

The Complexity of Grundy Coloring and its Variants

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Warm Up

Connected Grundy Coloring

Weak Grundy Coloring

Grundy Colorings

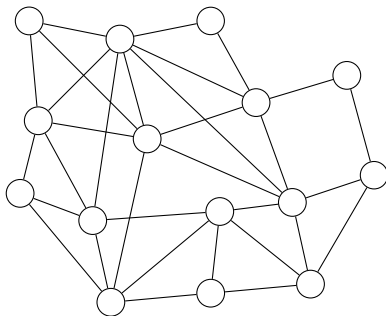
- ▶ Order the vertices $v_1, v_2 \dots v_n$ to maximize the number of colors used by the greedy coloring.
- ▶ That is, v_i is colored with $c(v_i)$ the first color that is *not* in its neighborhood.

Grundy Colorings

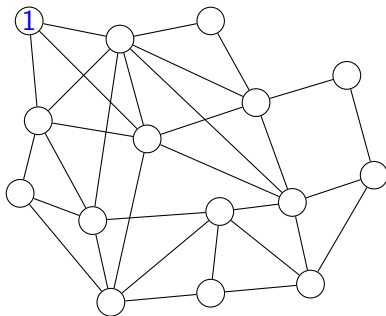
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Grundy Colorings

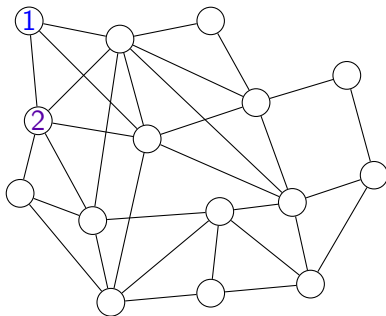
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- ▶ That is, v_i is colored with $c(v_i)$ the first color that is *not* in its neighborhood.
- ▶ Connected version: $\forall i, G[v_1 \cup \dots \cup v_i]$ is connected.
- ▶ Weak version: v_i can be colored with any color in $\{1, \dots, c(v_i)\}$.



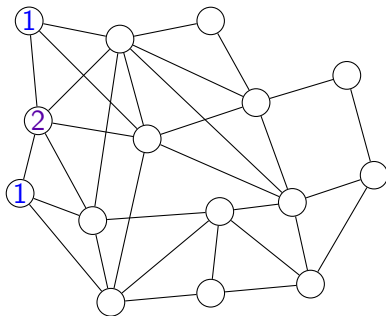
Can you achieve color 6?



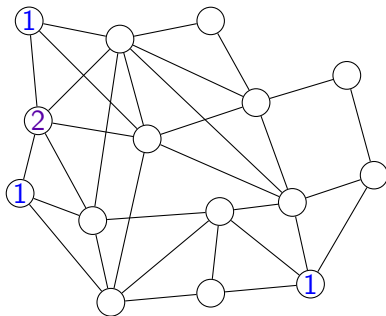
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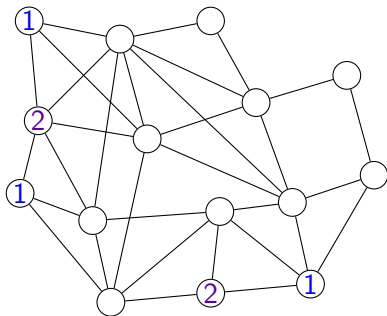
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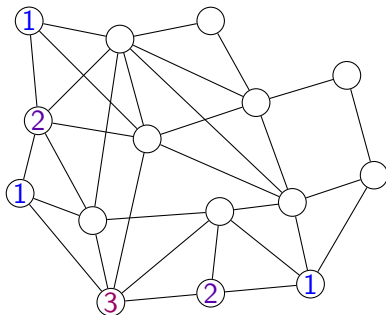
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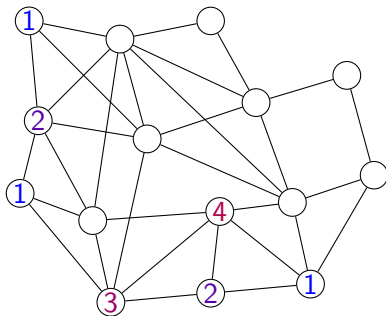
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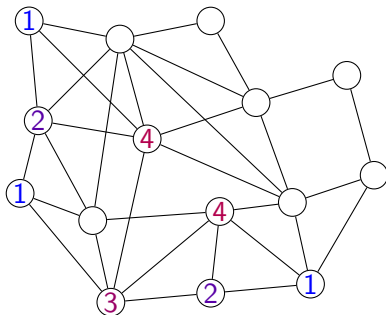
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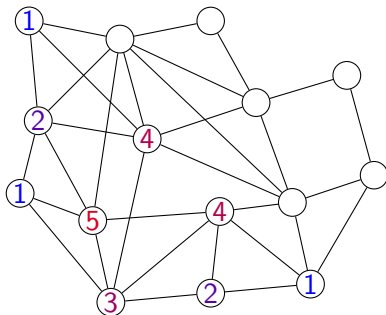
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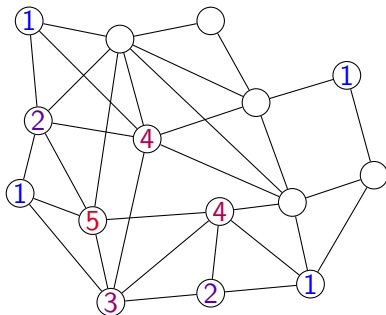
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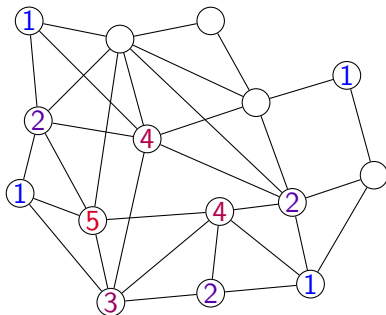
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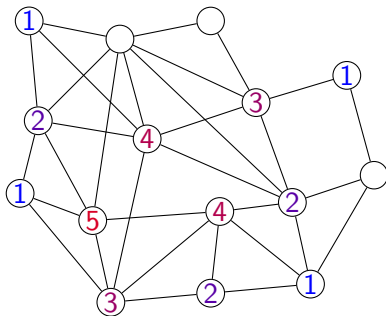
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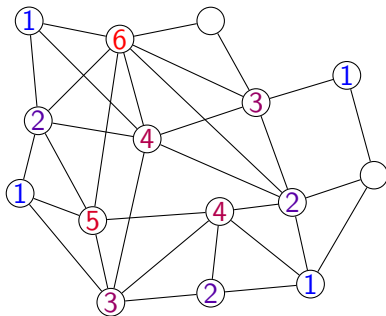
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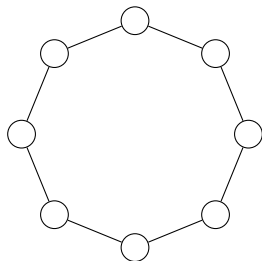
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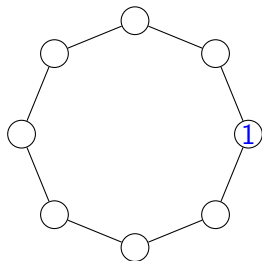
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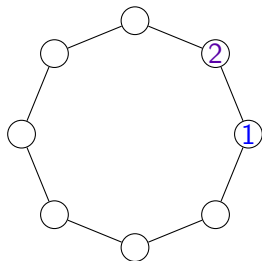
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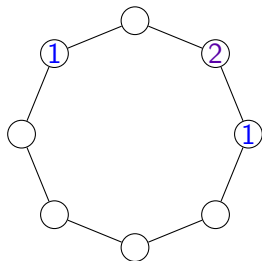
Grundy number = 3



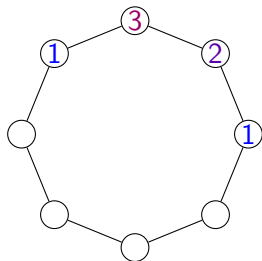
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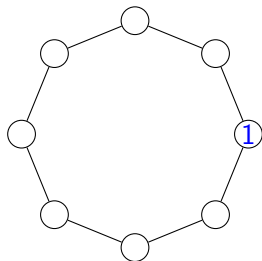
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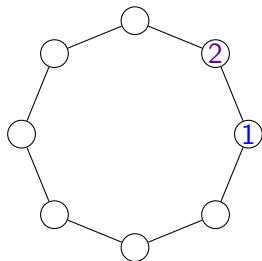
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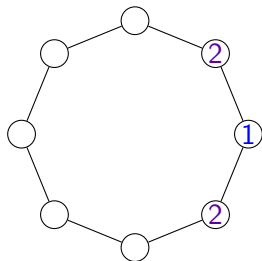
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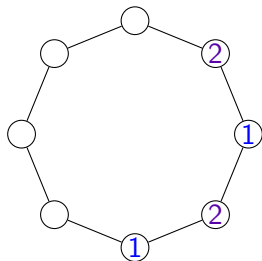
Connected Grundy number = 2



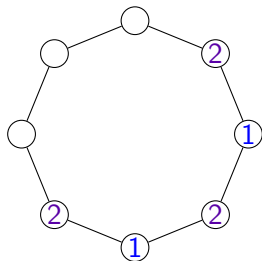
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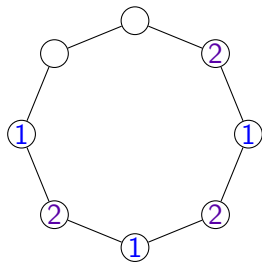
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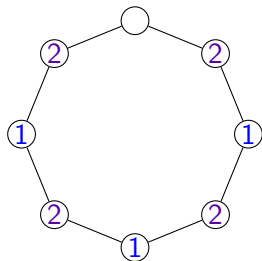
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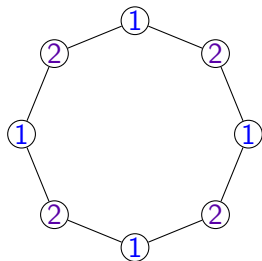
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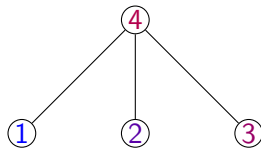
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How many vertices (at most) did we need to achieve color k ?

