

DETERMINISATION OF HISTORY-DETERMINISTIC AUTOMATA

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History-deterministic automata

- non-deterministic
- parity ω -word automata
- there exists $\sigma: A^* \rightarrow Q$ satisfying:
for every word $w \in L(\mathcal{A}) \subseteq A^\omega$
the sequence of states $\sigma(w) \in Q^\omega$
is an **accepting** run of \mathcal{A}

Open problem

What is the **state blow-up**
when **determinising**
a **history-deterministic automaton**?

Solution: exponential for **co-Büchi** automata
polynomial for **Büchi** automata

Applications

- synthesis (Büchi, Landweber '69)
- branching time verification (Emerson, Sistla '84)
- derived languages (Niwiński, Walukiewicz '98)
- symbolic representation (Henzinger, Piterman '06)
- quantitative models (Colcombet, Löding '10)
- ...

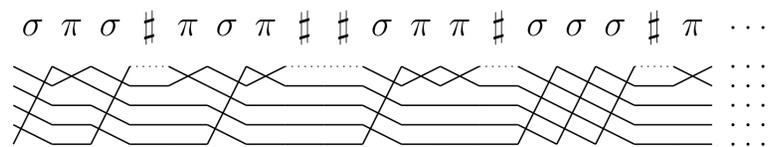
Co-Büchi automata

Theorem 1. There are **H-D** automata \mathcal{A}_n :

- $|\mathcal{A}_n| = \Theta(n)$ and \mathcal{A}_n is *co-Büchi*
- if \mathcal{B} is *deterministic* and $L(\mathcal{B}) = L(\mathcal{A}_n)$ then $|\mathcal{B}| \geq 2^n$

History-deterministic automata are **succinct**!

Proof scheme: Language of permutations:



+ **compactness** for **pumping** the limitary behaviour.

Is \mathcal{A} history-deterministic?

New complexity bounds:

- for **co-Büchi**: in **P**
- for **Büchi**: in **NP**
- for **parity**: at least as hard as **parity games**
(**NP** \cap **co-NP**)

Technical details at **arXiv.org** soon!

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Büchi automata

Theorem 2. Every **H-D** Büchi automaton \mathcal{A} admits \mathcal{B} :

- \mathcal{B} is *deterministic* and *Büchi*
- $L(\mathcal{A}) = L(\mathcal{B})$ and $|\mathcal{B}| \leq |\mathcal{A}|^2$

Polynomial determinisation procedure!

Proof scheme:

- **residual** languages and brutal powerset determinisation
- **rank signatures** of Walukiewicz
- iterative **normalization** of \mathcal{A}
- **dependency graph** over the automaton