

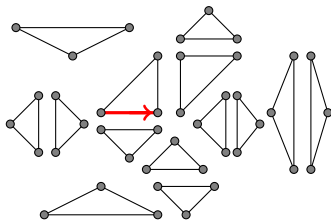
# Markovian triangulations and robust convergence to the UIPT

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Journée Combitop, ENS Lyon

- The UIPT is a natural "uniform" model of discrete, infinite planar geometry:
  - obtained by convergence of finite models relying on enumerative combinatorics;
  - nice to study because of its *Spatial Markov property*.
- Goals:
  - classify infinite objects exhibiting a similar Markov property;
  - use this to prove the convergence of finite models to the UIPT in a "robust" way, without precise enumeration.

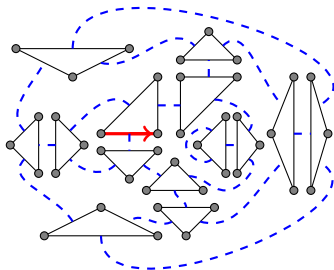
# Triangulations of the sphere



- A *triangulation of the sphere* with  $2n$  faces is a set of  $2n$  triangles whose sides have been glued two by two, in a way that is homeomorphic to the sphere.
- Our triangulations are *of type I* (we may glue two sides of the same triangle), and *rooted* (oriented root edge).
- Exact, explicit enumeration [Tutte 60s]:

$$\#\mathcal{T}_n = 2 \frac{4^n (3n)!!}{(n+1)!(n+2)!!}.$$

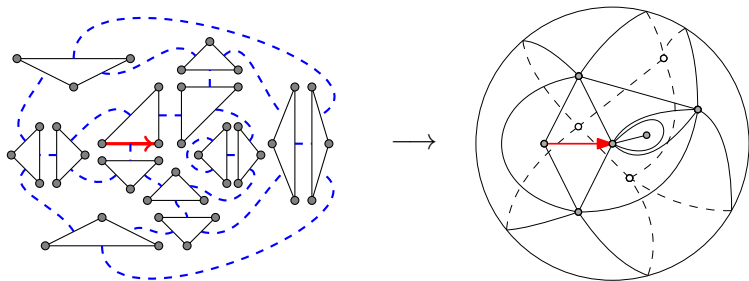
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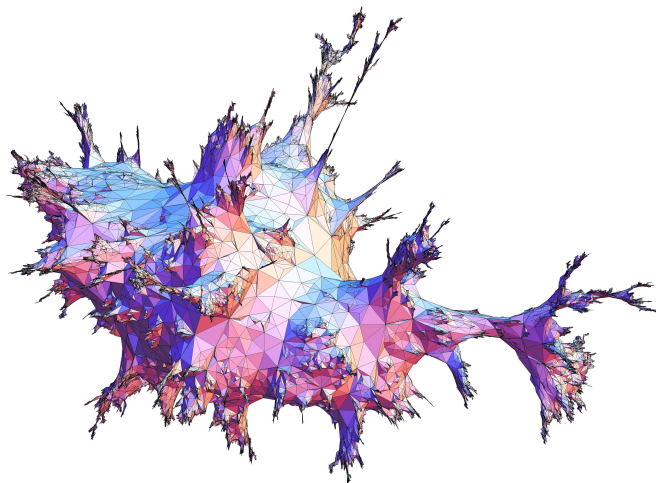


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# Random triangulations of the sphere

- Let  $T_n$  be a uniform triangulation of the sphere with  $2n$  faces. What does  $T_n$  look like?

