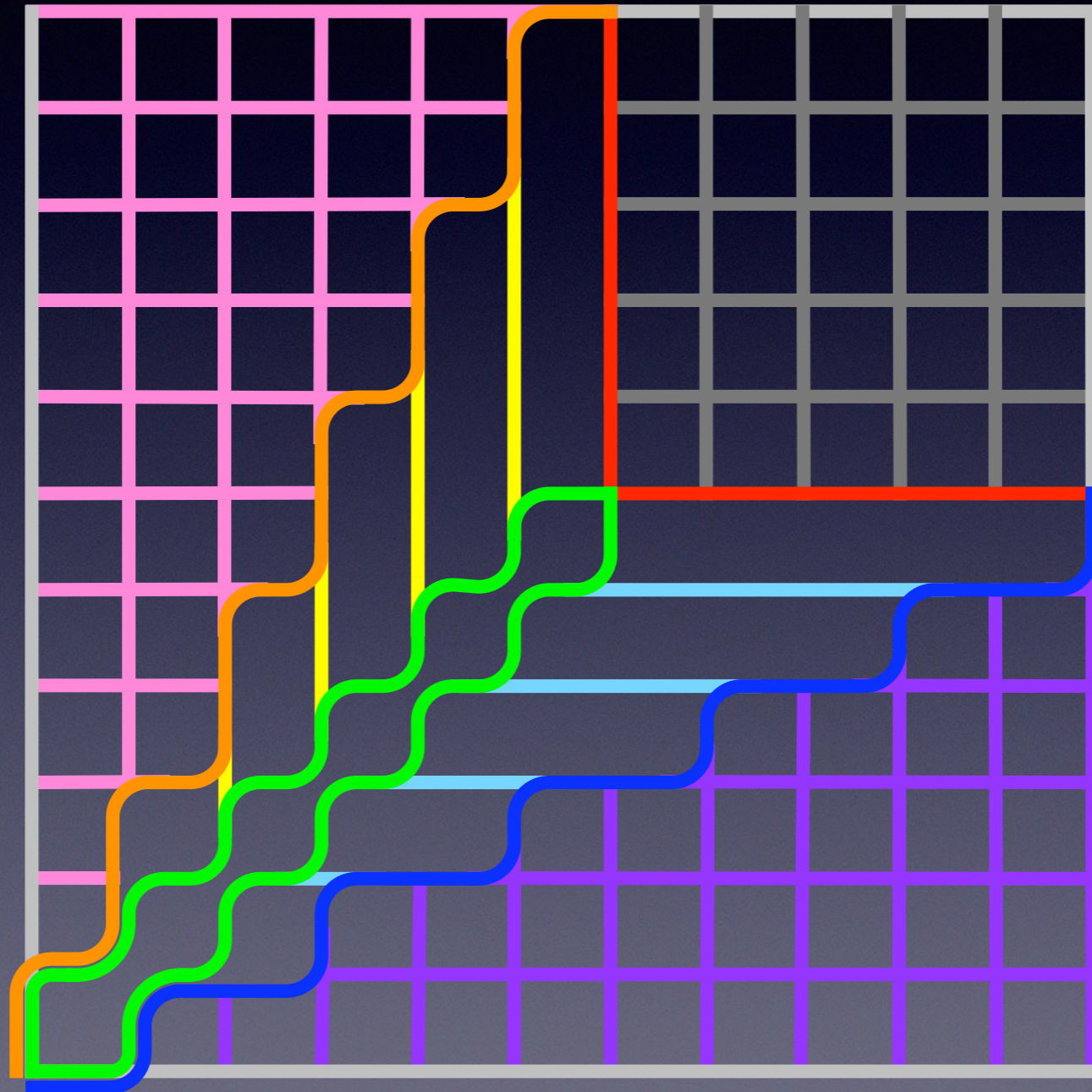
- 1) being able to guess the types of its successors from its own predecessors
- 2) synchronizing the two parts of the skeleton

