



LIP Laboratory
ENS-Lyon - France



Session Awareness Issues for Next Generation Cluster based Network Traffic Load Balancing Frameworks

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*The fifth ACS/IEEE International Conference on Computer Systems and Applications
AICCSA 2007 – Amman, Jordan, May 13-16, 2007*



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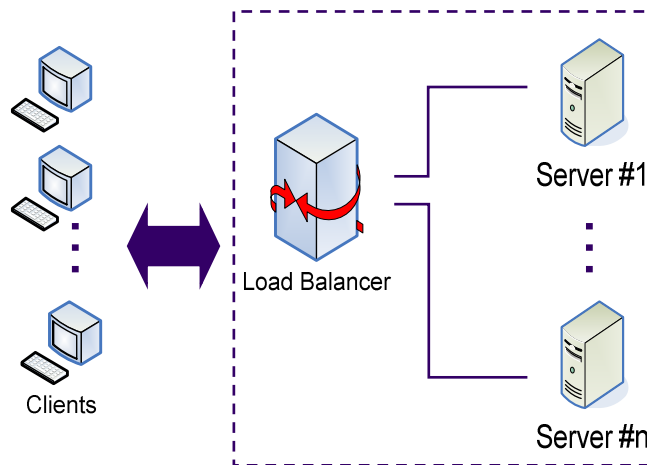
Agenda &

- ➔ Context
- ➔ Network traffic load balancing
- ➔ Session aware network traffic load balancing
- ➔ The requirements of NGN cluster based load balancing frameworks
- ➔ Conclusion & perspectives

➔ Clusters are a set of networked servers which

- Offer the same single system image to clients while providing additional processing capabilities
- Improve the scalability and the availability of the rendered service

➔ Operators are replacing single servers by clusters of servers



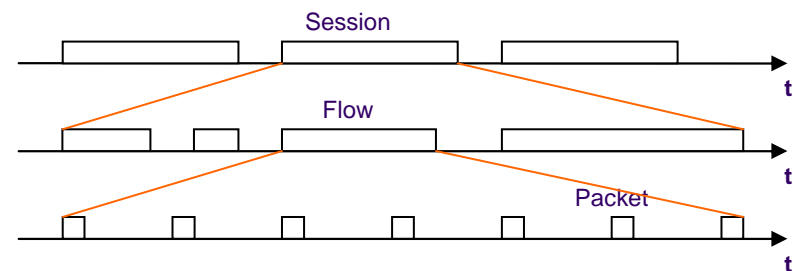
Context 2/2 &

- ➔ How to **efficiently spread** the offered network traffic in a cluster of servers while **preventing the interruption** and the **QoS degradation** of the rendered service?
- ➔ A special focus on the requirements of multiple flow based sessions

- A **session** is defined as an association between two communicating end points

- It can span over single or over multiple flows

- NGN services
 - *Video streaming, Voice over IP, etc.*
- Familiar services
 - *File transfer, etc.*



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Network traffic load balancing



The approaches 1/2

➡ Stateless versus stateful engines

➤ Stateless

- Less latency versus
 - *Unfair throughput*
 - *Unavailability and session unawareness*

➤ Statefull

- More resource consuming (memory & CPU) versus
 - *Means to achieve fairness and session awareness while distributing the offered traffic*
 - *Means to detect and react to failures*