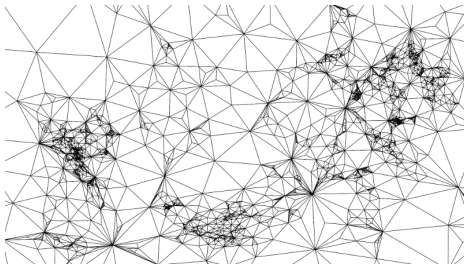


# Local limits of uniform triangulations

- *Local convergence:*

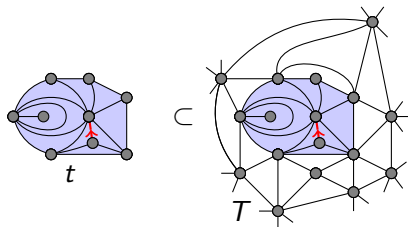
$$d_{\text{loc}}(t, t') = (1 + \max \{r \geq 0 \mid B_r(t) = B_r(t')\})^{-1}.$$

- Then  $T_n$  converges in distribution for the local topology to a random triangulation of the plane called the UIPT (Uniform Infinite Planar Triangulation)  $\mathbb{T}$  [Angel–Schramm 03].



# The argument of Angel and Schramm

- Let  $t$  be a small triangulation with a hole:



- Then  $\mathbb{P}(t \subset T_n) = \frac{\text{\#ways to fill the hole}}{\#T_n}$  depends on the perimeter and volume of  $t$ . It is explicit by the enumeration of Tutte, and converges as  $n \rightarrow +\infty$ .



# The argument of Angel and Schramm

- Before that:

- Tightness (control vertex degree): uses that  $\frac{\#\mathcal{T}_{n+1}}{\#\mathcal{T}_n}$  is bounded.
- One-endedness, i.e. in the limit there is no finite set separating two infinite regions: uses that

$$\sum_{\substack{k+\ell=n \\ k,\ell \gg 1}} \#\mathcal{T}_k \times \#\mathcal{T}_\ell \ll \#\mathcal{T}_n.$$

- Can be mimicked for many models, as long as exact enumeration is known.
- What if add "perturbations" that make the model too hard to count? For example, if we start with  $n$  triangles and  $o(n)$  quadrangles?

# The spatial Markov property of the UIPT $\mathbb{T}$

- By the Angel–Schramm argument  
 $\mathbb{P}(t \subset \mathbb{T}) = \lim_{n \rightarrow +\infty} \mathbb{P}(t \subset T_n)$  only depends on the perimeter and volume of  $t$ .
- Consequence: when we explore the UIPT "face by face" using *peeling explorations*, the perimeter and volume of the explored region follow a Markov chain with values in  $\mathbb{N}^2$ .
- Useful tool to study the fractal-like properties of  $\mathbb{T}$  (e.g. volume growth in  $r^4$  [Angel 04]).

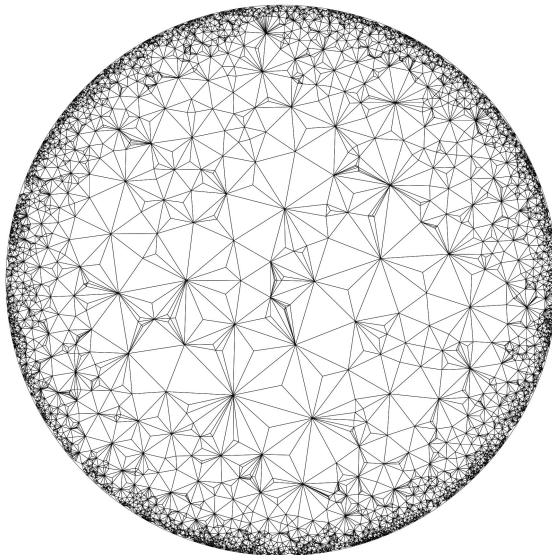
# Planar Stochastic Hyperbolic Triangulations

- The PSHT (Planar Stochastic Hyperbolic Triangulations) are the local limits of uniform triangulations of genus  $g$  with  $2n$  faces when  $g$  is proportional to  $n$  [B.–Louf 19].
- They form a one-parameter family  $(\mathbb{T}_\lambda)_{0 \leq \lambda \leq \lambda_c}$ , where  $\lambda_c = \frac{1}{12\sqrt{3}}$ . Characterized by a stronger version of the Markov property [Curien 14]:

$$\mathbb{P}(t \subset \mathbb{T}_\lambda) = C_{|\partial t|}(\lambda) \times \lambda^{|\partial t|}.$$

- What do they look like?
  - $\lambda = \lambda_c$ : the UIPT;
  - $0 < \lambda < \lambda_c$ : hyperbolic (mean degree  $> 6$ , exponential volume growth...);
  - $\lambda = 0$ : dual of a complete binary tree (degenerate object with infinite vertex degrees).