The resulting tiling set

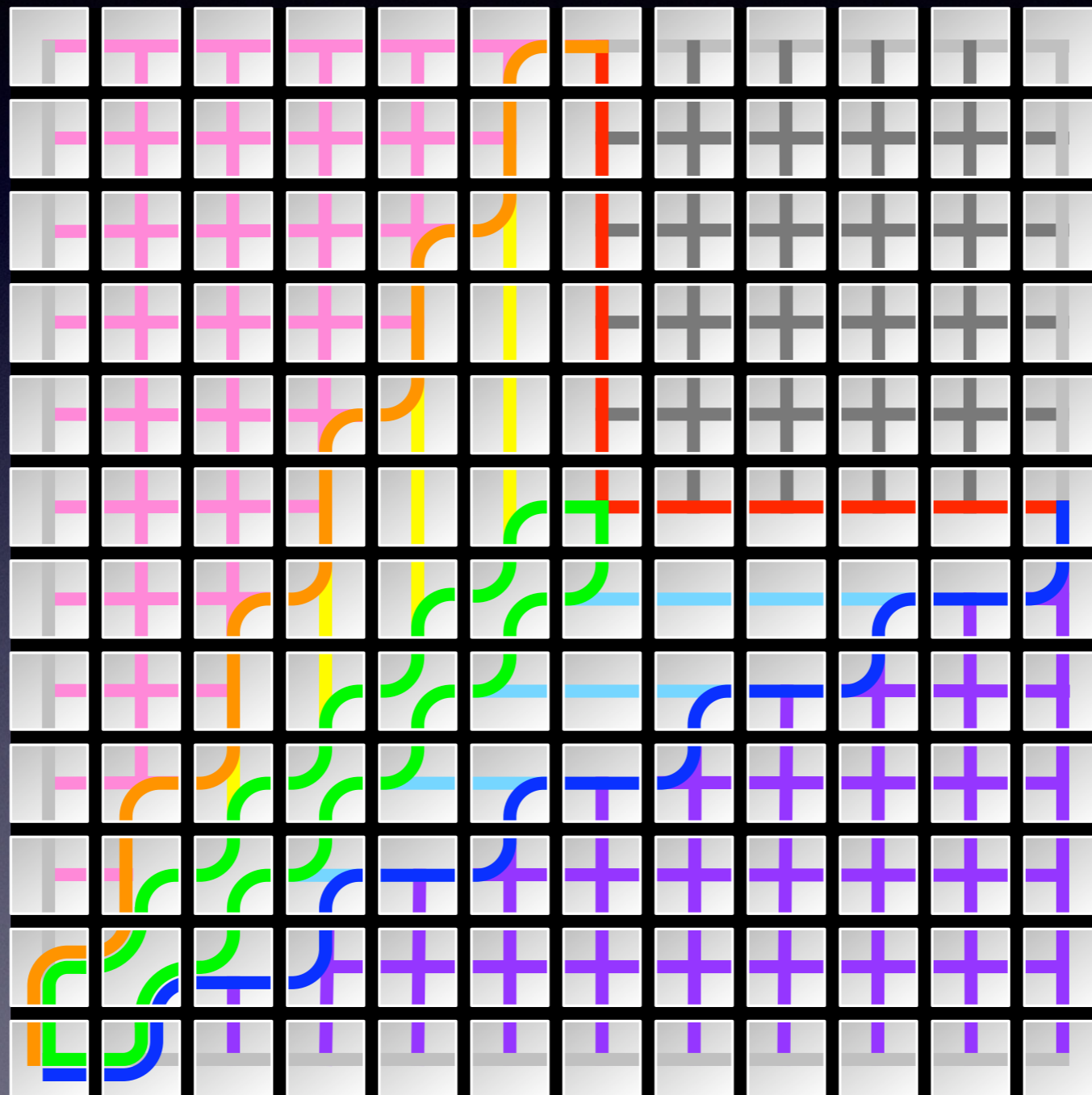
The resulting tiling set



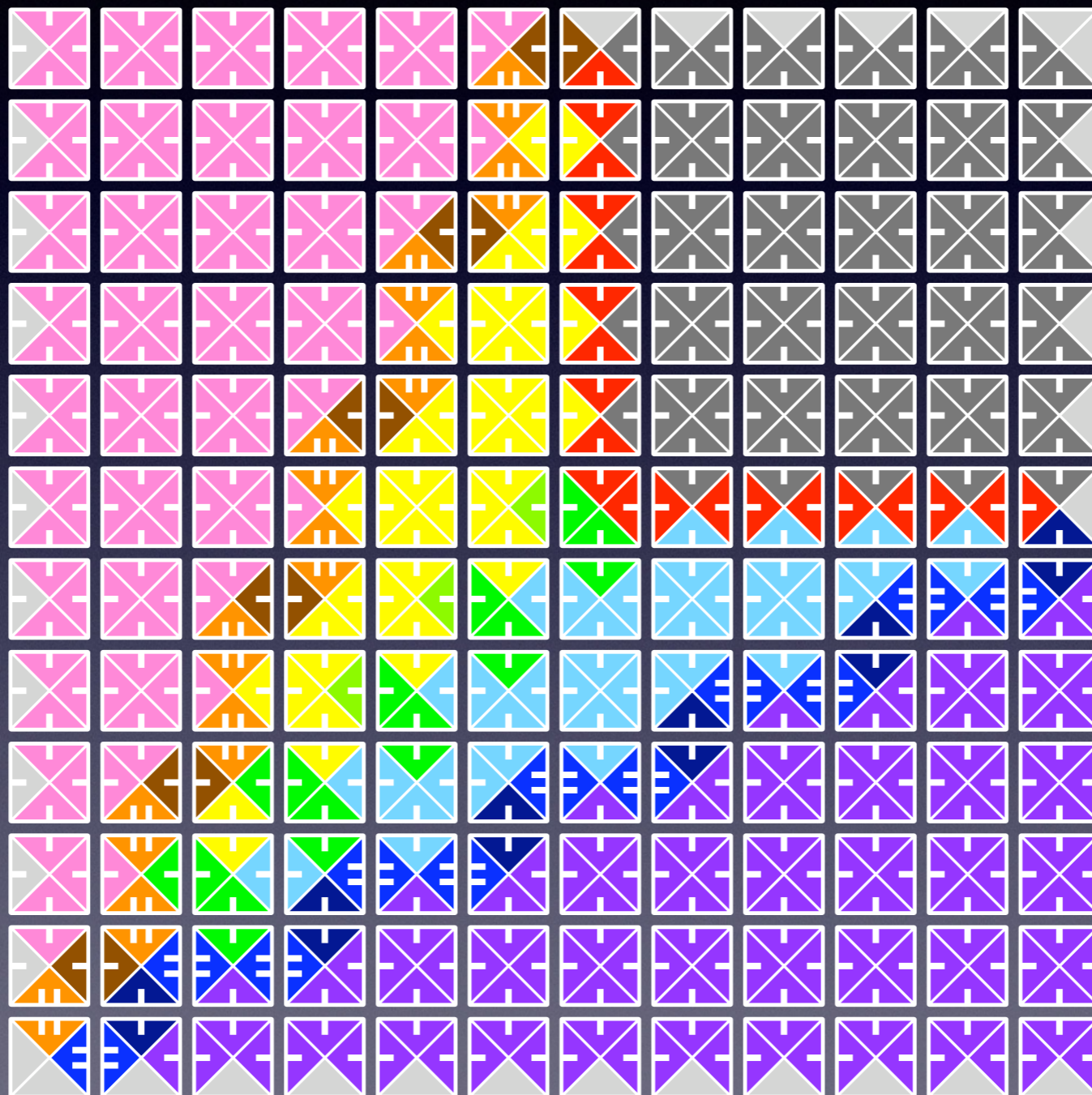
The resulting tiling set



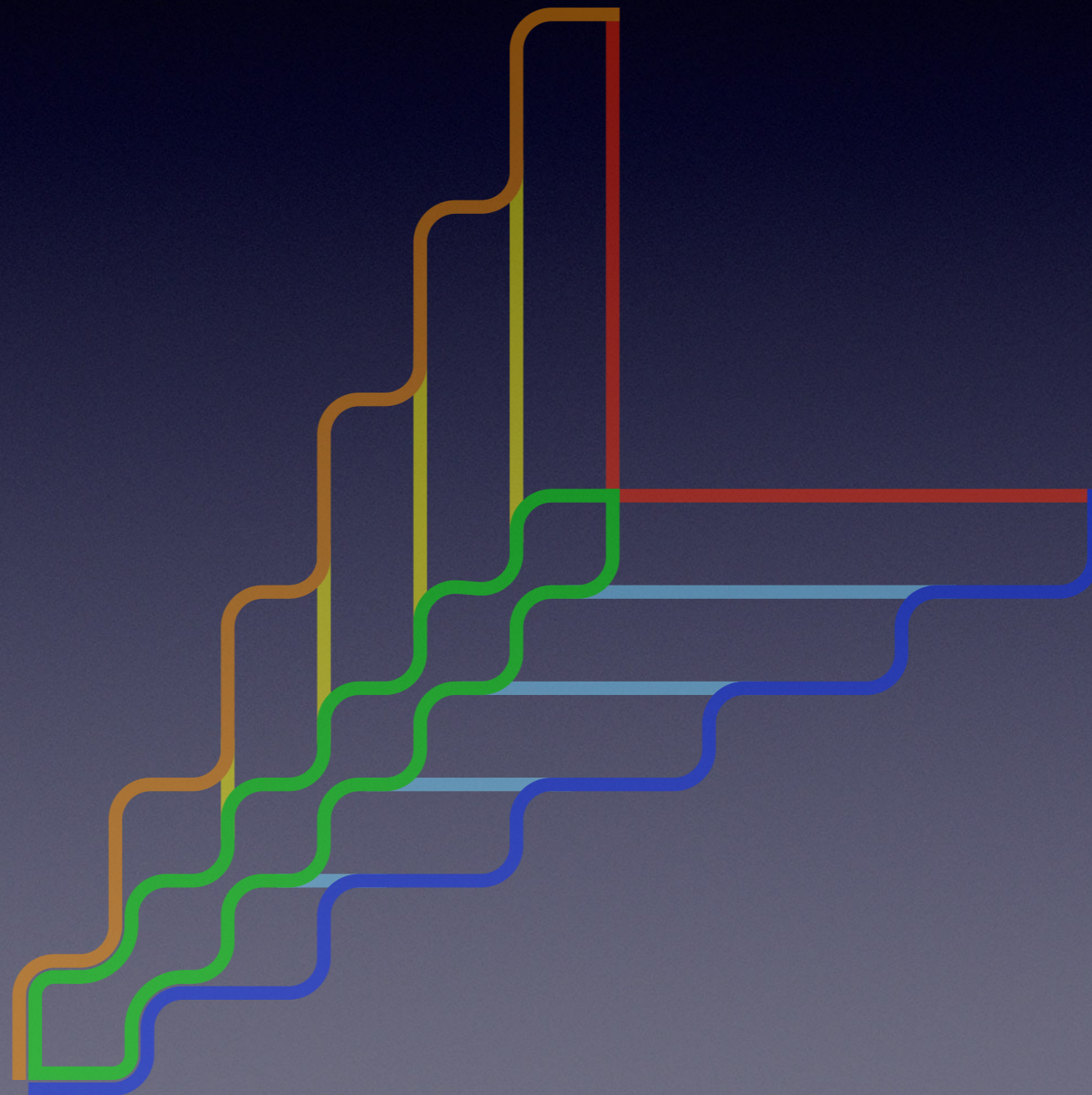
The resulting tiling set



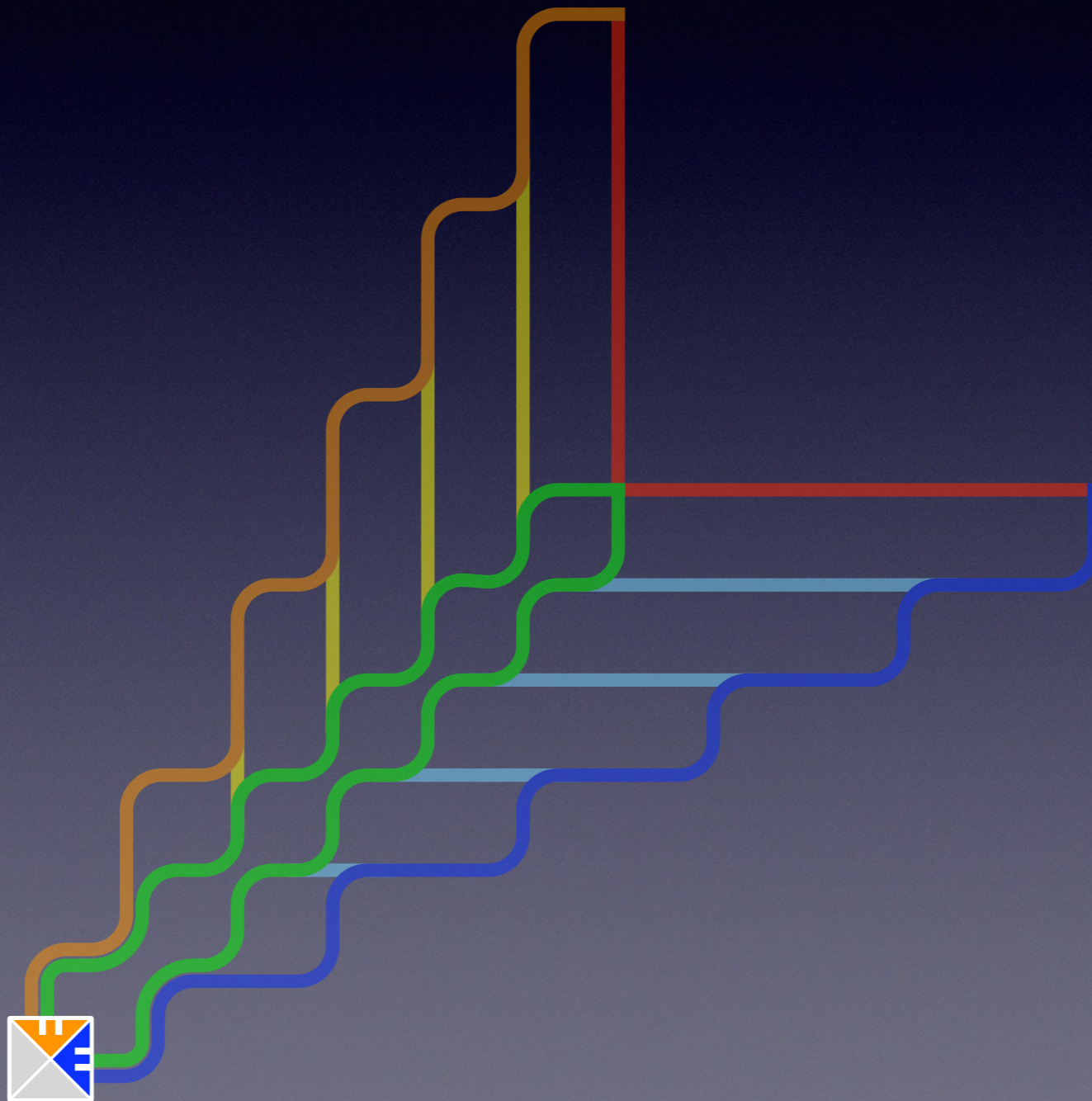
The resulting tiling set



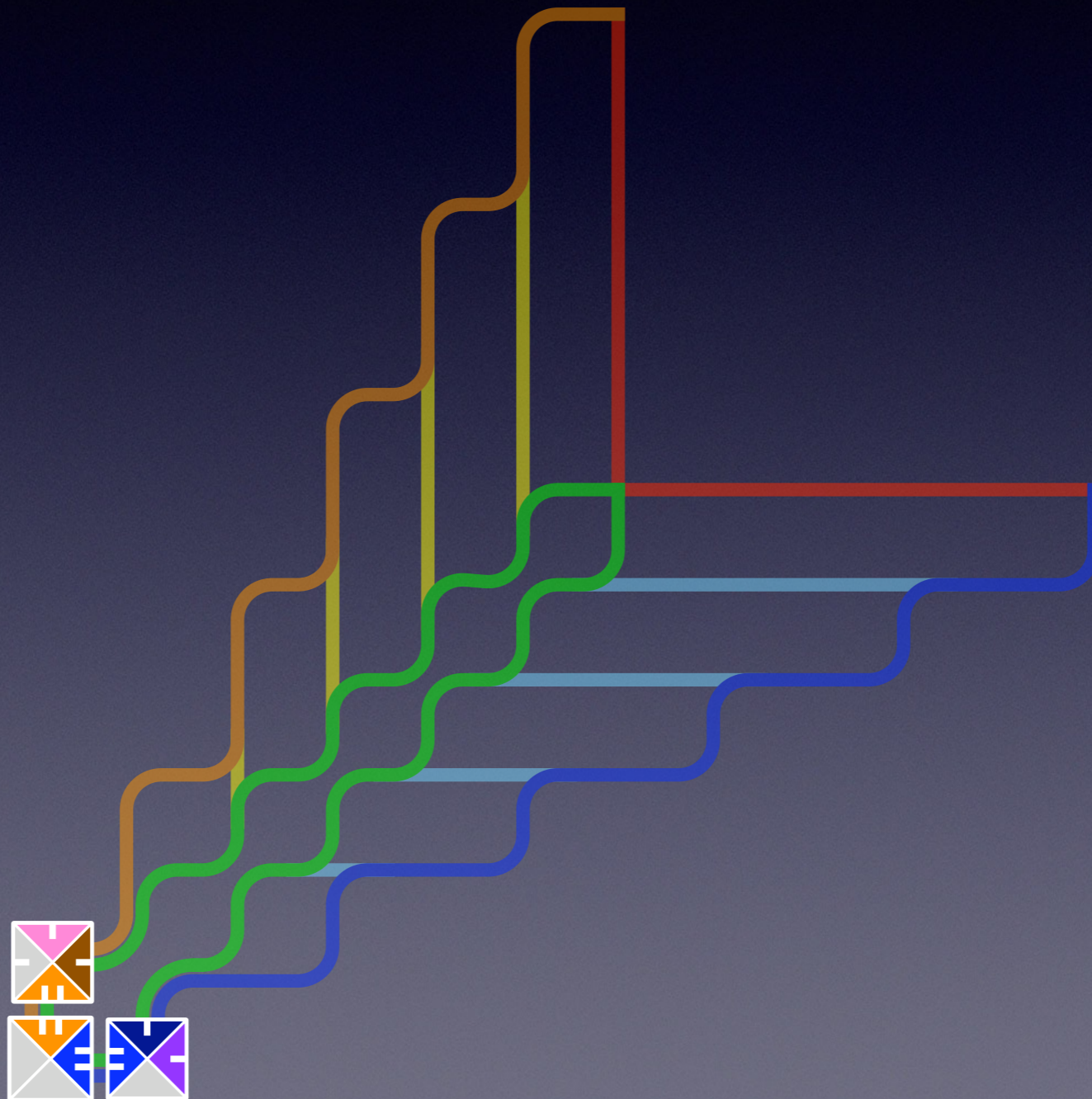