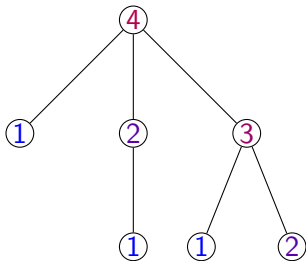
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④

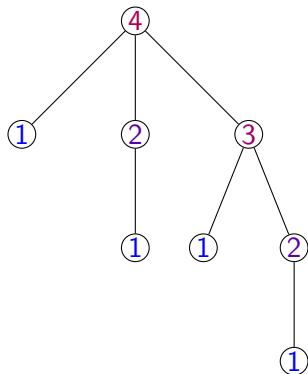
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$t_1 = 1$ and $t_k = \sum_{1 \leq i \leq k-1} t_i$.
So, $t_k = 2^{k-1}$.

Theorem (Zaker '05)

The Grundy number can be computed in $O(f(k)n^{2^{k-1}})$.

XP algorithm: $O(f(k)n^{g(k)})$

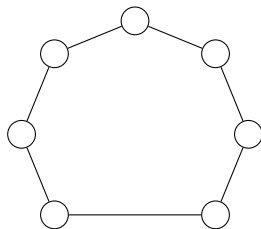
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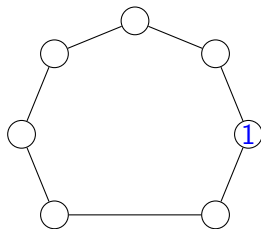
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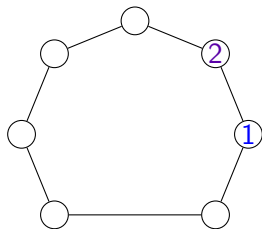
Can we do the same for the connected Grundy number?



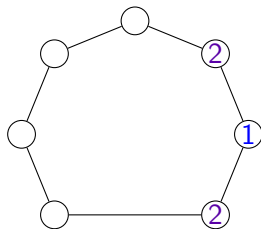
Connected Grundy number = 3, unbounded witness



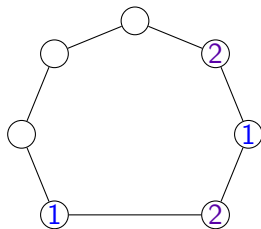
Connected Grundy number = 3, unbounded witness



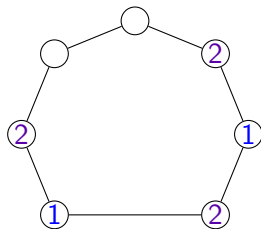
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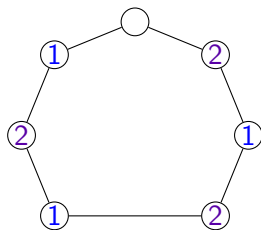
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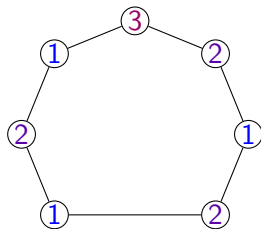
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Theorem (BCDGMSS '14)

CONNECTED GRUNDY COLORING *is NP-complete.*

Theorem

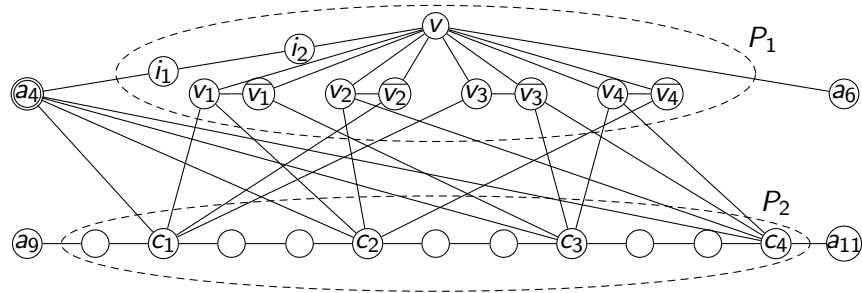
CONNECTED GRUNDY COLORING *is NP-complete even for $k = 7$.*

- ▶ Reduction from 3SAT-3OCC.

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coloring such a vertex by 3 \equiv setting the literal to true.

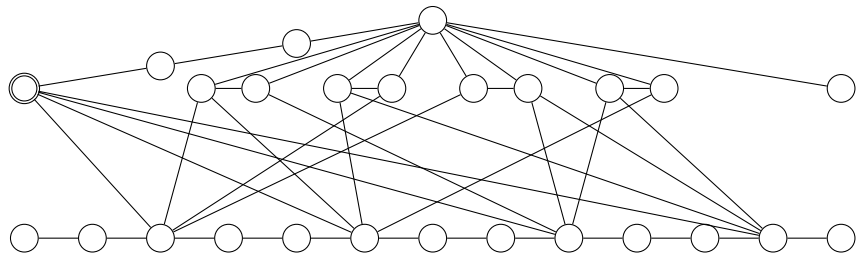
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- ▶ To achieve color 7, three special neighbors of the c_j s should be colored by 1, 2 and 3 respectively.

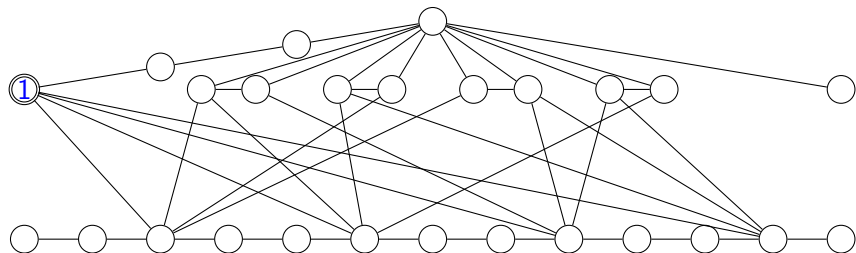


P_1 and P_2 for the instance

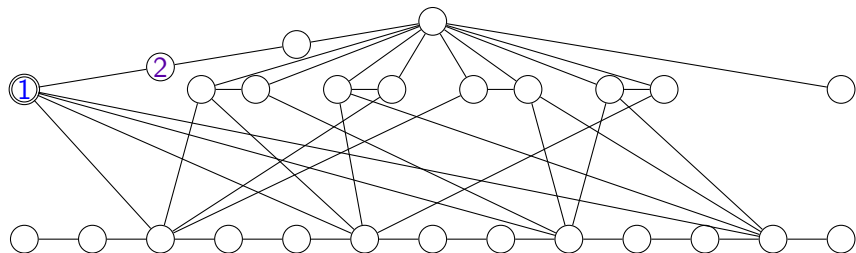
$$\{x_1 \vee \neg x_2 \vee x_3\}, \{x_1 \vee x_2 \vee \neg x_4\}, \{\neg x_1 \vee x_3 \vee x_4\}, \{x_2 \vee \neg x_3 \vee x_4\}.$$



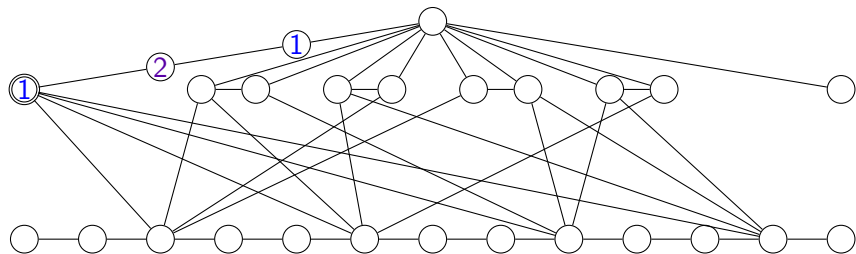
A connected Grundy coloring setting all the c_j s to 4.