A PSHT  $\mathbb{T}_\lambda$  with  $0 < \lambda < \lambda_c$



# Convergence of high genus triangulations

- The proof does not rely on asymptotic enumeration!
- Main ingredients:
  - Tightness: "robust" adaptation of the tightness argument of Angel–Schramm;
  - Planarity and one-endedness of subsequential limits: use the Goulden–Jackson recursion (coming from algebraic combinatorics);
  - Any limit must have the spatial Markov property;

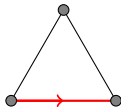
## Theorem (B.–Louf 19)

*Let  $T$  be a Markovian planar, one-ended, random infinite triangulation. Then  $T$  is of the form  $\mathbb{T}_\Lambda$ , where  $\Lambda$  is a random variable on  $[0, \lambda_c]$  (PSHT with a random parameter).*

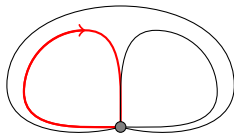
- Compare estimates on the PSHT and surgery arguments to prove that  $\Lambda$  is deterministic.
- Almost "enumeration-free" proof: if we want to replicate this sketch on planar models, the weak point is one-endedness.

# Infinite triangulations

- Infinite triangulation: family of countably many triangles glued along their edges and vertices. We do not assume one-endedness, nor finite vertex degrees.
- Examples of "degenerate" planar infinite triangulations:



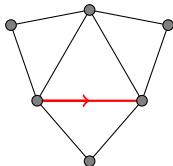
$T_0$  (PSHT of parameter 0)



$T_*$  (only one vertex)

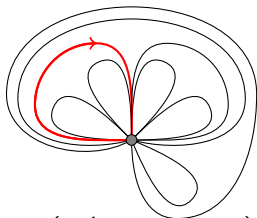
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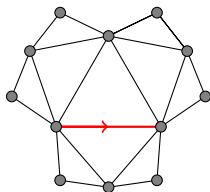
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$\mathbb{T}_\star$  (only one vertex)

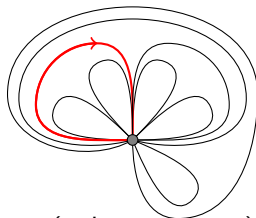
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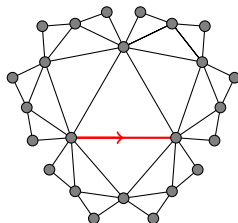


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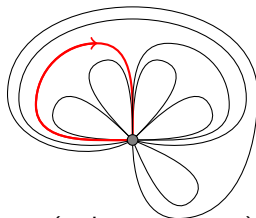
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# Spatial Markov property (general version)

## Definition

A random, infinite, planar triangulation  $T$  is *Markovian* if for any finite planar triangulation  $t$  with one or several holes, the probability  $\mathbb{P}(t \subset T)$  only depends on the perimeters of the holes of  $t$  and its total number of faces.

## Theorem (B.21+)

Let  $T$  be an infinite, planar, Markovian random triangulation. Then  $T$  is of the form  $\mathbb{T}_\Lambda$ , where  $\Lambda$  is a random variable with values in  $[0, \lambda_c] \cup \{\star\}$ .

- Consequences:
  - No nice notion of "uniform planar multi-ended triangulation".
  - The UIPT is the only Markovian planar triangulation where the expected inverse degree is  $1/6$ .

