

Mathematical foundations in deep learning

Part VII: Duality

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Fenchel-Rockafellar duality

Primal problem

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

Let $f: \mathcal{H} \rightarrow]-\infty, +\infty]$, $g: \mathcal{G} \rightarrow]-\infty, +\infty]$. Let $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$.

We want to

$$\underset{x \in \mathcal{H}}{\text{minimize}} \quad f(x) + g(Lx).$$

Dual problem

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

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We want to

$$\underset{v \in \mathcal{G}}{\text{minimize}} \quad f^*(-L^*v) + g^*(v).$$

Fenchel-Rockafellar duality

Weak duality

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

Let f be a proper function from \mathcal{H} to $]-\infty, +\infty]$, g be a proper function from \mathcal{G} to $]-\infty, +\infty]$, and $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$. Let

$$\mu = \inf_{x \in \mathcal{H}} f(x) + g(Lx) \quad \text{and} \quad \mu^* = \inf_{v \in \mathcal{G}} f^*(-L^*v) + g^*(v).$$

We have $\mu \geq -\mu^*$. If $\mu \in \mathbb{R}$, $\mu + \mu^*$ is called the **duality gap**.

Fenchel-Rockafellar duality

Weak duality

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Proof: According to Fenchel-Young inequality, for every $x \in \mathcal{H}$ and $v \in \mathcal{G}$,

$$f(x) + g(Lx) + f^*(-L^*v) + g^*(v) \geq \langle x \mid -L^*v \rangle + \langle Lx \mid v \rangle = 0.$$

Fenchel-Rockafellar duality

Strong duality

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

Let $f \in \Gamma_0(\mathcal{H})$, $g \in \Gamma_0(\mathcal{G})$, and $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$.

If $\text{int}(\text{dom } g) \cap L(\text{dom } f) \neq \emptyset$ or $\text{dom } g \cap \text{int}(L(\text{dom } f)) \neq \emptyset$, then

$$\mu = \inf_{x \in \mathcal{H}} f(x) + g(Lx) = -\min_{v \in \mathcal{G}} f^*(-L^*v) + g^*(v) = -\mu^*.$$

Fenchel-Rockafellar duality

Duality theorem (1)

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

Let $f \in \Gamma_0(\mathcal{H})$, $g \in \Gamma_0(\mathcal{G})$, and $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$.

$$\text{zer}(\partial f + L\partial g L^*) \neq \emptyset \quad \Leftrightarrow \quad \text{zer}((-L)\partial f^*(-L^*) + \partial g^*) \neq \emptyset.$$

Fenchel-Rockafellar duality

Duality theorem (1)

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

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$$\text{zer}(\partial f + L\partial g L^*) \neq \emptyset \quad \Leftrightarrow \quad \text{zer}((-L)\partial f^*(-L^*) + \partial g^*) \neq \emptyset.$$

Proof:

$$\begin{aligned} (\exists x \in \mathcal{H}) \quad 0 \in \partial f(x) + L^* \partial g(Lx) &\Leftrightarrow (\exists x \in \mathcal{H})(\exists v \in \mathcal{G}) \quad \begin{cases} -L^*v \in \partial f(x) \\ v \in \partial g(Lx) \end{cases} \\ &\Leftrightarrow (\exists x \in \mathcal{H})(\exists v \in \mathcal{G}) \quad \begin{cases} x \in \partial f^*(-L^*v) \\ Lx \in \partial g^*(v) \end{cases} \\ &\Leftrightarrow (\exists v \in \mathcal{G}) \quad 0 \in -L\partial f^*(-L^*v) + \partial g^*(v). \end{aligned}$$

Fenchel-Rockafellar duality

Duality theorem (2)

Let \mathcal{H} and \mathcal{G} be two real Hilbert spaces.

Let $f \in \Gamma_0(\mathcal{H})$, $g \in \Gamma_0(\mathcal{G})$, and $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$.

