# EXECUTION ANALYSIS OF DSM APPLICATIONS: A DISTRIBUTED AND SCALABLE APPROACH

Lionel Brunie, Laurent Lefèvre et Olivier Reymann Laboratoire de l'Informatique du Parallélisme Ecole Normale Supérieure de Lyon 69364 LYON Cedex 07, France (lbrunie, llefevre, oreymann)@lip.ens-lyon.fr

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# Plan

- What is a DSM system ?
- The DOSMOS system
- Why monitoring a DSM application ?
- A model for DSM application monitoring
  - Trace detection and collection : **Event Manager Process**
  - Trace management : *Meta-Object concept*
  - Analysis and visualization of execution traces:  ${\it Visualization\ process}$
- Implementation and architecture
  - Meta-Objects
  - System architecture
- Protocols : write, read, acquire, release operation protocols
- Experiments
- Conclusion and futur works

#### What is a DSM system ?

- Designed for distributed memory architectures
- Allows transparent accesses to shared data
  - Virtual shared memory systems
  - Object based distributed shared memory systems
- Hides the inter-process communication



Figure 1: A taxonomy of DSM systems

## The designing goals of DOSMOS

- Providing an easy-to-use programming environment
- Breaking the centralization of the memory accesses
- Scalability
- Portability

#### Basics of DOSMOS

- Dedicated processes (AP, MP...)
- Array allocation
- $\bullet$  Hierarchical structuring of application processes
- Optimized weak consistency protocols
- Mixing PVM and DOSMOS code
- From distributed systems to parallel machines

#### About Array splitting up

shared float A(14,20)[3,5];



Figure 2: Example of a split matrix

- Increase of the number of objects
- Improve the object distribution over the processors
- $\bullet$  reduce the false sharing and the access bottleneck

### Clustering using hierarchic groups

Processes that share a common set of variables are gathered.

- $\implies$  Limit the number of variable copies in the system
- $\implies$  Reduce the cost of the consistency maintenance



### DOSMOS architecture



Figure 4: DOSMOS system: an example of software configuration with two groups and three objects A, B and C

## Why monitoring a DSM application ?

- Informations :
  - DSM system administration
  - Shared data
- Detections :
  - Bottlenecks
  - Ping-Pong effects
  - No-sharing
  - Specific features

A model for DSM application monitoring

Trace detection and collection : **Event Manager Process** 

This process takes charge of the collection of the information concerning traces and its management.

- Do not overload the DSM system with the monitoring task
- Generate distributed trace files
- Easy to implement
- Definition of a communication protocol between the system and the monitoring tool

#### Trace management

- To store the events in a file on disk
  - For a post-mortem use
- $\bullet$  Use a dedicated structure : the  ${\bf Meta-Object\ structure}$

This is a data structure that contains trace information about the operations performed on a variable.

- Such a structure is linked to each variable
- Managed by an Event Manager Process
- Data stored in a judicious way
- Adapted for on-line analysis and optimization

#### DOSMOS-Trace architecture



Figure 5: DOSMOS-Trace: example of monitoring environment



Figure 6: Protocol implemented to collect the trace information about a write operation.



Figure 7: Protocol implemented to collect the trace information about a read operation.



Figure 8: Protocol implemented to collect the trace information about an Acquire operation.

## Release operations



Figure 9: Protocol implemented to collect the trace information about a Release operation.

Analysis and visualization of execution traces

The  $Visualization \ process$  provides two kinds of diagrams

- Diagrams showing variable accesses according to time
- Diagrams showing statistical results on variable accesses



Figure 10: DOSMOS-Trace: the Visualization Process (VP)  $\implies$  E.M.P.



Diagrams showing statistical results on variable accesses

- Number of readings during a time interval
- Number of writings during a time interval



Figure 11: Number and origin of the read accesses performed on an object vs execution time (in black: inter-group accesses)



Figure 12: Number and origin of the read accesses performed on an object vs execution time (note this execution does not include inter-group accesses)

# Histories

Diagrams showing variable accesses according to time

- History of accesses performed by the processes
- History of accesses performed on the variables



Figure 13: Object activity vs execution time



Figure 14: Process activity vs execution time

#### DOSMOS-Trace environmement



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Configuration	1	2	3	4		
Execution time	21.90	40.00 (+83%)	$48.20 \ (+120\%)$	$28.30 \ (+29\%)$		

Table 1: Execution time (in seconds) for several configurations

Ratio \ Configuration	1	2		2		3		4	
R=1	43.80	61.90	(41%)	70.10	(60%)	50.20	(15%)		
R=2	65.70	83.80	(28%)	92.00	(40%)	72.10	(10%)		
R=3	87.60	105.70	(21%)	113.90	(30%)	94.00	(7%)		

Table 2: Calculated execution time (in seconds) for different (computation/shared data access) ratios

# Conclusion

- Weak intrusion
- Flexibility
- System scalability
- Meta-object : good storage structure for on-line analysis
- User-orientation
- $\bullet$  Independence
- $\implies$  A distributed debugger for DSM systems

Adress: http://www.ens-lyon.fr/ llefevre/DOSMOS/Dosmos.html