# Robust convergence to the UIPT

- Sketch of a "combinatorics-free" proof of convergence of triangulations of the sphere to the UIPT:
  - "Dual local topology": use dual distance instead of graph distances.

$$d_{\text{loc}}^*(t, t') = (1 + \max \{r \geq 0 \mid B_r^*(t) = B_r^*(t')\})^{-1}$$

- Makes tightness immediate, but limits may have infinite vertex degrees;
  - The finite model is Markovian, so any subsequential limit is Markovian;
  - The expected inverse of the root degree in a triangulation of the sphere is  $1/6$ , and this passes to the limit;
  - So the UIPT is the only possible subsequential limit;
  - In particular the limit has finite vertex degrees, so convergence for the dual local topology implies convergence for the usual local topology.
- Robust argument: still works if we add a perturbation "small compared to the size  $n$ ".

# Examples of applications

- For example, this sketch allows to prove the convergence to the UIPT of:
  - "Triangulations with defects", i.e. maps with prescribed face degrees where triangles represent a proportion  $1 - o(1)$  of the edges;
  - High temperature Ising triangulations, i.e. triangulations of size  $n$  equipped with an Ising model on the faces with inverse temperature  $\beta_n \rightarrow 0$ ;
  - Uniform triangulations of size  $n$  and genus  $g_n = o(n)$ .
- "Meta-theorem": If we perturb the uniform measure by factors  $e^{o(n)}$ , we still have convergence to the UIPT.

# Large deviations for pattern occurrences in uniform triangulations

- For  $t_0$  a triangulation with a hole and  $T$  a triangulation of the sphere, let  $\text{occ}_{t_0}(T)$  be the number of occurrences of  $t_0$  in  $T$ .
- Fix  $t_0$ , and let  $T_n^{(\beta)}$  be a triangulation of the sphere of size  $n$ , picked with probability proportional to  $e^{\beta \text{occ}_{t_0}(T)}$ .
- The previous sketch shows that if  $\beta_n \rightarrow 0$ , then  $T_n^{(\beta_n)}$  converges locally to the UIPT.

## Corollary

Let  $T_n$  be a uniform triangulation of the sphere with  $2n$  triangles. Then for every  $\varepsilon > 0$ , the probability that

$$\left| \frac{\text{occ}_{t_0}(T_n)}{6n} - \mathbb{P}(t_0 \subset \text{UIPT}) \right| > \varepsilon$$

decreases exponentially in  $n$ .

# Sketch of proof in the one-ended case

- Let  $T$  be a one-ended, planar, infinite Markovian triangulation. For  $v \geq p \geq 1$ , let

$$a_v^p = \mathbb{P}(t \subset T)$$

for  $t$  a triangulation with perimeter  $p$  and  $v$  vertices in total.

- Peeling equations:

$$a_v^p = a_{v+1}^{p+1} + 2 \sum_{i=0}^{p-1} \sum_{j \geq 0} a_{v+j}^{p-i} \# \mathcal{T}_{i+1,j}.$$

- In particular, the law of  $T$  is determined by the numbers  $a_v^1$ .
- For a mixture of PSHT  $\mathbb{T}_\Lambda$ , we have  $a_v^1 = \mathbb{E}[\Lambda^{v-1}]$ .
- So we need to show that  $(a_v^1)_{v \geq 1}$  is the sequence of moments of some variable  $\Lambda$ .

# The Hausdorff moment problem

- Let  $\Delta$  be the discrete derivative operator:

$$(\Delta u)_n = u_n - u_{n+1}.$$

## Theorem (Hausdorff)

Let  $(u_n)$  be a sequence of real numbers. Then  $(u_n)$  is the sequence of moments of some  $[0, 1]$ -valued random variable if and only if

$$\forall k \geq 0, \forall n \geq 0, \Delta^k u_n \geq 0.$$

- By induction on  $k$ , we prove  $(\Delta^k a^p)_v \geq 0$  for all  $k \geq 0$  and  $v \geq p \geq 1$ .
- So there is a variable  $\Lambda \in [0, 1]$  such that  $a_v^1 = \mathbb{E}[\Lambda^{v-1}]$ .
- The convergence of the sum in the peeling equation implies  $\Lambda \in \left[0, \frac{1}{12\sqrt{3}}\right]$ .