Network traffic load balancing

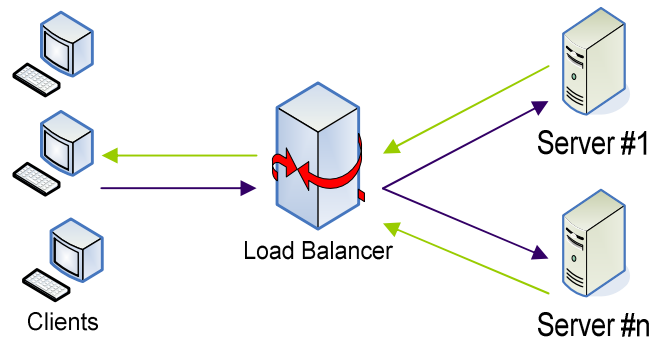


The approaches 2/2

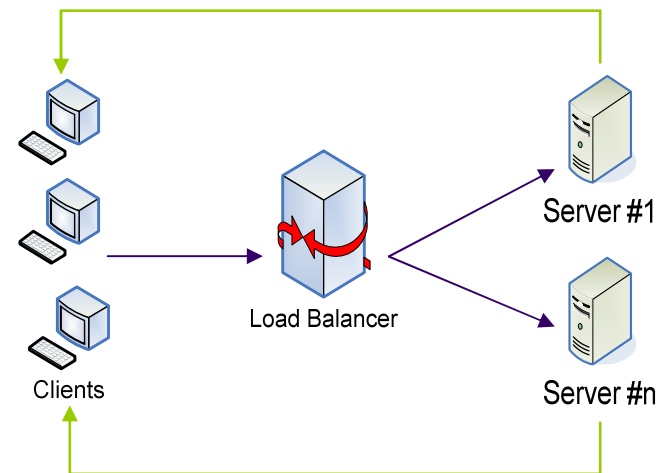
➔ One way versus two way based architectures

- Limited throughput and scalability issues when wasting the cluster entry point resources with the processing of the outgoing traffic

Two way architectures



One way architectures



— Incoming traffic
— Outgoing traffic

Network traffic load balancing



The mechanisms 1/8

➡ Network traffic load balancing falls into

- Flow aware engines
 - Provides flow level integrity
 - *Packet double rewriting*
 - *Packet forwarding*
 - *Packet tunnelling*
- Application aware engines
 - Uses DPI to provide application level integrity
 - *TCP Gateway*
 - *TCP Splicing variants*
 - *TCP handoff variants*

Network traffic load balancing

The mechanisms 2/8

Packet double rewriting

- NAT like operations
- Assumes that the cluster head as well as the cluster nodes share the same private network
- Suffers performance limitations
 - Scalability
- due to
 - The double processing overhead
 - Two way architecture based

Network traffic load balancing

The mechanisms 3/8

Packet forwarding

- Performs a link layer encapsulation of each incoming datagram
- The outgoing traffic will be directly forwarded to the clients
- Applies only to servers clustered within the same local network
- Assumes that the dispatcher statically maintains in-cluster ARP associations

Network traffic load balancing

The mechanisms 4/8

Packet tunneling

- Assumes that the entry point to the cluster as well as each node within the cluster supports tunneling
- Performs a network level encapsulation of each incoming datagram
- Applies to servers networked in different LANs
- Assumes that the entry point to the cluster as well as each node within the cluster supports tunneling

Network traffic load balancing

The mechanisms 5/8

➡ Application level load balancing involves more steps

- The traffic is forwarded in almost four steps
 - A connection between the client and the dispatcher is first established
 - Data is buffered for content inspection
 - Session states are established or updated
 - A cluster node is chosen and the traffic is forwarded to that node

