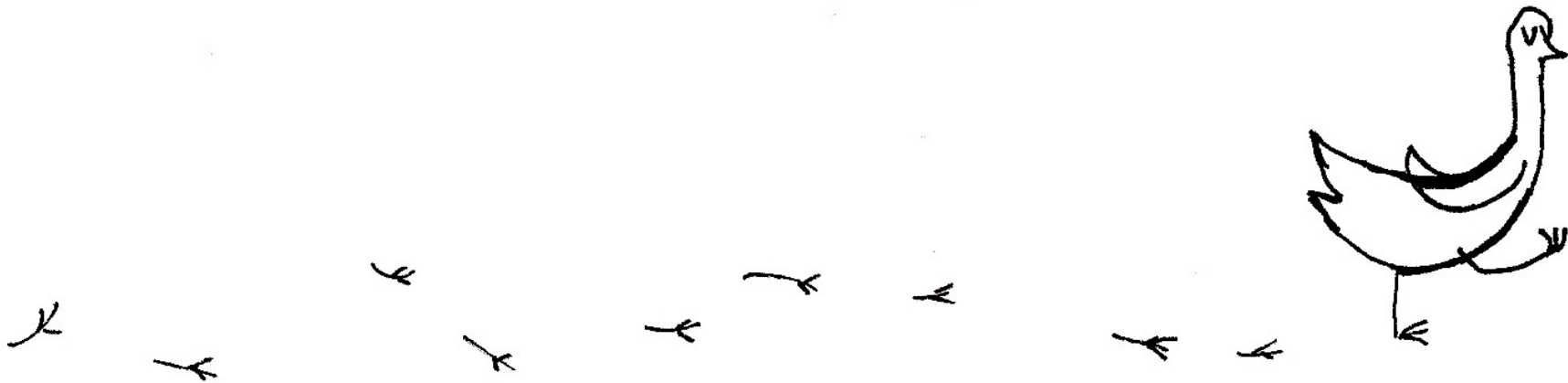


Q-categories as generalized domains

Paweł Waszkiewicz



Ralph Koppelman's advice

MR. MCQUIRE

Ben - I just want to say one word to
you - just one word -

BEN

Yes, sir.

MR. MCQUIRE

Are you listening?

BEN

Yes I am.

MR. MCQUIRE

(gravely)

Plastics.

They look at each other for a moment.

BEN

Exactly how do you mean?

