

Numerical optimization

$f: \vec{x} \in \mathbb{R}^N \rightarrow f(\vec{x}) \in \mathbb{R}$
 continuous variables

$f: \vec{\sigma} \rightarrow f(\vec{\sigma}) \in \mathbb{R}$

$\vec{\sigma} = (\sigma_1, \sigma_2, \dots, \sigma_N)$ (Ising spins)
 $\sigma_i = \pm 1$ Ising variable
 (binary variable)

Find the minimum of f under constraint $\sigma^2 = 1$

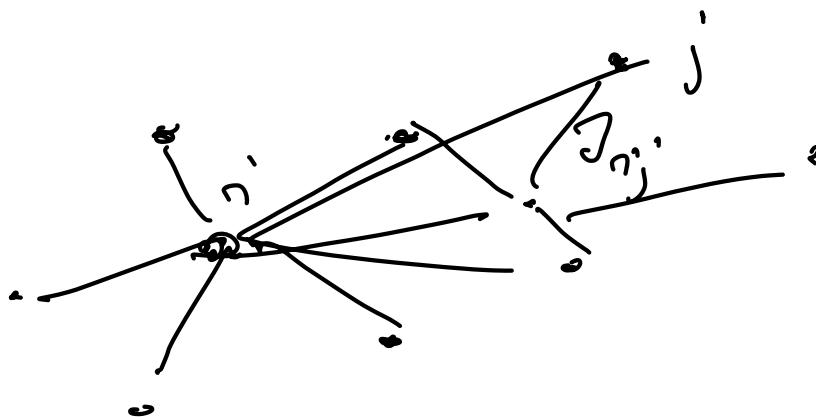
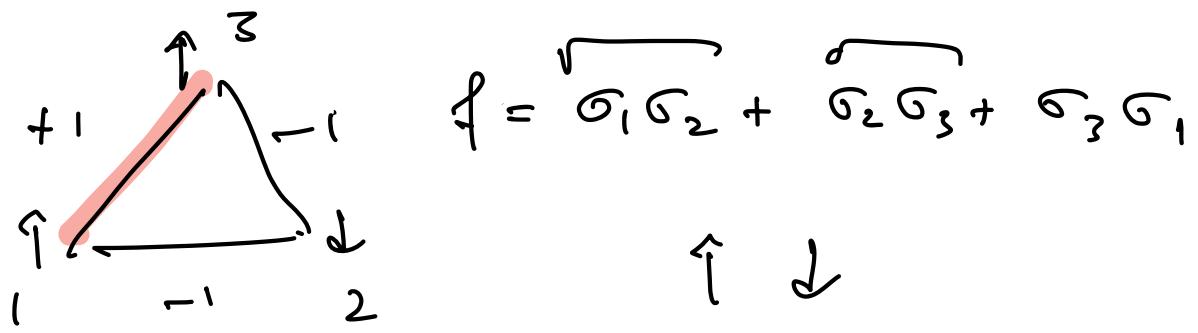
magnetic field

$$f(\vec{\sigma}) = -\sum_{i=1}^N h_i \sigma_i + \sum_{i,j} J_{ij} \sigma_i \sigma_j + \dots$$

two-spin int.

