

# Non-closure of the set of quantum correlations

•  $C_q(n, k)$  "quantum"

PVM's  $(P_{v,i})_{i=1}^k$  and  $(Q_{w,j})_{j=1}^k$  on fin. dim. Hilbert spaces  $H_A$  and  $H_B$

A unit vector  $\xi \in H_A \otimes H_B$

$$p(i,j|v,w) = \langle (P_{v,i} \otimes Q_{w,j}) \xi, \xi \rangle$$

•  $C_{qa}(n, k) = \overline{C_q(n, k)}$  "quantum approximate"

•  $C_{qc}(n, k)$  "quantum commuting"

PVM's  $(P_{v,i})_{i=1}^k$  and  $(Q_{w,j})_{j=1}^k$  on the same (possibly infinite dimensional) Hilbert space  $H$

satisfying  $P_{v,i} Q_{w,j} = Q_{w,j} P_{v,i} \quad \forall v, w, i, j$

A unit vector  $\xi \in H$

$$p(i,j|v,w) = \langle P_{v,i} Q_{w,j} \xi, \xi \rangle$$

Terminology: A tuple  $((P_{v,i})_i)_v, ((Q_{w,j})_j)_w, H, \xi$  is called a realization in  $C_{qc}$  of  $(p(i,j|v,w))$  if the above is satisfied.

Exercise:  $p \in C_q(n, k)$  iff it has a realization in  $C_{qc}(n, k)$  with  $H$  finite dimensional.

Rk:  $C_{qa}(n, k) = C_{qc}(n, k) \quad \forall n, k$  iff CEP is true      Rk:  $C_{qc}$  is closed

Goal of talk:  $C_q$  is not always closed.

•  $C_{ns}(n, k)$  "non-signaling"

Any  $(p(i,j|v,w)) \in [0,1]^{n^2 k^2}$  satisfying

$$\sum_{i=1}^k p(i,j|v,w) = 1, \quad \forall v, w$$

$$\sum_{j=1}^k p(i,j|v,w) = \sum_{j=1}^k p(i,j|v,w') \quad \forall i, v, w, w'$$

$$\sum_{i=1}^k p(i,j|v,w) = \sum_{i=1}^k p(i,j|v',w) \quad \forall j, v, v', w$$

Notation: Marginal densities

$$P_A(i|v) = \sum_{j=1}^k p(i,j|v,w)$$

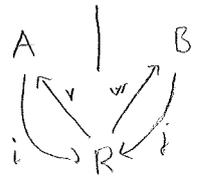
$$P_B(j|w) = \sum_{i=1}^k p(i,j|v,w)$$

answers ← questions

$$(p(i,j|v,w)) \in$$

$$C_r(n, k) \subseteq [0,1]^{n^2 k^2}$$

↑ #questions    #answers



Correlation sets and  $C^*$ -algebras

Thm:  $p \in C_{qc}(n, k)$  (resp.  $C_q(n, k)$ ) iff  $\exists$  a unital  $C^*$ -algebra  $\mathcal{A}$  (resp. a fin. dim. such) generated by families of projections  $(e_{v,i}), (f_{w,j}) \subset \mathcal{A}$  satisfying  $\sum_{i=1}^k e_{v,i} = 1_{\mathcal{A}} = \sum_{j=1}^k f_{w,j}$  and  $e_{v,i} f_{w,j} = f_{w,j} e_{v,i}$ , and a state  $\rho$  on  $\mathcal{A}$  such that

$$p(i, j | v, w) = \rho(e_{v,i} f_{w,j})$$

GNS construction

$(\mathcal{A}, \rho)$   $C^*$ -algebra w. a state

$\rightarrow (\pi, H, \xi)$

•  $\pi: \mathcal{A} \rightarrow \mathcal{B}(H)$   $*$ -hom.

•  $\xi$  is a cyclic unit vector

( $\pi(\mathcal{A})\xi$  is dense  $H$ )

•  $\rho(x) = \langle \pi(x)\xi, \xi \rangle$

Rk Marginal distributions:

•  $\rho_A(i | v) = \rho(e_{v,i})$

•  $\rho_B(j | w) = \rho(f_{w,j})$

Synchronous games

Def  $(p(i, j | v, w))$  is called synchronous if, for all  $v$ ,  $p(i, j | v, v) = 0$   $i \neq j$ .

Notation:  $C_r^S(n, k)$

Rk: Synchronous subsets are convex and we have the same line of inclusions.

