## HW 4: Error-correcting codes

(due December 4th, before tutorial)

- 1. Let  $A_q(n, d)$  be the largest k such that a code over alphabet  $\{1, \ldots, q\}$  of block length n, dimension k and minimum distance d exists (recall that this corresponds to the notation  $(n, k, d)_q$ ). Determine  $A_2(3, d)$  for all integers  $d \ge 1$ .
- 2. Suppose C is a  $(n, k, d)_2$ -code with d odd. Construct using C a code C' that is a  $(n + 1, k, d + 1)_2$ -code.
- 3. By constructing the columns of a parity check matrix in a greedy fashion, show that there exists a binary linear code  $[n, k, d]_2$  provided that

$$2^{n-k} > 1 + \binom{n-1}{1} + \dots + \binom{n-1}{d-2}.$$
(1)

This is a small improvement compared to the general Gilbert-Varshamov bound. In particular, it is tight for the  $[7, 4, 3]_2$  Hamming code.