Running the tilesset



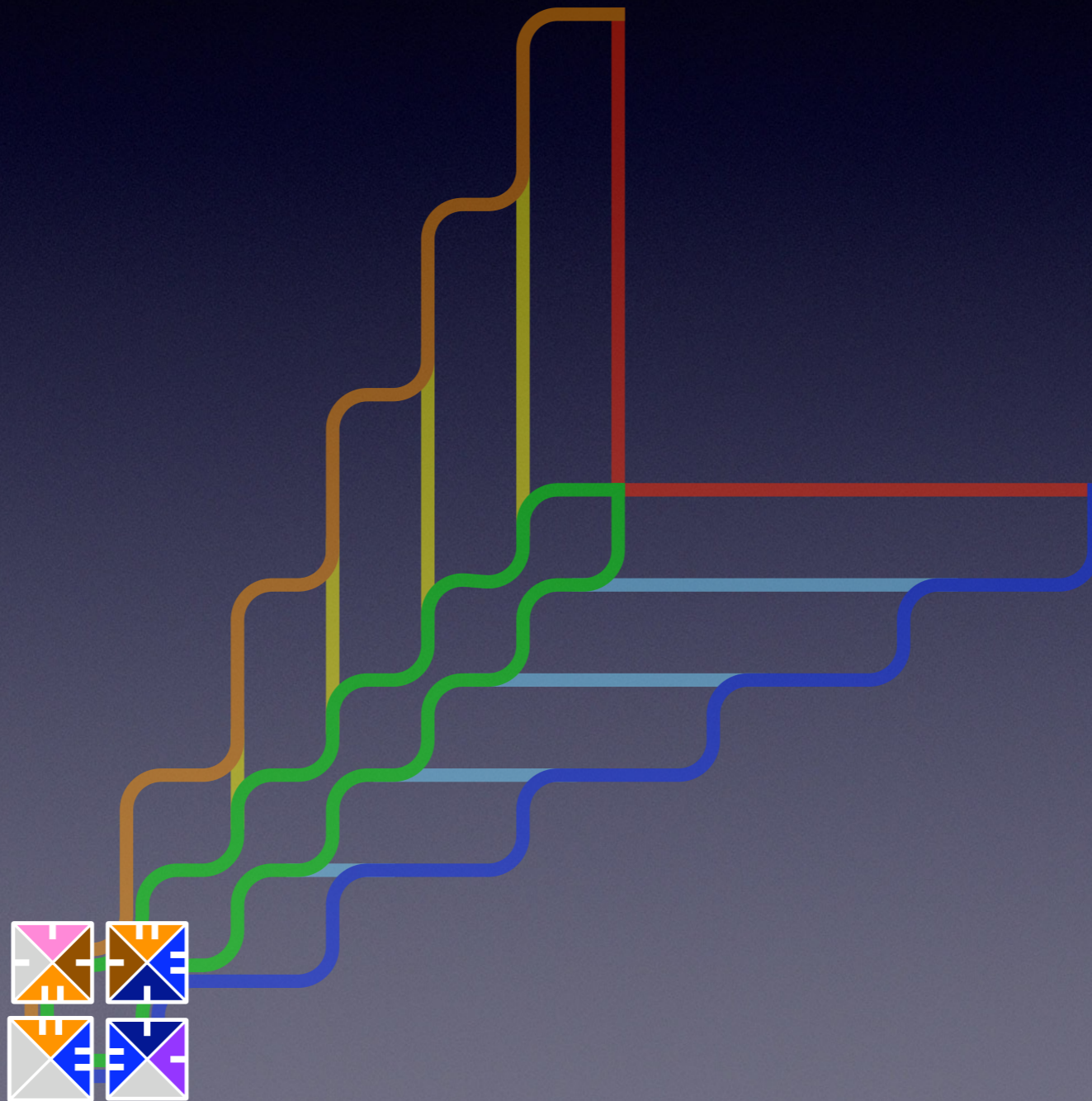
Running the tilesset



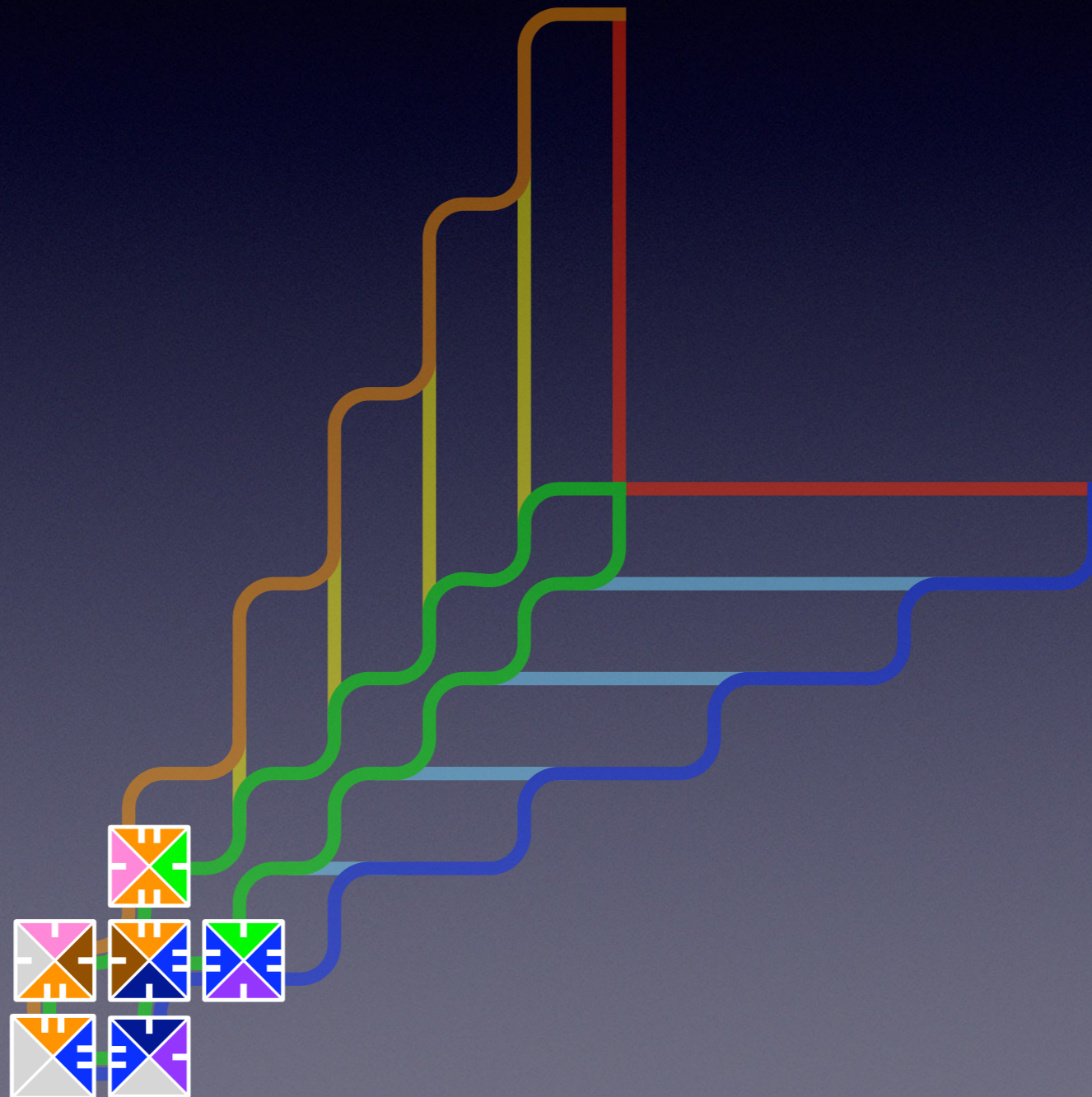
Running the tilename



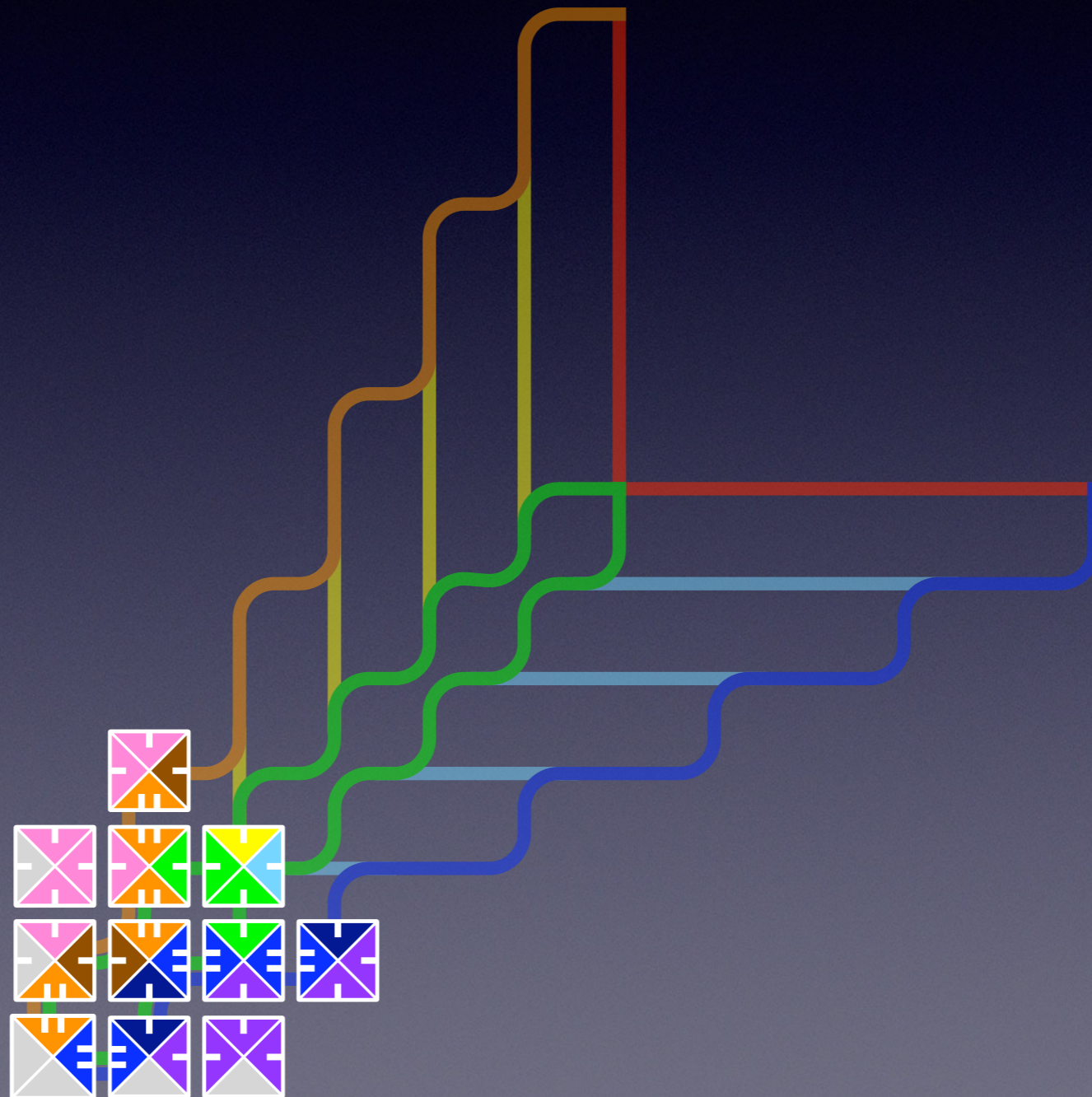
Running the tilename



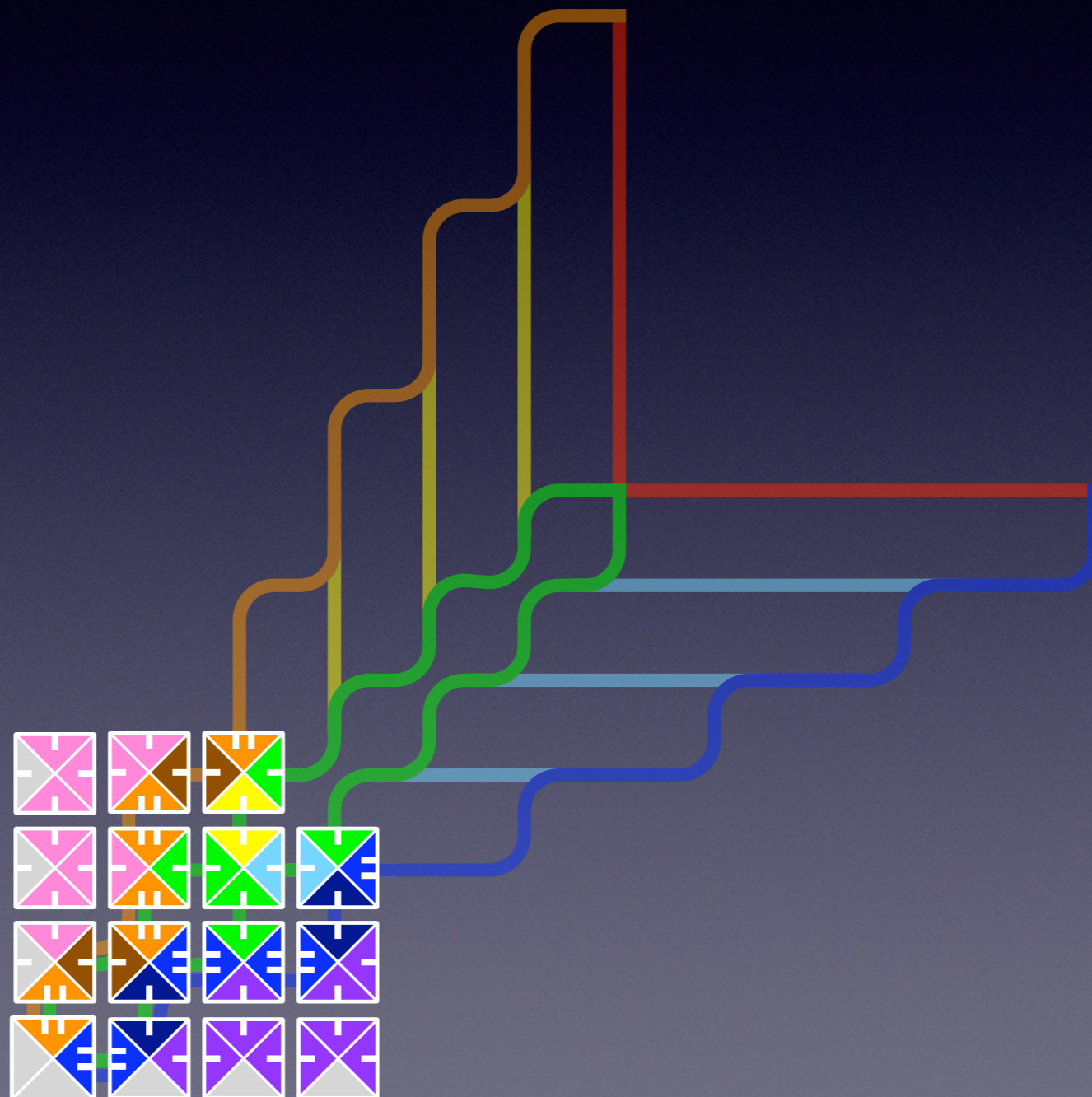
Running the tilename



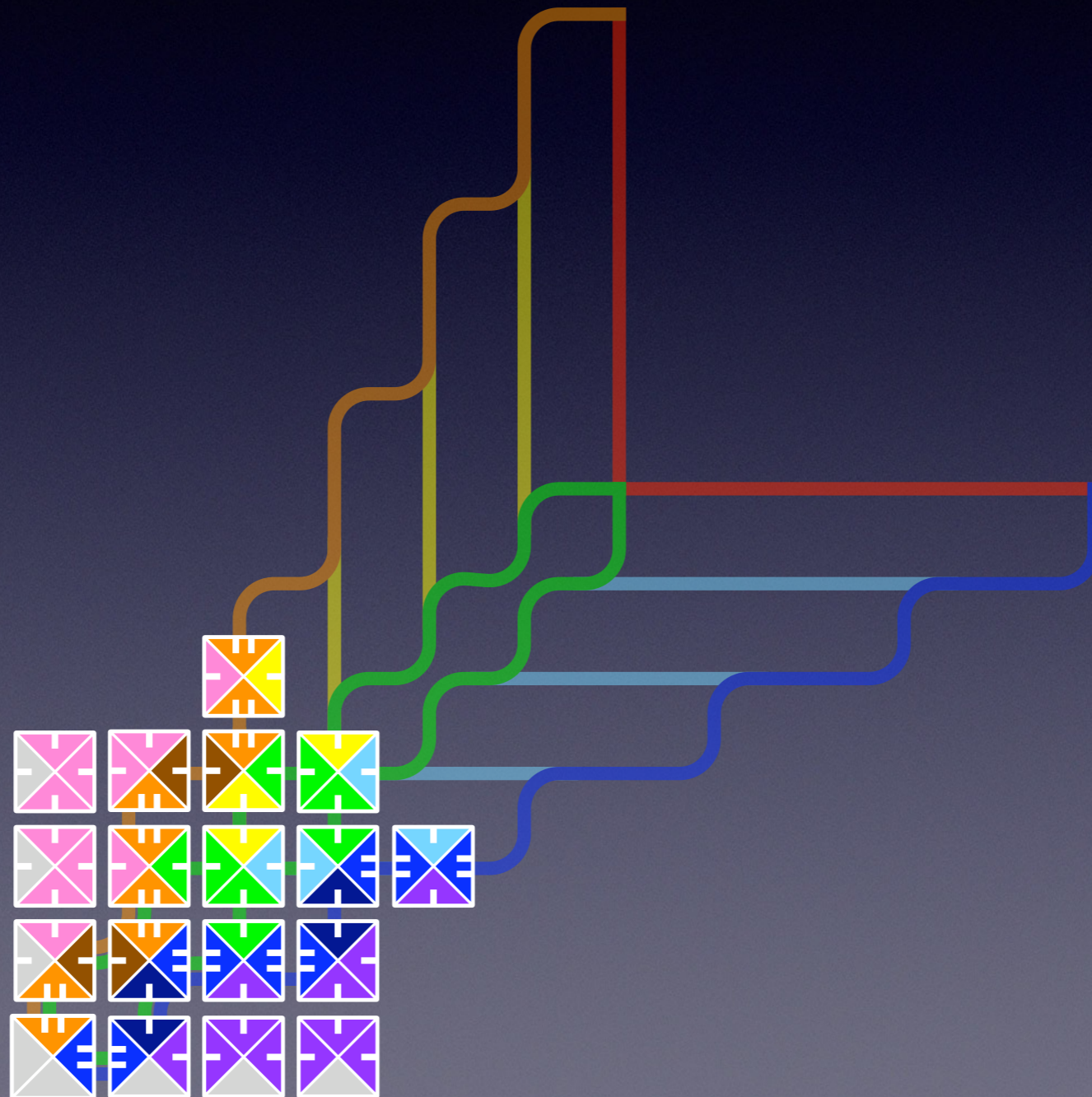
Running the tileset