Prop: If  $(p(i, j | v, w)) \in C_{qc}^S(n, k)$  is realized by  $((P_{v,i})_i)_v, ((Q_{w,j})_j)_w, H, \xi$  then

(i)  $P_{v,i}\xi = Q_{v,i}\xi \quad \forall v, i$  (OBS! does not force  $P_{v,i}$  and  $Q_{v,i}$  to be equal)

(ii)  $p(i, j | v, w) = \langle P_{v,i} P_{w,j} \xi, \xi \rangle = \langle Q_{w,j} Q_{v,i} \xi, \xi \rangle = p(j, i | w, v)$

Pf: Exercise (Hint for (i): C.S.)

Thm:  $p \in C_{qc}^S(n, k)$  (resp.  $C_q^S(n, k)$ ) iff  $\exists$  unital  $C^*$ -algebra  $\mathcal{A}$  (resp. a fin. dim. such) generated by a family of projections  $(e_{v,i}) \subset \mathcal{A}$  satisfying  $\sum_{i=1}^k e_{v,i} = 1_{\mathcal{A}}$ , and with a tracial state  $\tau$  such that

$$p(i, j | v, w) = \tau(e_{v,i} e_{w,j})$$

Pf: - If  $p \in C_{qc}^S(n, k)$  has realisation  $((P_{v,i})_i)_v, ((Q_{w,j})_j)_w, H, \xi$ ,

Let  $\mathcal{A}$  be the  $C^*$ -algebra generated by the  $P_{v,i}$ 's ( $Q_{w,j}$  commute with  $\mathcal{A}$ )

Define  $\tau = \langle \cdot, \xi, \xi \rangle$

Then  $p(i, j | v, w) = \tau(P_{v,i} P_{w,j})$  by prop.

$\tau$  is tracial: For  $x \in \mathcal{A}$

$$\tau(P_{v,i} x) = \langle P_{v,i} x \xi, \xi \rangle = \langle Q_{v,i} x \xi, \xi \rangle = \langle x P_{v,i} \xi, \xi \rangle = \tau(x P_{v,i}) \quad \text{+ induction + continuity}$$

- Conversely, suppose  $(\mathcal{A}, \tau)$  and  $(e_{v,i})$  are as in the theorem. For  $i \neq j, e_{v,i} e_{v,j} = 0$ , so the correlation is synchronous

Let  $(\pi, H, \xi)$  be the GNS construction and set  $P_{v,i} = \pi(e_{v,i})$

Then  $(P_{v,i})_i$  is a PVM,  $\forall v$ , and  $p(i, j | v, w) = \tau(e_{v,i} e_{w,j}) = \langle P_{v,i} P_{w,j} \xi, \xi \rangle$

For  $x$  finite linear combination of words on  $(P_{v,i})$ , set

$$Q_{w,j} x \xi = x P_{w,j} \xi$$

•  $Q_{w,j}$  is a well-defined contraction  $\rightarrow$  define  $Q_{w,j}$  on all of  $H$

•  $Q_{w,i}^2 = Q_{w,i}^* = Q_{w,i}$  and  $\sum_{j=1}^k Q_{w,j} = 1_H$

•  $Q_{w,j} P_{v,i} (x \xi) = P_{v,i} (x Q_{w,j} \xi) = P_{v,i} Q_{w,j} x \xi \rightarrow P_{v,i}$  and  $Q_{w,j}$  commute

## Synchronous binary answer games

Rk  $(p(i, j | v, w)) \in C_{ns}(n, 2)$

$$\bullet p(1, 2 | v, w) = p_A(1 | v) - p(1, 1 | v, w)$$

$$\bullet p(2, 1 | v, w) = p_B(1 | w) - p(1, 1 | v, w)$$

$$\bullet p(2, 2 | v, w) = 1 - p_A(1 | v) - p_B(1 | w) + p(1, 1 | v, w)$$

(knowing the marginals and each probability of both answering 1 gives everything)

Rk  $(p(i, j | v, w)) \in C_{ns}^s(n, 2)$

$$\text{That is, } p(1, 2 | v, v) = p(2, 1 | v, v) = 0$$

$$\leadsto p(1, 1 | v, v) = p_A(1 | v) = p_B(1 | v)$$

Def (A slice of  $C_r^s(n, 2)$ ) For each  $t \in [0, 1]$

$$\Gamma_r(t) = \{p \in C_r^s(n, 2) \mid p_A(1 | v) = p_B(1 | w) = t \forall v, w\}$$

Rk  $\Gamma_r(t)$  is non-empty and convex.