MR. MCQUIRE

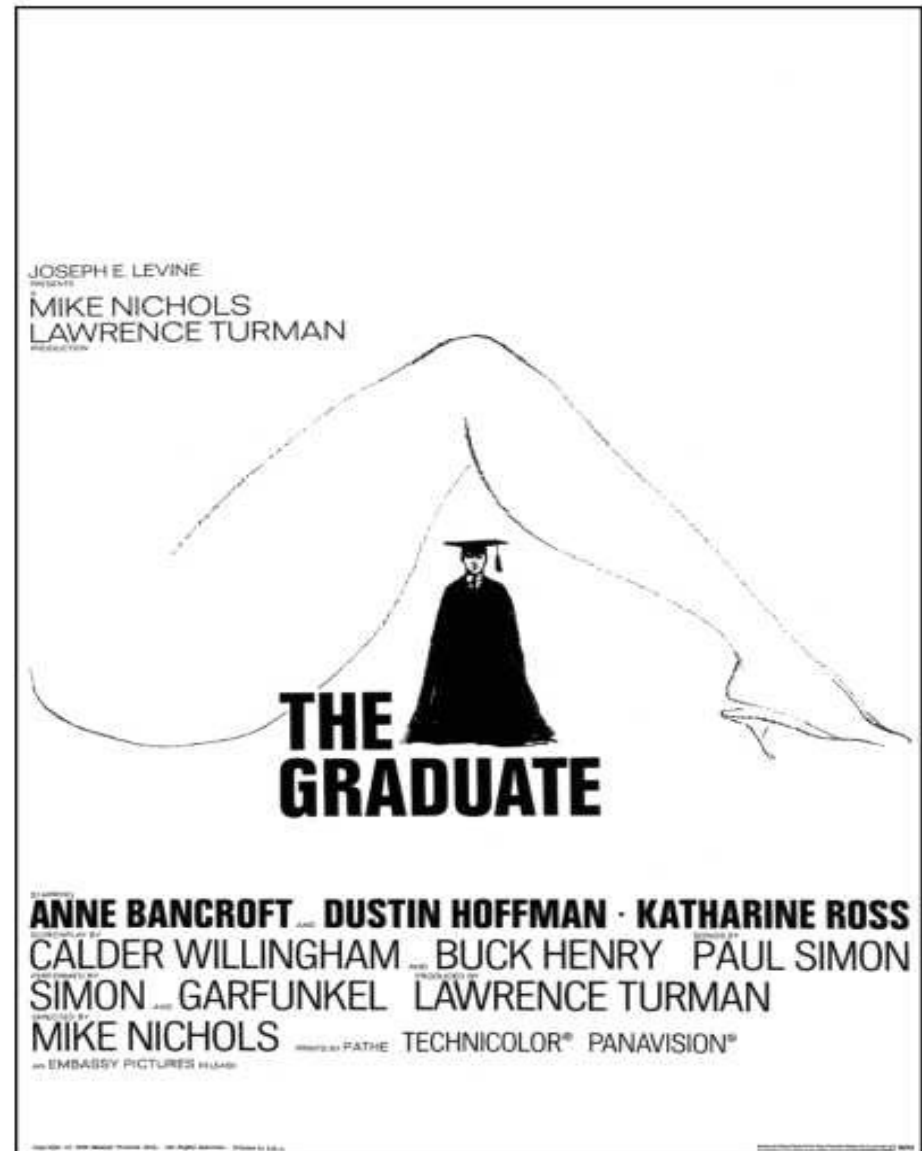
There is a great future in plastics.
Think about it. Will you think
about it?

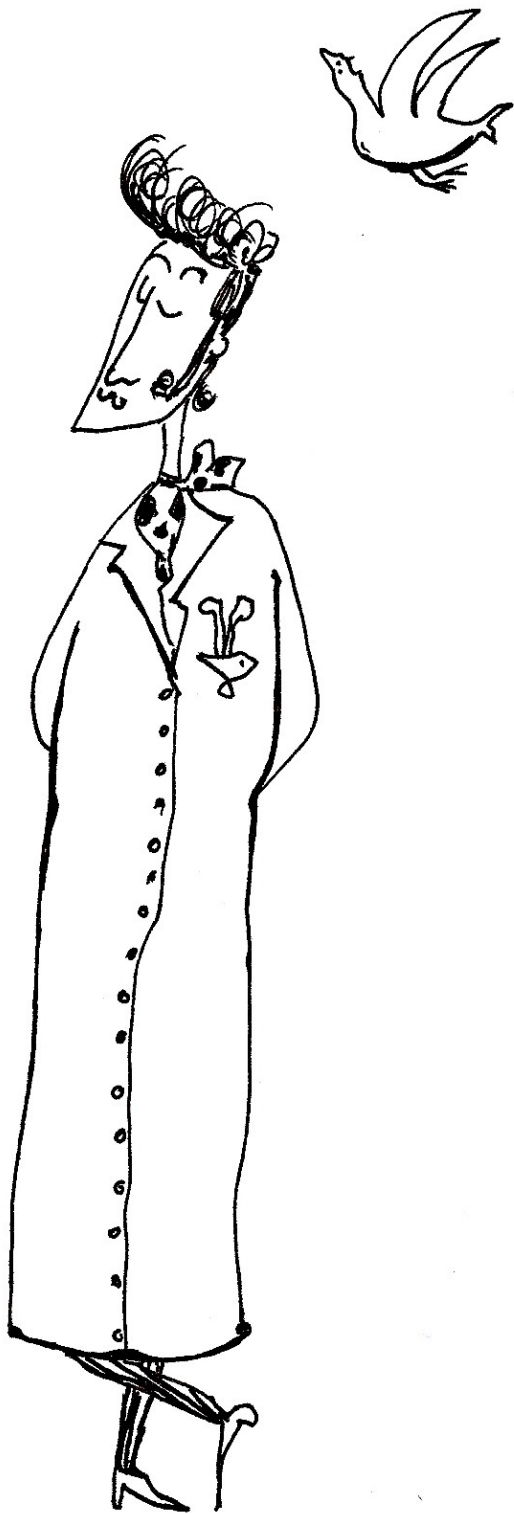
BEN

Yes, I will.

