

**Internship, Master 2**  
**Structure of classes of graphs defined by constraints on chords**  
**Proposed by Nicolas Trotignon, LIP, ENS de Lyon**

## Description

The goal of this internship is to get familiar with structural graph theory through the study of a particular class of graphs. Here graphs are non oriented and simple. A cycle in a graph is a *cluc* if it has length at least 5 and a unique chord (cluc stands for “Cycle, Long, with a Unique Chord”). We call  $\mathcal{C}$  the class of graphs that contains no cluc as an induced subgraph. We call  $\mathcal{B}$  the class of graphs such that no cycle has a unique chord. Note that  $\mathcal{B}$  is a subclass of  $\mathcal{C}$ . Note that if a graph is in  $\mathcal{C}$  and not in  $\mathcal{B}$ , then it must contain a cycle of length 4 with a unique chord. We denote by  $\chi(G)$  the chromatic number of  $G$  and by  $\omega(G)$  the maximum number of pairwise adjacent vertices in  $G$ .

The following two questions are open.

- Is there a polynomial time algorithm that decides whether an input graph is in  $\mathcal{G}$ ?
- Is there a polynomial  $f$  such that every graph  $G$  in  $\mathcal{C}$  satisfies  $\chi(G) \leq f(\omega(G))$ ?

The class under consideration is a generalisation of two known classes: the class  $\mathcal{B}$  defined above, and the class of chordal graphs (these received a lot of attention, a wikipedia page is devoted to them). The papers cited below all give hints toward the solutions of these two questions. First both questions are solved for  $\mathcal{B}$  in [5], and for chordal graphs (classical result). A function  $f$  such that  $\chi(G) \leq f(\omega(G))$  for all graphs in  $\mathcal{C}$  is known [4], but it is not a polynomial. The class  $\mathcal{B}$  and the class of chordal graphs share common structural properties [3, 1], and the so-called class of HHD-free graphs introduced in [2] might be important to understand the structure of the graphs that are in  $\mathcal{C} \setminus \mathcal{B}$ .

## Skills

Knowledge of graph theory is the main required skill. Basic knowledge of complexity theory and programming is appreciated.

## References

- [1] A. Brandstädt, F.F. Dragan, V.B. Le, and T. Szymczak. On stable cutsets in graphs. *Discrete Applied Mathematics*, 105(1-3):39–50, 2000.
- [2] C.T. Hoàng and N. Khouzam. On brittle graphs. *Journal of Graph Theory*, 12(3):391–404, 1988.
- [3] T. McKee. Independent separator graphs. *Utilitas Mathematica*, 73:217–224, 2007.
- [4] I. Penev. Amalgams and  $\chi$ -boundedness. Manuscript available at <http://perso.ens-lyon.fr/irena.penev/>, 2014.
- [5] N. Trotignon and K. Vušković. A structure theorem for graphs with no cycle with a unique chord and its consequences. *Journal of Graph Theory*, 63(1):31–67, 2010.