## Quantum Information Theory

## Exercize sheet 1

(Dated: January 18, 2022)

1. (a) Let  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C$ , with  $\mathcal{H}_A, \mathcal{H}_B, \mathcal{H}_C \simeq \mathbb{C}^2$ , and  $|\psi\rangle \in \mathcal{H}$  with

$$|\psi\rangle = \frac{1}{\sqrt{2}} \left(|0\rangle_A \left|0\rangle_B \left|0\rangle_C + \left|1\rangle_A \left|1\rangle_B \left|0\rangle_C\right)\right.\right.$$
(1)

Compute the reduced density matrices  $\rho_A = \text{Tr}_{BC}(|\psi\rangle \langle \psi|), \rho_C = \text{Tr}_{AB}(|\psi\rangle \langle \psi|)$ . Show that  $\rho_A$  is a mixed state, while  $\rho_C$  is pure.

(b) Let  $\rho$  be a density operator. Prove that

$$\operatorname{Tr}[\rho^2] \le 1 \,, \tag{2}$$

with the equality if and only if  $\rho$  is a pure state.

- 2. Let  $\{M_n\}$  with  $\sum_m M_m^{\dagger} M_m = 1$  define a measurement on a system with associated Hilbert space  $\mathcal{H}_S$ . Show that it is always possible to realize such measurement by means of a unitary transformation in an enlarged Hilbert space, followed by a projective measurement.
- 3. (a) Show that the transpose map is positive but not completely positive. Hint: consider the state  $\rho_{AB} = |\psi\rangle_{AB} \langle \psi|_{AB}$ , with

$$|\psi\rangle = \frac{1}{\sqrt{2}} (|0\rangle_A |0\rangle_B + |1\rangle_A |1\rangle_B).$$
(3)

(b) [*PPT criterion*] Given  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$ , a density operator  $\rho \in \mathcal{H}$  is called separable if there exist  $p_k \ge 0$ and density operators  $\rho_k^A \in \mathcal{H}_A$ ,  $\rho_k^B \in \mathcal{H}_B$  such that

$$\rho = \sum_{k} p_k \rho_A^k \otimes \rho_B^k \,. \tag{4}$$

Otherwise  $\rho$  is called an entangled state. Show that if  $\rho$  is separable, then  $\rho^{T_B}$  and  $\rho^{T_A}$  are density operators.

4. (a) [Bloch-sphere interpretation of trace distance] Let  $\rho$ ,  $\sigma$  be qubit density operators, with

$$\rho = \frac{1}{2} \left( \mathbb{1} + \vec{r} \cdot \vec{\sigma} \right) \,, \qquad \sigma = \frac{1}{2} \left( \mathbb{1} + \vec{s} \cdot \vec{\sigma} \right) \,, \tag{5}$$

and  $\vec{\sigma} = (\sigma^x, \sigma^y, \sigma^z)$ . Prove

$$\Delta(\rho,\sigma) = \frac{|\vec{r} - \vec{s}|}{2} \,. \tag{6}$$

(b) Prove the following formula for the trace distance of two density matrices

$$\Delta(\rho, \sigma) = \max_{P} \left\{ \operatorname{Tr} \left[ P(\rho - \sigma) \right] \right\} \,, \tag{7}$$

where the maximization is taken either over all positive operators  $P \leq 1$ , or over all projectors P. Hint: show that for any states  $\rho$ ,  $\sigma$  there exist positive operators Q, S with support on orthogonal spaces such that

$$\rho - \sigma = Q - S \,. \tag{8}$$

(c) Prove the triangle inequality

$$\Delta(\rho, \sigma) \le \Delta(\rho, \tau) + \Delta(\tau, \sigma), \qquad (9)$$

where  $\rho$ ,  $\sigma$ ,  $\tau$  are arbitrary density matrices.

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- (d) Show that  $\Delta(\rho, \sigma) = 1$  if and only if  $\rho$  and  $\sigma$  have orthogonal support.
- (e) [Operational interpretation of trace distance] We know that a quantum system is prepared in the state  $\rho_0$  or  $\rho_1$ , with classical probability p = 1/2. Suppose that we choose a binary POVM  $\Lambda = {\Lambda_0, \Lambda_1}$  and measure the system. In the case we measure 0 we guess that the state of the system is  $\rho_0$ , otherwise we guess it is  $\rho_1$ . What is the probability that our guess is correct as a function of  $\Lambda$ ? Show that our optimal probability of success is

$$p_{\rm succ} = \frac{1}{2} \left( 1 + \Delta(\rho, \sigma) \right) \,. \tag{10}$$

What is the optimal probability in the case the system is prepared in three possible density matrices, with probability p = 1/3?

5. [Trace-distance contractivity]. Prove that for any trace preserving positive linear map  $\mathcal{E}$ 

$$\Delta(\mathcal{E}(\rho), \mathcal{E}(\sigma)) \le \Delta(\rho, \sigma) \,. \tag{11}$$

In particular, it follows that the trace distance is contractive for all quantum channels.

6. Show that for any M-code (E, D), there exists an M/2-code (E', D') such that

$$p_{\rm err,max}(E',D') \le 2p_{\rm err}(E,D).$$
(12)