Def (The correlation function)  $f_r: [0, 1] \rightarrow \mathbb{R}$

$$f_r(t) = \inf \{F(p) \mid p \in \Gamma_r(t)\}$$

where  $F: C_{ns}(n, 2) \rightarrow \mathbb{R}$  is the function

$$F(p) = \sum_{v \neq w} p(1, 1 | v, w)$$

Rk If  $C_r^s(n, 2)$  is closed, the infimum is attained, for all  $t \in [0, 1]$

Rk  $f_q(t) = f_{qa}(t) \geq f_{qc}(t) \geq f_{ns}(t) \geq 0$

Prop  $f_r$  is a convex function

pf: By convexity of  $C_r^s(n, 2)$ ,

$$\lambda \Gamma_r(t_1) + (1-\lambda) \Gamma_r(t_2) \subset \Gamma_r(\lambda t_1 + (1-\lambda)t_2), \quad t_1, t_2, \lambda \in [0, 1]$$

we have

$$\begin{aligned} \lambda f_r(t_1) + (1-\lambda) f_r(t_2) &= \lambda \inf \{F(p) \mid p \in \Gamma_r(t_1)\} + (1-\lambda) \inf \{F(p) \mid p \in \Gamma_r(t_2)\} \\ &= \inf \{F(p) \mid p \in \lambda \Gamma_r(t_1) + (1-\lambda) \Gamma_r(t_2)\} \\ &\geq \inf \{F(p) \mid p \in \Gamma_r(\lambda t_1 + (1-\lambda)t_2)\} \\ &= f_r(\lambda t_1 + (1-\lambda)t_2) \end{aligned}$$

□

# The quantum correlation function

Lemma: Let  $t \in [0, 1]$  be irrational and suppose  $f_q(t)$  is attained in the infimum defining it.

Then there exist rational numbers  $r, s \in [0, 1]$  s.t.  $t \in [r, s]$  and the restriction of  $f_q$  to  $[r, s]$  is linear.

Pf: Take  $p \in I_q^+(t)$  s.t.  $f_q(t) = F(p)$ .

There is a unital fin. dim.  $C^*$ -algebra  $\mathcal{A}$  generated by projections  $\{P_v\}_{v=1}^n$  and  $w$ , a tracial state  $\tau: \mathcal{A} \rightarrow \mathbb{C}$  w.  $\tau(P_v) = t \ \forall v$  and such that

$$f_q(t) = \sum_{v \neq w} \tau(P_v P_w)$$

Write  $\mathcal{A} = \bigoplus_{L=1}^m M_{d_L}(\mathbb{C})$  and  $\tau = \bigoplus_{L=1}^m \lambda_L \text{tr}_{d_L}$  where  $\lambda_L > 0$  with  $\sum_{L=1}^m \lambda_L = 1$  (normalized trace)

Moreover, write  $P_v = \bigoplus_{L=1}^m P_{v,L}$  with  $P_{v,L} \in M_{d_L}(\mathbb{C})$  a projection

For each  $v$  and each  $1 \leq L \leq m$ , set

$$\tilde{P}_{v,L} = \bigoplus_{\sigma \in S_n} P_{\sigma(v),L} \in \bigoplus_{\sigma \in S_n} M_{d_L}(\mathbb{C}) =: \mathcal{A}_L$$

Define tracial state  $\tau_L: \mathcal{A}_L \rightarrow \mathbb{C}$  by  $\tau_L = \frac{1}{n!} \bigoplus_{\sigma \in S_n} \text{tr}_{d_L}$

Observe that, for any  $v$  and  $w$

$$\begin{aligned} \tau_L(\tilde{P}_{v,L}) &= \frac{1}{n!} \sum_{\sigma \in S_n} \text{tr}_{d_L}(P_{\sigma(v),L}) = \frac{1}{n!} \sum_{\sigma \in S_n} \text{tr}_{d_L}(P_{\sigma \circ \sigma^{-1}(v),L}) \\ &= \tau_L(\tilde{P}_{w,L}) =: r_L \end{aligned}$$

↑ this is rational

Summarizing: we have a unital  $C^*$ -algebra  $\mathcal{A}_L$  w. a tracial state  $\tau_L$  and all marginals equal to  $r_L$

Hence,  $\{\tilde{P}_{v,L}\}_{v=1}^n$  defines a synchronous quantum correlation in the slice  $I_q^+(r_L)$ , and thus

$$f_q(r_L) \leq \sum_{v \neq w} \tau_L(\tilde{P}_{v,L} \tilde{P}_{w,L})$$

Set  $\tilde{\mathcal{A}} = \bigoplus_{L=1}^m \mathcal{A}_L$ , define a tracial state  $\tilde{\tau}: \tilde{\mathcal{A}} \rightarrow \mathbb{C}$  by  $\tilde{\tau} = \bigoplus_{L=1}^m \lambda_L \tau_L$