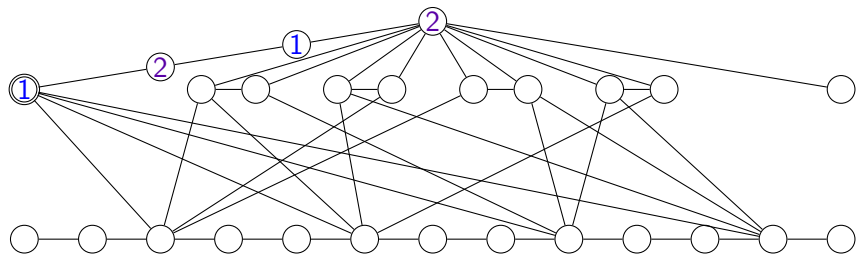
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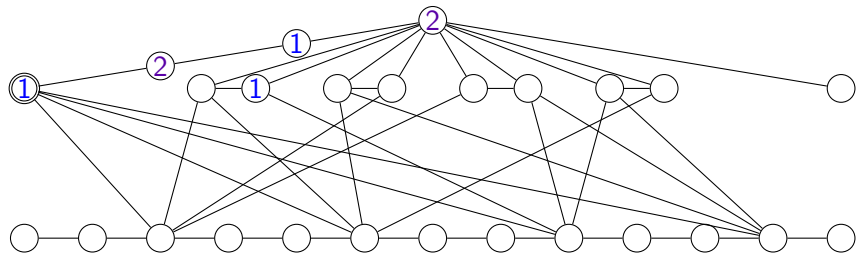
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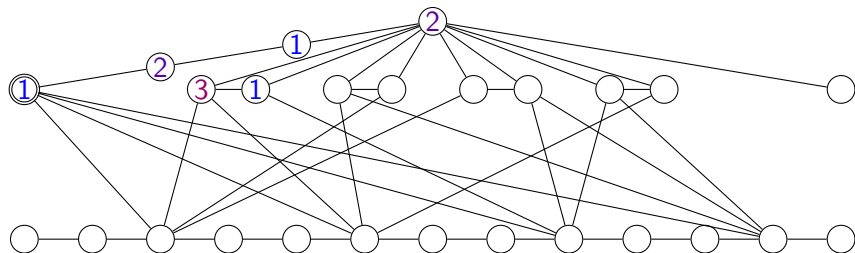
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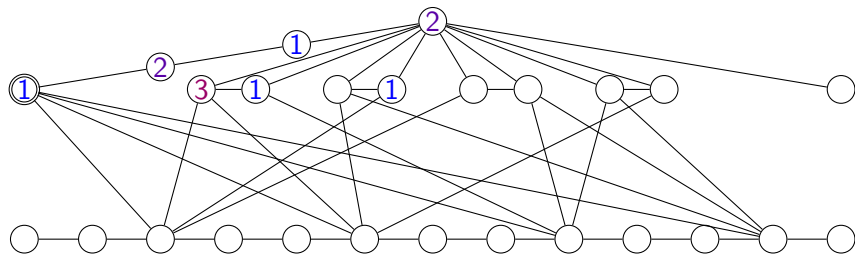
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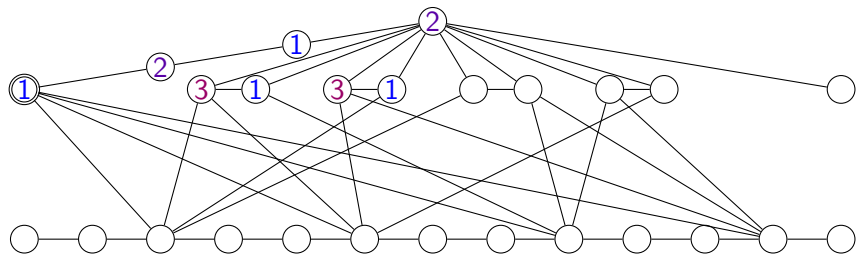
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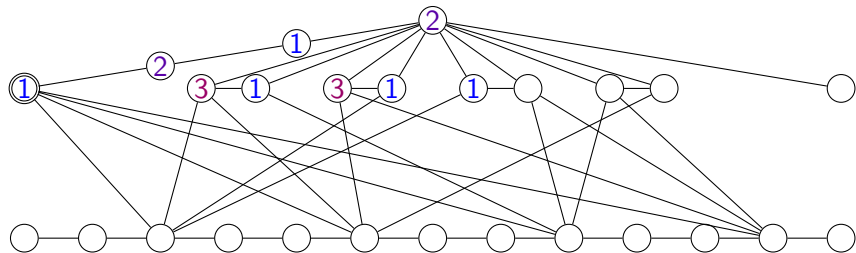
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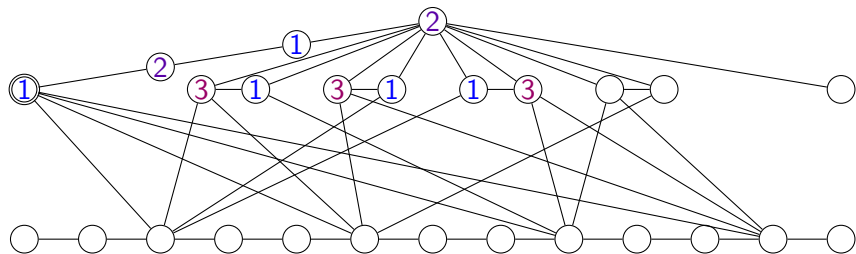
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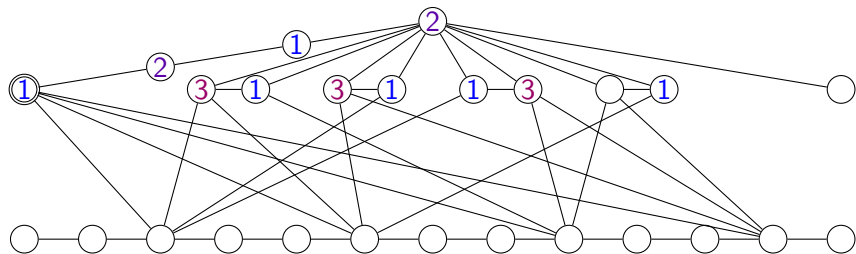
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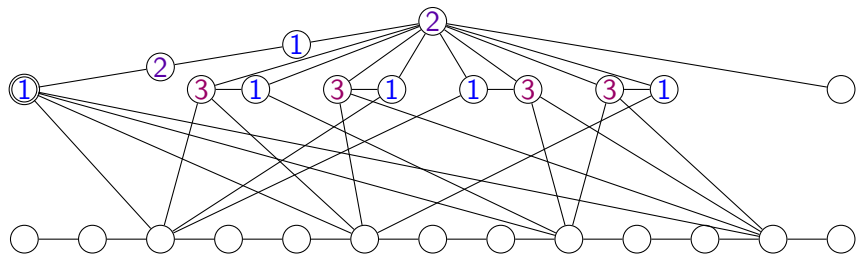
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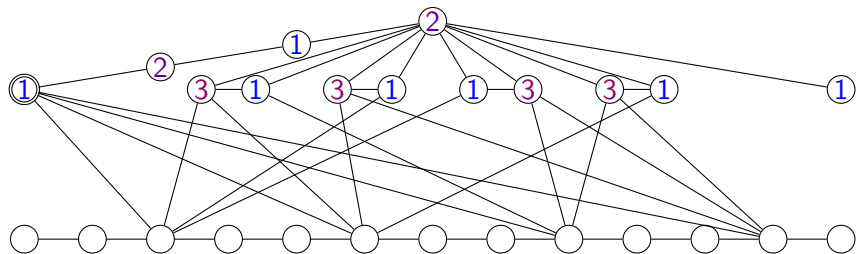
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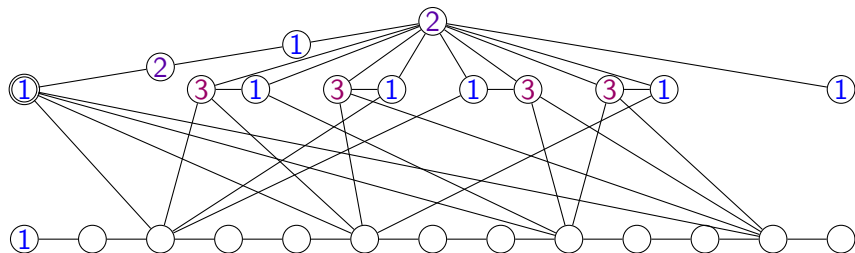
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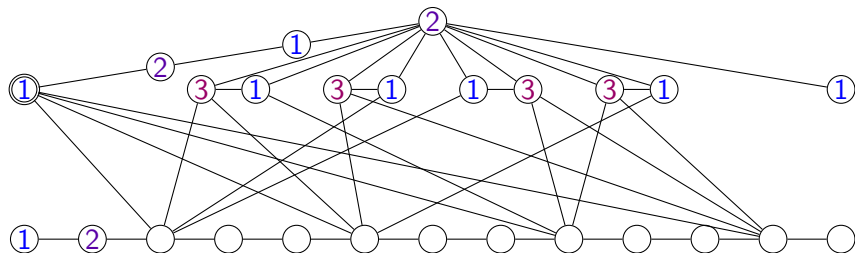
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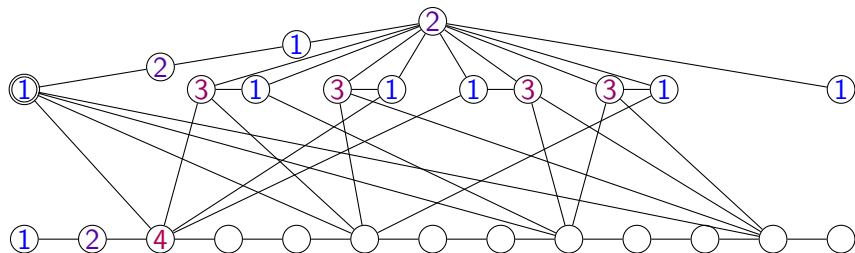
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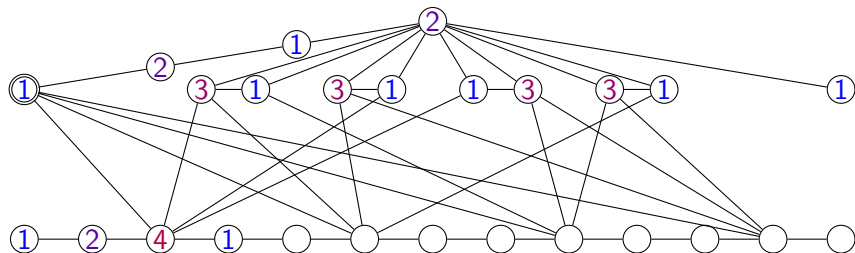
