

# Subject 2: Maps of event structures, and natural transformations

to be returned on the 25/10

## 1 Event structures and stable families

The required definitions for these exercises appear in the slides for lecture 3, available at <http://perso.ens-lyon.fr/pierre.clairambault/enseignement/CR07-slides-lecture3.pdf>.

**Question 1.** Let  $f : A \rightarrow B$  be a map of event structures.

(i) Show  $f$  locally reflects causal dependency, in the following sense. Let  $e, e' \in x$ , a configuration of  $A$ . Show if  $f(e) \leq_B f(e')$  (with both  $f(e)$  and  $f(e')$  defined) then  $e \leq_A e'$ .

(ii) Show  $f$  preserves the concurrency relation, *i.e.* if  $e \text{ co}_A e'$  and  $f(e)$  and  $f(e')$  are both defined, then  $f(e) \text{ co}_B f(e')$ .

**Question 2.** Let  $(A, \leq_A, \text{Con}_A), (B, \leq_B, \text{Con}_B)$  be event structures. Let  $f : A \rightarrow B$  be a partial function. Show  $f$  is a map of event structures,  $f : (A, \leq_A, \text{Con}_A) \rightarrow (B, \leq_B, \text{Con}_B)$ , iff

(i)  $\forall a \in A, b \in B. b \leq_B f(a) \implies \exists a' \in A. a' \leq_A a \ \& \ f(a') = b$ , and

(ii)  $\forall X \in \text{Con}_A. fX \in \text{Con}_B \ \& \ \forall a_1, a_2 \in X. f(a_1) = f(a_2) \implies a_1 = a_2$ .

**Question 3.** Verify that the finite configurations of an event structure form a stable family.

**Question 4.** Say an event structure  $A$  is *tree-like* when its concurrency relation is empty (so two distinct events are either causally related or inconsistent).

Suppose  $B$  is tree-like and  $f : A \rightarrow B$  is a total map of event structures (a map of event structures is *total* if the underlying function between events is total). Show  $A$  must also be tree-like, and moreover that the map  $f$  is rigid, *i.e.* preserves causal dependency.

**Question 5.** Show a map  $f : A \rightarrow B$  of event structures is *stable*, *i.e.*

$$x \uparrow y \implies f(x \cap y) = (fx) \cap (fy).$$

**Question 6.** Let  $\mathcal{F}$  be a nonempty family of finite sets satisfying the Completeness axiom in the definition of stable families. Show  $\mathcal{F}$  is coincidence-free iff

$$\forall x, y \in \mathcal{F}. x \subsetneq y \implies \exists x_1, e_1. x \overset{e_1}{\subset} x_1 \subseteq y.$$

By  $x \overset{e_1}{\subset} x_1$ , for configurations  $x, x_1$  and event  $e_1$ , is meant that  $e_1 \notin x \ \& \ x \cup \{e_1\} = x_1$ . [Hint: For ‘only if’ use induction on the size of  $y \setminus x$ .]

**Question 7.** Let  $f : \mathcal{F} \rightarrow \mathcal{G}$  be a map of stable families. Let  $e, e' \in x$ , a configuration of  $\mathcal{F}$ . Show if  $f(e) \leq_{fx} f(e')$  (with both  $f(e)$  and  $f(e')$  defined) then  $e \leq_x e'$ .

(the definition of  $\leq_x$  for  $x \in \mathcal{F}$  in a stable family  $\mathcal{F}$  appears on slide 17)

## 2 Natural transformations and equivalence of categories

**Question 8.** Which of the following functions correspond to natural transformations?

(a) Is this a natural transformation between functors from **Set** to **Set**?

$$\begin{aligned} \text{fst}_A : \quad \text{List}(A) &\rightarrow A + \{*\} \\ [a_1, \dots, a_n] &\mapsto (1, a_1) \\ [] &\mapsto (2, *) \end{aligned}$$

where  $A + B = (\{1\} \times A) \cup (\{2\} \times B)$  is the disjoint union.

(b) Is this a natural transformation between functors from **Set** to **Set**?

$$\begin{aligned} \text{equal}_A : \quad \text{List}(A) &\rightarrow \mathbb{B} \\ l &\mapsto \text{true} \quad \text{if all elements of } l \text{ are equal} \\ &\quad \text{false} \quad \text{otherwise} \end{aligned}$$

(note: an object  $A \in \mathcal{C}$  may be regarded as the constant functor  $\mathcal{C} \rightarrow \mathcal{C}$  which to any  $X \in \mathcal{C}$  associates  $A$ , and to any  $f : X \rightarrow Y$  associates  $\text{id}_A$ ).

(c) Is this a natural transformation between functors from **Set** to **Set**?

$$\begin{aligned} \text{length}_A : \quad \text{List}(A) &\rightarrow \mathbb{N} \\ [] &\mapsto 0 \\ [a_1, \dots, a_n] &\mapsto n \end{aligned}$$

(d) Is this a natural transformation between functors from **Ford** to **Set**?

$$\begin{aligned} \text{sort}_A : \quad \text{List}(A) &\rightarrow \text{List}(A) \\ l &\mapsto \text{sort}(l) \end{aligned}$$

where **Ford** is the category of totally ordered finite sets and order-preserving functions? What if the morphisms in **Ford** were not required to preserve the order?

**Question 9.** Let  $F, G : \mathcal{C} \rightarrow \mathcal{D}$  be functors, and  $\eta : F \Rightarrow G$  be a natural transformation. Show that  $\eta$  has an inverse iff for all  $A \in \mathcal{C}$ ,  $\eta_A$  has an inverse.

**Question 10.** Consider the category **Fam** having

- *Objects:* families of sets  $(A_i)_{i \in I}$  (where  $I \in \mathbf{Set}$ , and for all  $i \in I$ ,  $A_i \in \mathbf{Set}$ ).
- *Morphisms* from  $(A_i)_{i \in I}$  to  $(B_j)_{j \in J}$ : pairs  $(f, (g_i)_{i \in I})$  where  $f : I \rightarrow J$ , and for all  $i \in I$ ,  $g_i : A_i \rightarrow B_{f(i)}$ .

Show that **Fam** is equivalent to the presheaf category  $[1 \rightarrow 2, \mathbf{Set}]$ , where  $1 \rightarrow 2$  is the category having two objects 1, 2 and only one morphism from 1 to 2 (plus identities).

**Question 11.** Show that if  $F : \mathcal{C} \rightarrow \mathcal{D}$  is part of an equivalence of categories, then it is full, faithful and essentially surjective on objects. [Hint: prove faithful first]

**Question 12.** Show that equivalence of categories is an equivalence relation [you may use results in the lecture].