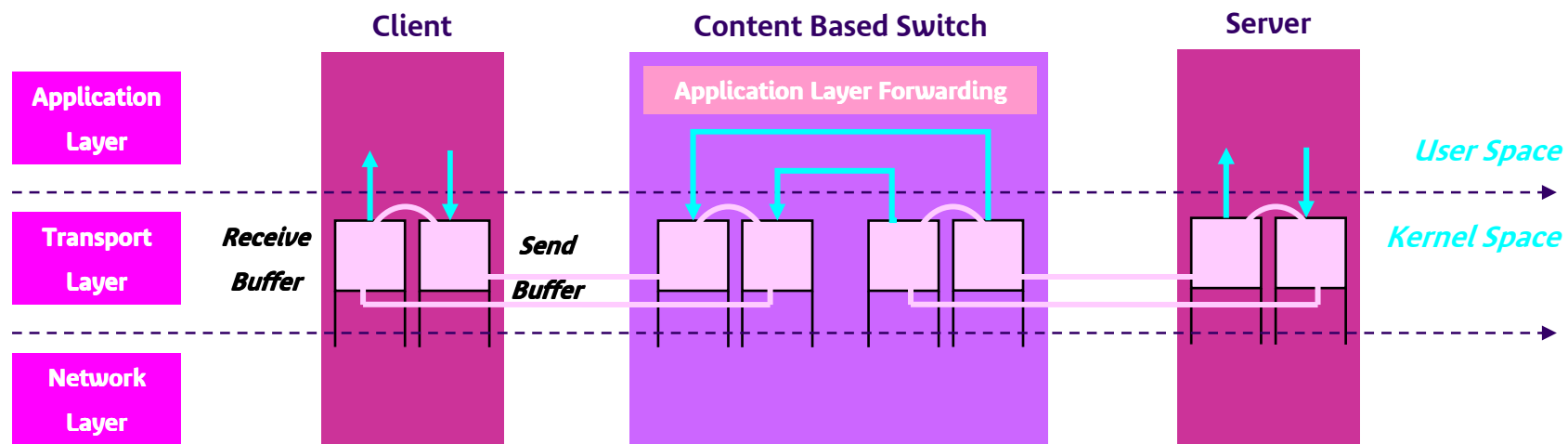
Network traffic load balancing &

The mechanisms 6/8

➔ TCP Gateway

➤ Cost effective

- Multiple copies and context switching

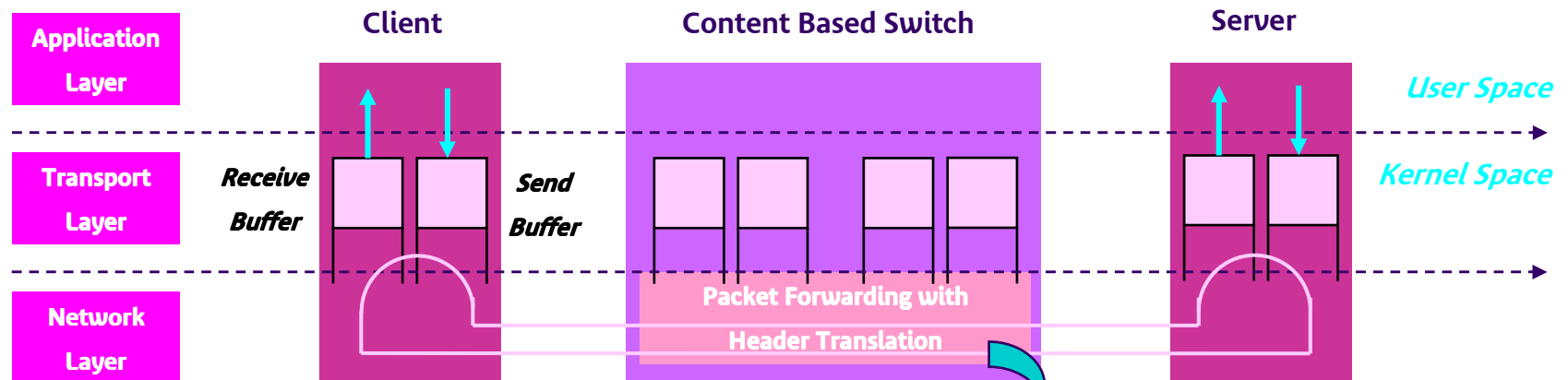


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Network traffic load balancing &

The mechanisms 7/8

➔ TCP Splicing



- Modifications also affect
 - IP pseudo Header
 - Socket options

SourcePort		DestPort	
SEQNber			
ACKNber			
Len	FLG	AdvWin	
Checksum		UrgPtr	
Options		Padding	

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Network traffic load balancing &