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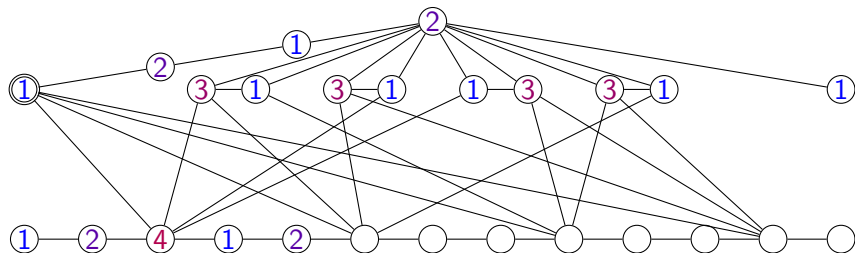
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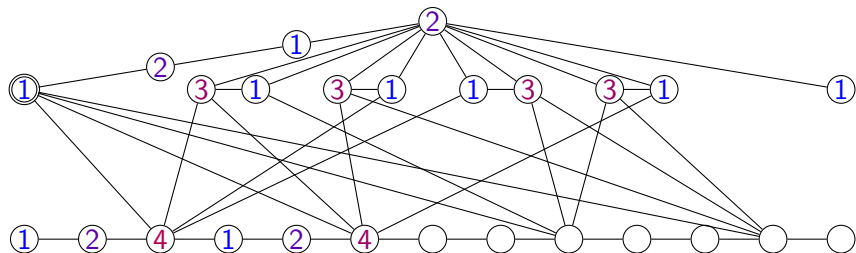
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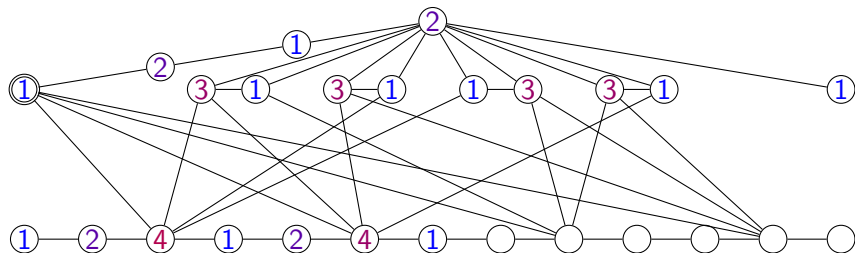
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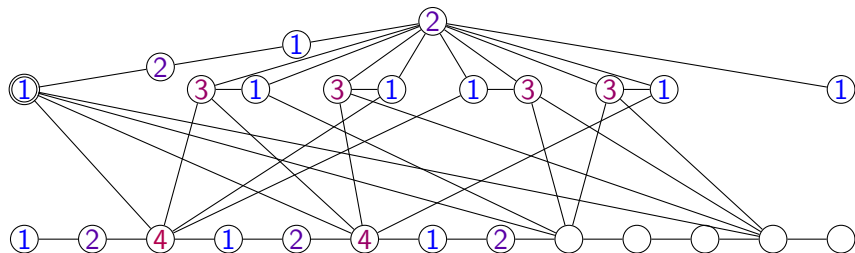
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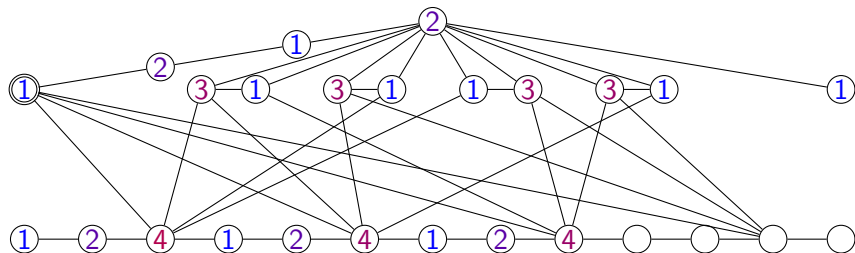
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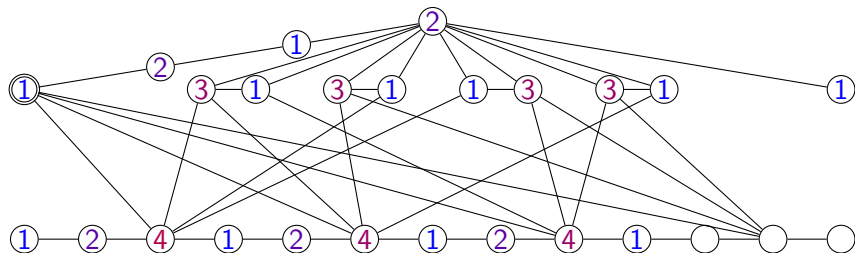
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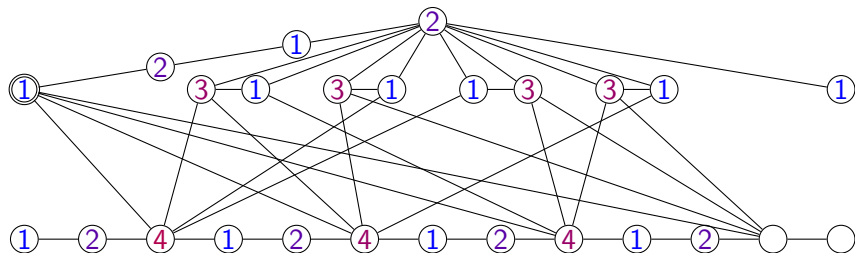
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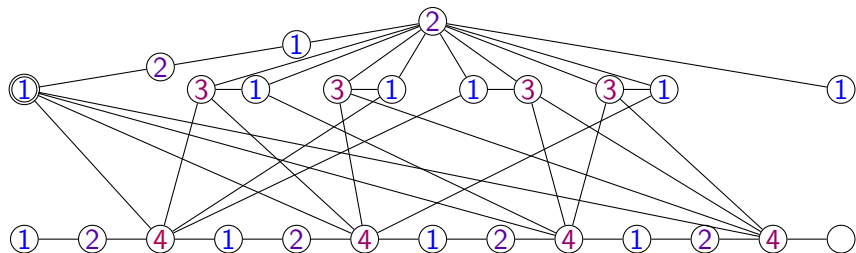
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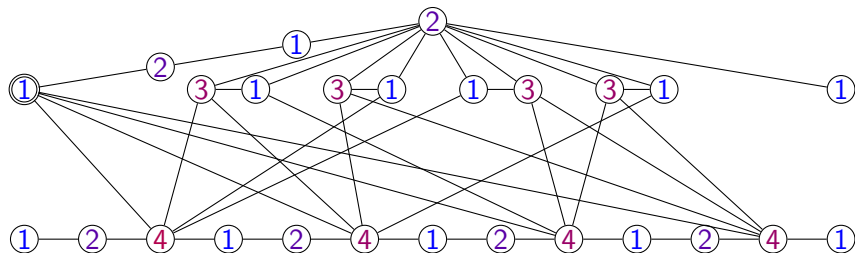
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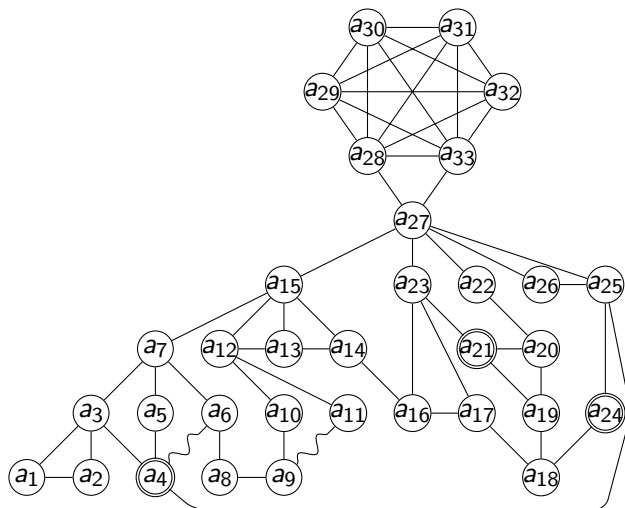
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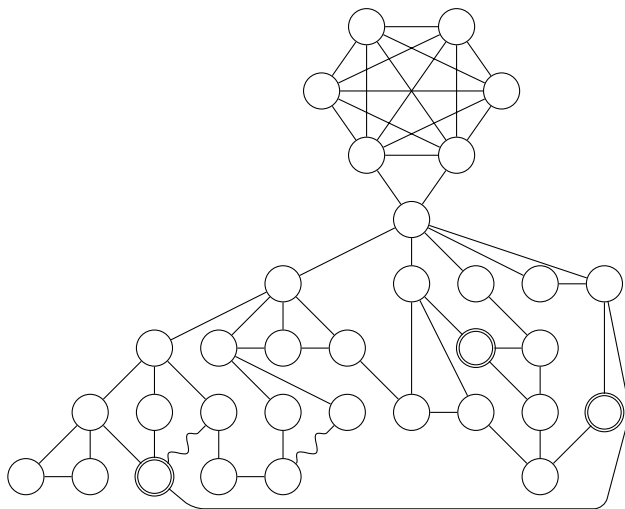
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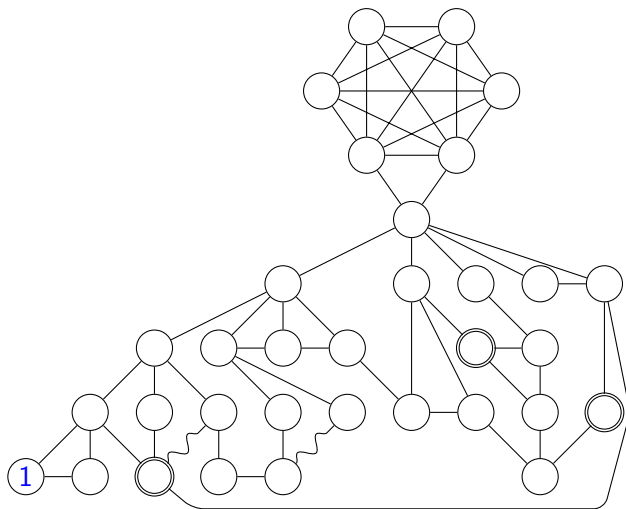
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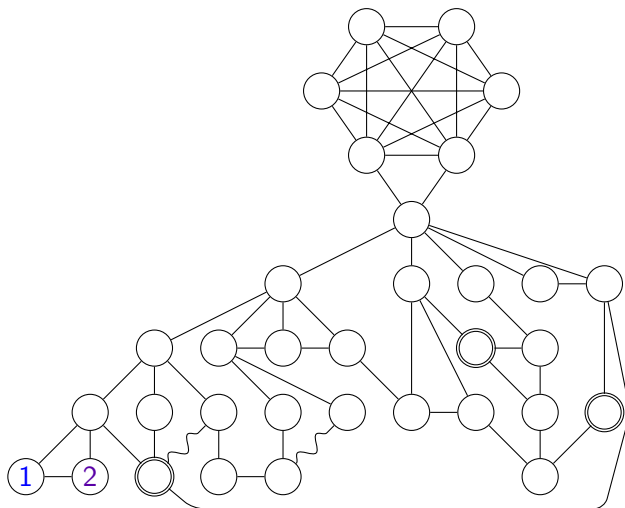
The doubly-circled vertices are linked to all the *clause* vertices c_j s.



