## The Calculus of Concurrent Systems Labelled Transition System and Behavioural equivalence

Names (channels) 
$$a, b, c, \dots$$
Processes  $P ::= a.P \mid \overline{a}.P \mid P_1 \mid P_2 \mid P_1 + P_2 \mid \mathbf{0} \mid (\nu a)P \mid !a.P$ 
Actions  $\mu ::= a \mid \overline{a} \mid \tau$ 

T-inp  $\overline{a.P \stackrel{a}{\to} P}$  T-out  $\overline{a.P \stackrel{\overline{a}}{\to} P}$  T-suml  $\overline{P \stackrel{\mu}{\to} P'}$  T-sumr  $\overline{Q \stackrel{\mu}{\to} Q'}$ 

T-com1  $\overline{P \stackrel{a}{\to} P' \quad Q \stackrel{\overline{a}}{\to} Q'}$  T-com2  $\overline{P \mid Q \mid P \mid Q \mid P' \mid Q'}$ 

T-parl  $\overline{P \mid P \mid P' \mid Q}$  T-parr  $\overline{Q \mid P \mid Q \mid P \mid Q'}$ 

T-res  $\overline{Q \mid P \mid P \mid P' \mid Q}$  T-rep  $\overline{Q \mid P \mid Q \mid P \mid Q'}$  T-rep  $\overline{Q \mid P \mid Q \mid P \mid Q'}$ 

**Definition:** A relation  $\mathcal{R}$  between processes is a bisimulation if whenever  $P\mathcal{R}Q$ , we have the following:

- 1. if  $P \xrightarrow{\mu} P'$ , then there exists Q' such that  $Q \xrightarrow{\mu} Q'$  and  $P'\mathcal{R}Q'$ , and
- 2. if  $Q \xrightarrow{\mu} Q'$ , then there exists P' such that  $P \xrightarrow{\mu} P'$  and  $P'\mathcal{R}Q'$ .

**Definition:** Bisimilarity, written  $\sim$ , is a relation between processes defined as follows: we have  $P \sim Q$  if there exists a relation  $\mathcal{R}$  such that

- 1.  $\mathcal{R}$  is a bisimulation, and
- 2. PRQ.