MR. MCQUIRE

Okay. Enough said. That's a deal.





Quasi-uniformities

$\mathcal{U} \subseteq \mathcal{P}(X \times X)$ is a quasi-uniformity if:

- $\forall V \in \mathcal{U} \quad \Delta \subseteq V,$
- $\forall U \in \mathcal{U} \quad \exists V \in \mathcal{U} \quad V \cdot V \subseteq U.$

EXAMPLES

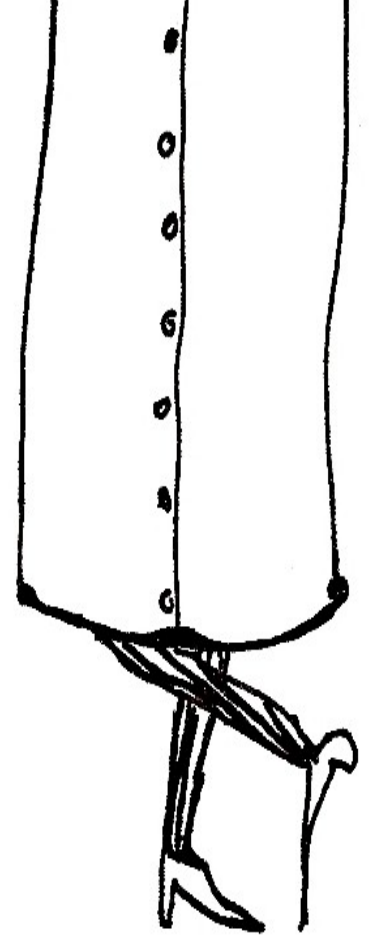
EXAMPLE 1. $\mathcal{U} = \{A\}$ iff A is a preorder.

EXAMPLE 2. In a metric space X , for

$$V_\varepsilon = \{(x, y) \mid X(x, y) < \varepsilon\}$$

define

$$\mathcal{U} = \{V_\varepsilon \mid \varepsilon > 0\}.$$



Quasi-uniform spaces as domains of computation

- [Smy88] Smyth, M.B. (1988) Quasi-uniformities: Reconciling Domains and Metric Spaces. *Lecture Notes in Computer Science* **298**, pp. 236–253.
- [Smy91] Smyth, M.B. (1991) Totally bounded spaces and compact ordered spaces as domains of computation. In G.M. Reed, A. W. Roscoe, and R. F. Wachter, editors, *Topology and Category Theory in Computer Science*, pp. 207–229. Clarendon Press.
- [Smy94] Smyth, M.B. (1994) Completeness of quasi-uniform and syntopological spaces. *Journal of the London Mathematical Society* **49**, pp. 385–400.