of configs: 2^N

"Frustrated" (competing) interactions

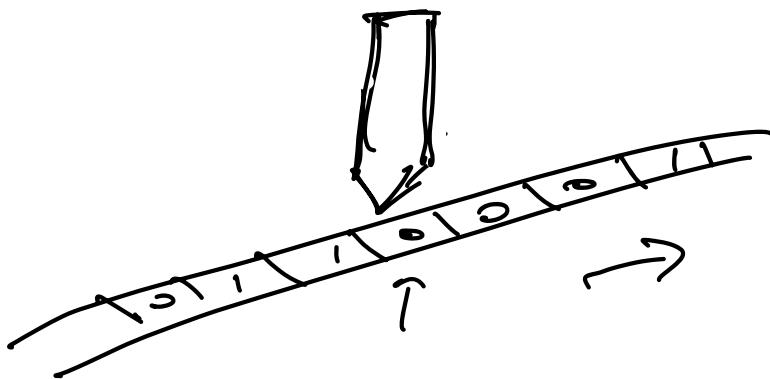


Curviness + Frustration \Rightarrow hardness of minimization



Primer in computational complexity

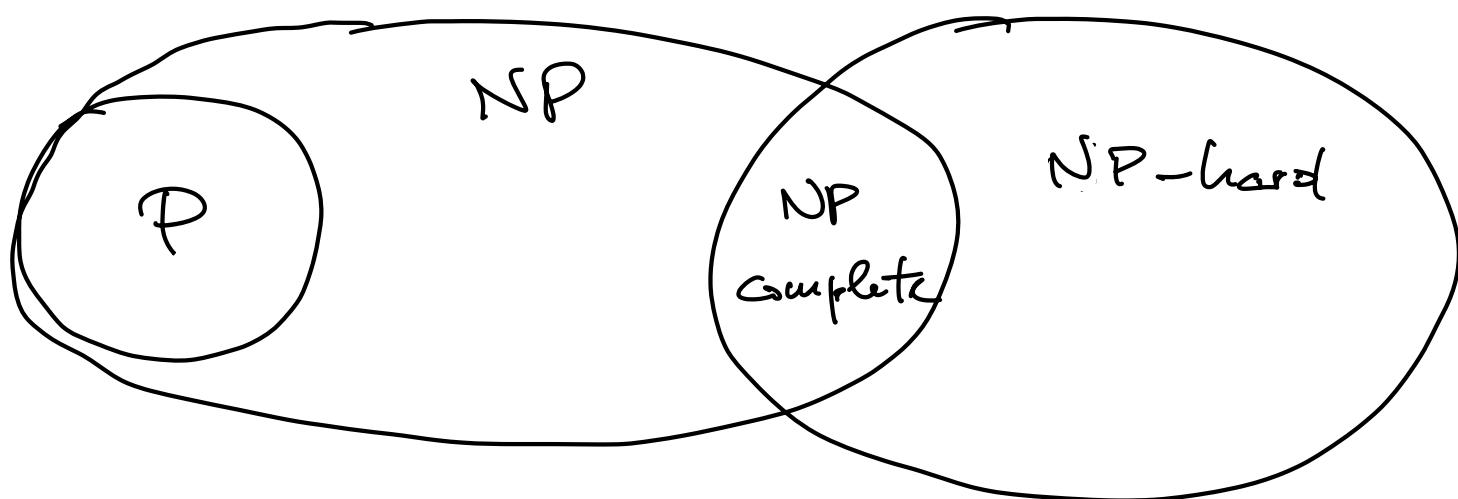
i) Deterministic Turing machine (DTM)



2) Non-deterministic Turing machine (NTM)

at every step \rightarrow write 1 with prob p
 \rightarrow write 0 with prob. $1-p$

Complexity classes



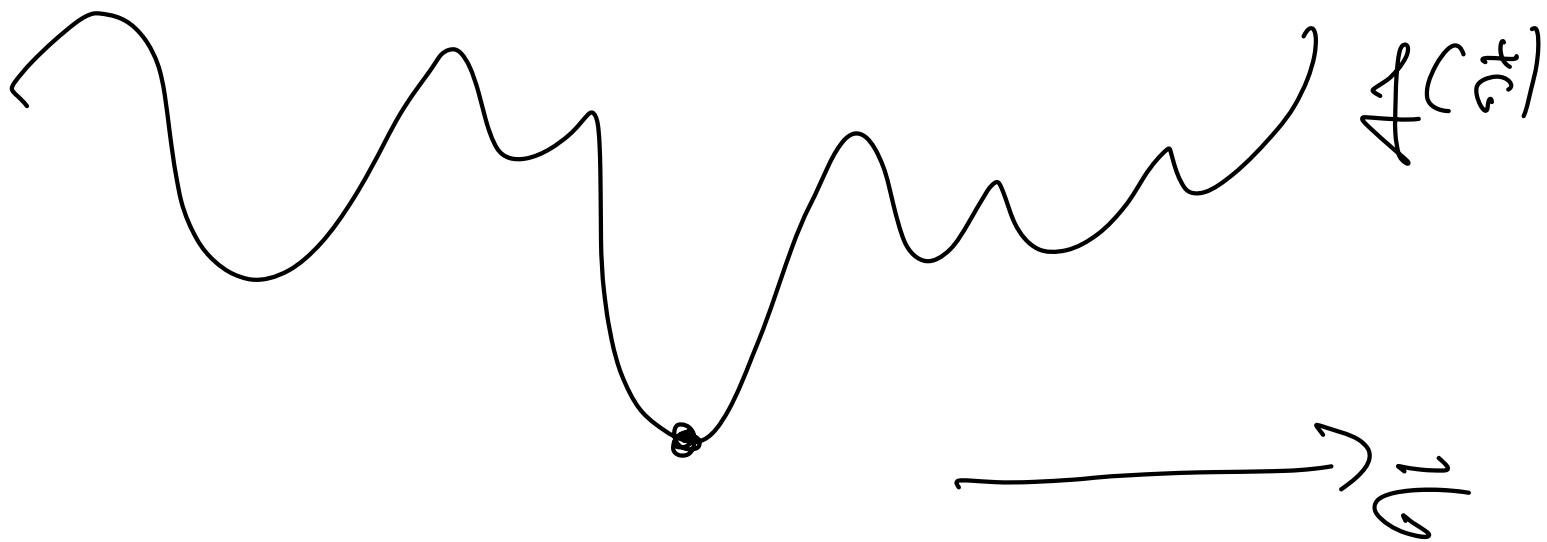
P : ensemble of problems solvable
in poly time by a DTM

