

# Commutative Diagrams

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## Exercise 1: Reasoning with Commutative Diagrams

Given a category  $\mathcal{C}$ , a *diagram* on  $\mathcal{C}$  is a directed graph with vertices labeled with objects of  $\mathcal{C}$  and edges labeled with morphisms of  $\mathcal{C}$  (with, if the source of edge  $e$  is labeled  $A$  and the target of  $e$  is labeled  $B$ , then the label of  $e$  is a morphism from  $A$  to  $B$ ). In other words, it is a directed graph equipped with a directed graph morphism into the underlying directed graph of  $\mathcal{C}$ .

A diagram is *commutative* if given any two vertices  $x$  and  $y$  and any two *non-empty* directed paths  $p$  and  $q$  from  $x$  to  $y$ , the two morphisms in  $\mathcal{C}$  obtained by composing the morphisms labeling the edges of  $p$  on one side, and those labeling the edges of  $q$  on the other side, are equal.

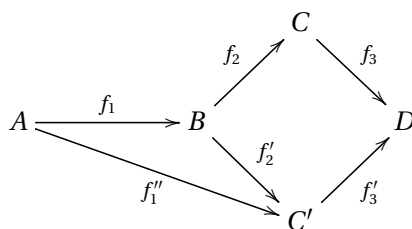
Building commutative diagrams gives a way to reason about equality of morphisms in categories which can be substituted to equational reasoning. As an example, here is a commutative square and the corresponding equation it describes on morphisms:

$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 g \downarrow & & \downarrow h \\
 A' & \xrightarrow{f'} & B'
 \end{array}
 \qquad
 \begin{array}{l}
 h \circ f = f' \circ g \\
 f ; h = g ; f'
 \end{array}$$

As another example, assuming we have morphisms  $f_1 : A \rightarrow B$ ,  $f_2 : B \rightarrow C$ ,  $f_3 : C \rightarrow D$ ,  $f'_2 : B \rightarrow C'$ ,  $f'_3 : C' \rightarrow D$ ,  $f''_1 : A \rightarrow C'$  satisfying:  $f_2 ; f_3 = f'_2 ; f'_3$  and  $f_1 ; f'_2 = f''_1$ . These two equations can be represented with the following two commutative diagrams:

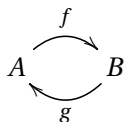
$$\begin{array}{ccc}
 B & \xrightarrow{f_2} & C \\
 f'_2 \downarrow & & \downarrow f_3 \\
 C' & \xrightarrow{f'_3} & D
 \end{array}
 \qquad
 \begin{array}{ccc}
 A & \xrightarrow{f_1} & B \\
 & \searrow f''_1 & \downarrow f'_2 \\
 & & C'
 \end{array}$$

By equational reasoning, we can deduce:  $f_1 ; f_2 ; f_3 = f_1 ; f'_2 ; f'_3 = f''_1 ; f'_3$ . The same result can be obtained by diagrammatic reasoning by considering the following commutative diagram:

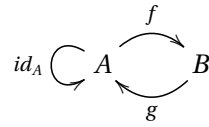


All this relies heavily on the associativity of composition in categories.

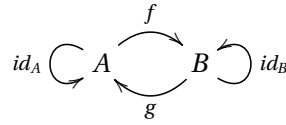
1. What are the equations induced by the following commutative diagram?