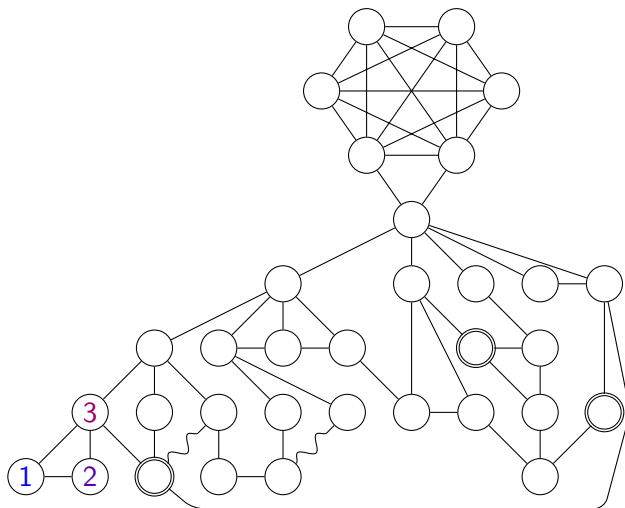
A connected Grundy coloring achieving color 7.



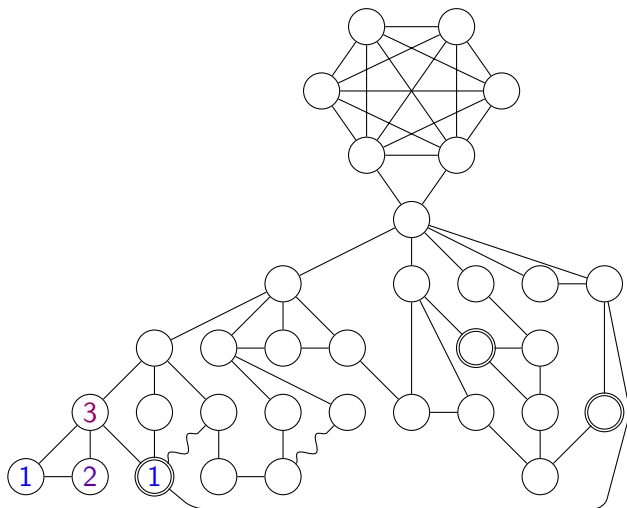
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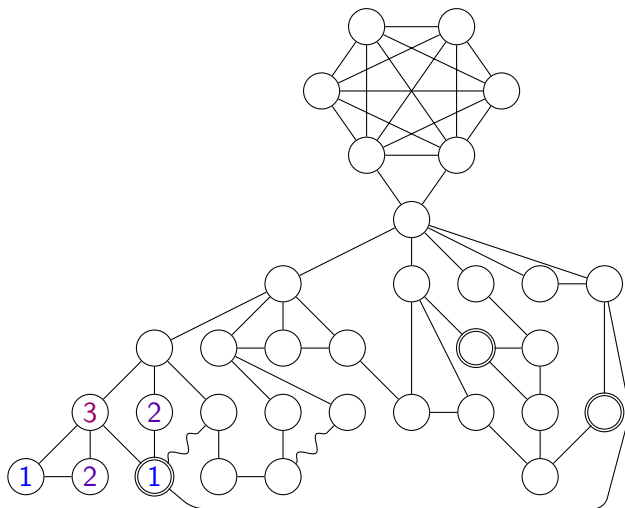
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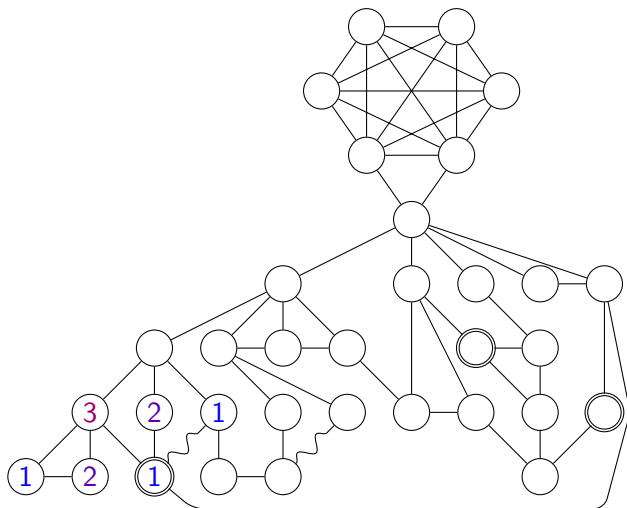
A connected Grundy coloring achieving color 7.