NP : ① ensemble of problems solvable
in poly time by a NTM -
② the solution can be checked in
poly time by a DTM

NP-hard : ① ~~②~~

all problems in NP can be mapped in polynomial time as problems in NP-Hard

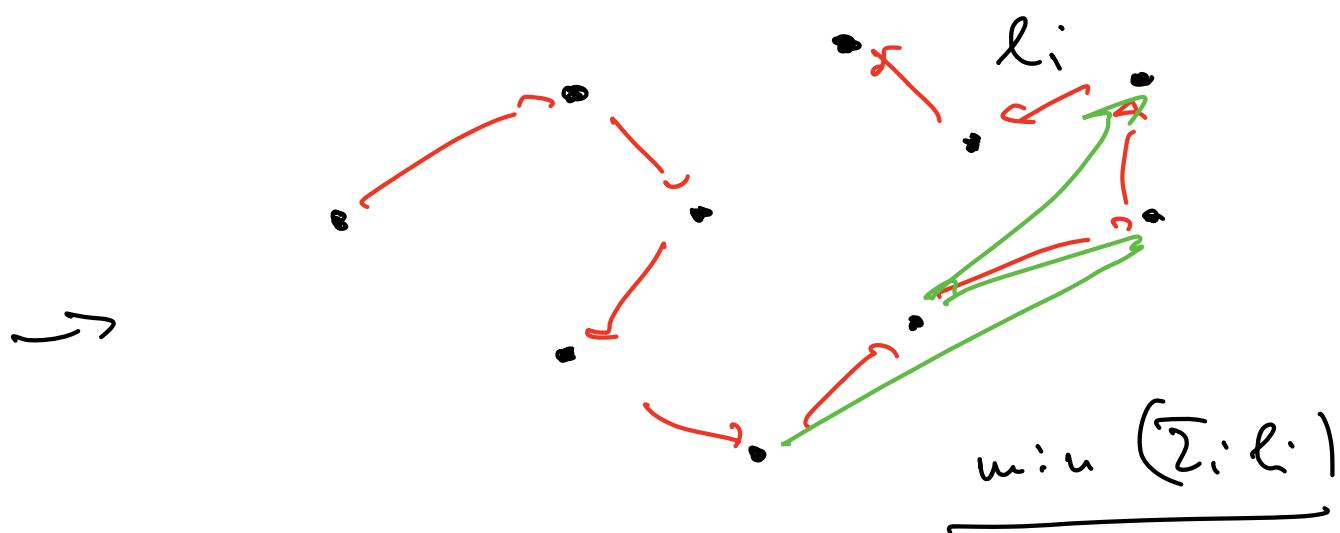
There are optimization problems that can be proven to be NP-hard