Lawvere's famous 1973 paper

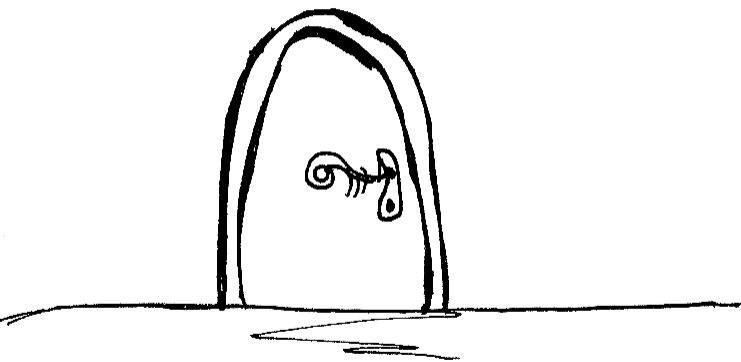


Ordered sets and metric spaces are \mathbf{Q} -enriched cats

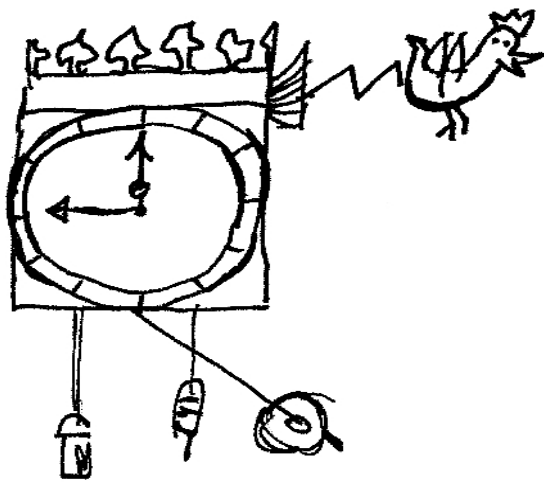
Lawvere, F.W. (1973) Metric spaces, generalized logic, and closed categories. Rend. Sem. Mat. Fis. Milano 43, pp. 135—166.

I WILL:

- Introduce $\mathcal{Q} = (Q, \leq, \otimes, 1)$
- Introduce $\mathcal{Q}\text{-Rel}$
- Introduce $\mathcal{Q}\text{-Cat}$



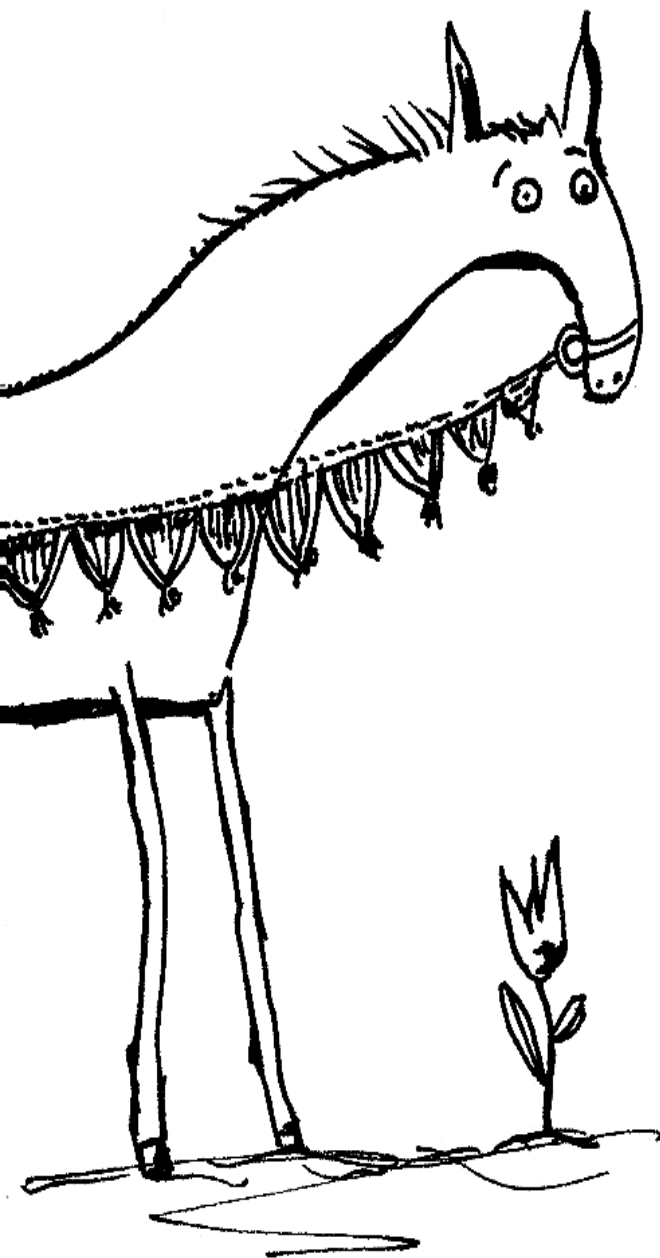
QUANTALES



A quantale is something that resembles non-negative real numbers.

$$\mathcal{Q} = (Q, \leq, \otimes, 1)$$

- Q is a complete lattice
- has addition \otimes
- addition has unit 1
- $a \otimes \bigvee S = \bigvee \{a \otimes s \mid s \in S\}$