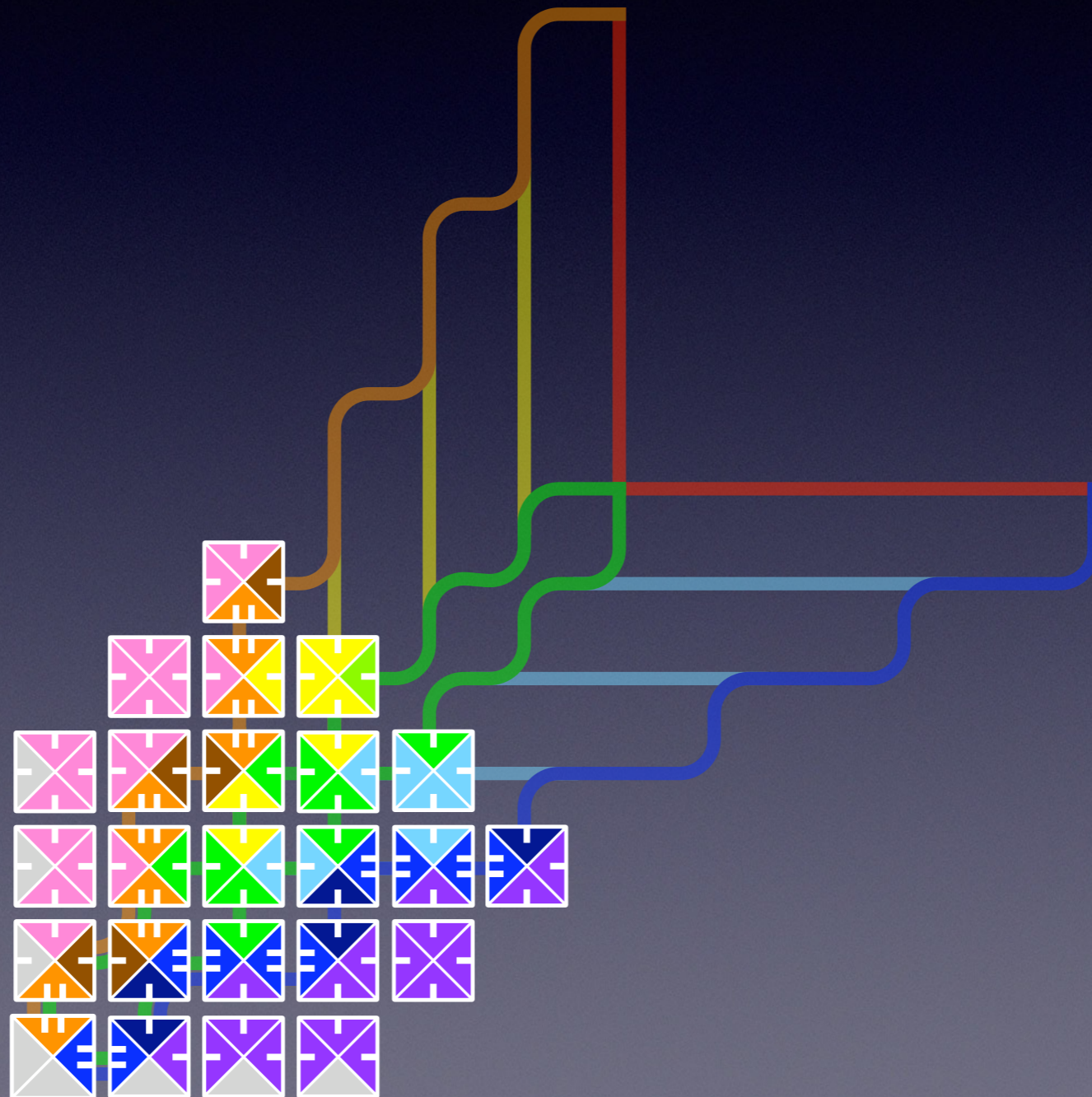
Running the tilename



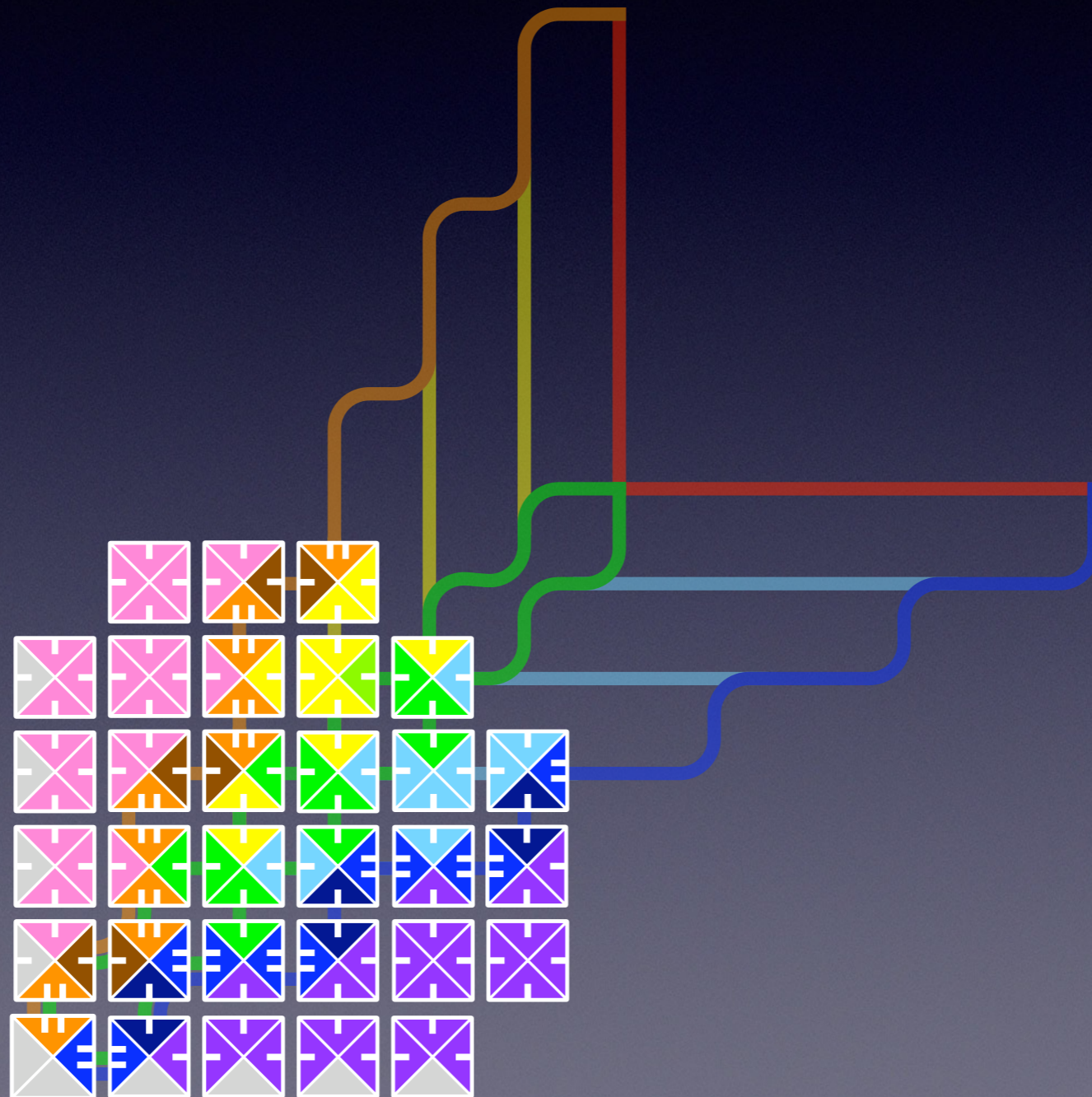
Running the tilename



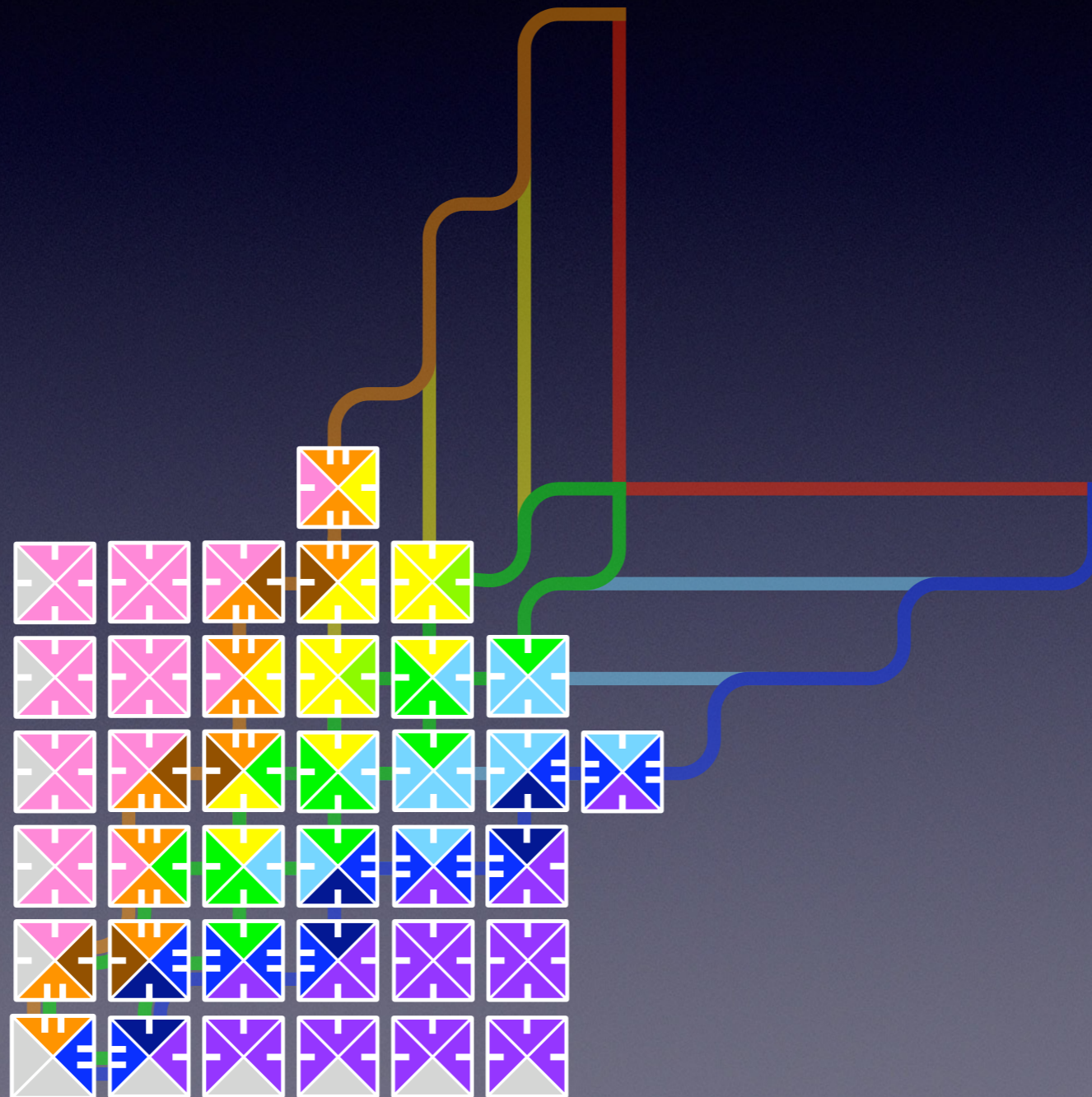
Running the tileset



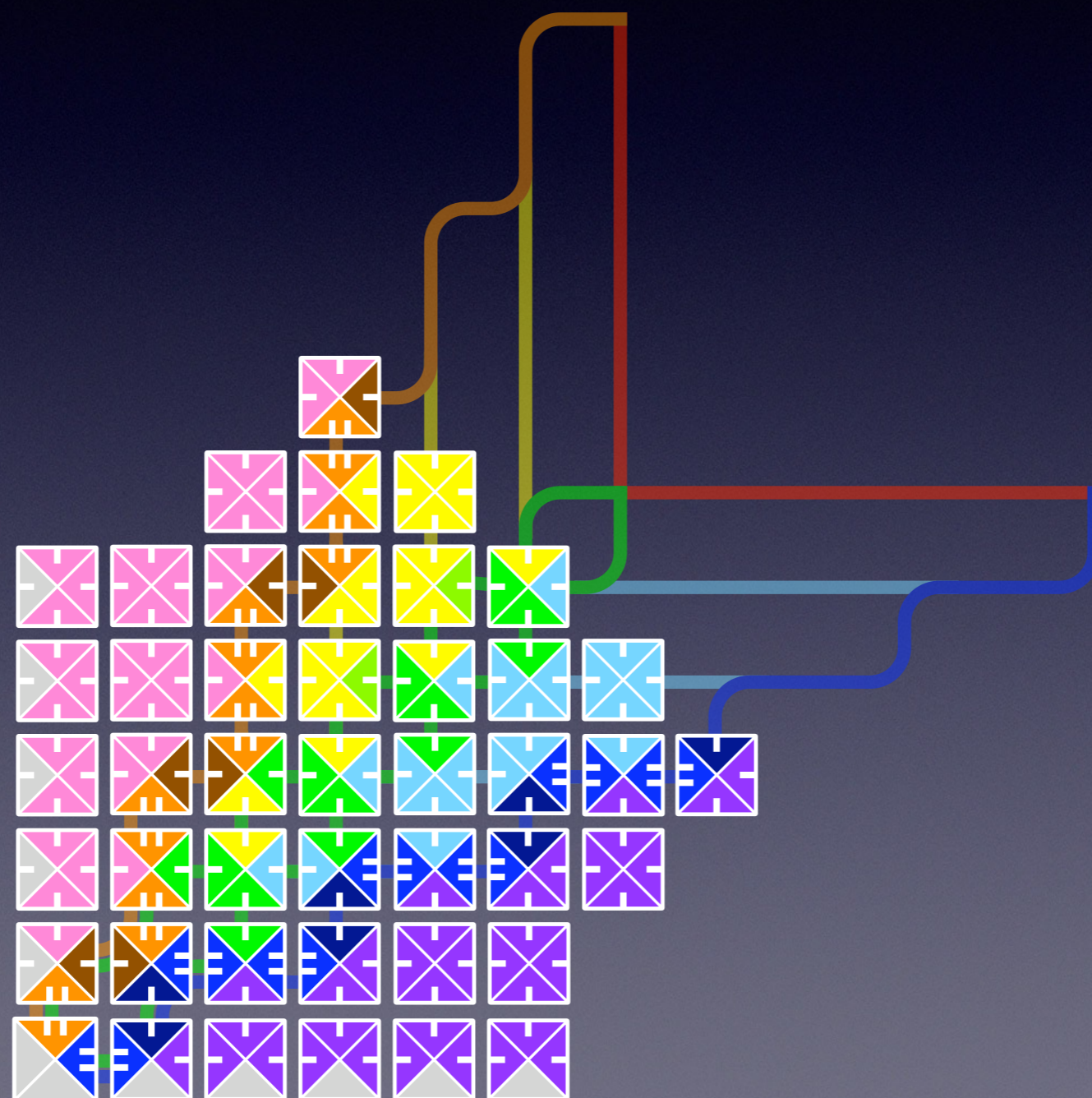
Running the tilename



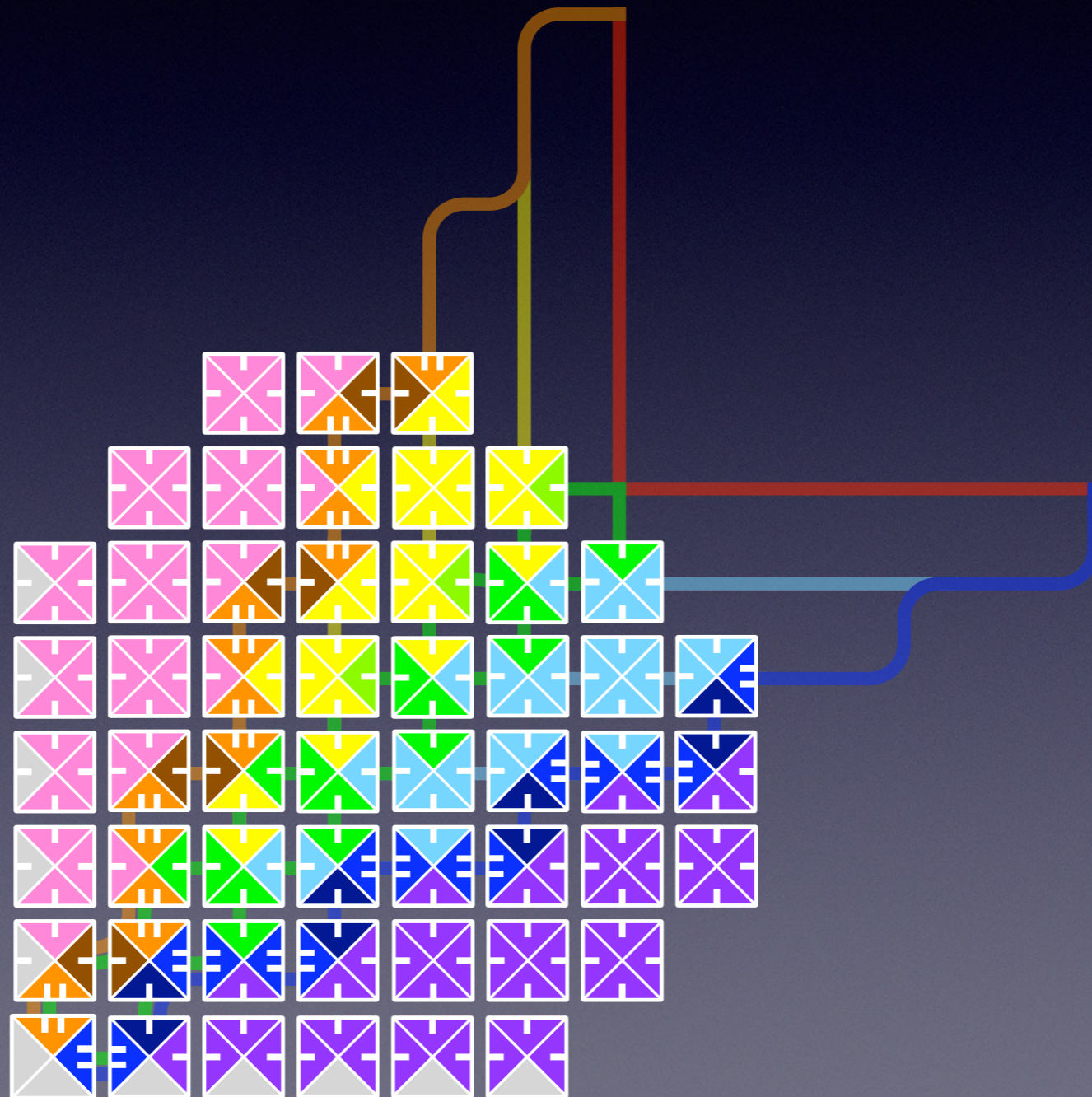
Running the tilename



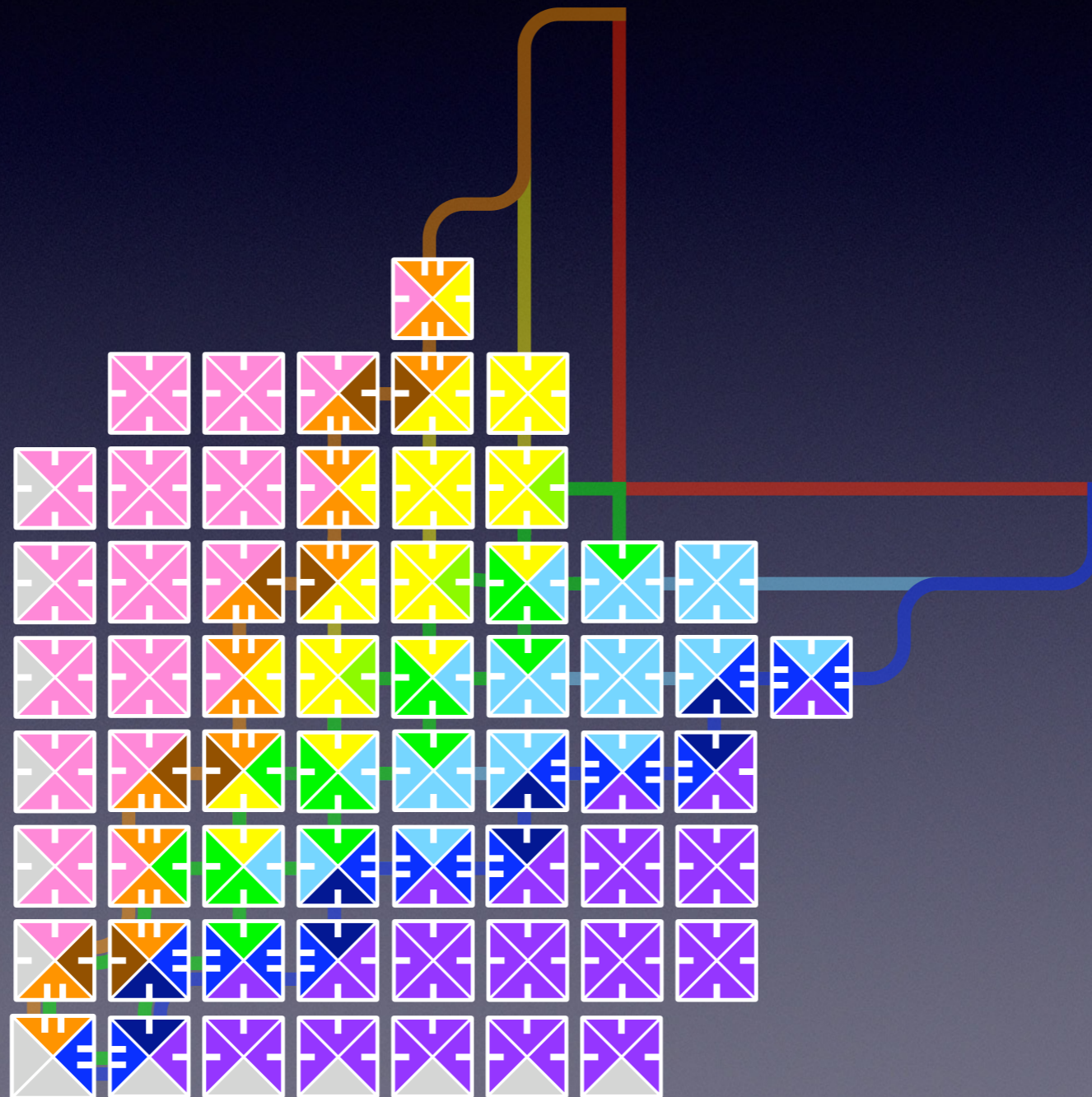