NP-complete : Boolean satisfiability

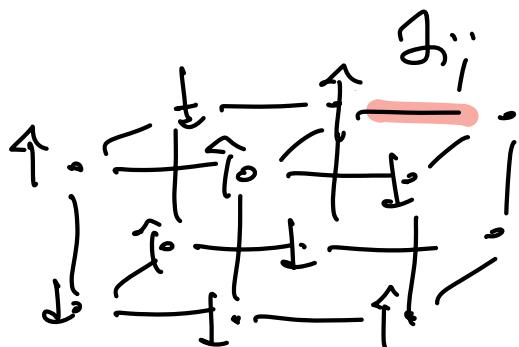
(a. or .. b). and. (d. or .. c) and . -

NP-hard : traveling salesman problem



Physics : minimization of the energy
of Ising models (spin glasses)

ex. :



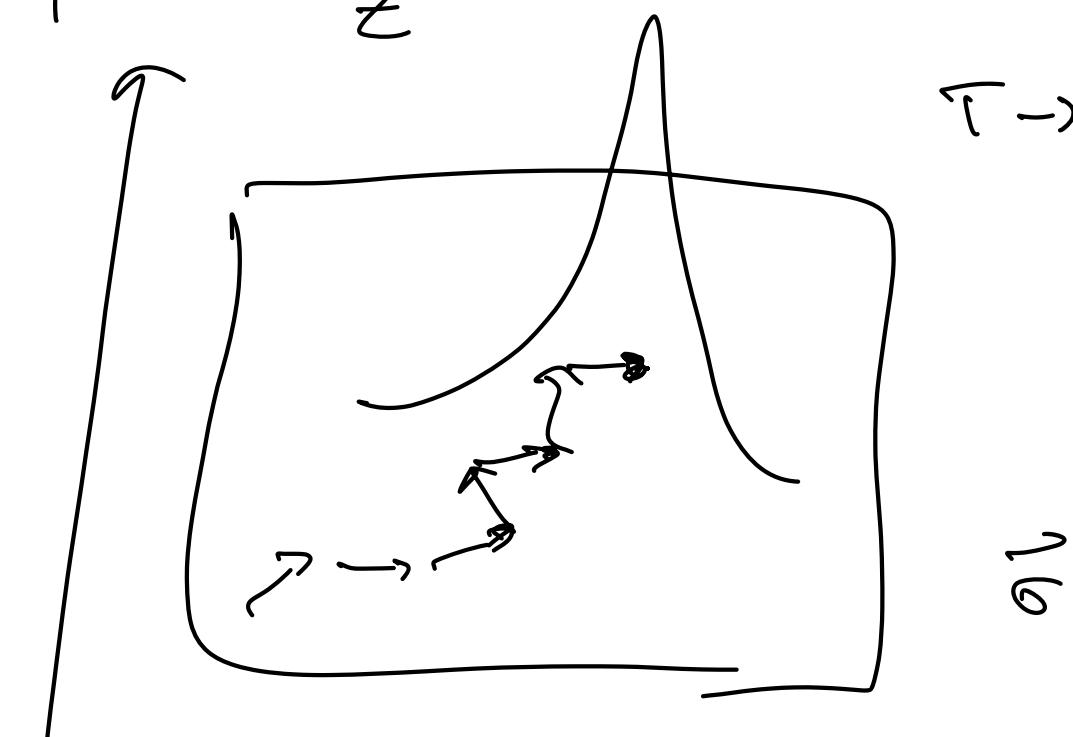
$$\Rightarrow J_{ij} = (0, -1, +1) \quad \text{scattered}$$

$$\rightarrow \sum_{\langle ij \rangle}^{(G)} = \sum_{\langle ij \rangle} J_{ij} G_i G_j$$