Define projections  $\tilde{P}_v = \bigoplus_{L=1}^m \tilde{P}_{v,L}$ . Then

$$\tilde{\tau}(\tilde{P}_v) = \sum_{L=1}^m \lambda_L \tau_L(\tilde{P}_{v,L}) = \frac{1}{n!} \sum_{\sigma \in S_n} \sum_{L=1}^m \lambda_L \text{tr}_{d_L}(P_{\sigma(v),L}) = \frac{1}{n!} \sum_{\sigma \in S_n} \tau(P_{\sigma(v)}) = t$$

$$t = \sum_{L=1}^m \lambda_L r_L$$

$$\tilde{\tau}(\tilde{P}_v \tilde{P}_w) = \sum_{L=1}^m \lambda_L \tau_L(\tilde{P}_{v,L} \tilde{P}_{w,L}) = \frac{1}{n!} \sum_{\sigma \in S_n} \tau(P_{\sigma(v)} P_{\sigma(w)})$$

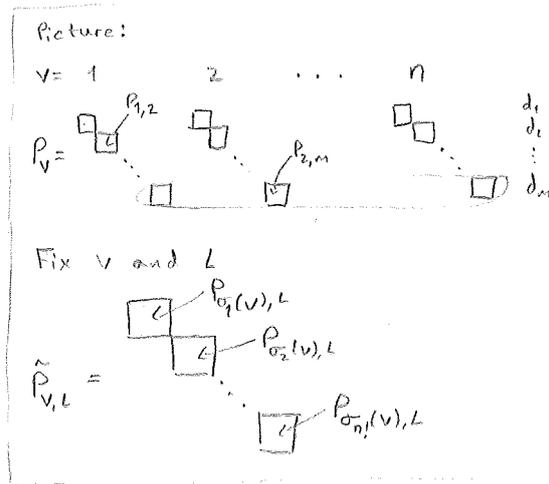
$$\rightarrow f_q(t) = \sum_{v \neq w} \tau(P_v P_w) = \sum_{v \neq w} \tau(P_{\sigma(v)} P_{\sigma(w)}) = \sum_{v \neq w} \tilde{\tau}(\tilde{P}_v \tilde{P}_w)$$

Thus,

$$f_q(t) = \sum_{L=1}^m \lambda_L \sum_{v \neq w} \tau_L(\tilde{P}_{v,L} \tilde{P}_{w,L}) \geq \sum_{L=1}^m \lambda_L f_q(r_L)$$

But  $f_q$  is convex, so this is an equality (by Jensen)

Since  $t$  is irrational and all the  $r_L$ 's are rational, it is not the case that all  $r_L$ 's are equal. So the equality in Jensen's inequality implies that  $f_q$  is linear on an interval containing the  $r_L$ 's.



# Talk: Non-closure of quantum correlations

Recap from Emilie's part:

Correlation sets:

$$C_q^S(n, k) \subseteq C_{qs}^S(n, k) \subseteq C_{qa}^S(n, k) \subseteq C_{qc}^S(n, k) \subseteq C_{qis}^S(n, k)$$

$\uparrow$   $\uparrow$   
closed?  $\frac{C_q^S(n, k)}{C_{qa}^S(n, k)}$   $\uparrow$  closed!

Define

$$\Gamma_r(\epsilon) = \left\{ (p(i, j | v, w)) \in C_r^S(n, 2) : p_A(a|v) = p_B(a|w) = \epsilon \right\}_{\forall v, w}$$

$$F(p) = \sum_{v, w} p(\begin{smallmatrix} 1 & 1 \\ 0 & 0 \end{smallmatrix} | v, w)$$

and

$$f_r(\epsilon) = \inf \left\{ F(p) \mid p \in \Gamma_r(\epsilon) \right\} \quad (*)$$

Proposition: Let  $t \in [0, 1]$  be irrational.

If  $f_q(t)$  is attained in  $(*)$ , then  $\exists r, s \in \mathbb{Q}$   $r < t < s$

s.t.

$f_q|_{[r, s]}$  linear.

Goal: Find  $t \in [0, 1]$  irrational s.t. ~~there is no interval~~  
there is no interval  $[r, t] \ni t$  where  
 $f_q|_{[r, t]}$  is linear!