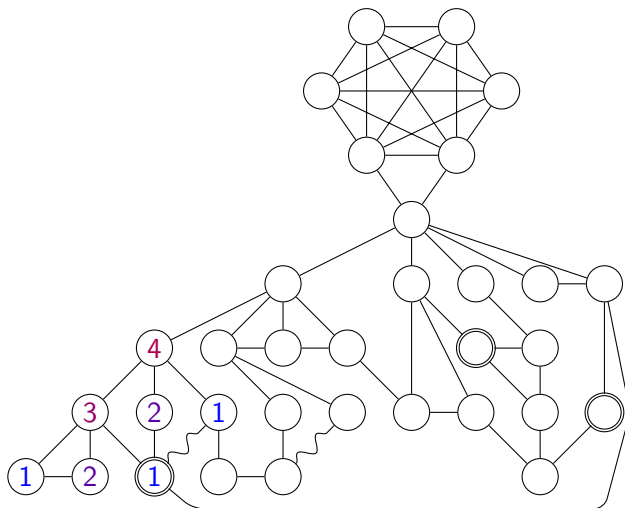
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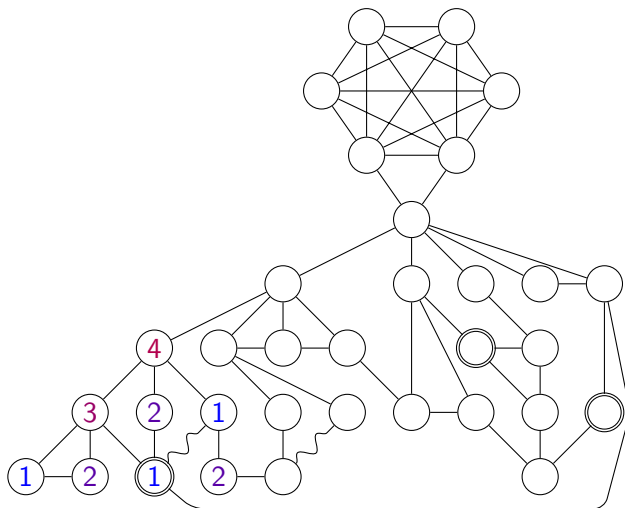
A connected Grundy coloring achieving color 7.



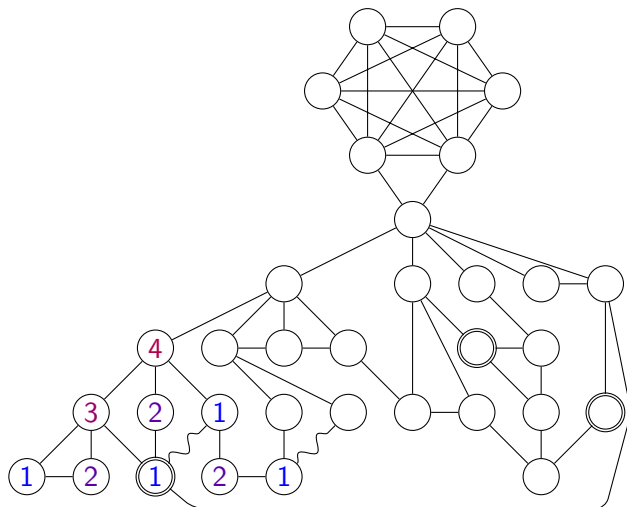
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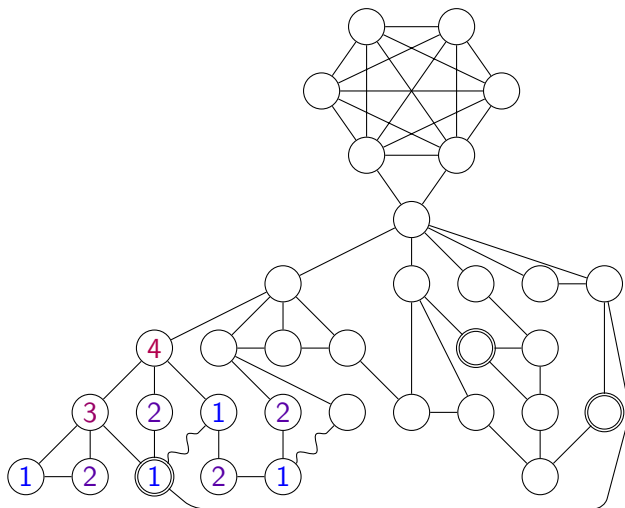
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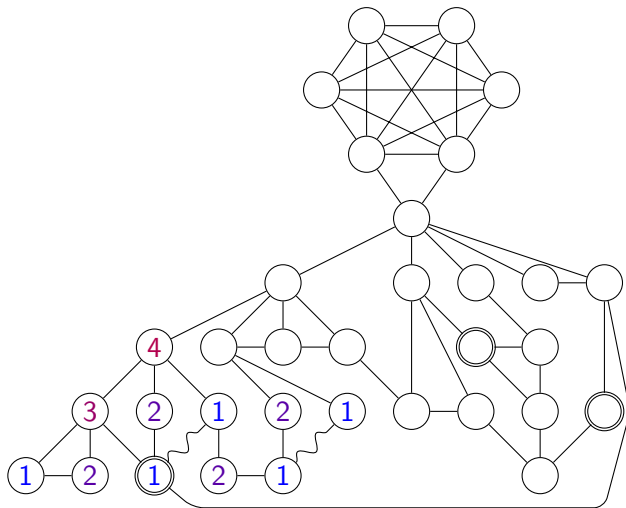
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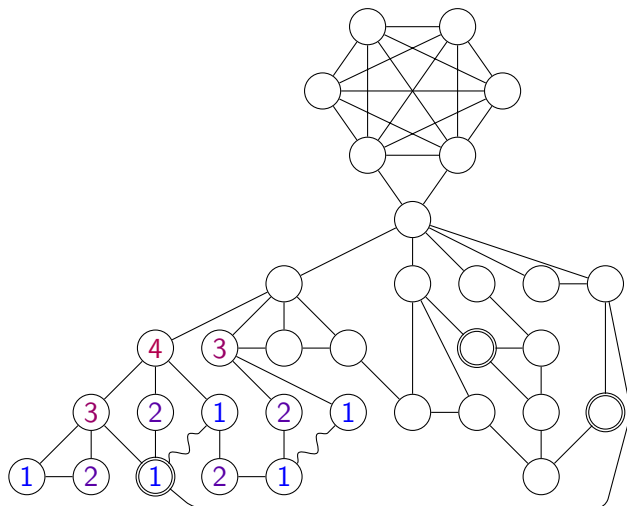
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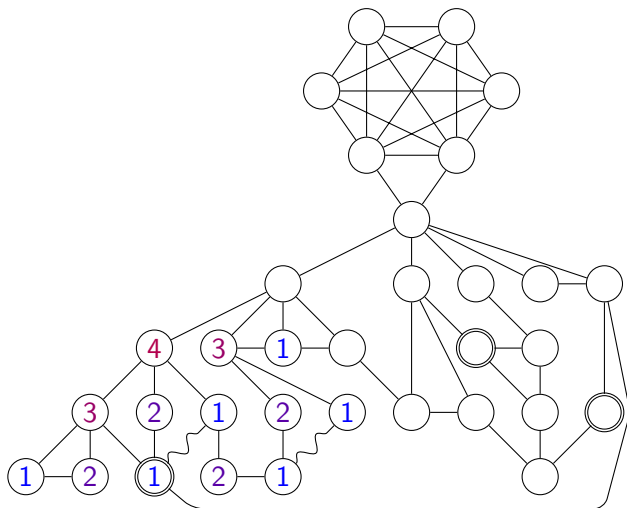
A connected Grundy coloring achieving color 7.



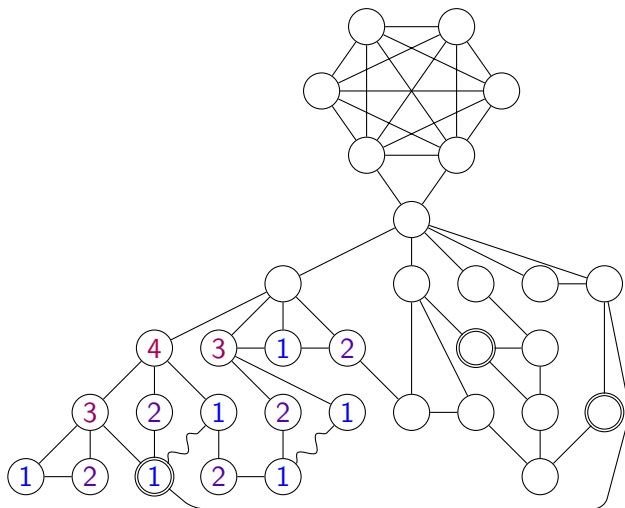
A connected Grundy coloring achieving color 7.



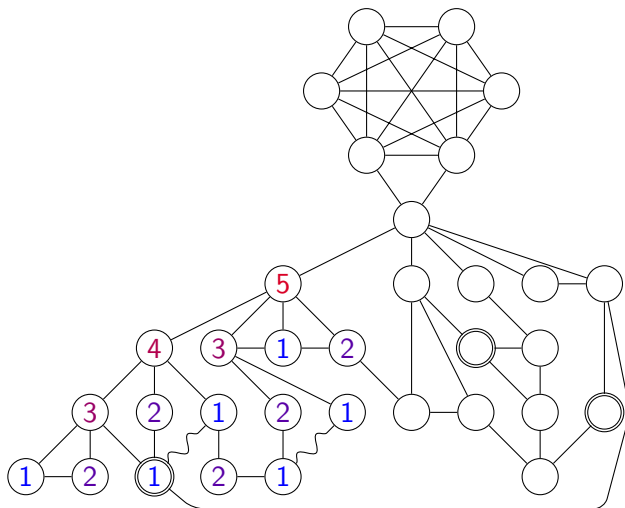
A connected Grundy coloring achieving color 7.