ED NP-hard

$$P(\vec{G}) = \frac{e^{-\frac{E(\vec{G})}{T}}}{Z}$$

$$Z = \sum_{\vec{G}} e^{\frac{-E(\vec{G})}{T}}$$

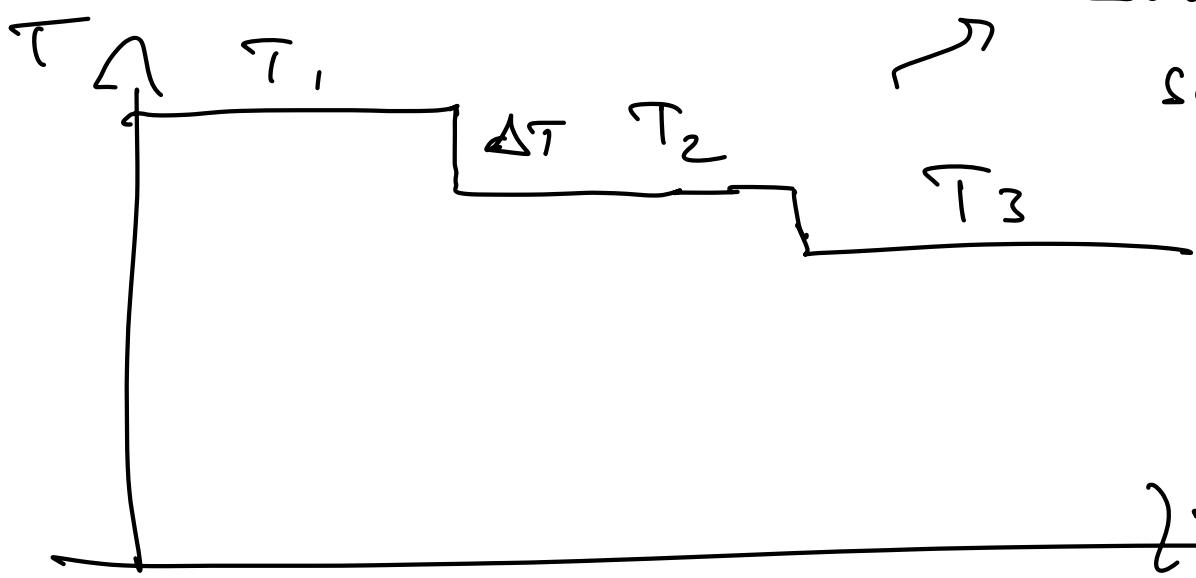


$$T \rightarrow 0$$

θ_1

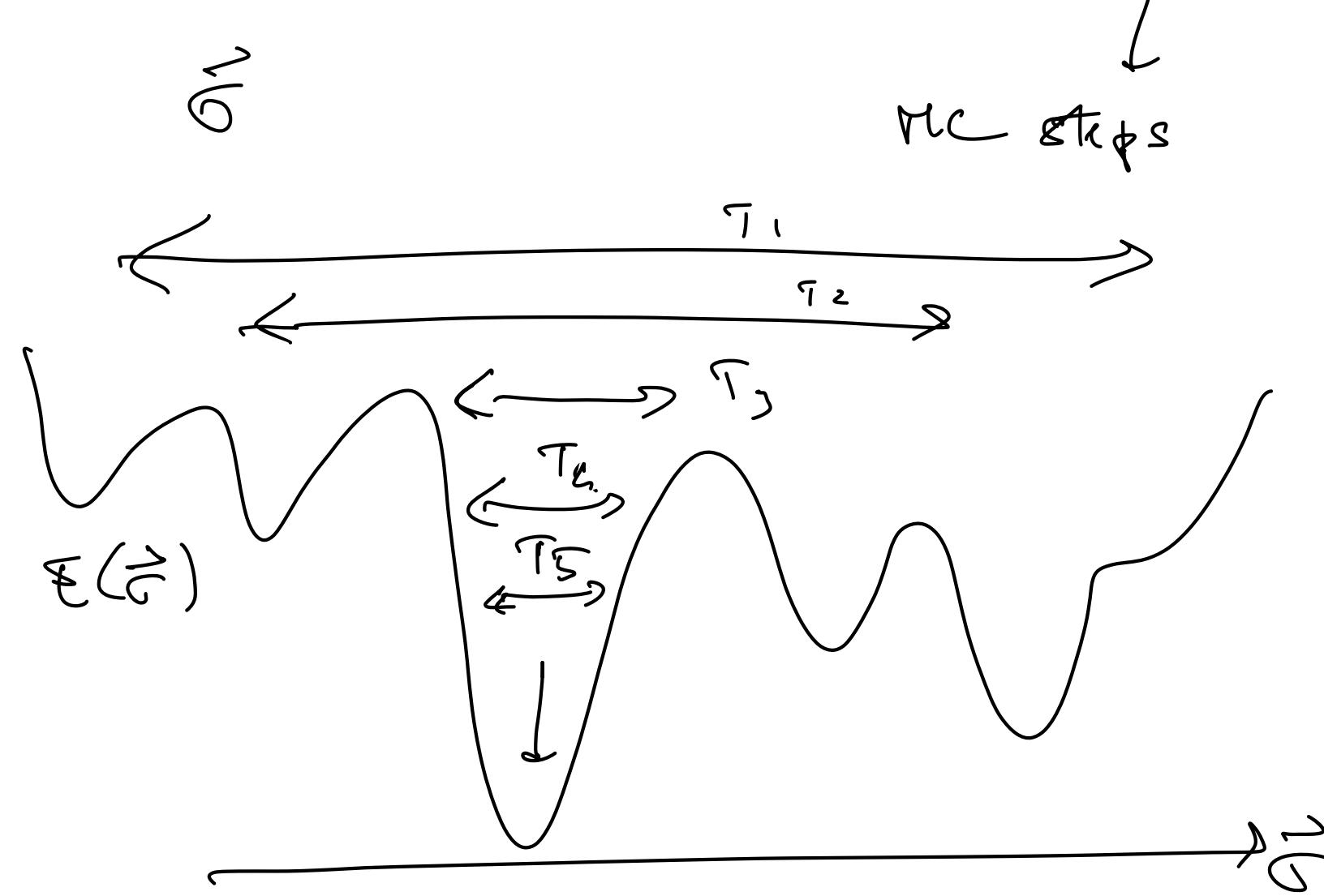
Simulated annealing

MC simulation



“cooling schedule”

\rightarrow \downarrow \downarrow \downarrow



Where $\epsilon(\vec{\sigma})$ is a model of an
Ising spin glass

$$E \rightarrow L = \sum_i l_i$$

$$- (\sum_i l_i) / T$$

$$\phi(\{l_i\}) = \frac{l}{2}$$

Probleem : $T \sim O(\exp(N))$

↓
auto correlation trees

Proof of efficiency for simulated annealing:

you are guaranteed to get to the
absolute minimum by using an annealing
schedule

$$T_n \approx \frac{c^N}{\log k}$$

$$k = \exp\left(\frac{c^N}{T}\right)$$



Numerical selection of ODEs

$$\Rightarrow \vec{x}(t) = (x_1(t), x_2(t), \dots, x_N(t))$$

↑ ↑ ↑

|

$$\vec{f}(\vec{x}, \dot{\vec{x}}, \ddot{\vec{x}}, \dots, \vec{x}^{(m)}; t) = 0$$