Def: (Vectorial correlations)

$(p(i,j | v,w))_{\substack{i,j \in \{1, \dots, n\} \\ v,w \in \{1, \dots, n\}}}$  is called a vectorial correlation

if  $\exists$  Hilbertspace  $H$  and vectors  $\begin{Bmatrix} x_{v,i} \\ y_{w,i} \\ h \end{Bmatrix} \in H$   
s.t.

- (1)  $\|h\|=1$
- (2)  $\langle x_{v,i}, x_{v,j} \rangle = 0 \quad \forall v \quad \forall i \neq j$
- (3)  $\langle y_{w,i}, y_{w,j} \rangle = 0 \quad \forall w \quad \forall i \neq j$
- (4)  $h = \sum_i x_{v,i} = \sum_j y_{w,j} \quad \forall v, w$
- (5)  $p(i,j | v,w) = \langle x_{v,i}, y_{w,j} \rangle \quad \forall v, w, i, j$

Write  $C_{vec}(n,k) = \left\{ \begin{array}{l} \text{set of all vectorial corr.} \\ \text{for } n, k \end{array} \right\}$ .

Exercise: Show  $C_{qc}(n,k) \subseteq C_{vec}(n,k) \subseteq C_{ns}(n,k)$

Exercise: Show that  $p \in C_{vec}^s(n,k)$  iff  $\exists$  H. space  $H$   
&  $\{x_{v,i}\} \subset H$   
 $h \in H$  s.t.

- (1)  $\|h\|=1$
- (2)  $\langle x_{v,i}, x_{v,j} \rangle = 0 \quad \forall v \quad \forall i \neq j$
- (3)  $h = \sum_i x_{v,i} \quad \forall v$
- (4)  $p(i,j | v,w) = \langle x_{v,i}, x_{w,j} \rangle \quad \forall v, w, i, j$

Computing  $f_{\text{rec}}(t)$ :

$$f_{\text{rec}}(t) = \inf \{ F(p) : p \in \Gamma_{\text{rec}}(t) \}$$

$$F(p) = \sum_{v \neq w} P(p, p | v, w)$$

$$\& \Gamma_{\text{rec}}(t) = \left\{ p \in C_{\text{rec}}^S(n, 2) : \begin{aligned} P_A(p | v) &= P_B(p | w) \\ &= t \end{aligned} \right\}$$

Lemma (Symmetrization):

$$f_{\text{rec}}(t) = \inf \{ F(p) : p \in \tilde{\Gamma}_{\text{rec}}(t) \}$$

where

$$\tilde{\Gamma}_{\text{rec}}(t) = \left\{ p \in \Gamma_{\text{rec}}(t) : \begin{aligned} P(p, p | v, w) &= P(p, p | x, y) \\ \forall v \neq w \quad \forall x \neq y \end{aligned} \right\}$$

Proof: For  $\pi \in S_n$  define  $\beta_\pi : C_{\text{rec}}^S \rightarrow C_{\text{rec}}^S$

$$\beta_\pi((P(i, j) | v, w)) = P(i, j | \pi^{-1}(v), \pi^{-1}(w))$$

Note:  $F = F \circ \beta_\pi$

$$\text{Define } \begin{cases} \text{Avg} : C_{\text{rec}}^S \rightarrow C_{\text{rec}}^S \\ \text{Avg}(p) = \frac{1}{n!} \sum_{\pi \in S_n} \beta_\pi(p) \end{cases}$$

Again:  $F = F \circ \text{Avg}$

Since  $\text{Avg}(\Gamma_{\text{rec}}(t)) = \tilde{\Gamma}_{\text{rec}}(t)$  we find

$$\begin{aligned} f_{\text{rec}}(t) &= \inf \{ F(p) : p \in \Gamma_{\text{rec}}(t) \} \\ &= \inf \{ F \circ \text{Avg}(p) : p \in \Gamma_{\text{rec}}(t) \} \\ &= \inf \{ F(p) : p \in \tilde{\Gamma}_{\text{rec}}(t) \}. \end{aligned}$$

□

How does  $p \in \tilde{\Pi}_{rec}(t)$  look like?