Running the tileset



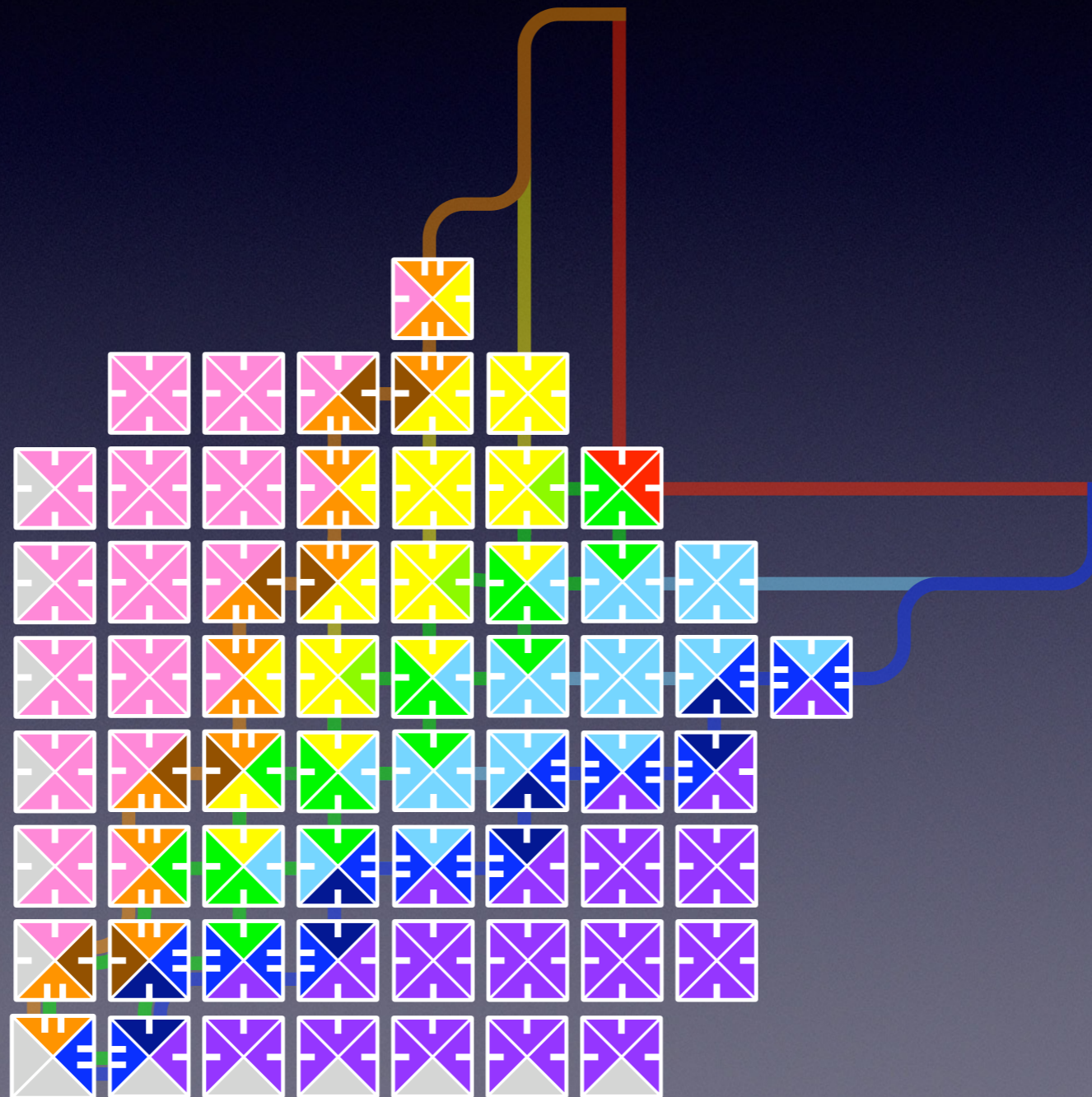
Running the tilesset



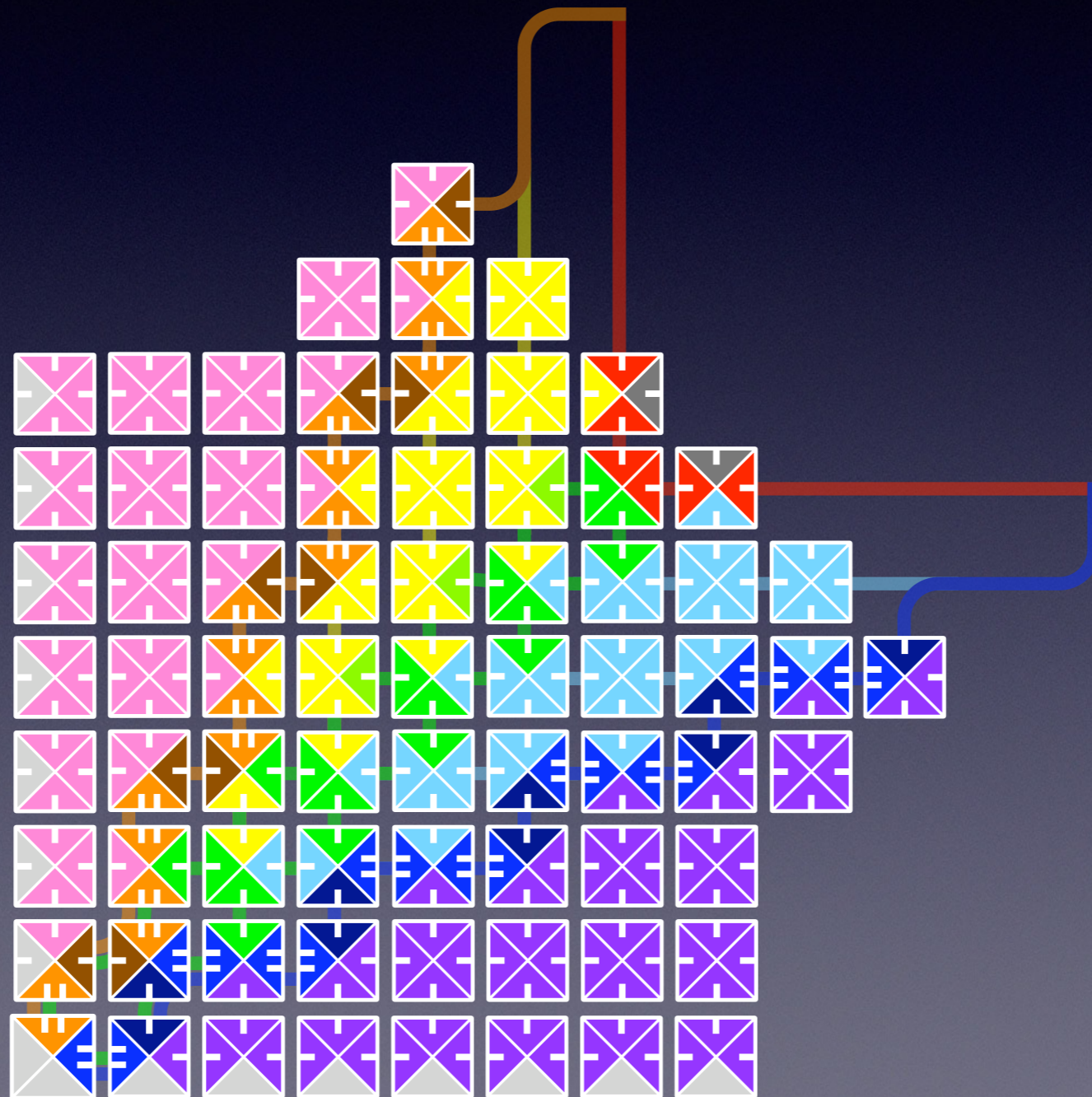
Running the tilename



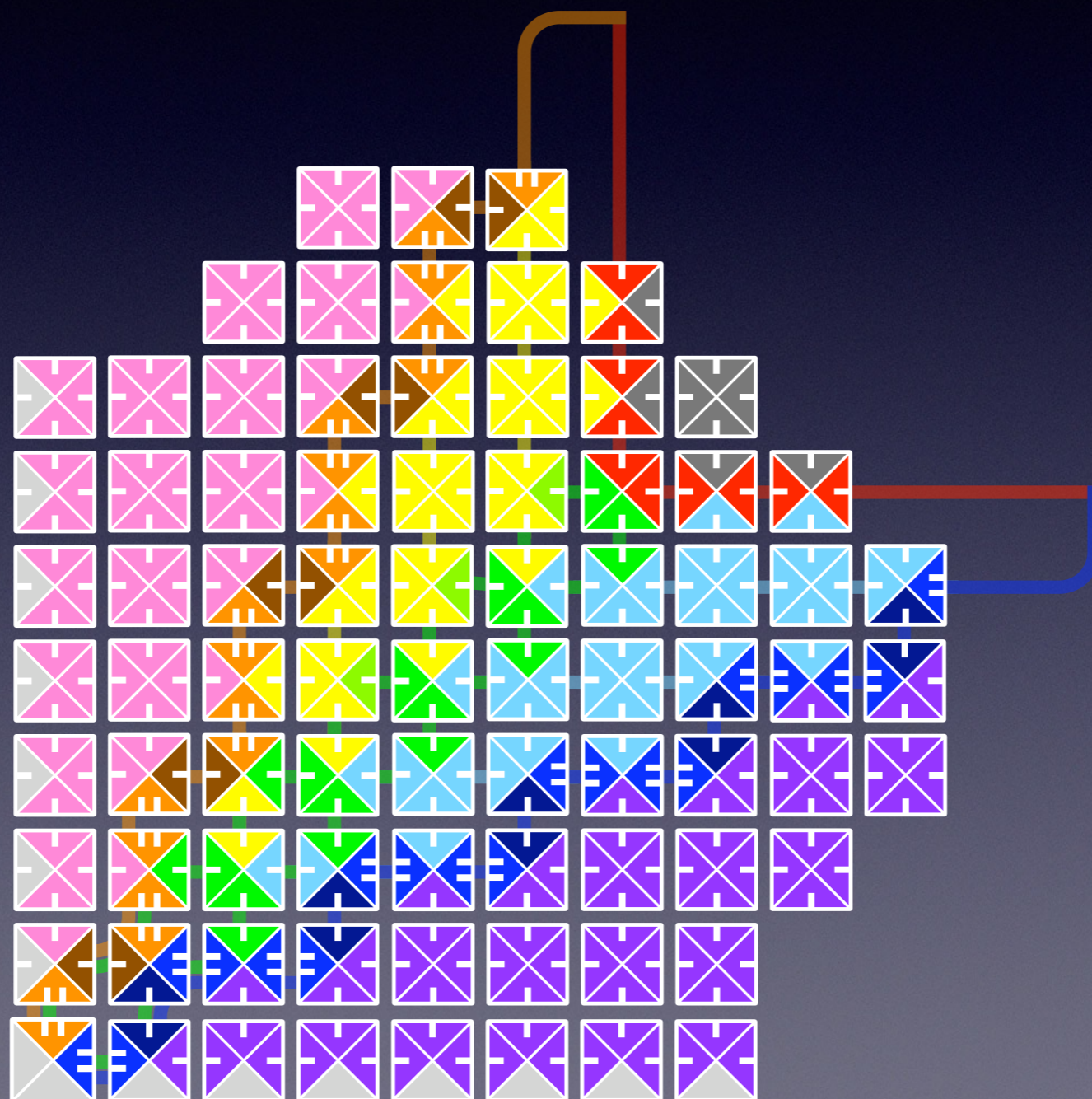
Running the tilename



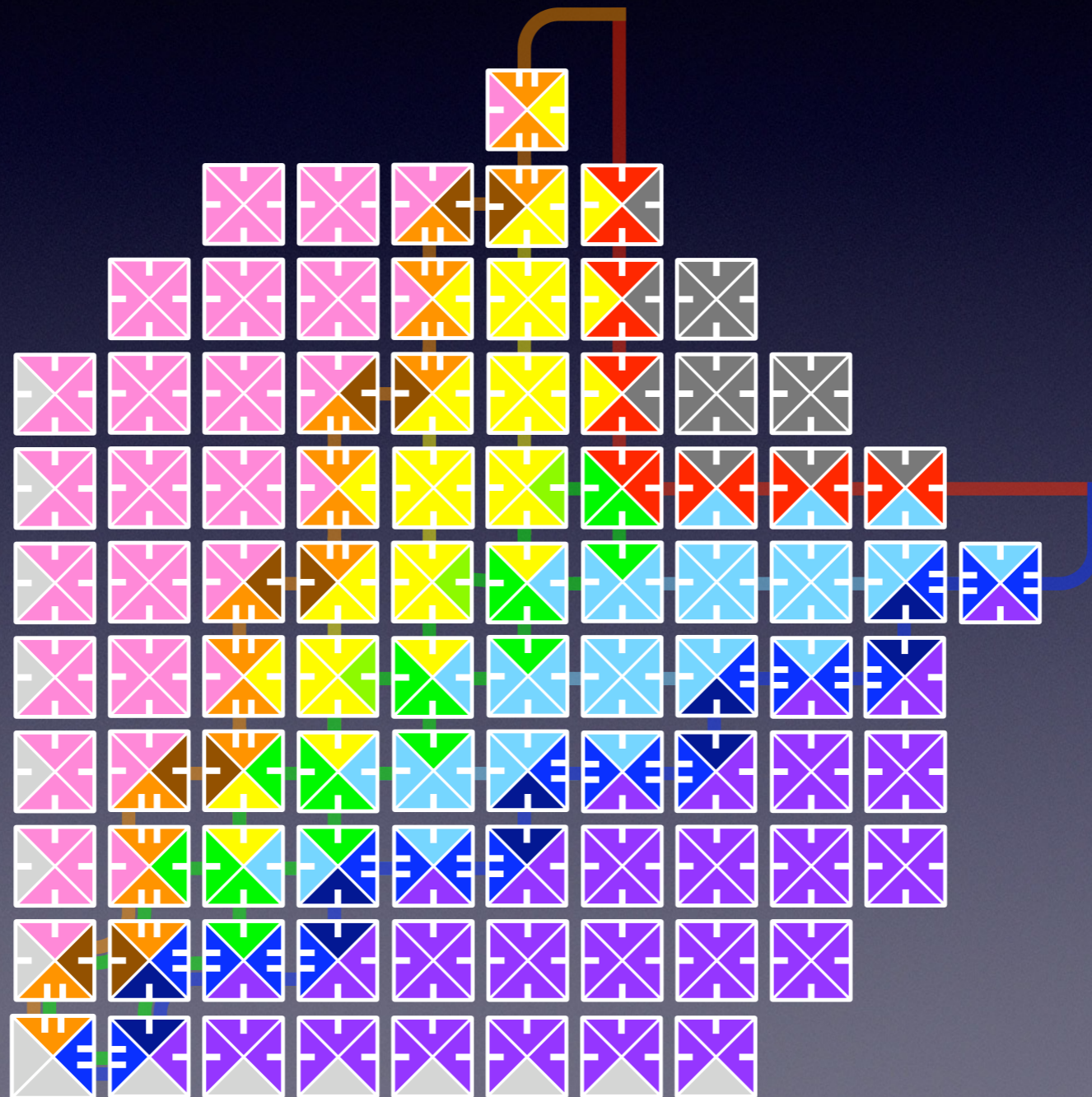
Running the tilename



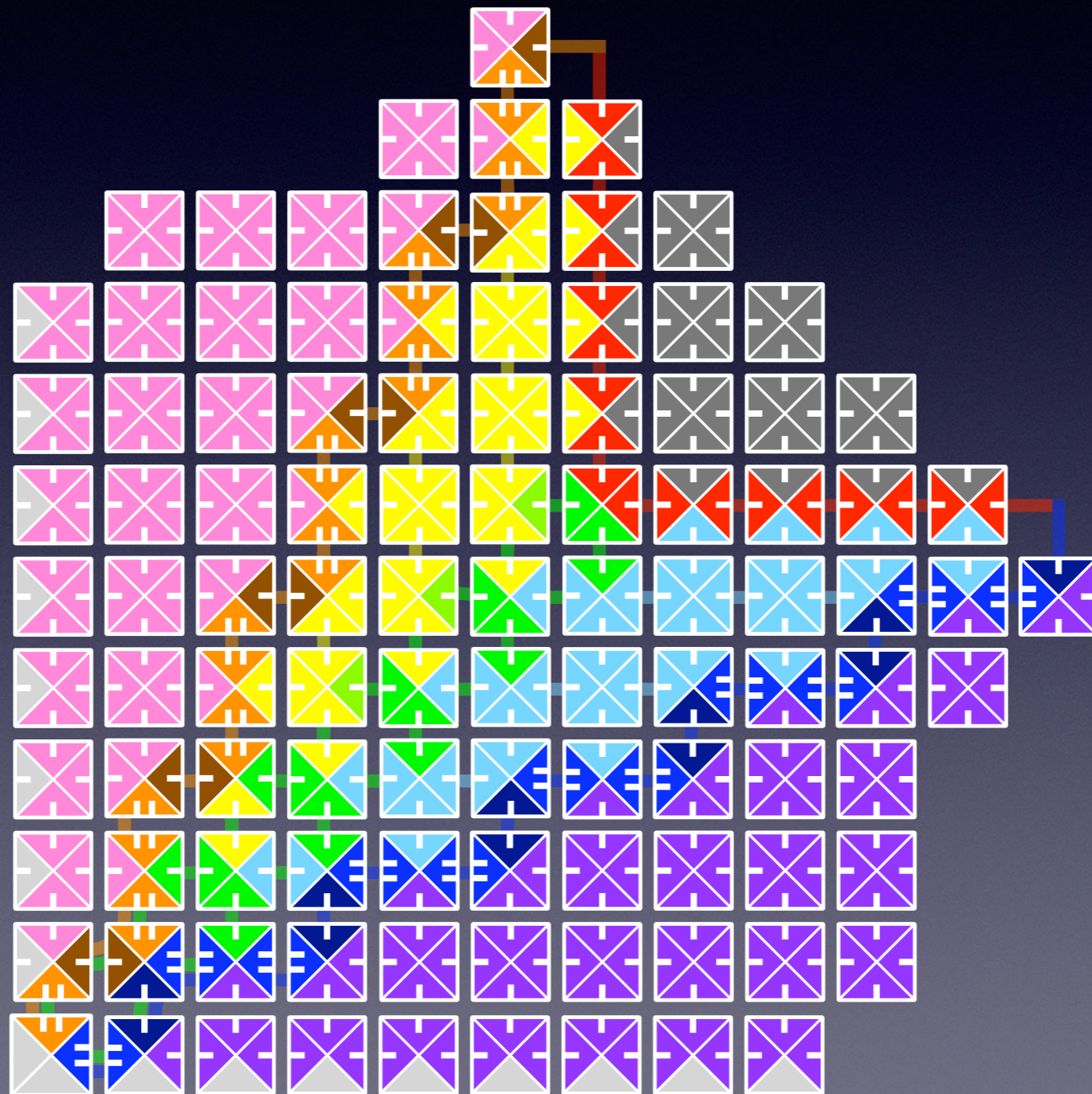