↗ m-th order ODE

↓ ordinary differential equation

1-st order ODE

$$\vec{q}_1(t) = \vec{x}(t)$$

$$\vec{q}_2 = \dot{\vec{x}}$$

$$\vec{q}_3 = \dot{\vec{q}}_2 = \ddot{\vec{x}}$$

⋮

$$\vec{q}_m = \vec{x}^{(m-1)} = \dot{\vec{q}}_{m-1}$$

$$\vec{G}(\vec{q}_1, \vec{q}_2, \vec{q}_3, \dots, \vec{q}_m; \dot{\vec{q}}_1, \dot{\vec{q}}_2, \dots, \dot{\vec{q}}_m; t) = 0$$

↙ \vec{q} ↗ $\dot{\vec{q}}$

$$\vec{G}(\vec{q}, \dot{\vec{q}}; t) = 0$$

Gaußig problem

$$\left\{ \begin{array}{l} \dot{\vec{q}} = \vec{f}(\vec{q}(t), t) \\ \vec{q}(0) = \vec{q}_0 \end{array} \right. \quad \rightarrow \text{initial conditions}$$

Ex. Newton's equations

$$\underline{m \ddot{\vec{x}}(t)} = \vec{f}(\vec{x}(t), t)$$

$$\vec{q}_1(t) = \vec{x}(t)$$

$$\vec{q}_2(t) = \dot{\vec{x}}(t)$$

$$\left\{ \begin{array}{l} \dot{\vec{q}}_1 = \vec{q}_2 \\ \dot{\vec{q}}_2 = \vec{f}(\vec{q}_1, t) \end{array} \right.$$