# The multi-ended case

- Let  $T$  be a multi-ended Markovian triangulation. For a triangulation  $t$  with volume  $v$  and  $k$  holes of perimeters  $p_1, \dots, p_k$ , we write

$$a_v^{p_1, p_2, \dots, p_k} = \mathbb{P}(t \subset T \text{ and } T \setminus t \text{ has } k \text{ infinite components}).$$

- New peeling equation:

$$\begin{aligned} a_v^{p_1, \dots, p_k} &= a_v^{p_1+1, p_2, \dots, p_k} + 2 \sum_{i=0}^{p_1-1} \sum_{j \geq 0} a_{v+i+j}^{p_1-i, p_2, \dots, p_k} \times \#\mathcal{T}_{i+1, j} \\ &\quad + \sum_{i=0}^{p_1-1} a_v^{i+1, p_1-i, p_2, \dots, p_k}. \end{aligned}$$

- The first term in the RHS has "larger" perimeters, so the law of  $T$  is determined by terms of the form  $a_v^{1, \dots, 1} = a_v^{k \otimes 1}$ .



# The multi-ended case

- The law of  $T$  is determined by terms of the form  $a_v^{1,\dots,1} = a_v^{k \otimes 1}$ .
- Hausdorff moment problem : there is  $(\Lambda, \Gamma)$  such that, for all  $k, v \geq 1$ :

$$a_v^{k \otimes 1} = \mathbb{E} \left[ \Lambda^{v-1} \Gamma^{k-1} \right].$$

- We want either  $\Gamma = 0$  (PSHT  $\mathbb{T}_\lambda$ ) or  $(\Lambda, \Gamma) = (0, 1)$  (degenerate  $\mathbb{T}_\star$ ).
- From here, we can solve completely the peeling equations:

$$a_v^{p_1, \dots, p_k} = \mathbb{E} \left[ \Lambda^{v-1} \Gamma^{k-1} \prod_{i=1}^k C_{p_i}(\Lambda, \Gamma) \right],$$

where  $C_p(\lambda, \gamma)$  is given by the induction

$$C_p = C_{p+1} + 2 \sum_{i=0}^{p-1} \lambda^i Z_{i+1}(\lambda) C_{p-i} + \gamma \sum_{i=0}^{p-1} C_{i+1} C_{p-i}.$$

# The multi-ended case

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- In particular, the generating function  $\sum_{p \geq 1} C_p(\lambda, \gamma) x^p$  is explicit.
- Behaviour of the generating function near its first singularity  
 $\rightarrow$  if  $\gamma > 0$  and  $(\lambda, \gamma) \neq (0, 1)$ , then  $C_p(\lambda, \gamma) < 0$  for some  $p$ .
- Recall  $a_v^{p, (k-1) \otimes 1} = \mathbb{E} [\Lambda^{v-1} \Gamma^{k-1} C_p(\Lambda, \Gamma)]$ .
- When  $k, v$  get large, this is close to  $\lambda_{\max}^{v-1} \gamma_{\max}^{k-1} C_p(\lambda_{\max}, \gamma_{\max})$ , which is negative for some  $p$  if  $\gamma_{\max} > 0$ .
- So almost surely  $\Lambda = 0$  or  $\Gamma = 0$ .
- Similar argument to exclude  $\Lambda = 0$  and  $0 < \Gamma < 1$ .

- Different kinds of face degrees ?
- What if we remove the planarity assumption? Conjectures:
  - There are no Markovian nonplanar triangulations with finite vertex degrees...
  - ...but there should be nonplanar, degenerate objects.

*THANK YOU !*