The mechanisms 8/8

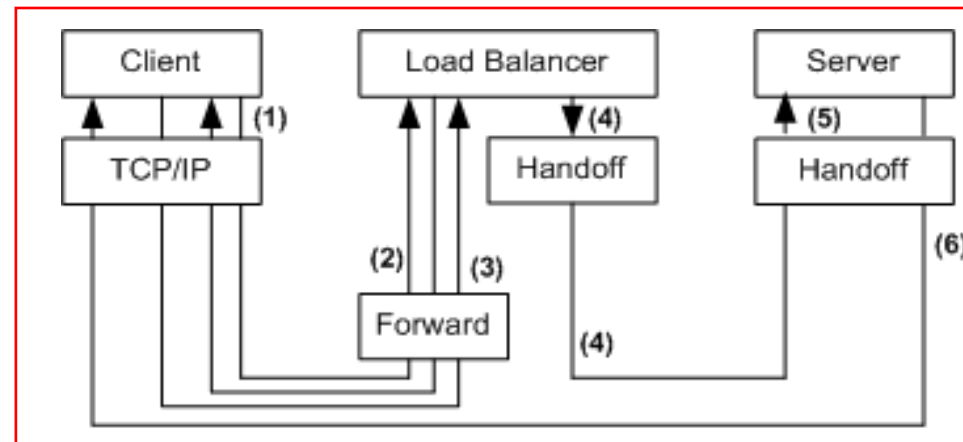
➔ TCP handoff

- Migrates the TCP connection from the front end to the back-end node using a Handoff protocol

Magic Nber	ConnMagic
Conn_Info	

Magic Nber	ConnMagic
Ack Msg	

- *The connection is created at the node without going through the three Way handshake procedure.*



Why session aware network traffic load & balancing?

➔ The session integrity issue

- Recalling that user level sessions can span over multiple flows
 - How to apply the same processing to flows pertaining to the same session?
- Basic transport level load balancing fails to achieve session integrity for multiple flow based services
 - Client persistency can 'force' session integrity but leads to unbalanced resource allocation among the cluster nodes