The category $\mathcal{Q}\text{-Rel}$:

- objects: sets X, Y, Z, \dots
- morphisms: $r : X \multimap Y$ is just a function $r : X \times Y \rightarrow \mathcal{Q}$
- composition:
for $r : X \multimap Y$ and $s : Y \multimap Z$

$$(s \cdot r)(x, z) = \bigvee_{y \in Y} r(x, y) \otimes s(y, z).$$

There is a functor $\mathbf{Set} \rightarrow \mathcal{Q}\text{-Rel}$ which maps objects identically and interprets a map $f : X \rightarrow Y$ as a \mathcal{Q} -relation $f : X \multimap Y$:

$$f(x, y) = \begin{cases} \mathbf{1} & \text{if } f(x) = y, \\ \perp & \text{otherwise.} \end{cases}$$

For example...

- 2-Rel is isomorphic to the category of relations
- $[0,1]$ -Rel is isomorphic to the category of fuzzy relations



A Q -category

$$X = (X, X(-, -))$$

is a set X with a Q -relation $X : X \multimap X$ satisfying:

$$1_X \leq X \text{ (reflexivity)}$$

$$X \cdot X \leq X \text{ (transitivity).}$$

A Q -functor $f : X \rightarrow Y$ must satisfy

$$f \cdot X \leq Y \cdot f.$$

2-Cat is the category **Ord** of (pre)orders.

$[0, \infty]$ -**Cat** is the category of (pre)metric spaces.

For the trivial quantale: **1-Cat** = **Set**.



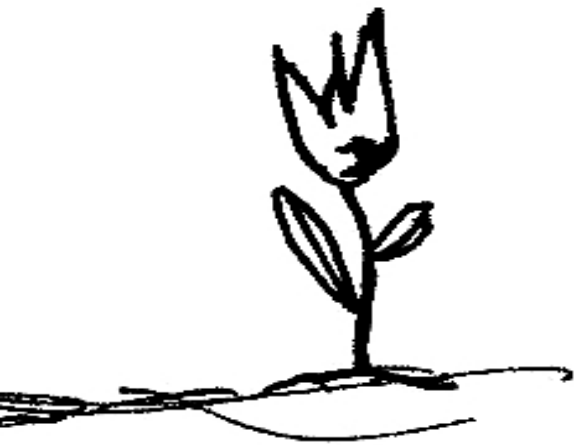
The category of \mathcal{Q} -categories

$\mathcal{Q}\text{-Cat}$ is symmetric monoidal closed with tensor product

$$X \otimes Y((x, y), (a, b)) = X(x, a) \otimes Y(y, b)$$

and internal hom:

$$Y^X(f, g) = \bigwedge_{x \in X} Y(fx, gx).$$



The internal hom describes the pointwise order if $\mathcal{Q} = \mathbf{2}$, and the non-symmetrized sup-metric if $\mathcal{Q} = [0, \infty]$ or $\mathcal{Q} = [0, 1]$.