$$1.) P(\mathbf{1}, \mathbf{1} | v, w) = P(\mathbf{1}, \mathbf{1} | x, y) =: \frac{s}{n(n-1)} \\ \forall v \neq w, \forall x \neq y$$

$$\hookrightarrow \underline{\underline{F(p) = s}}$$

$$2.) t = P_A(\mathbf{1} | v) = P(\mathbf{1}, \mathbf{1} | v, v) + P(\mathbf{1}, \mathbf{2} | v, v) \stackrel{= 0}{=} \text{by symm.}$$

$$\Rightarrow P(\mathbf{1}, \mathbf{1} | v, w) = \begin{cases} t & v = w \\ \frac{s}{n(n-1)} & v \neq w. \end{cases}$$

$$3.) t = P_A(\mathbf{1} | v) \stackrel{v \neq w}{=} P(\mathbf{1}, \mathbf{1} | v, w) + P(\mathbf{1}, \mathbf{2} | v, w) \\ \parallel \\ \frac{s}{n(n-1)}$$

$$\Rightarrow P(\mathbf{1}, \mathbf{2} | v, w) = P(\mathbf{2}, \mathbf{1} | v, w) = \begin{cases} 0 & v = w \\ t - \frac{s}{n(n-1)} & v \neq w. \end{cases}$$

$$4.) P(\mathbf{2}, \mathbf{2} | v, w) = 1 - \sum_{ij} P(i, j | v, w)$$

$$\Rightarrow P(\mathbf{2}, \mathbf{2} | v, w) = \begin{cases} 1 - t & v = w \\ 1 - 2t + \frac{s}{n(n-1)} & v \neq w \end{cases}$$

Theorem:

$$f_{\text{rec}}(t) = \min_{S \geq 0} \left\{ \begin{array}{l} \text{s.t.} \quad \frac{S}{n(n-1)} \leq t \\ (2t-1) \leq \frac{S}{n(n-1)} \\ \& \quad \begin{pmatrix} 1 & t & & t \\ t & t & & \frac{S}{n(n-1)} \\ & & \ddots & \\ t & \frac{S}{n(n-1)} & & t \end{pmatrix} \geq 0 \end{array} \right\} \quad (*)$$

Proof: Consider  $S \geq 0$  satisfying (\*).

Since every positive matrix is a Gram matrix

$$\exists \{h, x_{v,1}, x_{v,2}, \dots, x_{v,n}\} \subseteq \mathbb{C}^{n+1}$$

$$\text{s.t.} \quad \left\{ \begin{array}{l} \langle h, h \rangle = 1 \\ \langle h, x_{v,i} \rangle = t \quad \forall i \in \{2, \dots, n\} \\ \langle x_{v,i}, x_{v,i} \rangle = \frac{S}{n(n-1)} \quad \forall i \neq n \quad \& \quad \langle x_{v,i}, x_{v,n} \rangle = t \end{array} \right.$$

Set  $x_{v,2} = h - x_{v,1}$ . Then:

$$\langle x_{v,1}, x_{v,2} \rangle = \langle x_{v,1}, h \rangle - \langle x_{v,1}, x_{v,1} \rangle = t - t = 0.$$

$$\& \quad h = x_{v,1} + x_{v,2}$$

Set

$$P(i,j | v,w) = \langle x_{vi} | x_{wj} \rangle$$

$$\Rightarrow p \in \text{vec}(t) \Rightarrow f_{\text{vec}}(t) \leq \min(s \geq 0: \text{(*)})$$
$$F(p) = s$$

For converse assume that  $p \in \text{vec}(t)$  satisfies  $F(p) = s$ .

$$\text{Lemma} \Rightarrow \exists \tilde{p} \in \tilde{\text{vec}}(t) \text{ s.t. } F(\tilde{p}) = s.$$

Previous discussion

$$\Rightarrow \frac{s}{n(n-1)} \leq t$$

$$\& (2t-1) \leq \frac{s}{n(n-1)}$$

Build Gram matrix from realization of  $\tilde{p}$

$$\Rightarrow \begin{pmatrix} 1 & t & \dots & t \\ t & t & \dots & \frac{s}{n(n-1)} \\ \vdots & \vdots & \ddots & \vdots \\ t & \frac{s}{n(n-1)} & \dots & t \end{pmatrix} \geq 0.$$

$$\leadsto \min(s \geq 0 \text{ s.t. (*)}) \leq f_{\text{vec}}(t).$$

□

Corollary:

For  $n=5$  we have

$$f_{\text{rec}}(t) = \begin{cases} 0 & 0 \leq t \leq \frac{1}{5} \\ 5t(5t-1) & \frac{1}{5} \leq t \leq \frac{4}{5} \\ 20(2t-1) & \frac{4}{5} \leq t \leq 1 \end{cases} \quad \leftarrow \text{quadratic.}$$

Thm:  $C_q^S(5,2)$  is not closed.