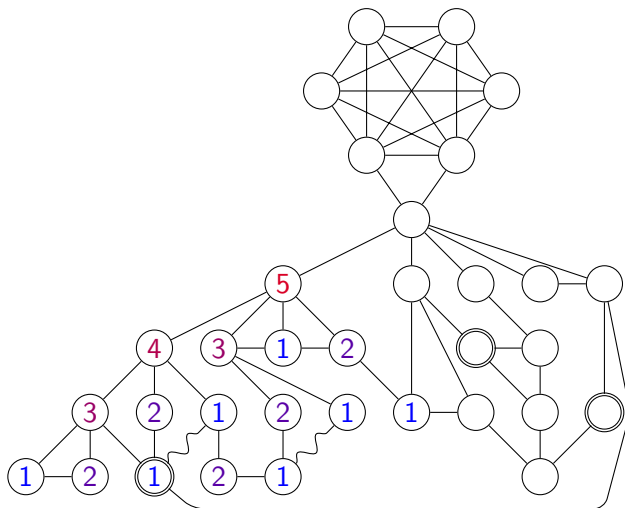
A connected Grundy coloring achieving color 7.



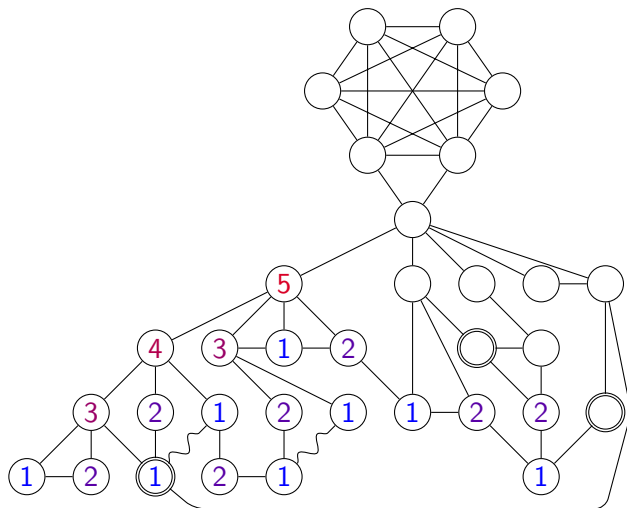
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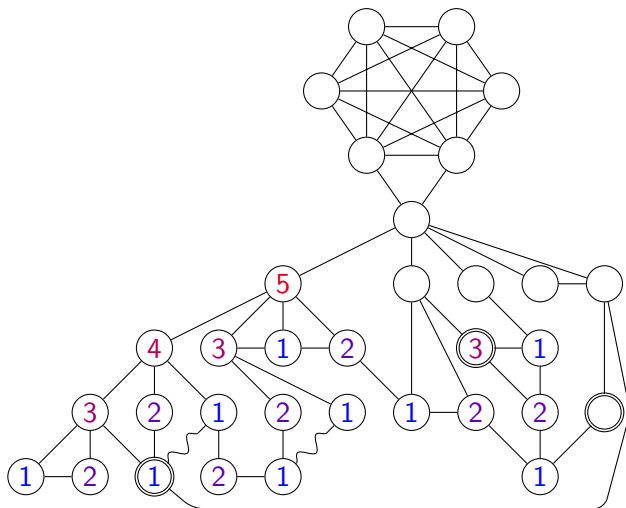
A connected Grundy coloring achieving color 7.



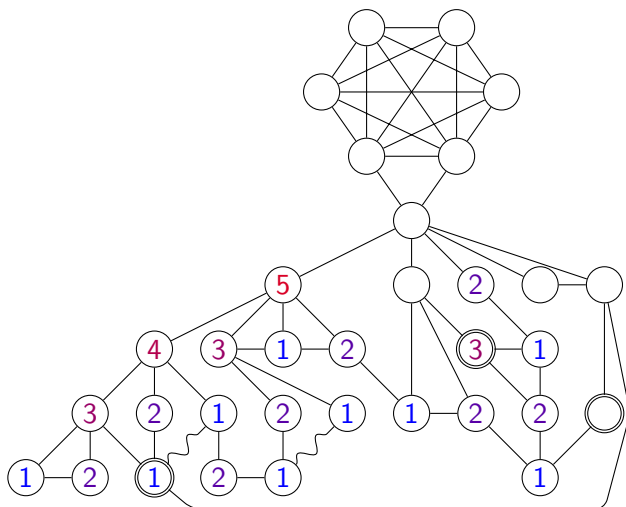
A connected Grundy coloring achieving color 7.



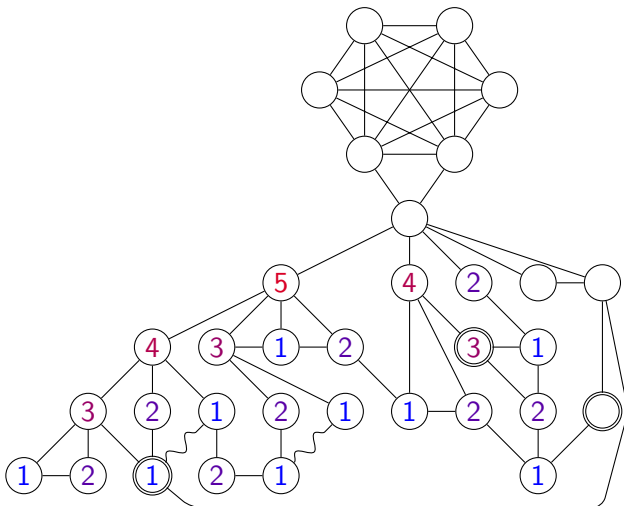
A connected Grundy coloring achieving color 7.



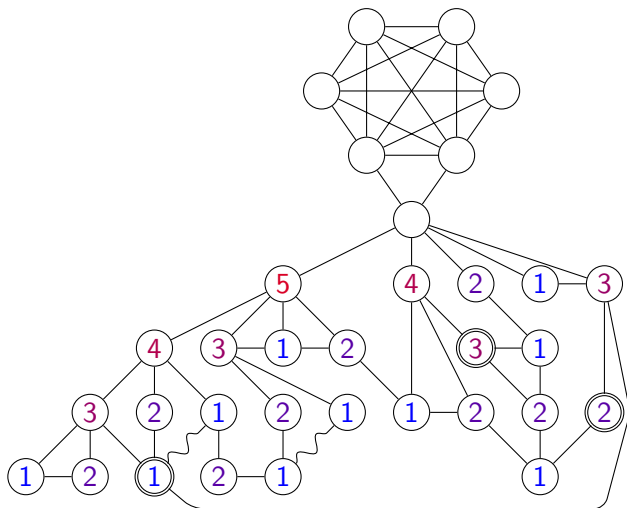
A connected Grundy coloring achieving color 7.



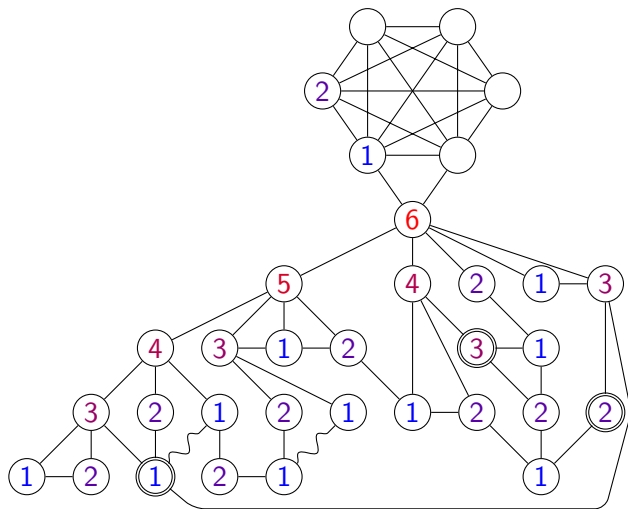
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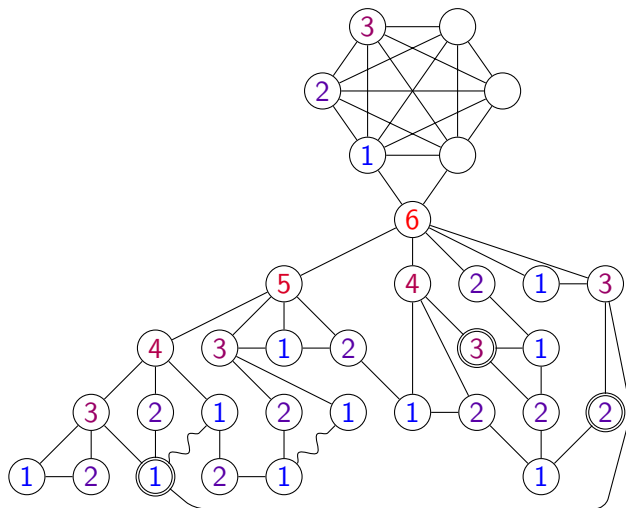
A connected Grundy coloring achieving color 7.