The Amsterdam Group

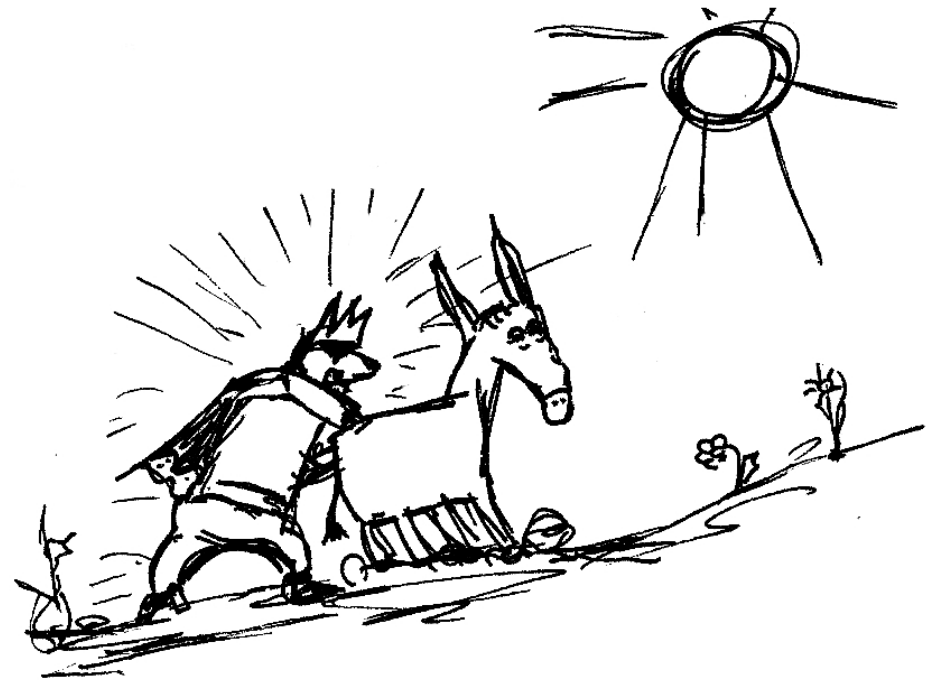
- [AR89] America, P. and Rutten, J.J.M.M. (1989) Solving reflexive domain equations in a category of complete metric spaces, *Journal of Computer and System Sciences* **39**(3), pp. 343–375.
- [BvBR96] Bonsangue, M. M., van Breugel, F. and Rutten, J. J. M. M. (1996) Alexandroff and Scott Topologies for Generalized Metric Spaces, *Annals of the New York Academy of Sciences*, pp. 49–68.
- [BvBR98] Bonsangue, M.M., van Breugel, F. and Rutten, J.J.M.M. (1998) Generalized Metric Spaces: Completion, Topology, and Powerdomains via the Yoneda Embedding, *Theoretical Computer Science* **193**(1-2), pp. 1–51.
- [Rut96] Rutten, J.J.M.M. (1996) Elements of generalized ultrametric domain theory, *Theoretical Computer Science* **170**, pp. 349–381.
- [Rut98] Rutten, J.J.M.M. (1998) Weighted colimits and formal balls in generalized metric spaces, *Topology and its Applications* **89**, pp. 179–202.

My research idea is to...

- phrase domain theory in the language of 2 -relations and...
- ...change 2-relations to arbitrary Q -relations...
- ...to obtain results for the general case.
- There are a lot of things that can go wrong, but in most cases it works!



Joint work



- (with Dirk Hofmann)

We introduce
continuous Q-cats

- (with Mateusz Kostanek)

The Bilimit Theorem: every expanding sequence
of Q-cats has a bilimit

- Theorem. $\mathcal{Q}\text{-Cat}$ is cartesian closed iff $\otimes = \wedge$

A \mathcal{Q} -distributor $\varphi: X \multimap Y$ is a \mathcal{Q} -relation $\varphi: X \multimap Y$ with

$$\varphi \cdot X = \varphi = Y \cdot \varphi.$$

The \mathcal{Q} -distributor $X: X \multimap X$ plays the role of the identity in the category $\mathcal{Q}\text{-Dist}$ of sets and \mathcal{Q} -distributors.

For a \mathcal{Q} -functor $f: X \rightarrow Y$, both

$$f_* = Y \cdot f \quad \text{and} \quad f^* = f^\circ \cdot Y$$

are \mathcal{Q} -distributors.



The following are equivalent for \mathcal{Q} -relations $\varphi: X \multimap Y$ between \mathcal{Q} -categories.

- $\varphi: X \multimap Y$ is a \mathcal{Q} -distributor.
- $\varphi: X^{\text{op}} \otimes Y \rightarrow \mathcal{Q}$ is a \mathcal{Q} -functor.

EXAMPLE. For $\mathcal{Q} = \mathbf{2}$, a \mathcal{Q} -distributor from X to 1 corresponds to a monotone map from X^{op} to $\mathbf{2}$, that is, to a lower set in X .