Proof: Consider

$$I_5 := \left[ \frac{\sqrt{5}-1}{2\sqrt{5}}, \frac{\sqrt{5}+1}{2\sqrt{5}} \right] \subset \left[ \frac{1}{5}, \frac{4}{5} \right]$$

Claim:  $f_{\text{rec}}(t) = f_q(t) \quad \forall t \in \mathbb{Q} \cap I_5$

Since  $f_{\text{rec}}(t) = 5t(5t-1)$  ~~not~~ linear on  $I_5$

Claim  $\Rightarrow f_q(t) = \inf \{ F(p) : p \in \Gamma_q(t) \}$

not attained in any  
irrational  $t \in I_5$ .

(otherwise  $\exists r, s \in \mathbb{Q}$  s.t.  $t \in [r, s]$   
&  $f_q|_{[r, s]}$  linear)

$\Rightarrow C_q^S(5,2)$  not closed.

To prove the claim, note  $f_{\text{rec}}(t) \leq f_q(t)$ .  
 $\parallel$   
 $5t(5t-1) \quad \forall t.$

Show: 
$$\left. \begin{array}{l} f_{\text{vec}}(t) \cong f_q(t) \quad \forall t \in I_s \cap \mathbb{Q} \\ \parallel \\ st(st-1) \end{array} \right\} (*)$$

Magical Lemma (Kruglyak, Rabanovich, Samoilenko 2003)

If  $t \in I_s \cap \mathbb{Q}$ , then  $\exists k \in \mathbb{N}$  &  $P_1, \dots, P_s \in M_k$   
projections s.t.

$$P_1 + \dots + P_s = st \mathbb{1}_k$$

Proof of (\*):

For  $t \in I_s \cap \mathbb{Q}$  consider projections as in lemma.

Set 
$$\left\{ \begin{array}{l} \tilde{P}_v = P_v \oplus \dots \oplus P_{v+1} \in M_k \oplus \dots \oplus M_k \\ \subseteq M_{sk} \\ \text{for } v \in \{0, \dots, s-1\} \end{array} \right.$$

Note:  $\sum_{v=0}^{s-1} \tilde{P}_v = st \mathbb{1}_{sk}$

Let  $\text{tr}_{sk}$  denote normalized trace (state) on  $M_{sk}$ .

$$\begin{aligned} 1.) \text{tr}_{sk}(\tilde{P}_v) &= \frac{1}{sk} \text{Tr}(\tilde{P}_v) \\ &= \frac{1}{sk} \sum_{v=0}^{s-1} \text{Tr}(P_v) = \frac{t}{sk} \text{Tr}(\mathbb{1}_{sk}) = \underline{\underline{t}} \end{aligned}$$

$$2.) \sum_{v,w} \tilde{P}_v \tilde{P}_w = 25t^2 \mathbb{1}_{5k}$$

$$\parallel$$

$$\sum_v \tilde{P}_v + \sum_{v \neq w} \tilde{P}_v \tilde{P}_w$$

$$\implies \sum_{v \neq w} \tilde{P}_v \tilde{P}_w = (25t^2 - 5t) \mathbb{1}_{5k}$$

$$\implies \sum_{v \neq w} \text{tr}_{5k} (\tilde{P}_v \tilde{P}_w) = 5t(5t-1)$$

~~QED~~

Finally define

$$P(i,j | v,w) = \text{tr}_{5k} (Q_{v,i} Q_{w,j})$$

with

~~$$Q_{v,i} = \begin{cases} \tilde{P}_v & i=0 \\ \mathbb{1}_{5k} - \tilde{P}_v & i=1 \end{cases}$$~~

$$Q_{v,i} = \begin{cases} \tilde{P}_v & i=0 \\ (\mathbb{1}_{5k} - \tilde{P}_v) & i=1 \end{cases}$$

$$\implies f_7(t) \leq 5t(5t-1)$$

□

Corollary:

$C_q(5,2)$  and  $C_{qs}(5,2)$  are not closed.

$$C_{qs}(5,2) \neq C_{qa}(5,2) = \overline{C_q(5,2)}.$$

## A few words about the magic:

We will not prove the magical Lemma. But we can clarify a few constructions.

Suppose:  $\alpha \mathbb{1}_k = P_1 + \dots + P_n$  (\*) ( $\& P_i \neq 0$ )

$$\Rightarrow \alpha k = \sum_i \text{rk}(P_i) \in \mathbb{N}, \& \alpha \geq 1.$$

Two tricks lead to new decompositions:

(1) (\*)  $\Rightarrow \exists Q_1, \dots, Q_n$  projections s.t.

$$\bullet (n-\alpha) \mathbb{1}_k = Q_1 + \dots + Q_n$$

( $\hookrightarrow$  Choose  $Q_i = \mathbb{1} - P_i$ ).