- If there exists $\hat{x} \in \mathcal{H}$ such that $0 \in \partial f(\hat{x}) + L^* \partial g(L\hat{x})$, then \hat{x} is a solution to the primal problem. Moreover, there exists a solution \hat{v} to the dual problem such that $-L^* \hat{v} \in \partial f(\hat{x})$ and $L\hat{x} \in \partial g^*(\hat{v})$.
- If there exists $(\hat{x}, \hat{v}) \in \mathcal{H} \times \mathcal{G}$ such that $-L^* \hat{v} \in \partial f(\hat{x})$ and $L\hat{x} \in \partial g^*(\hat{v})$ then \hat{x} (resp. \hat{v}) is a solution to the primal (resp. dual) problem.

If $(\hat{x}, \hat{v}) \in \mathcal{H} \times \mathcal{G}$ is such that $-L^* \hat{v} \in \partial f(\hat{x})$ and $L\hat{x} \in \partial g^*(\hat{v})$, then (\hat{x}, \hat{v}) is called a **Kuhn-Tucker point**.

Augmented Lagrangian method

ADMM algorithm (*Alternating-direction method of multipliers*)

⇒ **Lagrangian interpretation**

$$\underset{x \in \mathcal{H}}{\text{minimize}} \quad f(x) + g(Lx) \quad \Leftrightarrow \quad \underset{\substack{x \in \mathcal{H}, y \in \mathcal{G} \\ Lx=y}}{\text{minimize}} \quad f(x) + g(y)$$

Lagrange function:

$$(\forall (x, y, v) \in \mathcal{H} \times \mathcal{G}^2) \quad \mathcal{L}(x, y, z) = f(x) + g(y) + \langle z \mid Lx - y \rangle$$

where $z \in \mathcal{G}$ denotes the Lagrange multiplier.

Alternating-direction method of multipliers

Idea: iterations for finding a saddle point $(\hat{x}, \hat{y}, \hat{z})$:

$$(\forall n \in \mathbb{N}) \quad \begin{cases} x_n \in \operatorname{Argmin} \mathcal{L}(\cdot, y_n, z_n) \\ y_{n+1} \in \operatorname{Argmin} \mathcal{L}(x_n, \cdot, z_n) \\ v_{n+1} \text{ such that } \mathcal{L}(x_n, y_{n+1}, v_{n+1}) \geq \mathcal{L}(x_n, y_{n+1}, v_n). \end{cases}$$

But the convergence is not guaranteed in general !

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But the convergence is not guaranteed in general !

Solution: introduce an **Augmented Lagrange function**.

Let $\gamma \in]0, +\infty[$, we define

$$\begin{aligned} (\forall (x, y, z) \in \mathcal{H} \times \mathcal{G}^2) \quad \tilde{\mathcal{L}}(x, y, z) = & f(x) + g(y) + \gamma \langle z \mid Lx - y \rangle \\ & + \frac{\gamma}{2} \|Lx - y\|^2 \end{aligned}$$

The Lagrange multiplier is $v = \gamma z$.

Alternating-direction method of multipliers

Algorithm for finding a saddle point:

$$(\forall n \in \mathbb{N}) \quad \begin{cases} x_n = \operatorname{argmin}_{x \in \mathcal{H}} \tilde{\mathcal{L}}(x, y_n, z_n) \\ y_{n+1} = \operatorname{argmin}_{y \in \mathcal{G}} \tilde{\mathcal{L}}(x_n, y, z_n) \\ z_{n+1} \text{ such that } \tilde{\mathcal{L}}(x_n, y_{n+1}, z_{n+1}) \geq \tilde{\mathcal{L}}(x_n, y_{n+1}, z_n). \end{cases}$$