—

Numerical solution to a Gaußig problem
comes with a finite accuracy

$$|\vec{y}_{\text{num}}(t) - \vec{y}_{\text{exact}}(t)| \leq \epsilon$$

$$t \in [t_0, t_0]$$

↑

time to solution with accuracy ϵ

$$T = \frac{A}{\epsilon^\alpha} N\left(\frac{t_0}{\epsilon}\right)$$

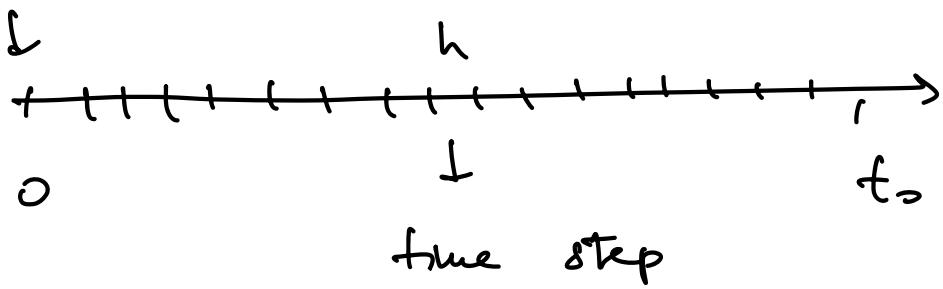
↓
general scaling

$$\omega \sim \omega(1)$$

=

Numerical scheme of integration

in general it goes via discretization
of time



$$\vec{q}(t) \rightarrow \vec{q}_n \approx \vec{q}(t_n)$$

$$t_n = nh$$

Numerical scheme : explicit

$$\vec{q}_{n+1} = \vec{g}_n(\vec{q}_n, \vec{q}_{n-1}, \vec{q}_{n-2}, \dots, \vec{q}_0; h)$$

$$\vec{q}_{n+1} = \vec{g}_n(\vec{q}_{n+1}, \vec{q}_n, \vec{q}_{n-1}, \dots, \vec{q}_0; h)$$

implicit

Special case : true-independent linear
set of ODEs

$$\dot{\vec{q}}(t) = A \vec{q}$$

$$\vec{q}(0) = \vec{q}_0$$

$$\Rightarrow \vec{q}(t) = e^{At} \vec{q}_0$$

$$A \text{ symmetric} \rightarrow U^T A U = \text{diag}(\lambda_1, \dots, \lambda_n)$$

$$e^{At} = U^T \text{diag}(e^{\lambda_1 t}, e^{\lambda_2 t}, \dots, e^{\lambda_n t}) U$$

↗ ↘

flow do you build an integration scheme ?

$$\dot{\vec{q}}(t) = \vec{f}(\vec{q}; t)$$

$$\int_t^{t+h} \dot{\vec{q}}(t) dt = \vec{q}(t+h) - \vec{q}(t)$$

$$= \int_t^{t+h} dt' \vec{f}(\vec{q}(t'); t')$$

$\underbrace{\quad}_{\text{if } h \text{ is small}}$

$t \quad t+h$ if h is small

expand $\vec{q}(t')$ around $\vec{q}(t)$

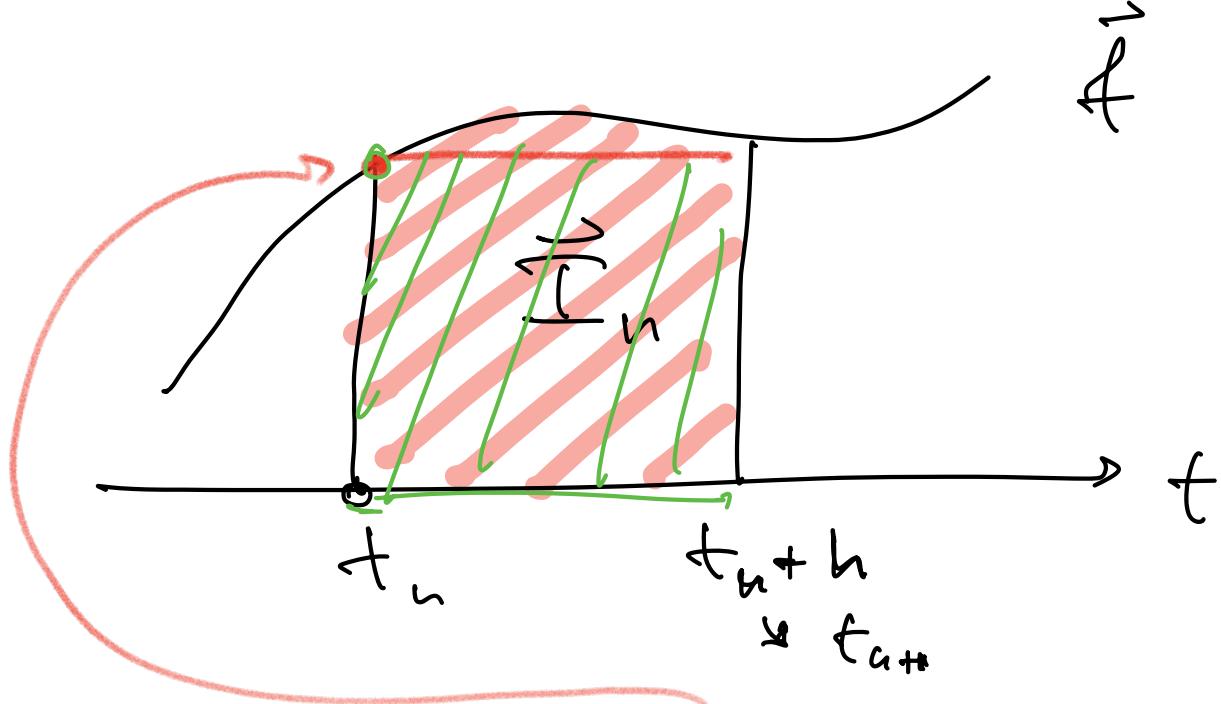
$$\approx \int_t^{t+h} dt' \left[\vec{f}(\vec{q}(t); t) + (t' - t) \frac{d\vec{f}}{dt} + \dots \right]$$