Let \widehat{X} denote the \mathcal{Q} -category of
all \mathcal{Q} -distributors $X \multimap 1$
(i.e. all \mathcal{Q} -functors $X^{op} \rightarrow \mathcal{Q}$):

$$\widehat{X}(\phi, \psi) = \bigwedge_{x \in X} \mathcal{Q}(\phi x, \psi x).$$

X embeds into \widehat{X} via the
Yoneda \mathcal{Q} -functor:

$$X \ni x \mapsto yx := x^* \in \widehat{X}$$

DEFINITION. A \mathcal{Q} -category X
is **COCOMplete**
iff $y: X \rightarrow \widehat{X}$ has a left adjoint
 $S: \widehat{X} \rightarrow X$:

$$S \dashv y.$$



Relative Cocompleteness

We consider any subcategory J of $\mathcal{Q}\text{-Dist}$ such that for all \mathcal{Q} -functors f :

$$f^* \in J$$

and

$$(X \xrightarrow{\varphi} Y \xrightarrow{y^*} 1) \in J \Rightarrow (X \xrightarrow{\varphi} Y) \in J.$$

Define

$$J(X) = \{\varphi: X \multimap 1 \mid \varphi \in J\}$$

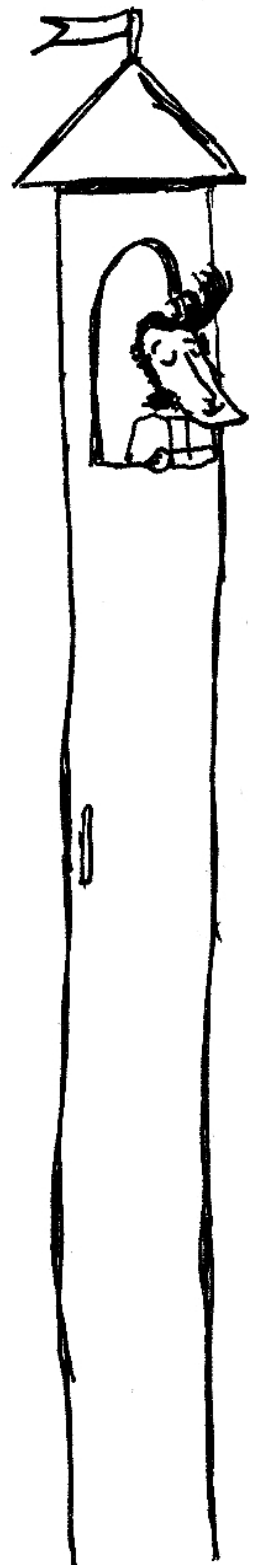
and

$$J_S(X) = \{\varphi \in J(X) \mid \varphi \text{ has a supremum}\}$$

A \mathcal{Q} -category X is

iff **J-COCOMplete**

$$\mathcal{S}: J_S(X) \rightarrow X \dashv y: X \rightarrow J_S(X).$$



EXAMPLE 1. Choose $J = \mathcal{Q}\text{-Dist}$.
 J -cocomplete means cocomplete.

EXAMPLE 2. For $\mathcal{Q} = \mathbf{2}$, choose

$$J = \{\phi: X \multimap 1 \mid \phi \text{ is an ideal}\}.$$

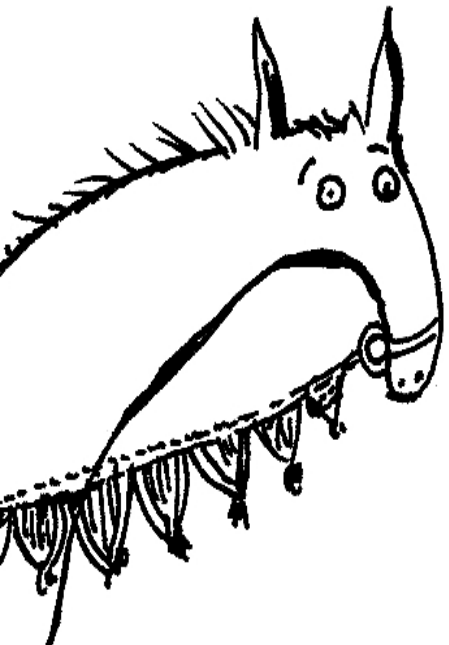
J -cocomplete means directed-complete.