By performing a gradient ascent on the Lagrange multiplier,

$$\begin{aligned} (\forall n \in \mathbb{N}) \quad & \begin{cases} x_n = \operatorname{argmin}_{x \in \mathcal{H}} f(x) + \gamma \langle z_n | Lx - y_n \rangle + \frac{\gamma}{2} \|Lx - y_n\|^2 \\ y_{n+1} = \operatorname{argmin}_{y \in \mathcal{G}} g(y) + \gamma \langle z_n | Lx_n - y \rangle + \frac{\gamma}{2} \|Lx_n - y\|^2 \\ z_{n+1} = z_n + \frac{1}{\gamma} \nabla_z \tilde{\mathcal{L}}(x_n, y_{n+1}, z_n) \end{cases} \\ \Leftrightarrow (\forall n \in \mathbb{N}) \quad & \begin{cases} x_n = \operatorname{argmin}_{x \in \mathcal{H}} \frac{1}{2} \|Lx - y_n + z_n\|^2 + \frac{1}{\gamma} f(x) \\ y_{n+1} = \operatorname{prox}_{\frac{g}{\gamma}}(z_n + Lx_n) \\ z_{n+1} = z_n + Lx_n - y_{n+1}. \end{cases} \end{aligned}$$

Augmented Lagrange method

ADMM algorithm (*Alternating-direction method of multipliers*)

Let \mathcal{H} and \mathcal{G} be two Hilbert spaces. Let $f \in \Gamma_0(\mathcal{H})$ et $g \in \Gamma_0(\mathcal{G})$.

Let $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$ such that L^*L is an isomorphism and let $\gamma \in]0, +\infty[$.

$$(\forall n \in \mathbb{N}) \quad \begin{cases} x_n = \operatorname{argmin}_{x \in \mathcal{H}} \frac{1}{2} \|Lx - y_n + z_n\|^2 + \frac{1}{\gamma} f(x) \\ s_n = Lx_n \\ y_{n+1} = \operatorname{prox}_{\frac{g}{\gamma}}(z_n + s_n) \\ z_{n+1} = z_n + s_n - y_{n+1}. \end{cases}$$

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Let $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$ such that L^*L is an isomorphism and let $\gamma \in]0, +\infty[$.

We assume that $\text{int}(\text{dom } g) \cap L(\text{dom } f) \neq \emptyset$ or $\text{dom } g \cap \text{int}(L(\text{dom } f)) \neq \emptyset$ and that $\text{Argmin}(f + g \circ L) \neq \emptyset$. Let

$$(\forall n \in \mathbb{N}) \quad \begin{cases} x_n = \underset{x \in \mathcal{H}}{\text{argmin}} \quad \frac{1}{2} \|Lx - y_n + z_n\|^2 + \frac{1}{\gamma} f(x) \\ s_n = Lx_n \\ y_{n+1} = \text{prox}_{\frac{g}{\gamma}}(z_n + s_n) \\ z_{n+1} = z_n + s_n - y_{n+1}. \end{cases}$$

We have:

- $x_n \rightharpoonup \hat{x}$ where $\hat{x} \in \text{Argmin}(f + g \circ L)$
- $\gamma z_n \rightharpoonup \hat{v}$ where $\hat{v} \in \text{Argmin}(f^* \circ (-L^*) + g^*)$.

Augmented Lagrangian method

ADMM algorithm (*Alternating-direction method of multipliers*)

≡ **Douglas-Rachford for the dual problem**

Let \mathcal{H} and \mathcal{G} be two Hilbert spaces. Let $f \in \Gamma_0(\mathcal{H})$ and $g \in \Gamma_0(\mathcal{G})$.
Let $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$ such that L^*L is an isomorphism and let $\gamma \in]0, +\infty[$.
The Douglas-Rachford iterations to minimize $f^* \circ (-L^*) + g^*$ are

$$(\forall n \in \mathbb{N}) \quad \begin{cases} v_n = \text{prox}_{\gamma g^*} u_n \\ w_n = \text{prox}_{\gamma f^* \circ (-L^*)} (2v_n - u_n) \\ u_{n+1} = u_n + w_n - v_n \end{cases}$$