2. Same question with:

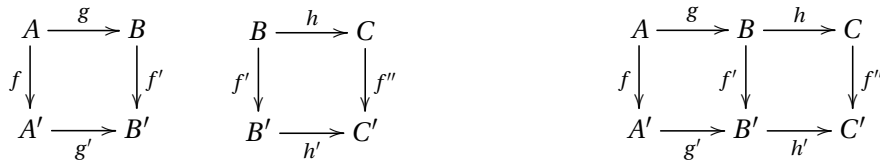


3. Same question with:

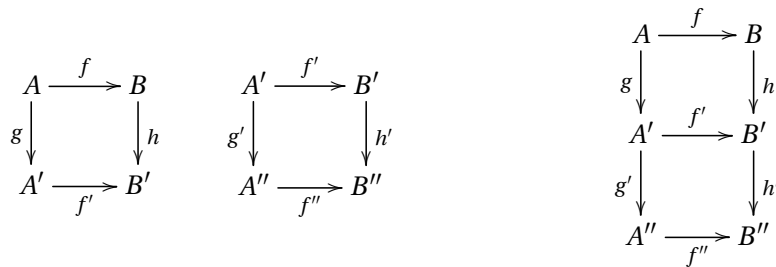


What property is it defining about  $f$  and  $g$ ?

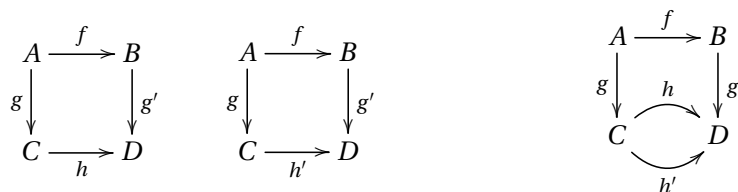
4. Prove that, in any commutative diagram  $\mathbb{D}$ , adding a *new* vertex  $v$  and a *new* edge  $e$  with source  $v$  and target a vertex of  $\mathbb{D}$  which is not in a cycle (resp. with target  $v$  and source a vertex of  $\mathbb{D}$  which is not in a cycle) gives another commutative diagram.
5. Prove that, in any commutative diagram with an edge  $e$  labeled  $f; g$  (and with source and target not belonging to cycles), adding a new vertex  $v$  and splitting  $e$  into an edge toward  $v$  labeled  $f$  and an edge from  $v$  labeled  $g$  gives another commutative diagram.
6. Prove that, in any commutative diagram, removing an edge gives another commutative diagram.
7. Prove that, in any commutative diagram, removing a node and all its adjacent edges gives another commutative diagram.
8. Prove that, in any commutative diagram containing a non-empty directed path  $p$  whose associated morphism is  $f$  (composition of the labels of the edges of  $p$ ), adding a new edge with same source as  $p$ , same target as  $p$  and label  $f$  gives another commutative diagram.
9. Prove that if the first two diagrams are commutative, then the third one is commutative as well:



10. Prove that if the first two diagrams are commutative, then the third one is commutative as well:



11. By choosing objects and morphisms in **Set**, prove that it is possible to have the first two diagrams commutative but not the third one:



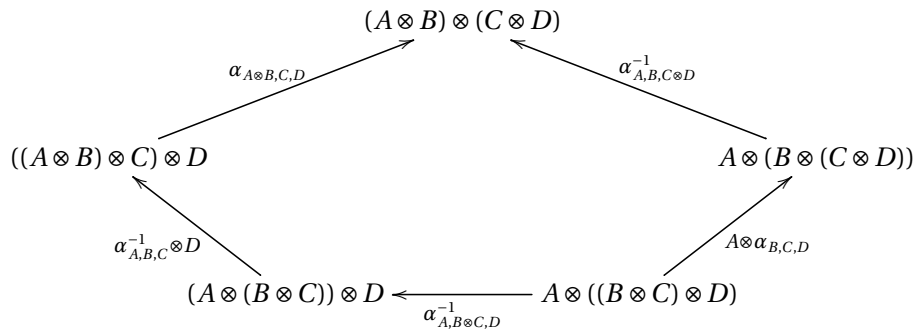
12. Given a commutative diagram in which all the labels of edges are isomorphisms, prove that the diagram obtained by reversing all edges and by turning each edge label  $\varphi$  into  $\varphi^{-1}$  is also commutative.
13. By using morphisms from  $\{0, 1\}$  to  $\{0, 1\}$  in **Set**, prove that given an isomorphism  $\varphi$ , the commutation of the first diagram does not imply the commutation of the second one:



14. Given a commutative diagram obtained exactly by identifying the two sources and the two targets of two disjoint non-empty directed paths (no other vertex or edge), and an edge  $e$  which is labeled by an isomorphism  $\varphi$ , prove that replacing  $e$  with the reverse edge labeled  $\varphi^{-1}$  gives a commutative diagram.