Running the tileset



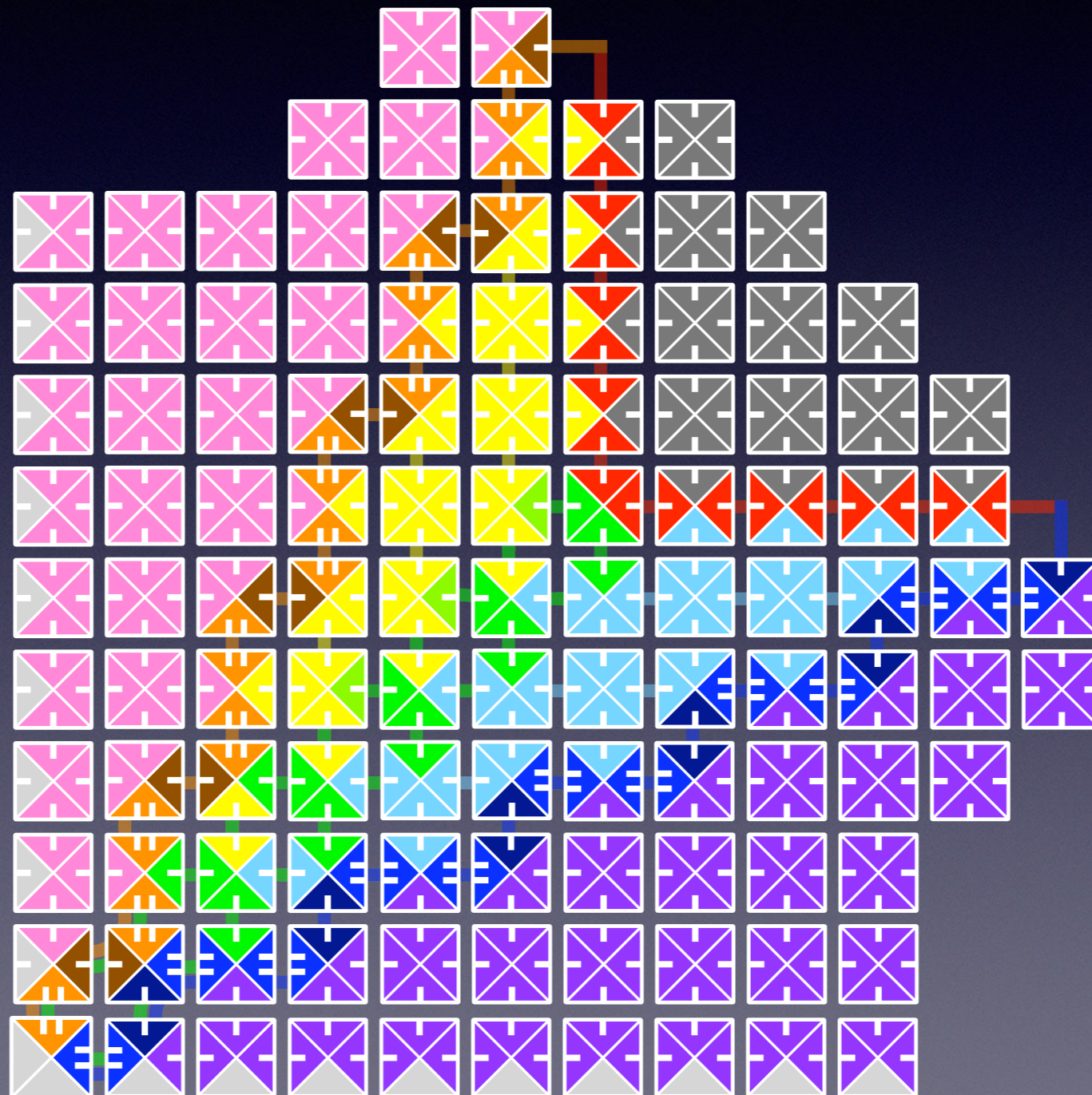
Running the tileset



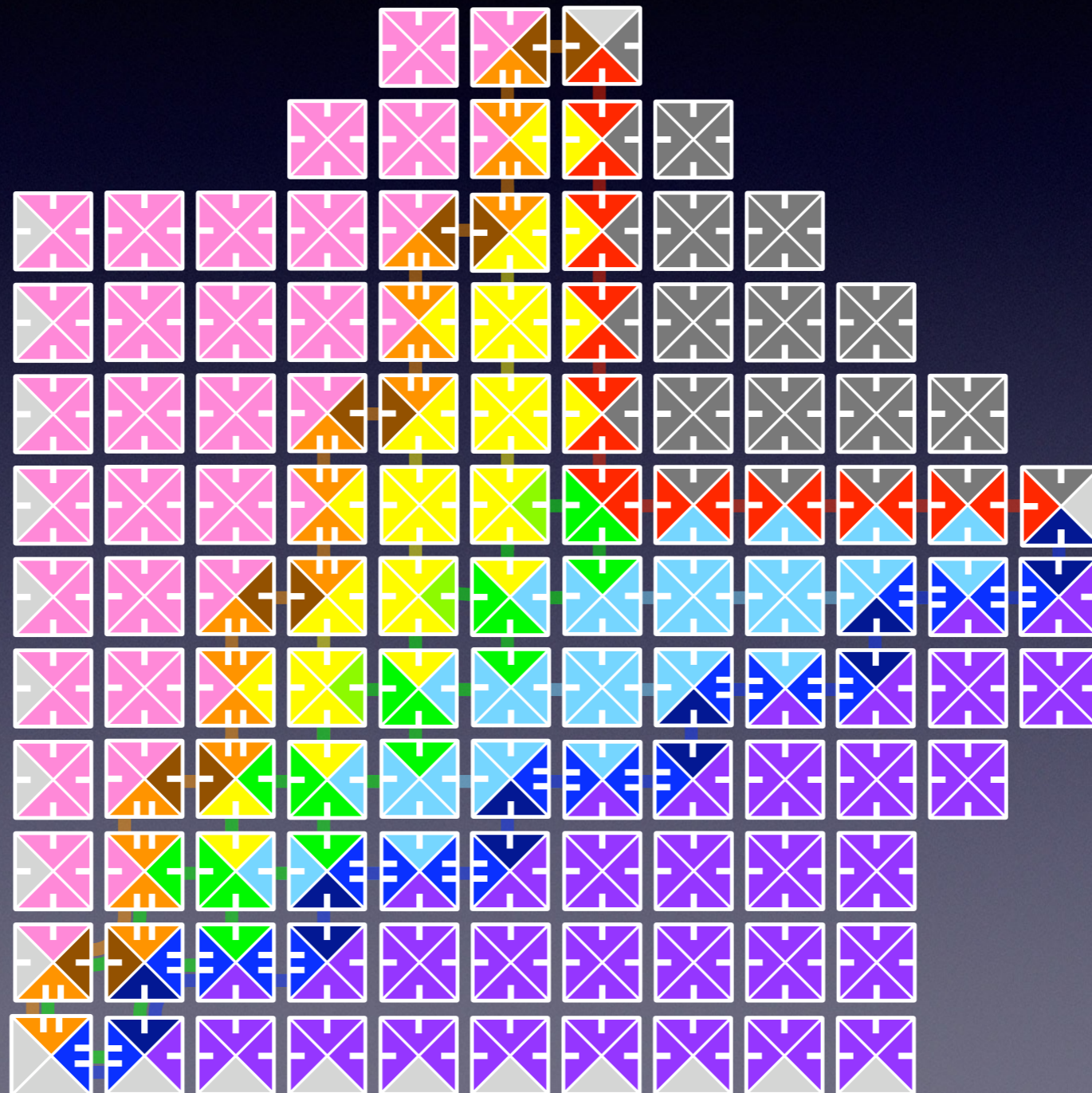
Running the tiling set



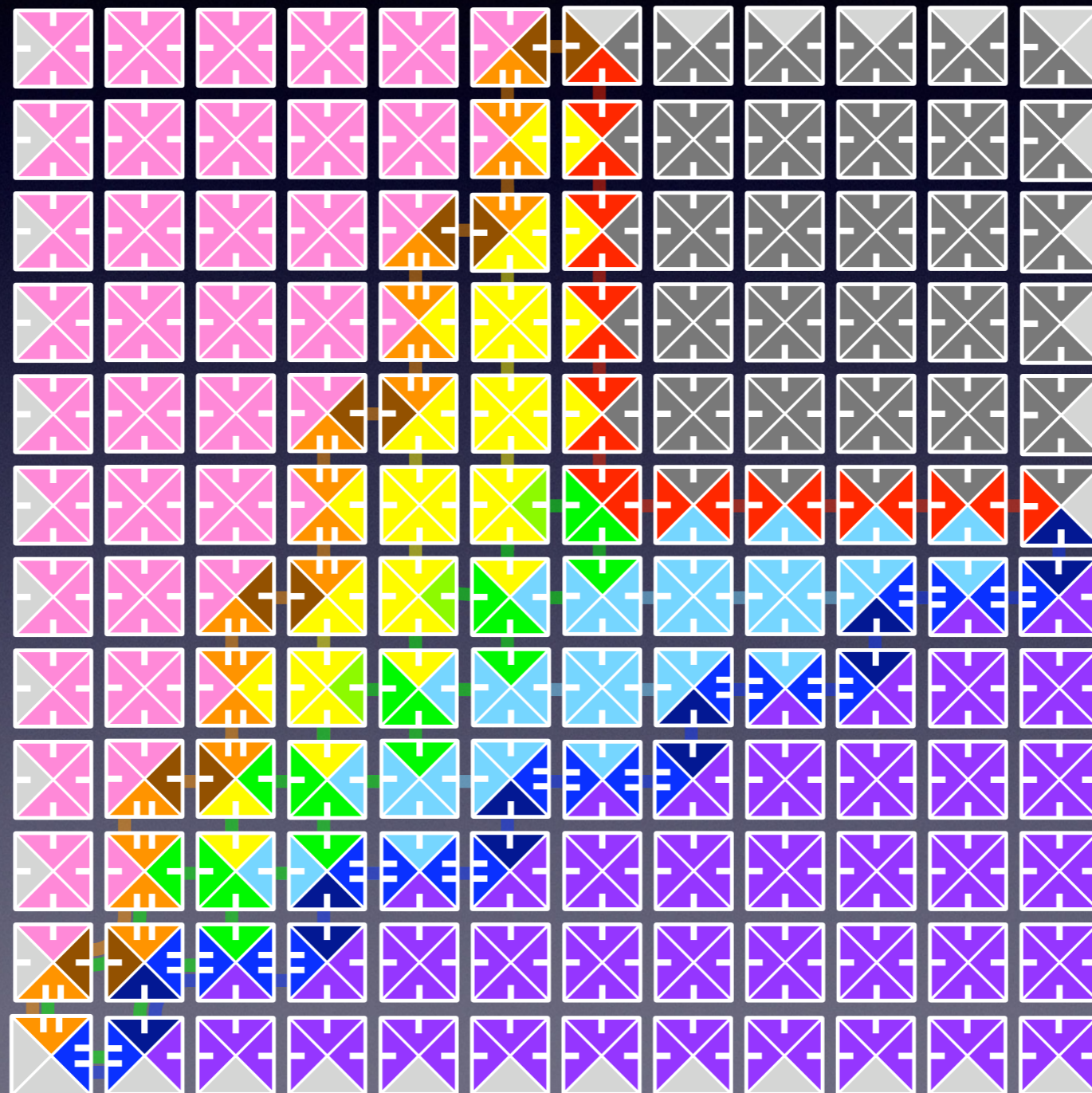
Running the tiling set



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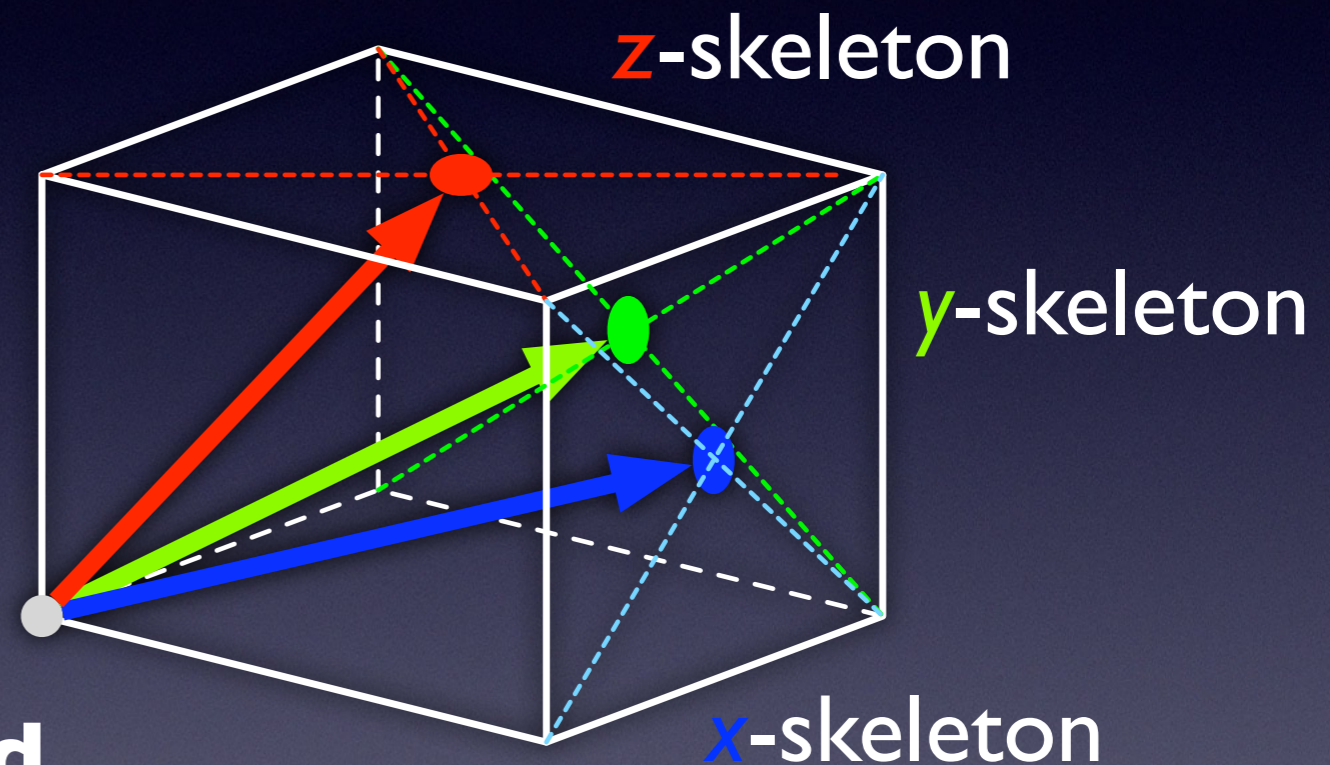


Assembling cubes in real time