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The Douglas-Rachford iterations to minimize $f^* \circ (-L^*) + g^*$ are

$$(\forall n \in \mathbb{N}) \quad \begin{cases} v_n = u_n - \gamma \operatorname{prox}_{\gamma^{-1}g}(\gamma^{-1}u_n) \\ 2v_n - u_n - w_n \in \gamma \partial(f^* \circ (-L^*)) w_n \\ u_{n+1} = u_n + w_n - v_n \end{cases}$$

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$$(\forall n \in \mathbb{N}) \quad \begin{cases} v_n = u_n - \gamma \operatorname{prox}_{\gamma^{-1}g}(\gamma^{-1}u_n) \\ 2v_n - u_n - w_n \in -\gamma L \circ \partial f^* \circ (-L^*)w_n \\ u_{n+1} = u_n + w_n - v_n \end{cases}$$

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$$(\forall n \in \mathbb{N}) \quad \begin{cases} v_n = u_n - \gamma \operatorname{prox}_{\gamma^{-1}g}(\gamma^{-1}u_n) \\ L^*(u_n - 2v_n - \gamma Lx_n) \in \partial f(x_n) \\ u_{n+1} = u_n + w_n - v_n \end{cases}$$

using $y_n = \gamma^{-1}(u_n - v_n)$ and $z_n = \gamma^{-1}v_n$

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$$(\forall n \in \mathbb{N}) \quad \begin{cases} y_n = \text{prox}_{\gamma^{-1}g}(\gamma^{-1}u_n) \\ L^*(y_n - z_n - Lx_n) \in \frac{1}{\gamma}\partial f(x_n) \\ u_{n+1} = \gamma z_n + \gamma Lx_n \end{cases}$$

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Primal-dual optimization algorithm : $\min_x f(x) + h(x) + g(Lx)$

Condat-Vũ algorithm:

$$(\forall n \in \mathbb{N}) \quad \begin{cases} p_n = \text{prox}_{\tau f}(x_n - \tau(\nabla h(x_n) + L^* v_n)) \\ q_n = \text{prox}_{\sigma g^*}(v_n + \sigma(L(2p_n - x_n))) \\ (x_{n+1}, v_{n+1}) = (x_n, v_n) + \lambda_n((p_n, q_n) - (x_n, v_n)). \end{cases}$$

- Remark:

- * $h \in C_{\zeta}^{1,1}(\mathcal{H})$

- * No operator inversion.

- * Allow the use of proximable or/and differentiable functions.

- * Convergence when $\frac{1}{\tau} - \sigma\|L\|^2 > \frac{\zeta}{2}$.

- * $x_n \rightharpoonup \hat{x} \in \text{Argmin}(f + h + g \circ L)$

Primal-dual optimization algorithm : $\min_x f(x) + h(x) + g(Lx)$

Condat-Vũ algorithm: \Rightarrow **Chambolle-Pock algorithm**

$$(\forall n \in \mathbb{N}) \quad \begin{cases} p_n = \text{prox}_{\tau f}(x_n - \tau(\nabla h(x_n) + L^* v_n)) \\ q_n = \text{prox}_{\sigma g^*}(v_n + \sigma(L(2p_n - x_n))) \\ (x_{n+1}, v_{n+1}) = (x_n, v_n) + \lambda_n((p_n, q_n) - (x_n, v_n)). \end{cases}$$

- Remark:

* When $h = 0$, $\lambda_n \equiv 1$ and $\sigma\tau\|L\|^2 < 1$, this yields the Chambolle-Pock algorithm.

