

Graphs for Relation Algebra

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Binary relations

$$R, S \in \mathcal{P}(X^2)$$

$$0 \triangleq \emptyset$$

$$1 \triangleq \{\langle x, x \rangle \mid x \in X\}$$

$$\top \triangleq X^2$$

$$R^\circ \triangleq \{\langle x, y \rangle \mid y R x\}$$

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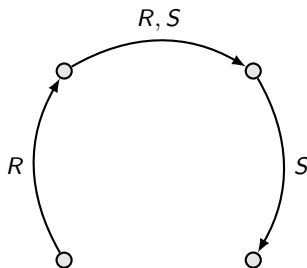
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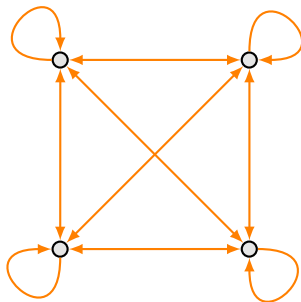
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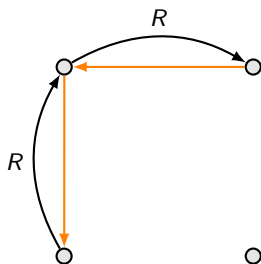
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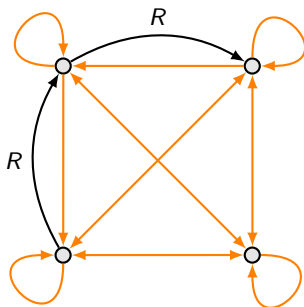
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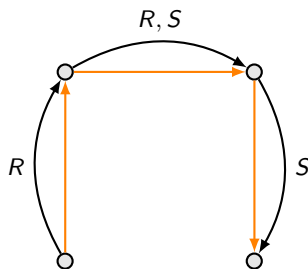
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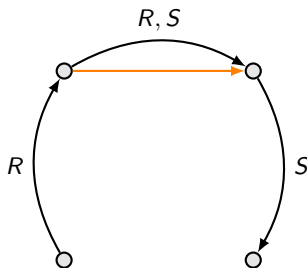
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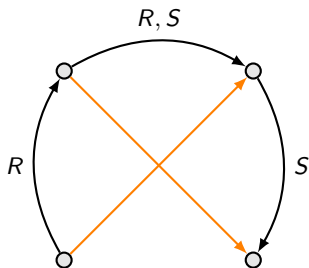
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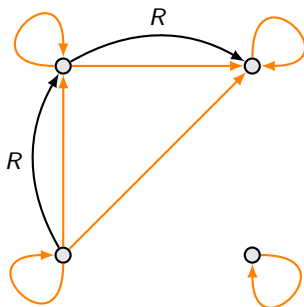
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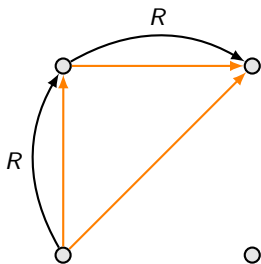
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Abstract properties

Many properties of relations can be expressed algebraically, without mentioning points:

$$1 \subseteq R$$

$$R \cdot R \subseteq R$$

$$R^+ \cap 1 = 0$$

$$R^0 \cdot R \subseteq R^* \cdot R^{0*}$$

Abstract properties

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local confluence

Laws

Many (in)equations hold universally:

- Intuitive ones:

$$R \cup S = S \cup R$$
$$R \cdot (S \cdot P) = (R \cdot S) \cdot P$$

...

$\langle \cup, \cap, \cdot^c, 0, \top \rangle$ form a Boolean lattice
 $\langle \cdot, \cup, 1, 0 \rangle$ form an idempotent semiring

- Stranger ones:

$$R \cdot S \cap P \subseteq (R \cap P \cdot S^{\circ}) \cdot (S \cap R^{\circ} \cdot P)$$
$$(R \cap S \cdot T) \cdot P = R \cdot P \cap S \cdot T$$
$$(R \cup S)^* = R^* \cdot (S \cdot R^*)^*$$

...