(2) (\*)  $\Rightarrow \exists Q_1, \dots, Q_n$  projections s.t.  
&  $\alpha > 1$

$$\frac{\alpha}{\alpha-1} \mathbb{1}_{k(\alpha-1)} = Q_1 + \dots + Q_n$$

## Construction for (2):

Suppose (\*).

in particular  
 $P_i \neq 0$ .

Write  $P_i = V_i V_i^*$  for  $V_i: \mathbb{C}^{\text{rk}(P_i)} \rightarrow \mathbb{C}^k$   
isometry.

and define

$$V = (V_1, \dots, V_n) : \mathbb{C}^{\alpha k} \rightarrow \mathbb{C}^k.$$

Note that

$$VV^* = \sum_i V_i V_i^* = \alpha \mathbb{1}_n$$

$\Rightarrow \frac{1}{\sqrt{\alpha}} V^*$  is an isometry.

$\hookrightarrow$  Extend to unitary

$$U = \begin{matrix} \text{Ker} \\ \hline \left( \begin{array}{ccc} \frac{1}{\sqrt{\alpha}} V_1 & \dots & \frac{1}{\sqrt{\alpha}} V_n \\ \sqrt{\frac{\alpha-1}{\alpha}} W_1 & \dots & \sqrt{\frac{\alpha-1}{\alpha}} W_n \end{array} \right) \\ \text{Ker} \end{matrix}$$

$$U^* U = \begin{pmatrix} \frac{1}{\alpha} V_1^* V_1 + \frac{\alpha-1}{\alpha} W_1^* W_1 & & * \\ & \ddots & \\ * & & \ddots \end{pmatrix} = \begin{pmatrix} \mathbb{1}_{\text{rk}(P_1)} & & \\ & \ddots & \\ & & \ddots \end{pmatrix}$$

$$\begin{aligned} \Rightarrow \frac{1}{\alpha} V_i^* V_i + \frac{\alpha-1}{\alpha} W_i^* W_i &= \mathbb{1}_{\text{rk}(P_i)} \\ \parallel \\ \frac{1}{\alpha} \mathbb{1}_{\text{rk}(P_i)} \end{aligned}$$

$$\Rightarrow W_i^* W_i = \mathbb{1}_{\text{rk}(P_i)}$$

$\hookrightarrow W_i$  isometry.

$$U^* U = \begin{pmatrix} \frac{\alpha-1}{\alpha} \sum_i W_i W_i^* & & * \\ & \ddots & \\ * & & \ddots \end{pmatrix} = \begin{pmatrix} \mathbb{1}_{K(\alpha-1)} & & \\ & \ddots & \\ & & \ddots \end{pmatrix}$$

$$\Rightarrow Q_i = W_i W_i^* \text{ satisfy } \sum_i Q_i = \frac{\alpha}{\alpha-1} \mathbb{1}_{K(\alpha-1)}$$

projections

Lemma: TFAE

$$(1) \alpha \mathbb{1}_k = P_1 + \dots + P_n \quad \text{for proj. } P_i$$

$$(2) \frac{\alpha}{\alpha-1} \mathbb{1}_{k(\alpha-1)} = \tilde{P}_1 + \dots + \tilde{P}_n \quad \text{for proj. } \tilde{P}_i$$

$$(3) (n-\alpha) \mathbb{1}_k = P'_1 + \dots + P'_n \quad \text{for proj. } P'_i$$

$$(4) \frac{(n-1)\alpha-n}{\alpha-1} \mathbb{1}_{k(\alpha-1)} = Q_1 + \dots + Q_n \quad \text{for proj. } Q_i$$

Interval in theorem

$$I_5 = \left[ \frac{\sqrt{5}-1}{2\sqrt{5}}, \frac{\sqrt{5}+1}{2\sqrt{5}} \right]$$

$$\Rightarrow 5I_5 = \left[ \frac{5-\sqrt{5}}{2}, \frac{5+\sqrt{5}}{2} \right]$$

$$\parallel$$
$$2+\phi$$

golden section  
1.618...

Recall: Fibonacci numbers

$$\left\{ \begin{array}{l} f_0 = f_1 = 1 \\ f_{n+1} = f_n + f_{n-1} \end{array} \right.$$

$$\left\{ \begin{array}{l} \phi = \lim_{n \rightarrow \infty} \frac{f_{n+1}}{f_n} \end{array} \right.$$