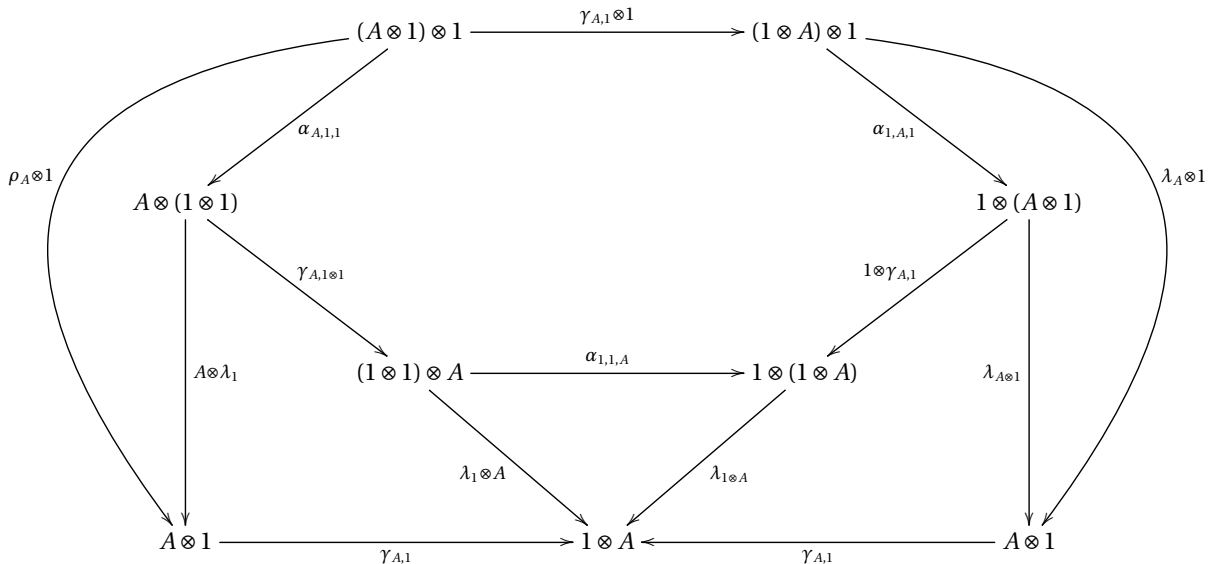
Note that it is in particular the case for any commutative acyclic triangle.

15. Given a monoidal category  $(\mathcal{C}, \lambda, \rho, \alpha)$ , prove the following diagram is commutative:



And give also an equational proof.

16. Justify the commutativity of the following diagram in any symmetric monoidal category  $(\mathcal{C}, \lambda, \rho, \alpha, \gamma)$ :



17. Deduce  $\rho_A = \gamma_{A,1}$ ;  $\lambda_A$  for any object  $A$  in any symmetric monoidal category  $(\mathcal{C}, \lambda, \rho, \alpha, \gamma)$ .

## Exercise 2:   **Preservation of Commutative Diagrams through Functors**

1. Prove the image of a commutative diagram under a functor is also a commutative diagram.
2. A morphism  $f$  is called an *epimorphism* if for any  $g$  and  $h$ ,  $f ; g = f ; h$  implies  $g = h$ .
  - (a) Prove that any isomorphism is an epimorphism.
  - (b) What are the epimorphisms of **Set**?
3. If  $F$  is an endofunctor and there exists a natural transformation from  $F$  to  $Id$  whose elements are epimorphisms, prove that the converse of Question 1 also holds: if the image under  $F$  of a diagram is commutative then the original diagram is commutative as well.
4. In any *monoidal* category, prove that  $f \otimes 1 = g \otimes 1$  implies  $f = g$ .