How do DPI helps in achieving session awareness



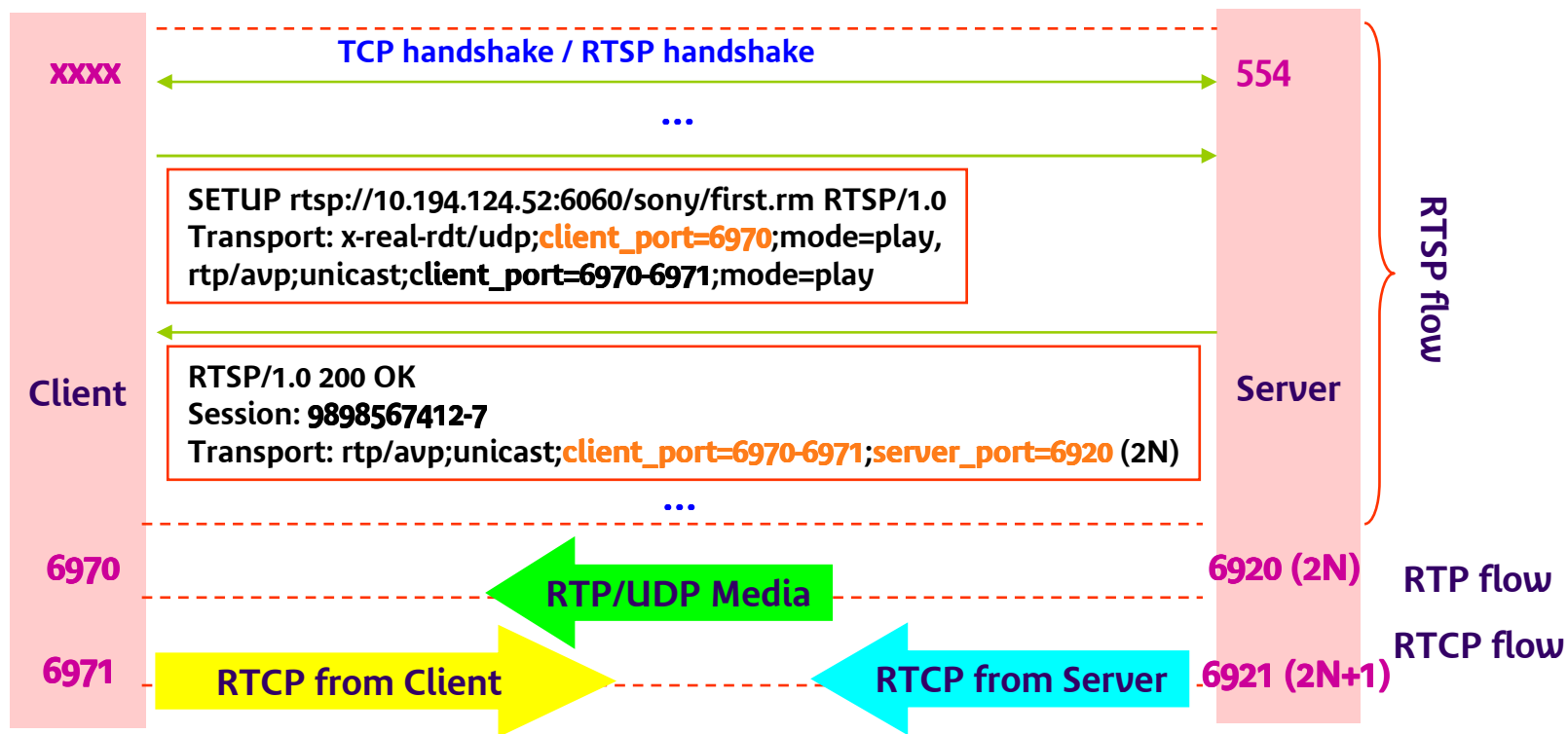
A unicast streaming media use case

➡ Can be built upon an RTSP/RTP/RTCP association

- RTSP is used to establish the signaling flow
 - Control commands (DESCRIBE, SETUP, PLAY, TEARDOWN, etc.) are transported over the RTSP channel
 - The media flow identifiers are negotiated on the control channel
- RTP and RTCP are used for the data and control data exchange

DPI for session awareness

A unicast streaming media use case



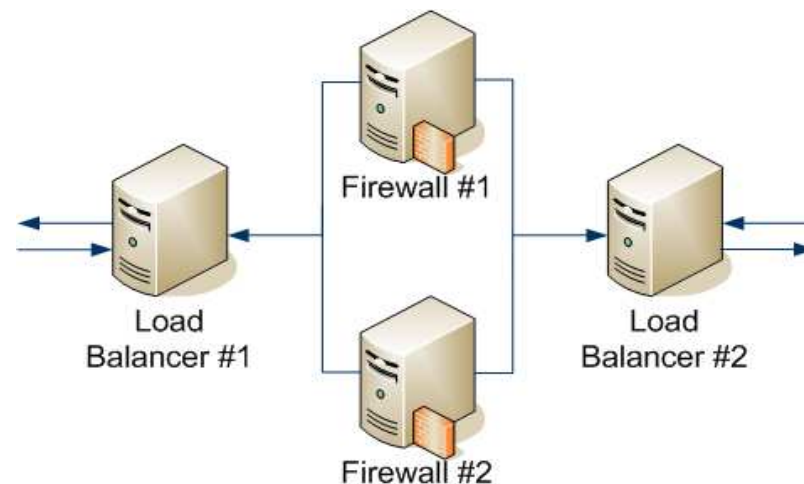
DPI for NAT & firewall traversal &

➡ NAT traversal issues

- How to scale servers when clients or servers are behind NAT devices?

➡ Firewall traversal issues

- How to scale firewalls while avoiding the interruption of the legitimate sessions?
 - Involved flows use non standard ports



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The requirements of NGN cluster & based load balancing frameworks

⇒ Performance

- Improved throughput in terms of number of **sessions completed** per unit of time
- Improved scaling factor inside the cluster
 - The system throughput grows as a function of the number of the processing servers inside the cluster

⇒ High availability

- The continuous processing of
 - the new incoming sessions
 - the already established sessions

in case of failure of

- the legitimate **entry point** to the cluster
- the legitimate **processing server** inside the cluster

Key features for NGN C-LB Frameworks

Adaptive and **session aware** processing of the offered network traffic to the cluster of servers

- Awareness of the usage and availability of the cluster resources
- Awareness of the session level constraints for **any** processing of the offered network traffic
 - Load balancing
 - Admission control
 - High availability
 - etc.

Conclusion and Perspectives &

- ➔ We questioned the appropriateness of flow aware load balancing in a cluster of servers
- ➔ We stated the requirements of NGN cluster based load balancing architectures in terms of
 - Performance
 - High availability
- ➔ We proposed an NGN architecture of IP traffic offered to a cluster of servers
 - Adaptive and session aware processing (European PATENT)
 - Provides means to avoid the interruption of the already flows when the legitimate server goes down (European PATENT)
- ➔ Our short term objectives focus on
 - Evaluating the session aware processing provided within this framework
 - Admission control,
 - Load balancing,
 - High availability.



Thanks
Any Questions?



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