EXAMPLE 3. For $\mathcal{Q} = [0, \infty]$, choose

$$J = \{\inf_i \sup_{j \geq i} X(-, x_j) \mid (x_i) \text{ is Cauchy}\}$$

J -cocomplete means metrically complete.

etcetera



A \mathcal{Q} -functor $f: X \rightarrow Y$ is *Scott-continuous* if for all $\varphi \in J_S(X)$,

$$f(\mathcal{S}\varphi) = \mathcal{S}(\varphi \cdot f^*).$$

PROPOSITION. Let Y be a J -cocomplete category. Then any \mathcal{Q} -functor $f: X \rightarrow Y$ uniquely extends to a Scott-continuous \mathcal{Q} -functor $F: J(X) \rightarrow Y$.

THEOREM. The inclusion functor

$$\mathbf{J-Cocont} \rightarrow \mathbf{\mathcal{Q}\text{-Cat}}$$

has a left adjoint which sends a \mathcal{Q} -category X to $J(X)$ and a \mathcal{Q} -functor f to $(-)\cdot f^*$.

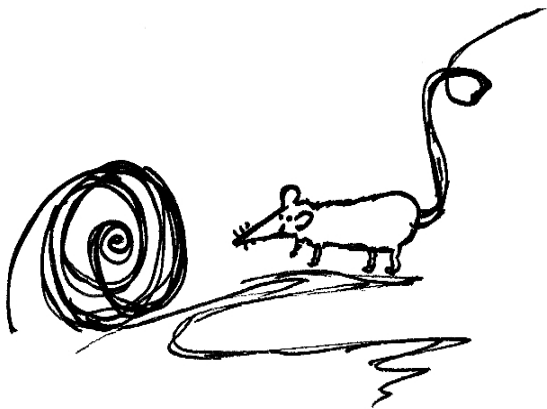


A fixpoint theorem

THEOREM. Let X be J -cocomplete, where $J = \{\bigvee_i \bigwedge_{j \geq i} X(-, x_j) \mid (x_i) \text{ is Cauchy}\}$. Let $f: X \rightarrow X$ be a \mathcal{Q} -functor. If

$$\phi \cdot f^* = \phi$$

for some \mathcal{Q} -distributor $\phi: X \multimap 1$ such that $\phi \in J$, then f has a fixed point.



PROOF: Define $x \leq y$ iff $1 = X(x, y)$. By J -cocompleteness, $\mathcal{S}\phi$ exists. Thus:

$$\mathcal{S}\phi = \mathcal{S}(\phi \cdot f^*) \leq f(\mathcal{S}\phi) \leq f^2(\mathcal{S}\phi) \leq \dots$$

However, by the choice of J , (X, \leq) is directed-complete, and f is monotone wrt \leq , so f has a fixpoint by the usual argument using transfinite induction.

Banach fixpoint theorem

Suppose X is a complete (quasi-)metric space. Let $f: X \rightarrow X$ be a contraction. Let $x \in X$. Then:

$$\phi = \bigvee_i \bigwedge_{j \geq i} X(-, f^j x)$$

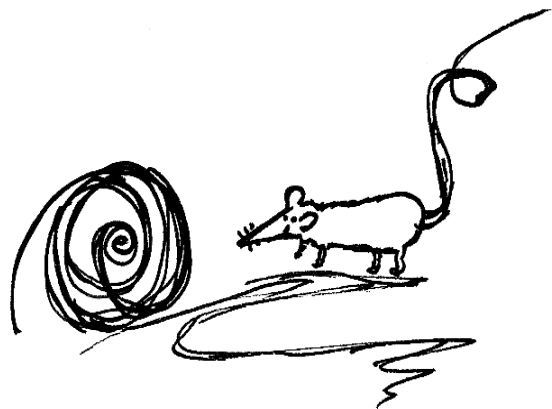
is an element of J , and satisfies:

$$\phi \cdot f^* = \phi.$$

Moreover,

$$f(\mathcal{S}\phi) = \mathcal{S}(\phi \cdot f^*) = \mathcal{S}\phi$$

is the fixed point. It is unique, since f is a contraction.



Corollary. A contraction on a complete metric space has a unique fixed point.