$$= \frac{h \vec{f}(\vec{q}(t), t)}{t \rightarrow t_0 \approx uh} + \frac{1}{2} h^2 \frac{d\vec{f}}{dt}(\vec{q}(t); t) + \dots$$

$$+ h^3$$

$$\vec{q}_{n+1} - \vec{q}_n = \int_{t_n}^{t_{n+1}} dt' \vec{f}(\vec{q}(t'); t')$$

$$= \sum_u$$



$$\vec{q}_{n+1} - \vec{q}_n \approx h \vec{f}(\vec{q}_n, t_n) + \frac{h^2}{2} \frac{d\vec{f}}{dt}(\vec{q}_n, t_n) + \dots$$

numerical integration scheme = good approximation
of \vec{I}_n^1

Euler method (18th century)

$$\vec{q}_{n+1} \approx \vec{q}_n + h \vec{f}(\vec{q}_n, t_n) + \mathcal{O}(h^2)$$

↑

Order of consistency of a numerical scheme

numerical scheme

$$\vec{q}_{n+1} = \vec{g}_n((\vec{q}_{n+1}), \vec{q}_n, \dots, \vec{q}_0; h)$$

"local" error

$$\vec{\epsilon}_{n+1} = \vec{y}(t_{n+1}) - \vec{g}_n(\vec{y}(t_{n+1}), \vec{y}(t_n), \dots, \vec{y}_0; h)$$

exact sol.

method is consistent of order ρ over an interval $[t_0, t_1]$

$$\text{if } \max_{n \in [1, t_1/h]} |\vec{\epsilon}_{n+1}| = h^\rho \quad \underline{\hspace{10cm}}$$

$$\lim_{h \rightarrow 0} \epsilon(h) \sim h^\rho \quad \underline{\hspace{1cm}}$$

$$\text{error} \sim \mathcal{O}(h^{p+1})$$

$$\text{Euler : error} \sim \mathcal{O}(h^2)$$

$$\Rightarrow p = 1$$



Convergence of a method

global error

$$\tilde{\xi}_n = \tilde{q}_n - \tilde{q}(t_n)$$

method is convergent of order p if

$$\lim_{h \rightarrow 0} \left(\max_{n \in [0, t/h]} |\tilde{\xi}_n| \right) \sim \underline{\mathcal{O}(h^p)}$$

Simple idea:

$$\tilde{q}_n = \tilde{q}(t_n) + \mathcal{O}(h^{p+1})$$

$$\tilde{y}_2 = \tilde{y}(t_2) + 2 O(h^{p+1})$$

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$$\tilde{y}_n = \tilde{y}(t_n) + n O(h^{p+1})$$

$$n = \frac{t_0/t_n}{\tau}$$
$$= \tilde{y}(t_n) + t_0 O(h^p)$$

Euler can be proven to be consistent
and convergent of order 1