Relation algebra

1. Can we decide the validity of an equation?
2. At which cost?

With complement

1. No
2. -

Merci Gödel

Kleene algebra

Consider the fragment with operations \cdot , \cup , \cdot^* , 1 , 0 ,
i.e., regular expressions

Theorem. *An equation is valid iff the corresponding regular expressions recognise the same language.*

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Theorem. *An equation is valid iff the corresponding regular expressions recognise the same language.*

\leftrightarrow The problem is thus PSPACE-complete in this fragment

Allegories

Now consider the fragment with operations \cdot , \cap , \cdot° , 1 , \top ,
called **allegories in the literature**

The previous theorem fails

Valid laws include

$$a \cdot b \cap c \subset (a \cap c \cdot b^\circ) \cdot b$$

$$a \cdot (b \cap c) \subset a \cdot b \cap a \cdot c$$

$$(a \cap b \cdot \top) \cdot c = a \cdot c \cap b \cdot \top$$

Graphs, at last

Call **graph** a directed edge-labelled graph with two distinguished vertices called input and output

Define the following operations on such graphs

$$G \cdot H \triangleq \text{---} \circ \text{---} G \text{---} \circ \text{---} H \text{---} \circ \text{---} \text{---}$$

$$G \cap H \triangleq \text{---} \circ \begin{array}{c} \nearrow G \\ \searrow H \end{array} \circ \text{---}$$

$$G^\circ \triangleq \text{---} \circ \text{---} G \text{---} \circ \text{---}$$

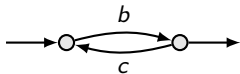
$$\underline{1} \triangleq \text{---} \circ \text{---}$$

$$\underline{\perp} \triangleq \text{---} \circ \quad \circ \text{---}$$

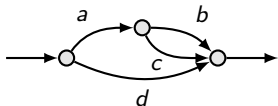
$$\underline{a} \triangleq \text{---} \circ \xrightarrow{a} \circ \text{---}$$

Graphs of expressions

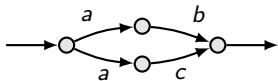
$$G(b \cap c^{\circ})$$



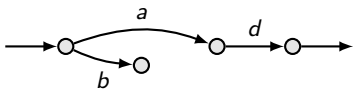
$$G(a \cdot (b \cap c) \cap d)$$



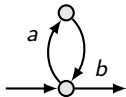
$$G(a \cdot b \cap a \cdot c)$$



$$G((a \cap b \cdot \top) \cdot d)$$



$$G(a \cdot b \cap 1)$$



Graph homomorphisms

Theorem. *An inequation $u \subseteq v$ is valid iff there exists a graph homomorphism from $G(v)$ to $G(u)$.*

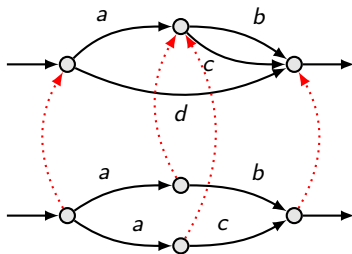
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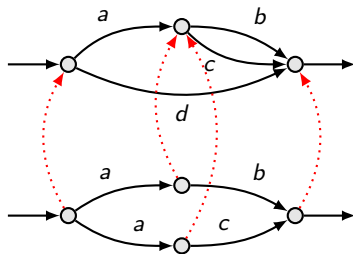
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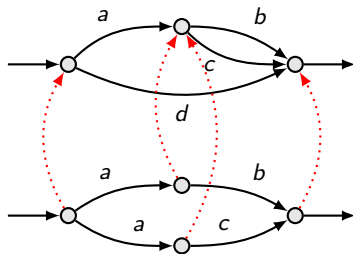
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- is it NP-hard?
- is it in P?

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Theorem. *An equation is valid in Kleene allegories iff the corresponding languages of graphs coincide.*

- “Petri automata” for recognising those languages
- Comparison of such automata is EXPSPACE-complete
... in a fragment where graphs are series-parallel