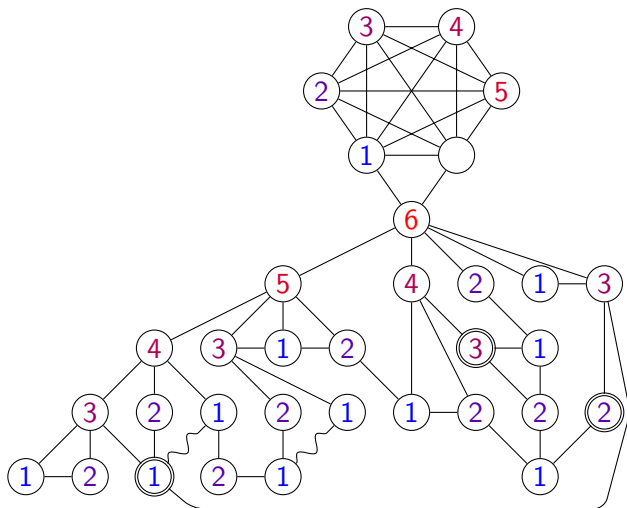
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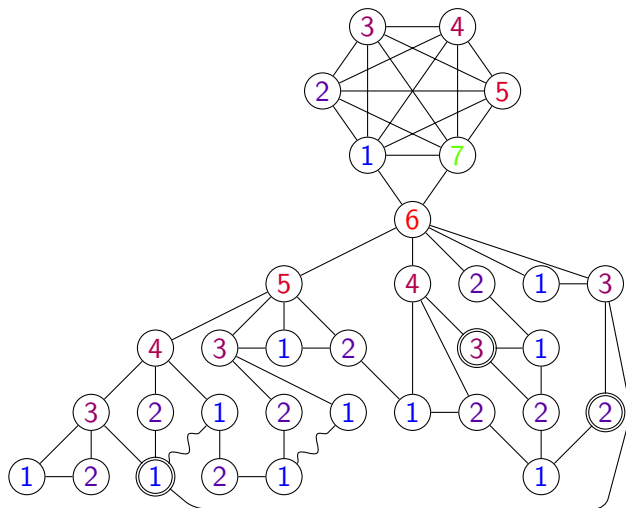
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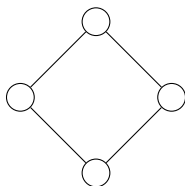
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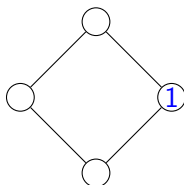
Are there graphs where weak Grundy exceeds Grundy number?

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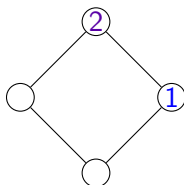
Weak Grundy number = 3, (connected) Grundy number = 2.

Are there graphs where weak Grundy exceeds Grundy number?



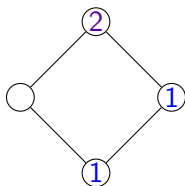
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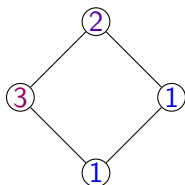
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Color Coding

- ▶ Add *colors* at random to the instance such that the colors *enhance* an optimal solution with probability $p(k)$.

Color Coding

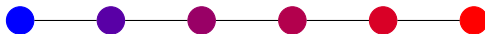
- ▶ Add *colors* at random to the instance such that the colors *enhance* an optimal solution with probability $p(k)$.
- ▶ If we try $\frac{100}{p(k)}$ times, we always fail with probability $(1 - \frac{1}{e})^{100}$ (that's not happening).

Color Coding

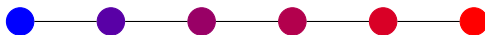
- ▶ Add *colors* at random to the instance such that the colors *enhance* an optimal solution with probability $p(k)$.
- ▶ If we try $\frac{100}{p(k)}$ times, we always fail with probability $(1 - \frac{1}{e})^{100}$ (that's not happening).
- ▶ Solving the instance is *easier* with this extra information.

Now, can you propose an algorithm in $O(f(k)n^c)$ for finding a path of length k in a graph?

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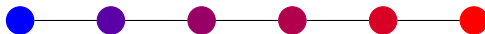


Now, can you propose an algorithm in $O(f(k)n^c)$ for finding a path of length k in a graph?



- ▶ How many times do you repeat the random coloring?

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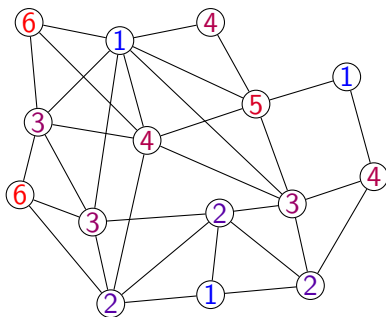
- ▶ How many times do you repeat the random coloring?
- ▶ How do find a path of length k in colored instances?

Theorem

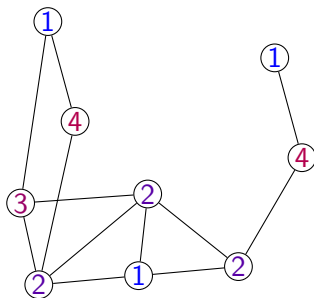
WEAK GRUNDY COLORING *is in* FPT.

FPT algorithm: $O(f(k)n^c)$.

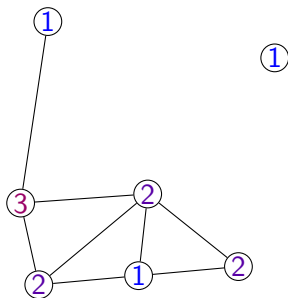
Guess #1



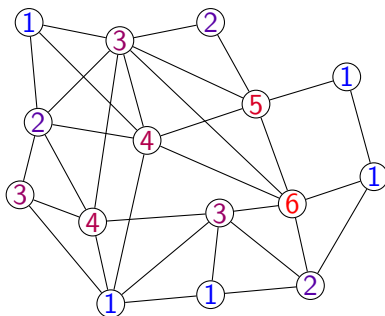
Guess #1



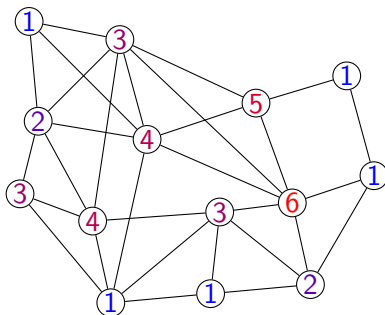
Guess #1



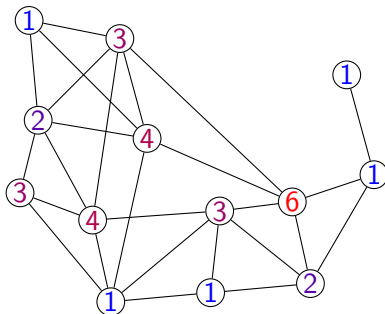
Guess #2



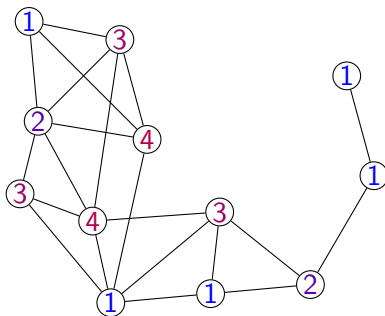
Guess #2



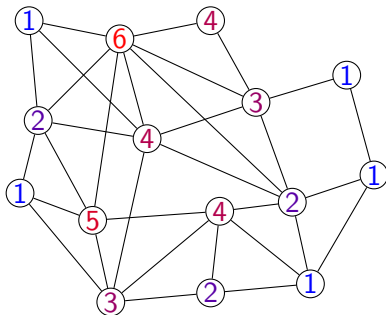
Guess #2



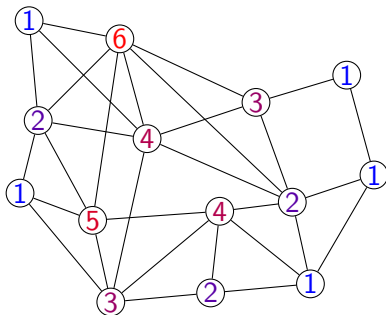
Guess #2



... $O(k^{2^k})$ unsuccessful guesses later ...



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Open Problems

- ▶ Is GRUNDY COLORING solvable in $O(f(k)n^c)$?
- ▶ What is the complexity of CONNECTED GRUNDY COLORING for $k = 4$, $k = 5$ and $k = 6$?
- ▶ What is the complexity of WEAK CONNECTED GRUNDY COLORING for k constant?