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$$(\forall n \in \mathbb{N}) \quad \begin{cases} x_{n+1} = \text{prox}_{\tau f}(x_n - \tau L^* v_n) \\ y_n = 2x_{n+1} - x_n \\ v_{n+1} = \text{prox}_{\sigma g^*}(v_n + \sigma L y_n). \end{cases}$$

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Condat-Vũ algorithm: \Rightarrow Forward-backward algorithm

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- Remark:

- * When $h = 0$, $\lambda_n \equiv 1$ and $\sigma\tau\|L\|^2 < 1$, this yields the Chambolle-Pock algorithm.

- * When $g = 0$ and $L = 0$, this yields the forward-backward algorithm.

Primal-dual optimization algorithm : $\min_x f(x) + h(x) + g(Lx)$

Condat-Vũ algorithm: \Rightarrow Forward-backward algorithm

$$(\forall n \in \mathbb{N}) \quad \begin{cases} p_n = \text{prox}_{\tau f}(x_n - \tau \nabla h(x_n)) \\ x_{n+1} = x_n + \lambda_n(p_n - x_n). \end{cases}$$

- Remark:

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Primal-dual optimization algorithm : $\min_x f(x) + h(x) + g(Lx)$

Condat-Vũ algorithm: \Rightarrow Douglas-Rachford algorithm

$$(\forall n \in \mathbb{N}) \quad \begin{cases} p_n = \text{prox}_{\tau f}(x_n - \tau(\nabla h(x_n) + L^* v_n)) \\ q_n = \text{prox}_{\sigma g^*}(v_n + \sigma(L(2p_n - x_n))) \\ (x_{n+1}, v_{n+1}) = (x_n, v_n) + \lambda_n((p_n, q_n) - (x_n, v_n)). \end{cases}$$

- Remark:

- * When $h = 0$, $\lambda_n \equiv 1$ and $\sigma\tau\|L\|^2 < 1$, this yields the Chambolle-Pock algorithm.

- * When $g = 0$ and $L = 0$, this yields the forward-backward algorithm.

- * In the limit case when $h = 0$, $\lambda_n \equiv 1$, $L = \text{Id}$ and $\sigma = 1/\tau$, this yields the Douglas-Rachford algorithm.

Primal-dual optimization algorithm : $\min_x f(x) + h(x) + g(Lx)$

Condat-Vũ algorithm: \Rightarrow Douglas-Rachford algorithm

$$(\forall n \in \mathbb{N}) \quad \begin{cases} x_{n+1} = \text{prox}_{\tau f}(x_n - \tau v_n) \\ s_n = \text{prox}_{\tau g}(2x_{n+1} - (x_n - \tau v_n)) \\ x_{n+1} - \tau v_{n+1} = (x_n - \tau v_n) + s_n - x_{n+1} \end{cases}$$

- Remark:

- * When $h = 0$, $\lambda_n \equiv 1$ and $\sigma\tau\|L\|^2 < 1$, this yields the Chambolle-Pock algorithm.

- * When $g = 0$ and $L = 0$, this yields the forward-backward algorithm.

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Optimization algorithms

Forward-Backward	$f_1 + f_2$	f_1 grad. Lipschitz prox_{f_2}	[Combettes,Wajs,2005]
ISTA	$f_1 + f_2$	f_1 grad. Lipschitz $f_2 = \lambda \ \cdot\ _1$	[Daubechies et al, 2003]
Douglas-Rachford	$f_1 + f_2$	prox_{f_1} prox_{f_2}	[Combettes,Pesquet, 2007]
PPXA	$\sum_i f_i$	prox_{f_i}	[Combettes,Pesquet, 2008]
ADMM	$\sum_i f_i \circ L_i$	prox_{f_i} $(\sum_{i=1}^m L_i^* L_i)^{-1}$	[Eckstein, Yao, 2015]
Chambolle-Pock	$f_1 + f_2 \circ L_2$	prox_{f_1} prox_{f_2}	[Chambolle, Pock, 2011]
Condat-Vũ	$f_1 + f_2 \circ L_2 + f_3$	prox_{f_1} prox_{f_2} f_3 grad. Lipschitz	[Condat, 2013][Vũ, 2013]