The skeleton & its rank function

The skeleton.

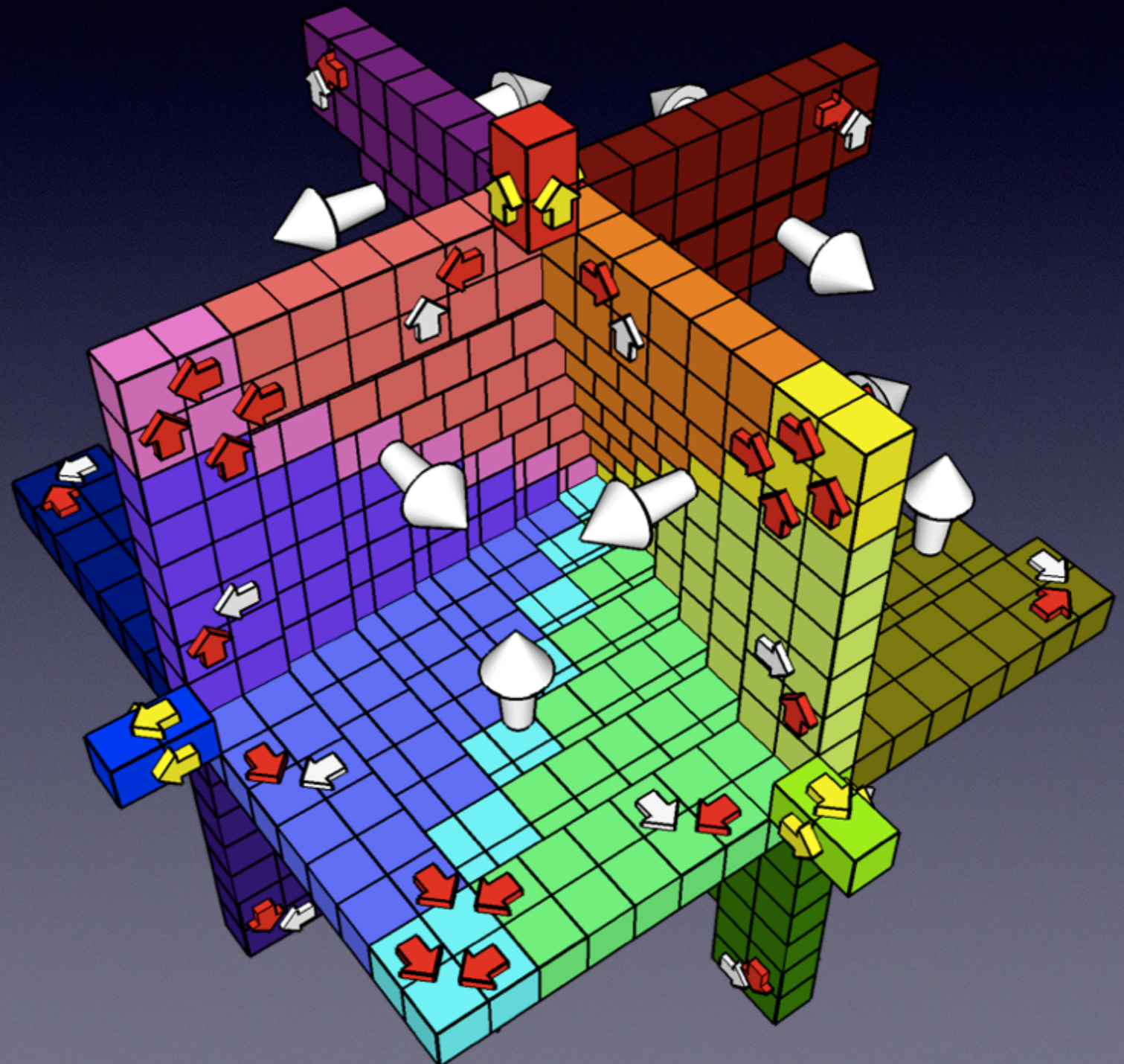
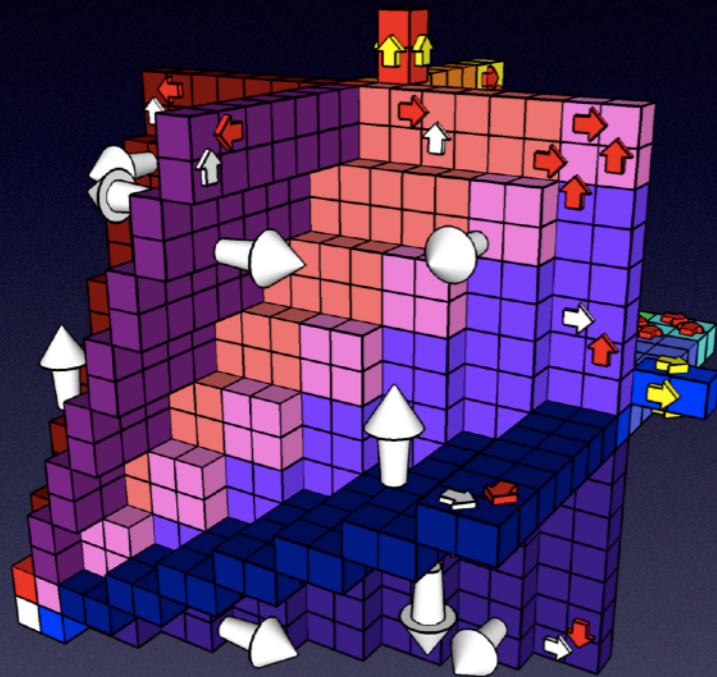
$$\begin{cases} a_i = (i, \lfloor i/2 \rfloor, \lfloor i/2 \rfloor) \\ b_j = (\lfloor j/2 \rfloor, j, \lfloor j/2 \rfloor) \\ c_k = (\lfloor k/2 \rfloor, \lfloor k/2 \rfloor, k) \end{cases}$$



The rank function induced.

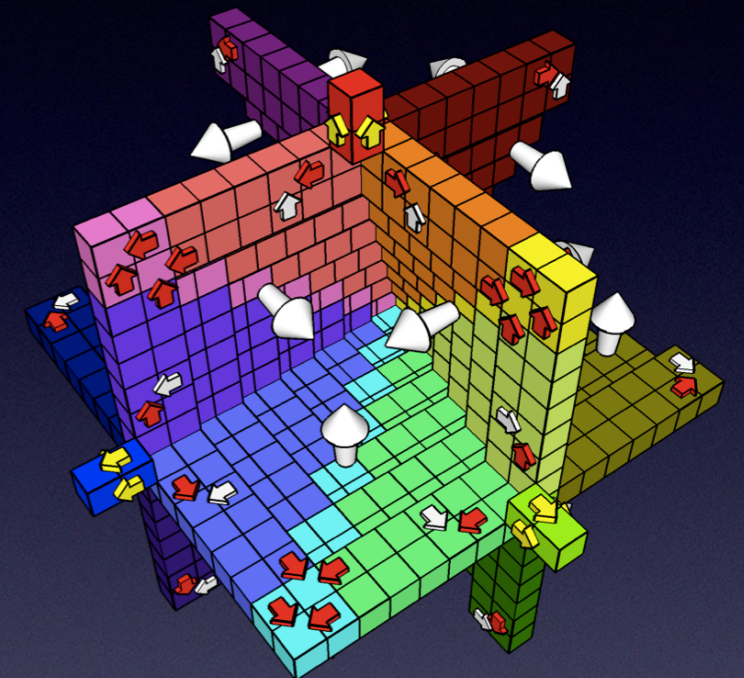
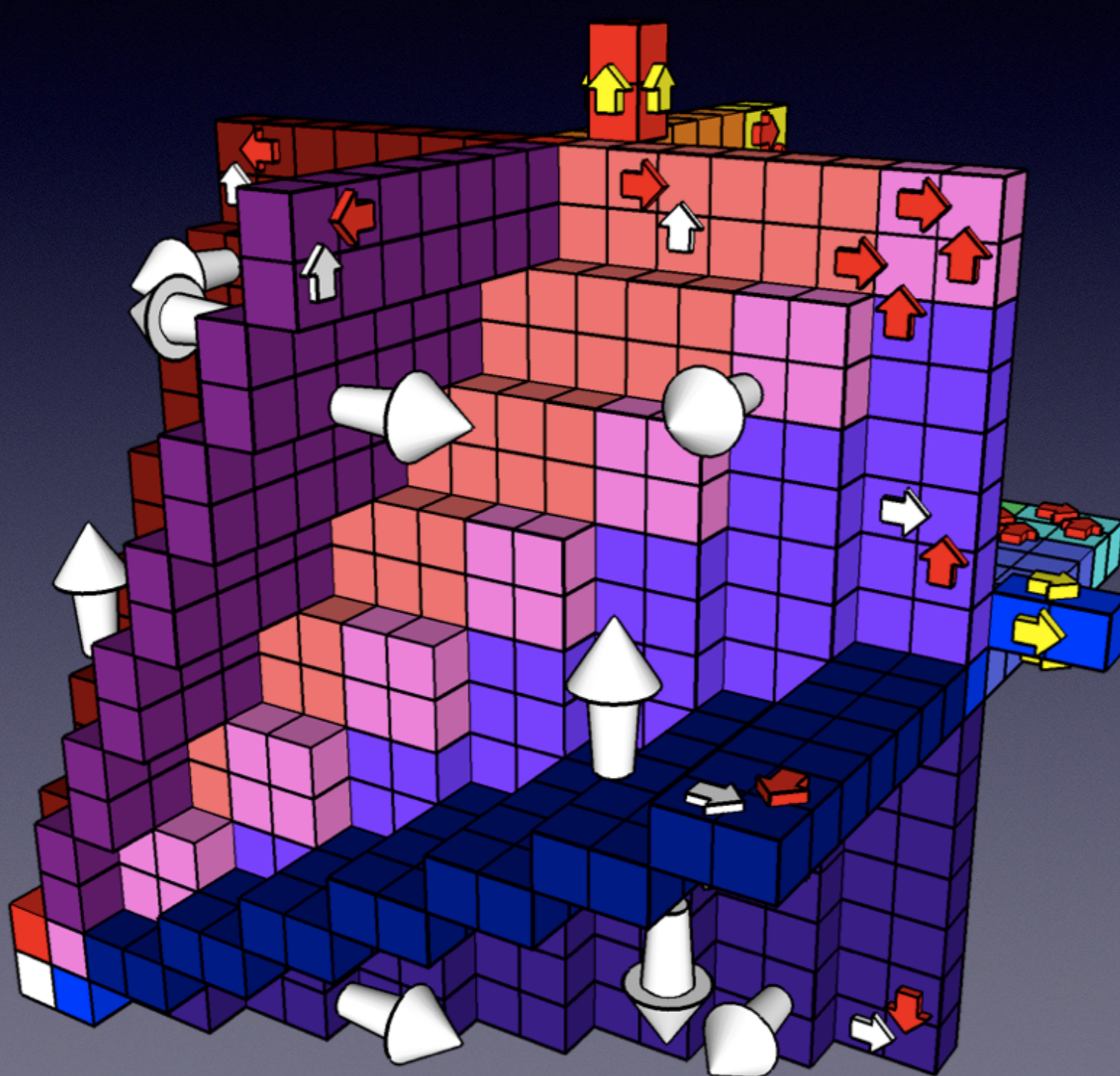
$$\text{rank}(u) = \max \begin{cases} \|a_i\|_1 + \|u - a_i\|_1 \\ \|b_j\|_1 + \|u - b_j\|_1 \\ \|c_k\|_1 + \|u - c_k\|_1 \end{cases}$$

Variations of the rank function



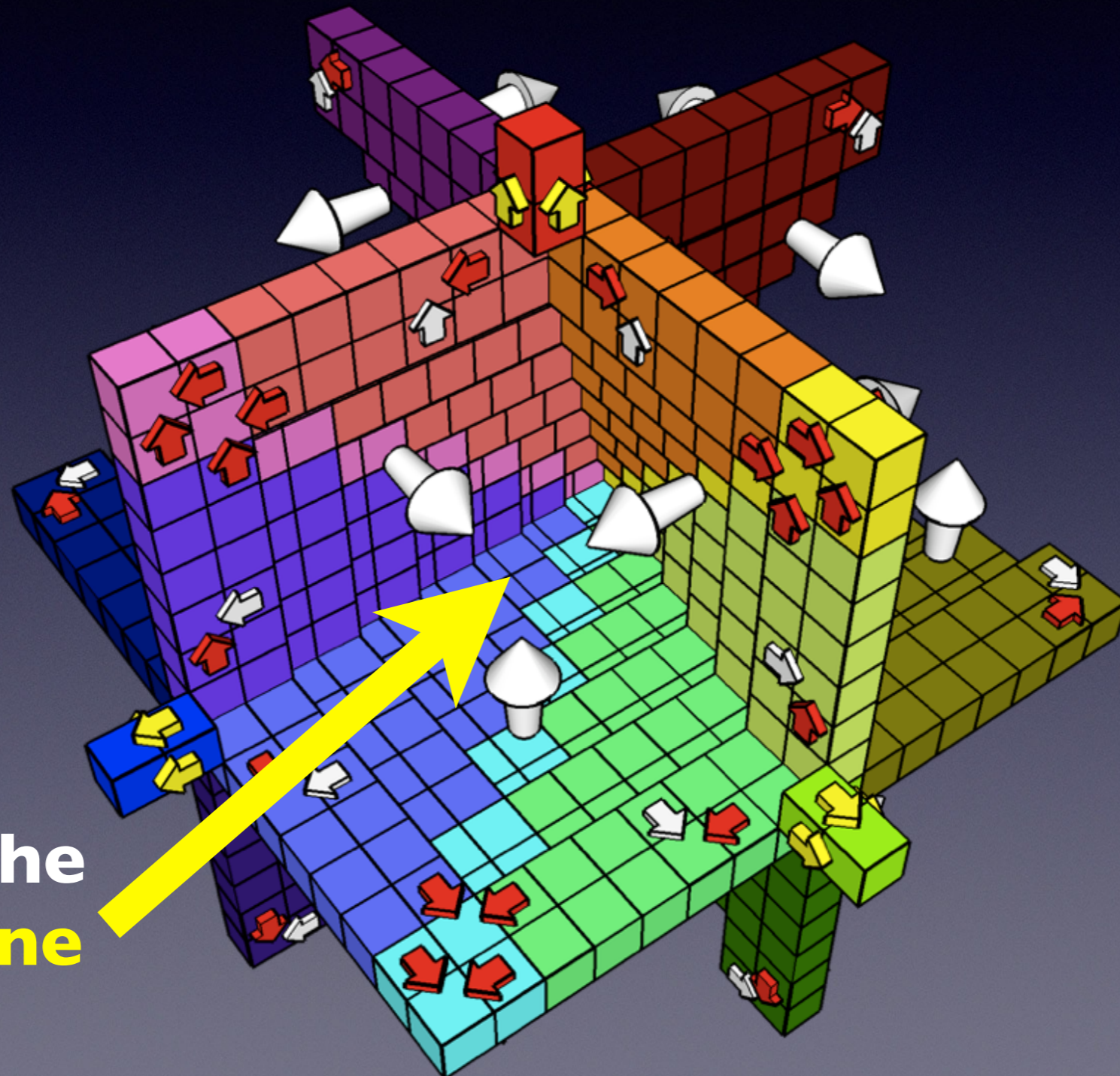
Key step:
determining the
successors from
the predecessors

Variations of the rank function



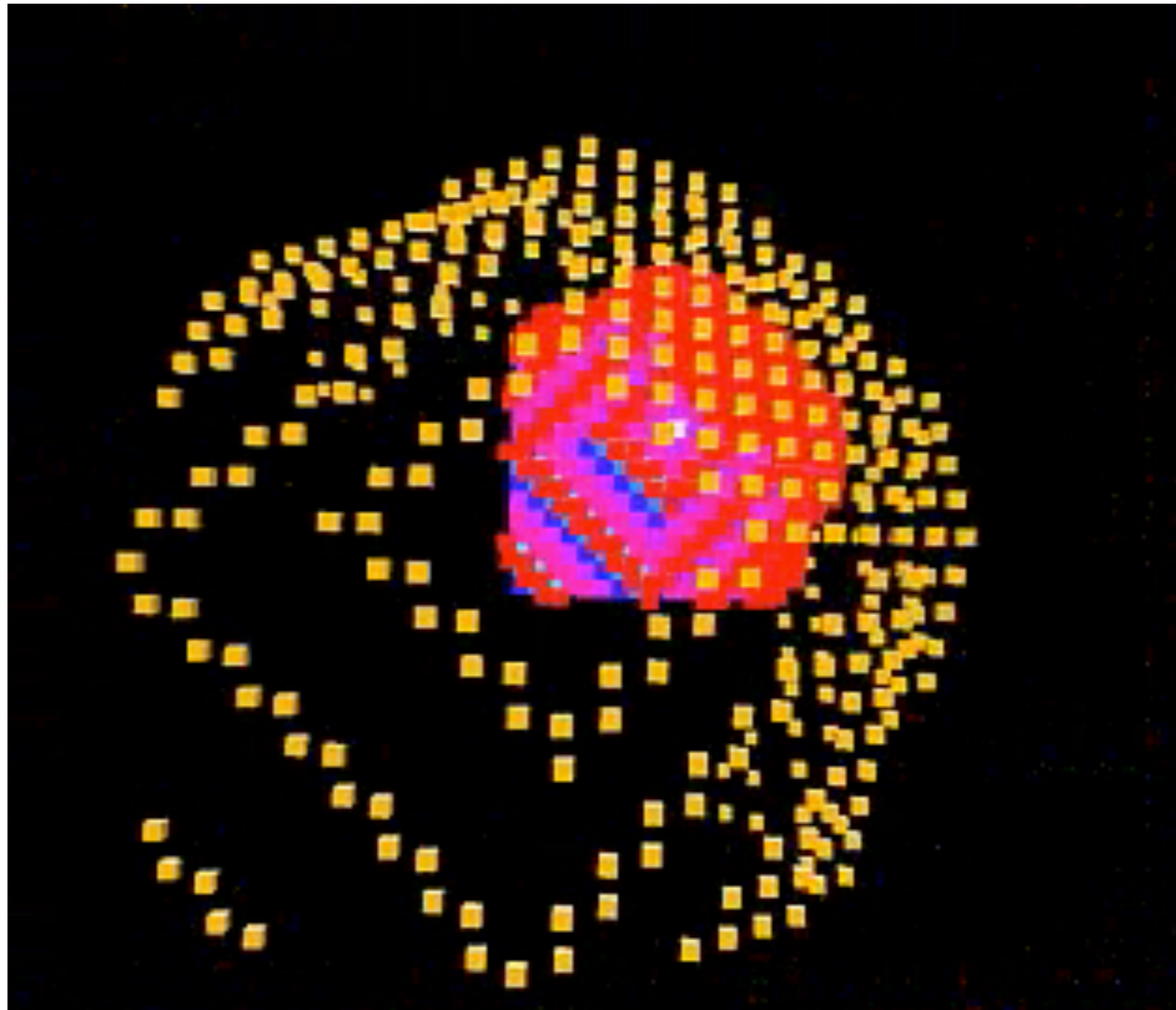
Key step:
determining the
successors from
the predecessors

Synchronizing the 3 skeleton branches



Use again the
on-time zone

Time optimal cube assembly



Time optimal cube assembly

